

CSE216 Programming Abstractions

Type Inference

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Type Inference

- Type inference
 - Determining the type of an expression
- HM Type Inference algorithm
 - Assign **type variables** to subexpressions
 - Build **type constraints** using the type variables
 - Find a **unifier** that solves the type constraints
- TODO
 - Download **reci_tiny_type_inter.zip** and implement **constr** and **unifier** functions

Type Inference

```
(*constr: type constraints for kvar = expr
   kvar: type variable for expr
   expr: expression (kvar = expr)
   env : variable name to kind variable map
   return: list of constraints (kexpr = kexpr)...
*)
let rec constr kvar kexpr env =
...
  match kexpr with
    (*TODO: make constraints for literals and variables*)
    | BOOL b -> ret [(kvar, KB)]
    | NUM n  ->
    | VAR v  ->
```

Type Inference

```
(*constr: type constraints for kvar = expr
   kvar: type variable for expr
   expr: expression (kvar = expr)
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   return: List of constraints (kexpr = kexpr)...
*)
let rec constr kvar kexpr env =
...
  match kexpr with
  (*TODO: make constraints for literals and variables*)
  | BOOL b -> ret [(kvar, KB)]
  | NUM n -> ret [(kvar, KN)]
  | VAR v -> ret [(kvar, lookup v env)]
```

Type Inference

(**TODO: make constraints for primitive operators**)

```
| ADD (a, b) | SUB (a, b) ->
    newvar ()      >>= fun kv1 ->
    newvar ()      >>= fun kv2 ->
    constr kv1 a env >>= fun c1 ->
    constr kv2 b env >>= fun c2 ->
    (kvar, KN)::(kv1, KN)::(kv2, KN)::c2@c1 |> ret

| EQ (a, b) | GE (a, b) ->
    newvar ()      >>= fun kv1 ->
    newvar ()      >>= fun kv2 ->

| AND (a, b) | OR (a, b) ->
    newvar ()      >>= fun kv1 ->
    newvar ()      >>= fun kv2 ->

| NOT a ->
    newvar ()      >>= fun kv1 ->
```

(**TODO*: make constraints for primitive operators*)

```
| ADD (a, b) | SUB (a, b) ->
    newvar ()          >>= fun kv1 ->
    newvar ()          >>= fun kv2 ->
    constr kv1 a env >>= fun c1 ->
    constr kv2 b env >>= fun c2 ->
    (kvar, KN)::(kv1, KN)::(kv2, KN)::c2@c1 |> ret

| EQ (a, b) | GE (a, b) ->
    newvar ()          >>= fun kv1 ->
    newvar ()          >>= fun kv2 ->
    constr kv1 a env >>= fun c1 ->
    constr kv2 b env >>= fun c2 ->
    (kvar, KB)::(kv1, KN)::(kv2, KN)::c2@c1 |> ret

| AND (a, b) | OR (a, b) ->
    newvar ()          >>= fun kv1 ->
    newvar ()          >>= fun kv2 ->
    constr kv1 a env >>= fun c1 ->
    constr kv2 b env >>= fun c2 ->
    (kvar, KB)::(kv1, KB)::(kv2, KB)::c2@c1 |> ret

| NOT a ->
    newvar ()          >>= fun kv1 ->
    constr kv1 a env >>= fun c1 ->
    (kvar, KB)::(kv1, KB)::c1 |> ret
```

Type Inference

(**TODO: make constraints for conditional expressions**)

```
| IF (c, t, f) ->
  newvar ()      >>= fun kv1 ->
  newvar ()      >>= fun kv2 ->
  newvar ()      >>= fun kv3 ->
```

Type Inference

```
(*TODO: make constraints for conditional expressions*)
| IF (c, t, f) ->
  newvar ()      >>= fun kv1 ->
  newvar ()      >>= fun kv2 ->
  newvar ()      >>= fun kv3 ->
  constr kv1 c env >>= fun c1 ->
  constr kv2 t env >>= fun c2 ->
  constr kv3 f env >>= fun c3 ->
  (kvar, kv2)::(kv1, KB)::(kv2, kv3)::c3@c2@c1 |> ret
```

Type Inference

(**TODO: make constraints for function definitions**)

```
| FUN (v, e) ->
  newvar ()          >>= fun kv1 ->
  newvar ()          >>= fun kv2 ->
```

Type Inference

```
(*TODO: make constraints for function definitions*)
| FUN (v, e) ->
  newvar ()      >>= fun kv1 ->
  newvar ()      >>= fun kv2 ->
  (v, kv1)::env |> fun ev ->
  constr kv2 e ev >>= fun c1 ->
  (kvar, KF (kv1, kv2))::c1 |> ret
```

Type Inference

(**TODO: make constraints for function applications**)

```
| APP (f, a) ->
  newvar ()          >>= fun kv1 ->
  newvar ()          >>= fun kv2 ->
```

Type Inference

(**TODO: make constraints for function applications**)

```
| APP (f, a) ->
  newvar ()          >>= fun kv1 ->
  newvar ()          >>= fun kv2 ->
  constr kv1 f env >>= fun c1 ->
  constr kv2 a env >>= fun c2 ->
  (kv1, KF (kv2, kvar))::c2@c1 |> ret
```

Type Inference

```
let rec unifier clist =
  match clist with
  | [] -> fun x -> x (*id substitution: no substitution*)
  | hd::tl ->
    match hd with
    | (KV v, ke) | (ke, KV v) ->
      if contains (KV v) ke
      then assert false (*recursive def is not supported*)
      else
        (*TODO: make a substitution for h*)
        let sub_h =                                     in
        let sub_t = tl
          (*TODO: apply sub_h to the remaining constraints*)
          |> List.map (fun (a, b) ->
                         )
          |> unifier in
        (*unifier is the composed substitutions*)
        fun e -> e |> sub_h |> sub_t

      (*TODO: unify function definitions*)
      | (KF (a, b), KF (c, d)) -> |> unifier
      ...
      ...
```

Type Inference

```
let rec unifier clist =
  match clist with
  | [] -> fun x -> x (*id substitution: no substitution*)
  | hd::tl ->
    match hd with
    | (KV v, ke) | (ke, KV v) ->
      if contains (KV v) ke
      then assert false (*recursive def is not supported*)
      else
        let sub_h = substitution (KV v) ke in
        let sub_t = tl
          |> List.map (fun (a, b) -> (sub_h a, sub_h b))
          |> unifier in
        (*unifier is the composed substitutions*)
        fun e -> e |> sub_h |> sub_t
    | (KF (a, b), KF (c, d)) -> (a, c)::(b, d)::tl |> unifier
      |> unifier
```

...