

HW 2 Key:

Q4.

Let $S[i,j]$ be the shortest string which is a super-sequence of $B1[1 \dots i]$ and $B2[1 \dots j]$.

$S[0,0] = 0$, $S[i,0] = i$, and $S[0,j] = j$.

Recurrence Formula:

$$S[i,j] = \text{Min} \left\{ \begin{array}{l} S[i-1,j-1] + 1 ; \text{ if } B1[i] = B2[j] \\ S[i-1,j] + 1 \\ S[i,j-1] + 1 \end{array} \right\}$$

Runtime: $O(nm)$, where length of $B1 = n$ and length of $B2 = m$.

Q5.

Similar to Q4, consider 3 sequences at a time.

Runtime: $O(nml)$, where length of $B1 = n$, length of $B2 = m$, and length of $B3 = l$.

For K sequences, generalize the above formula for K sequences, Runtime: $O(n^K)$

Q6.

Let $P[i,j]$ be the minimum number of inserts for string $S (S_i \dots S_j)$ to become a palindrome.

Initialize:

$$\begin{array}{l} P[i,i] = 0 \text{ for } i=0 \text{ to } n \\ P[0,i] = i \end{array}$$

Recurrence Formula:

$$P[i,j] = \text{Min} \left\{ \begin{array}{l} P[i+1,j-1] + 1, \text{ if } S[i] = S[j] \\ P[i,j-1] + 1 \\ P[i+1,j] + 1 \end{array} \right\}$$

Runtime: $O(n^2)$