CSE 613: Parallel Programming

Lecture 4
(Greedy Scheduling)

(inspiration for some slides comes from lectures given by Charles Leiserson)

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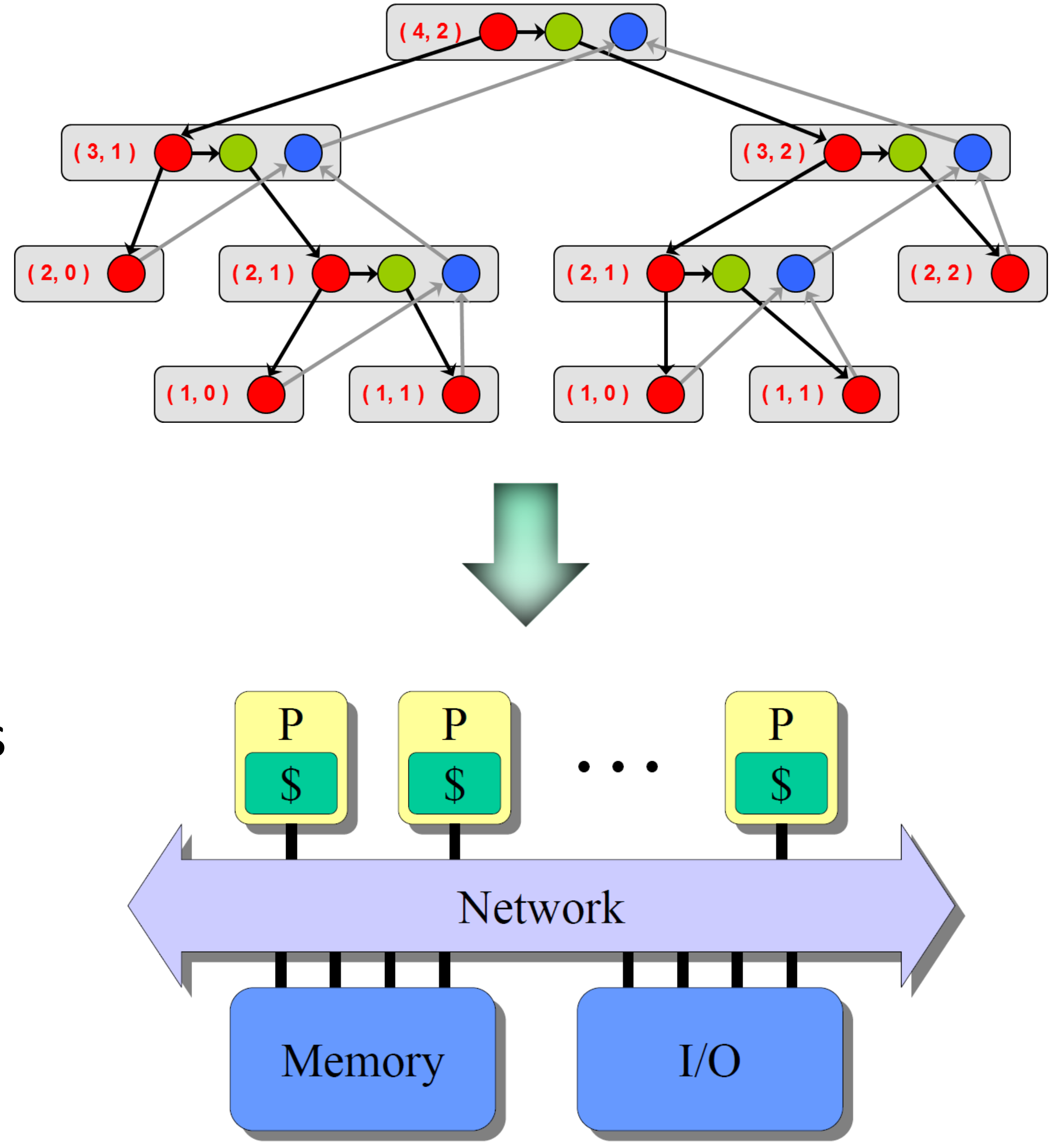
Spring 2019

Scheduler

A *runtime/online scheduler* maps tasks to processing elements dynamically at runtime.



The map is called a *schedule*.

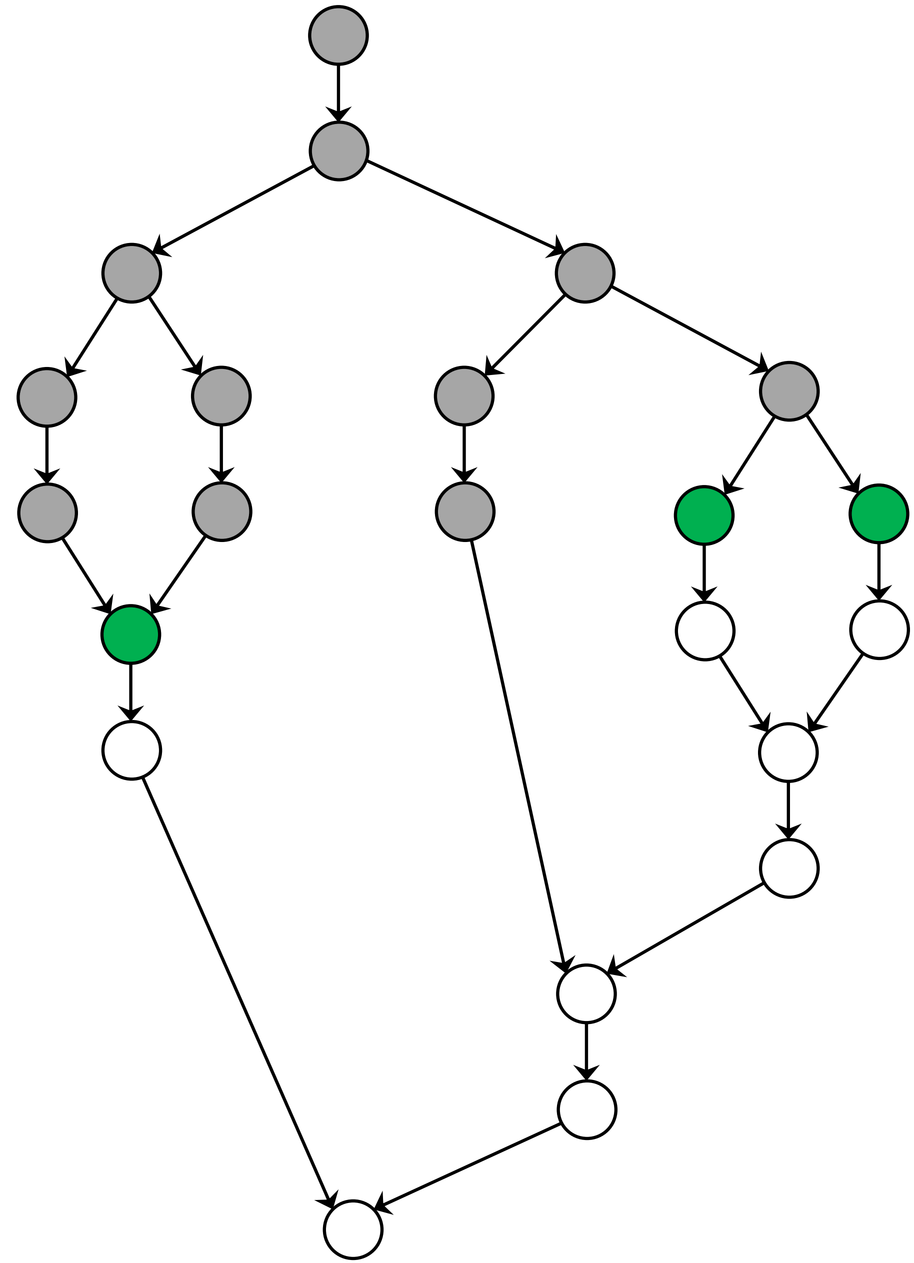
An *offline scheduler* prepares the schedule prior to the actual execution of the program.

Greedy Scheduling

A strand / task is called *ready* provided all its parents (if any) have already been executed.

- executed task
- ready to be executed
- not yet ready

A *greedy scheduler* tries to perform as much work as possible at every step.



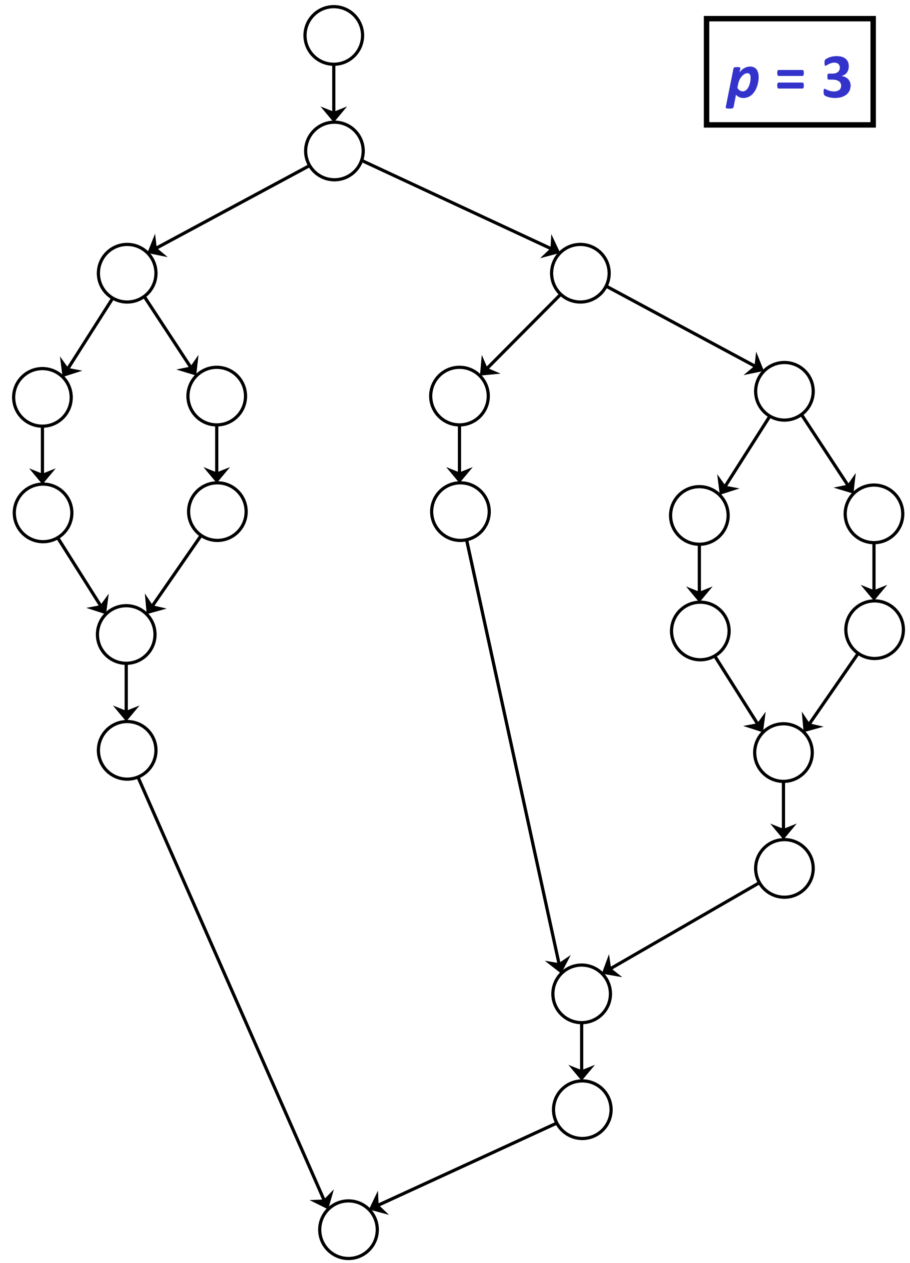
A Centralized Greedy Scheduler

Let p = number of cores

At every step:

– if $\geq p$ tasks are ready:
  execute any p of them
  (complete step)

– if $< p$ tasks are ready:
  execute all of them
  (incomplete step)



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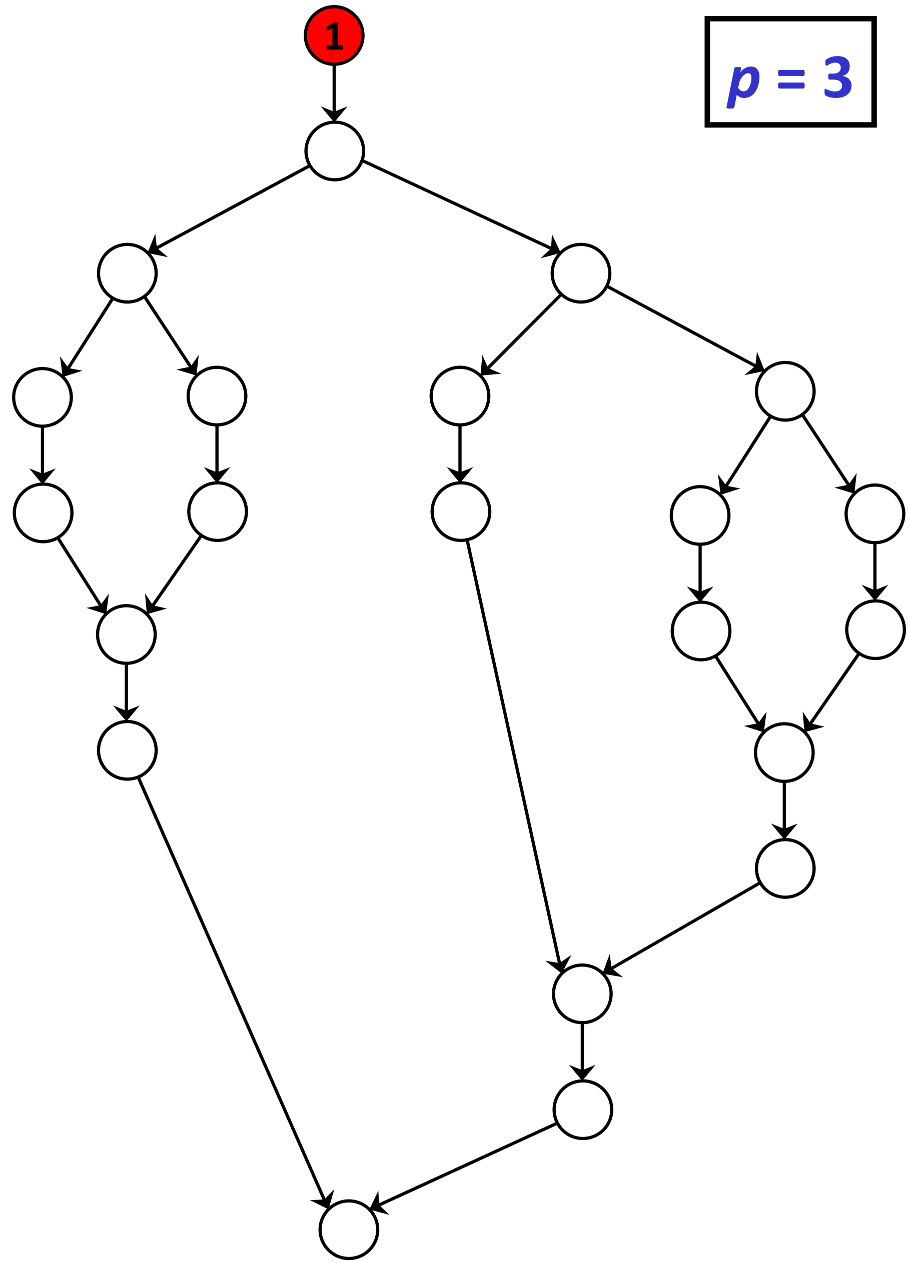

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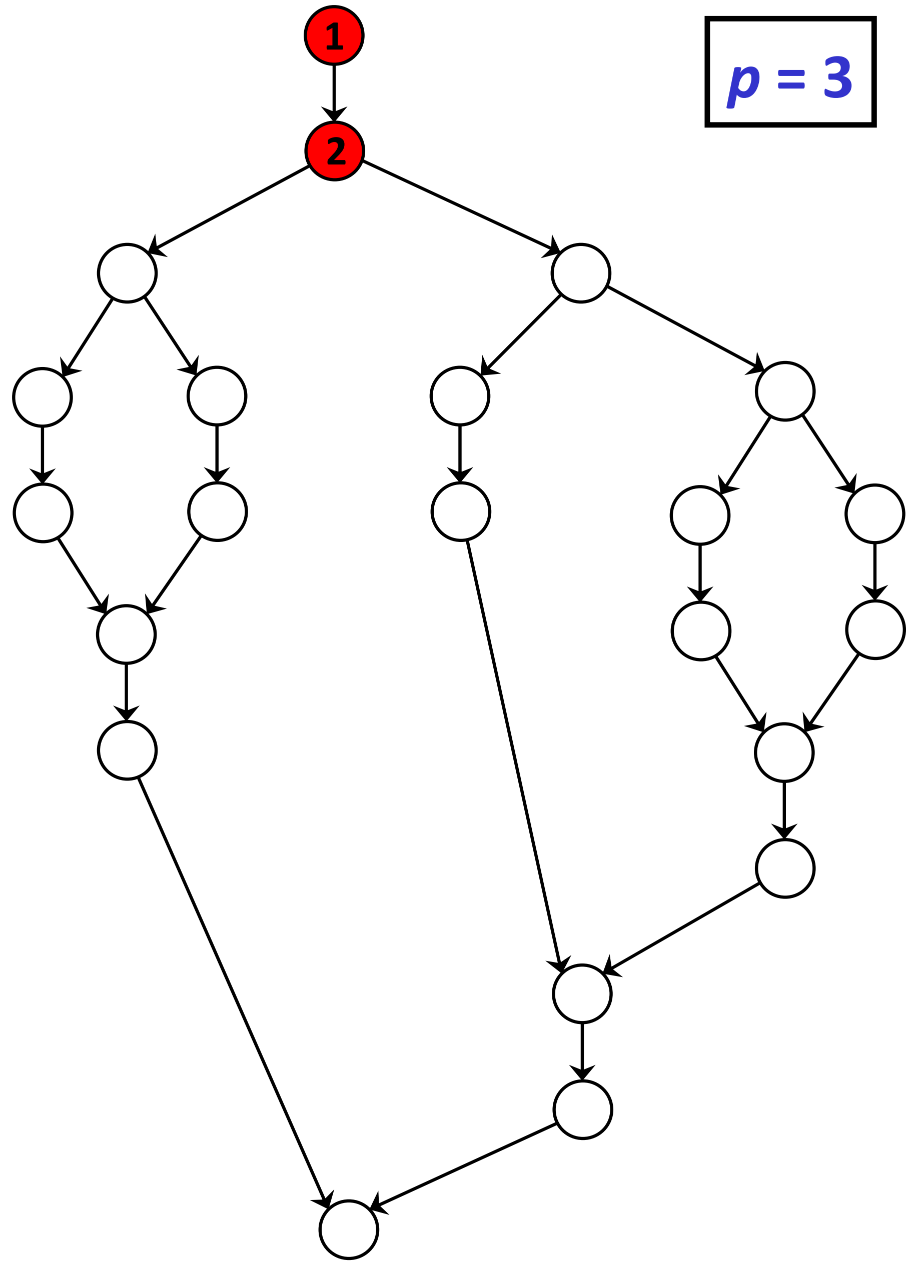

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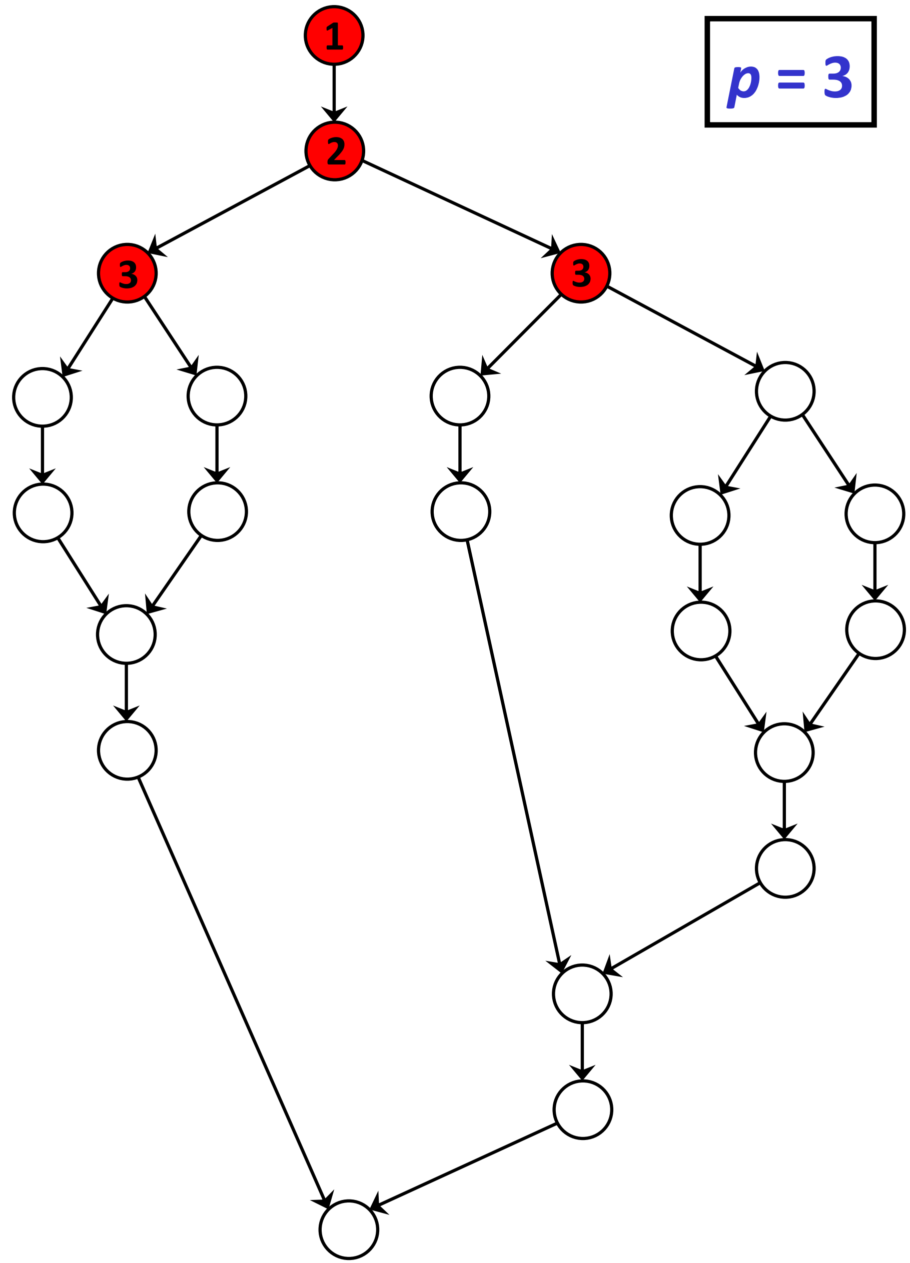

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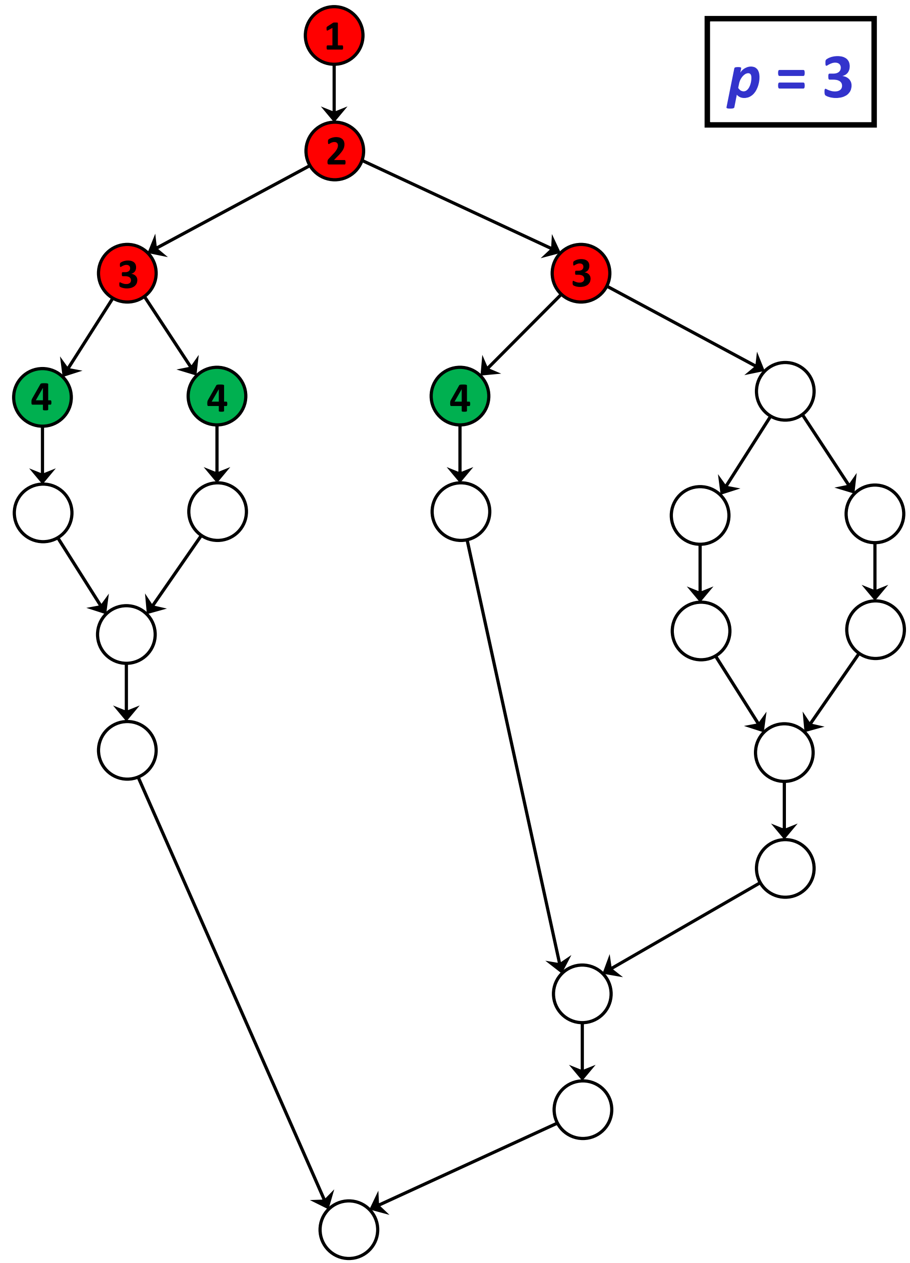

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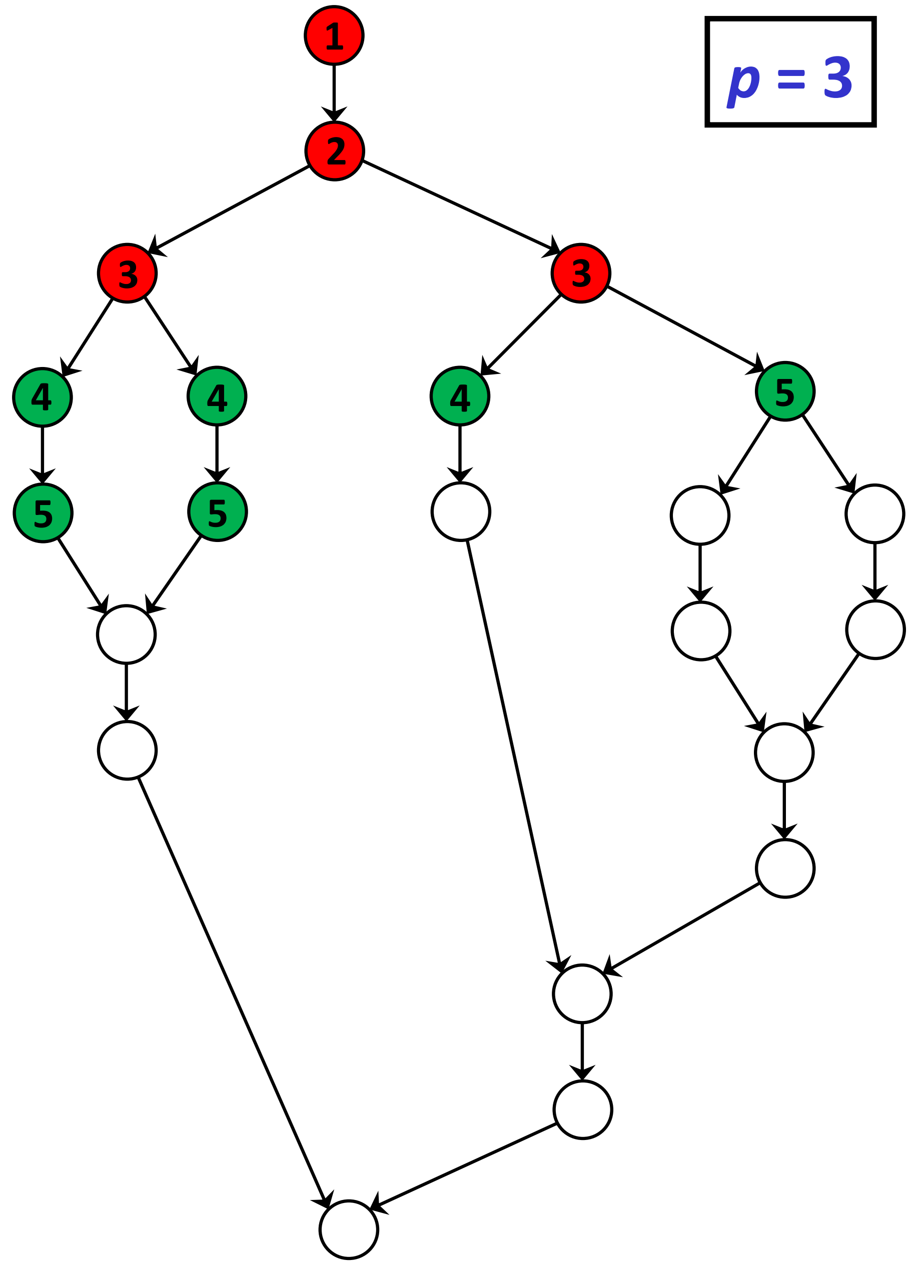

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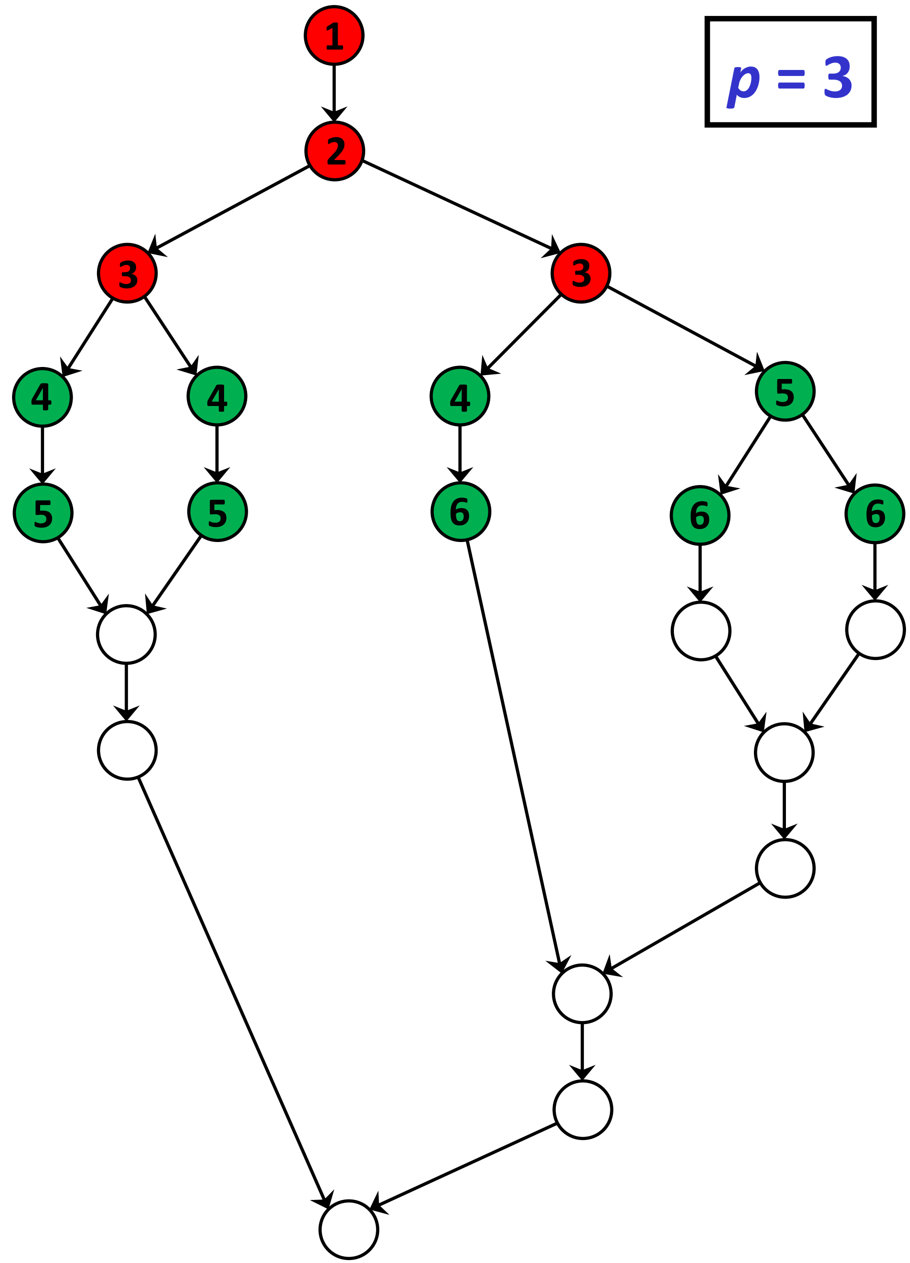

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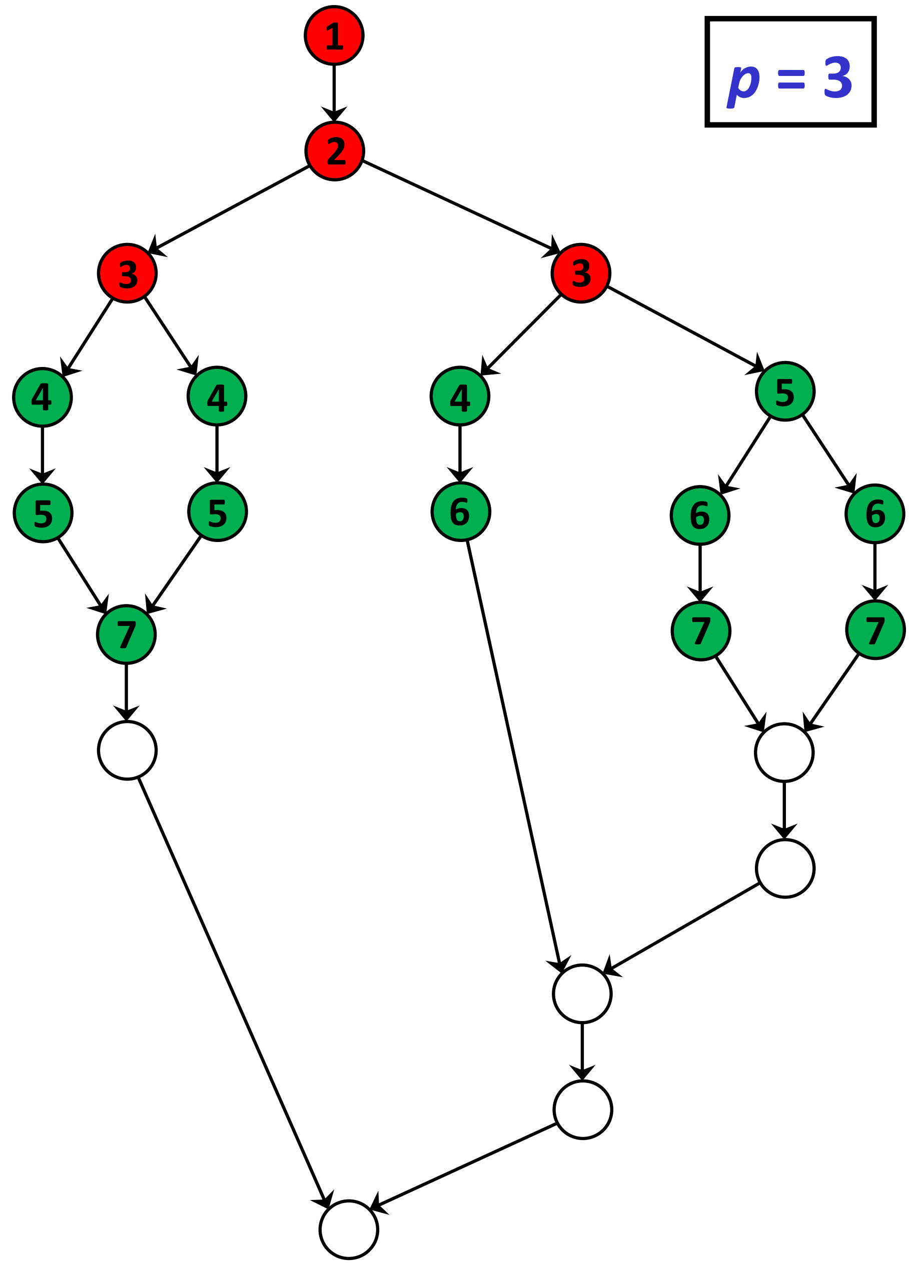

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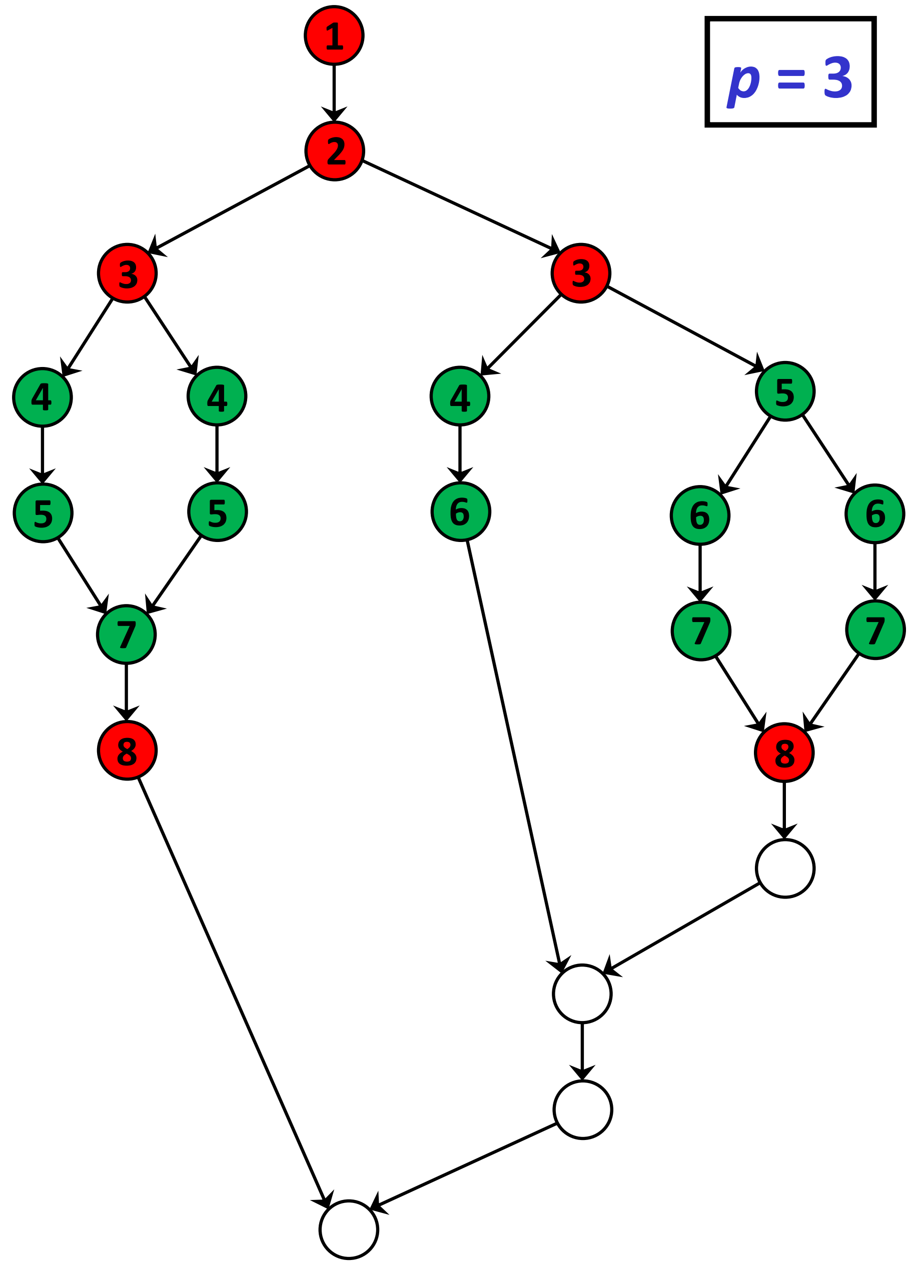

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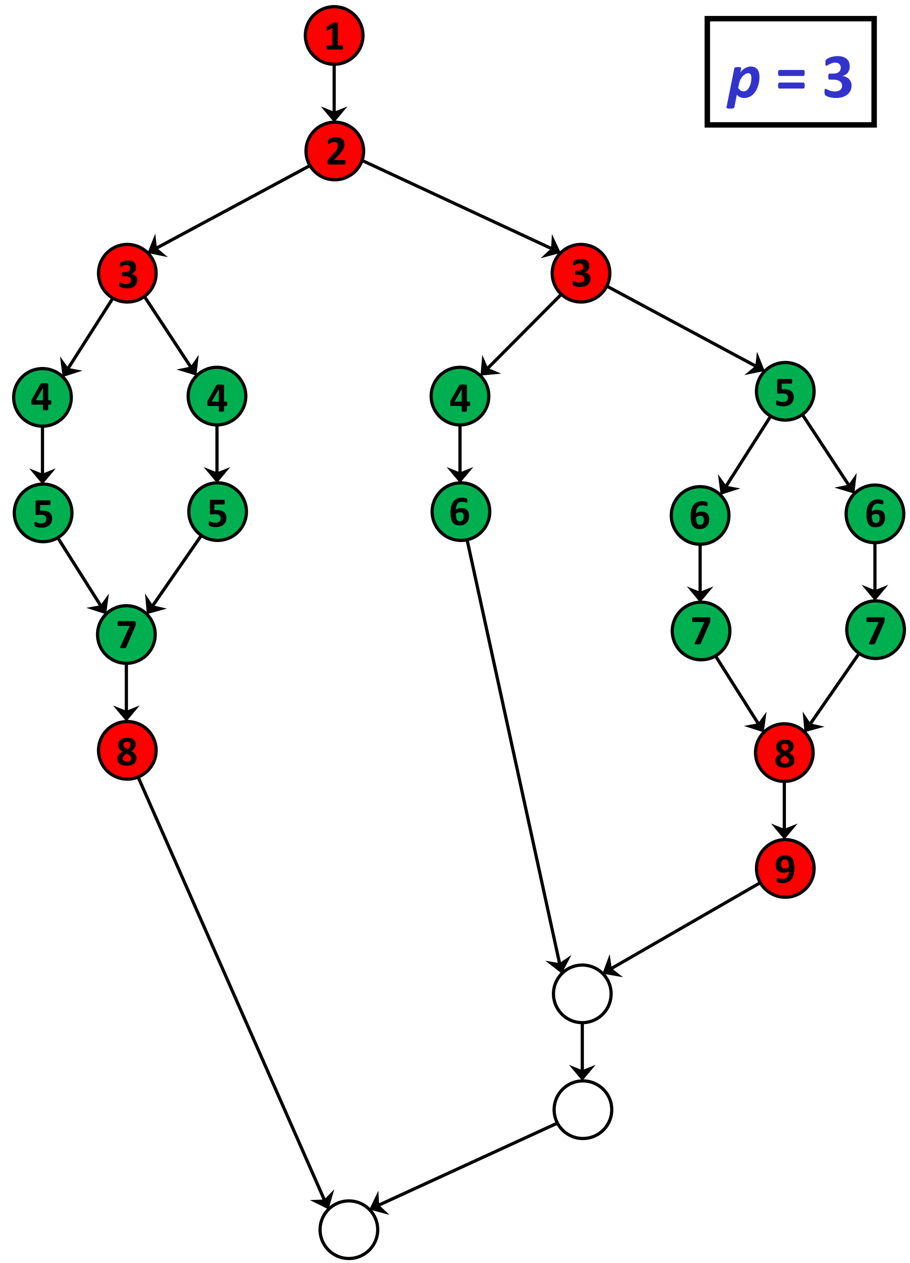

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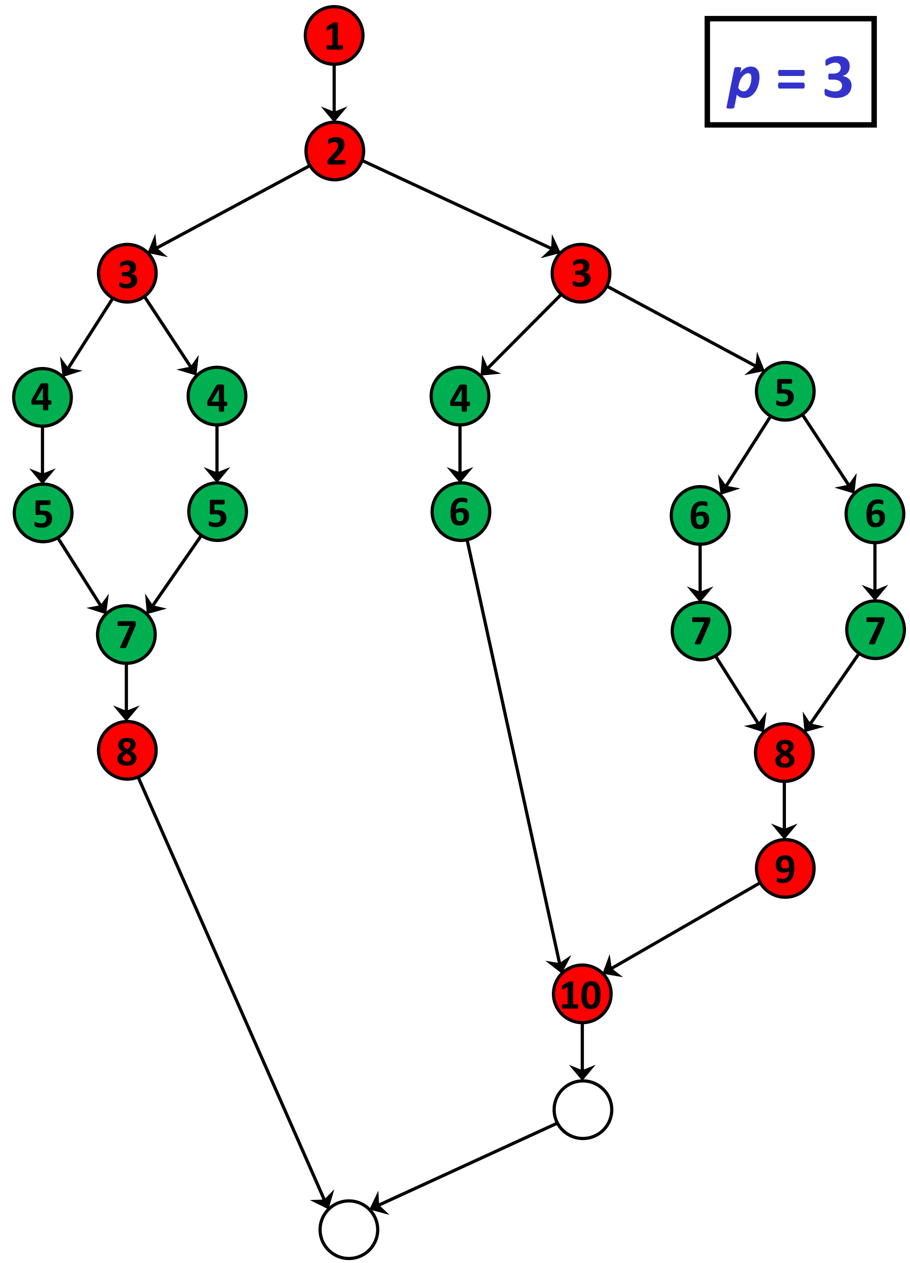

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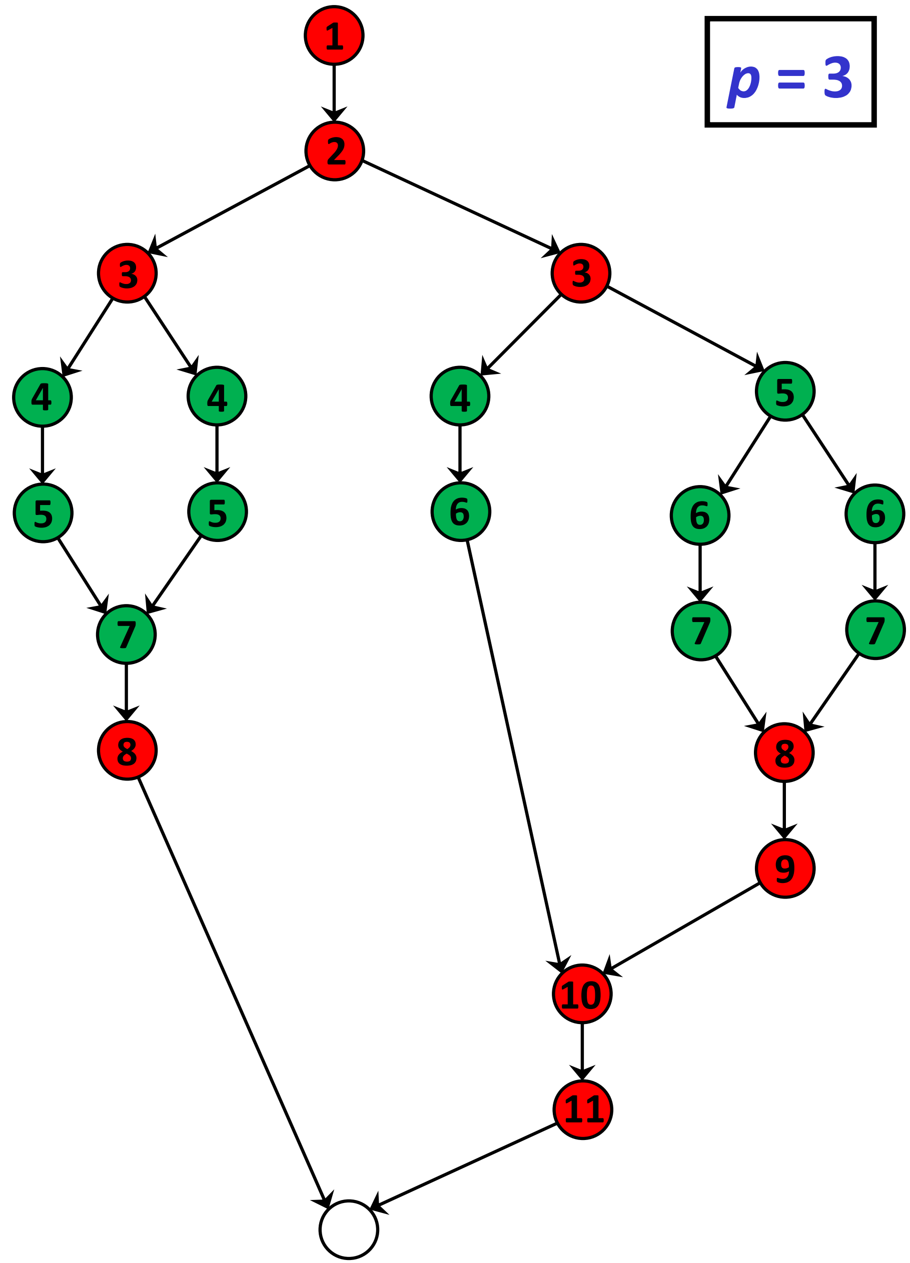

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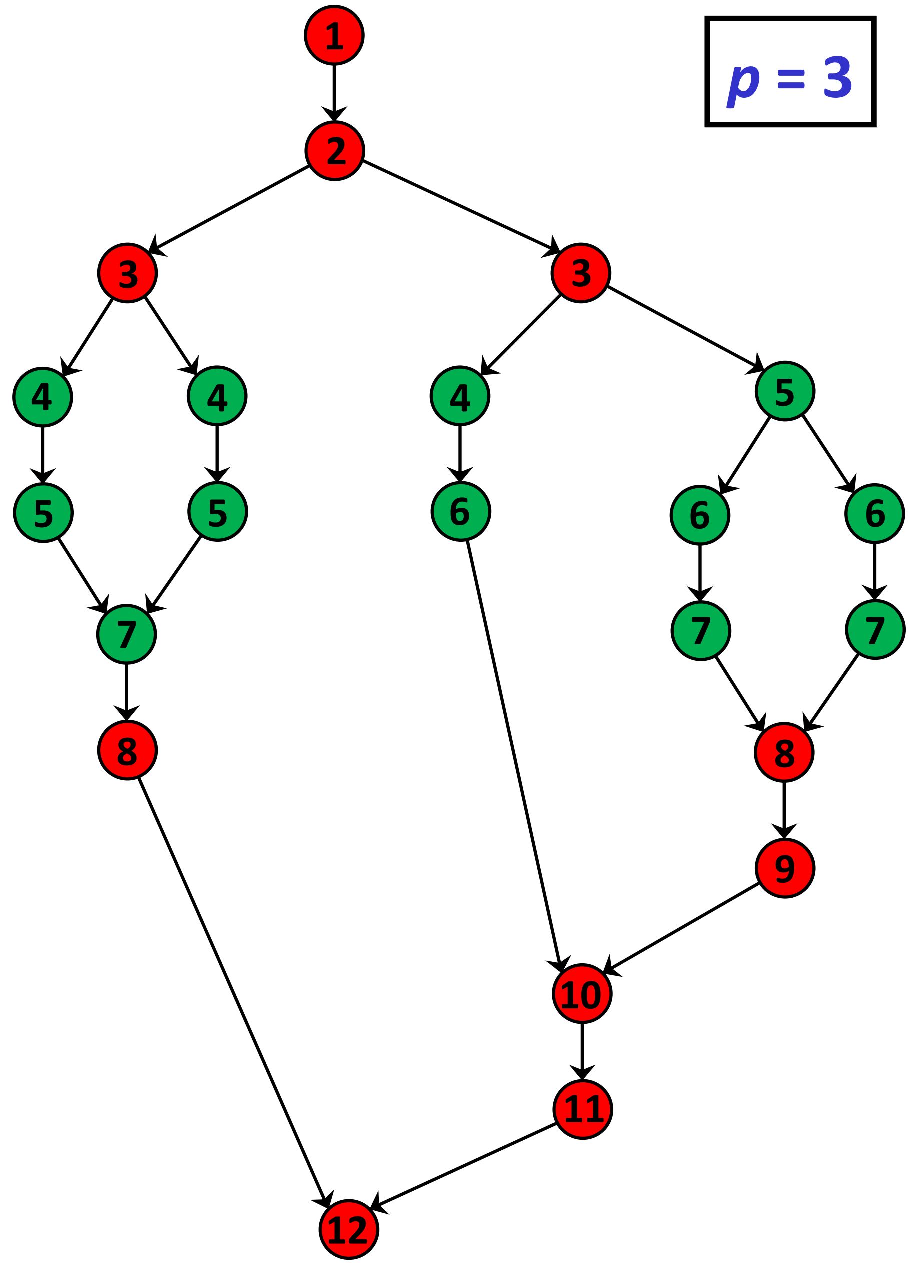

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Greed Scheduling Theorem

Theorem [Graham'68, Brent'74]:

For any greedy scheduler,

$$T_p \leq \frac{T_1}{p} + T_\infty$$

Proof:

T_p = #complete steps

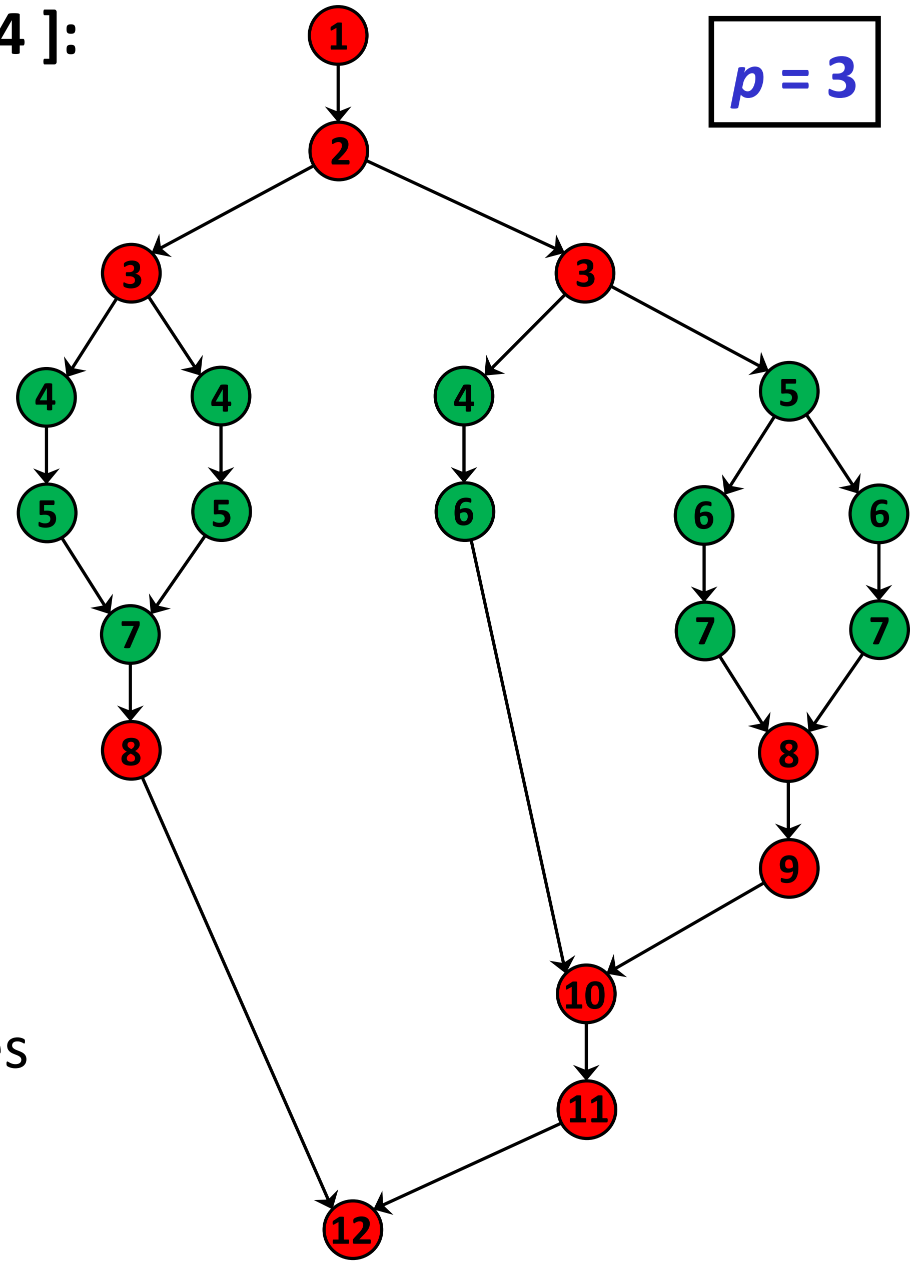
+ #incomplete steps

– Each complete step performs *p* work:

#complete steps $\leq \frac{T_1}{p}$

– Each incomplete step reduces the span by 1:

#incomplete steps $\leq T_\infty$



Optimality of the Greedy Scheduler

Corollary 1: For any greedy scheduler $T_p \leq 2T_p^*$, where T_p^* is the running time due to optimal scheduling on *p* processing elements.

Proof:

Work law: $T_p^* \geq \frac{T_1}{p}$

Span law: $T_p^* \geq T_\infty$

$\therefore$ From Graham-Brent Theorem:

$$T_p \leq \frac{T_1}{p} + T_\infty \leq T_p^* + T_p^* = 2T_p^*$$

Optimality of the Greedy Scheduler

Corollary 2: Any greedy scheduler achieves $S_p \approx p$ (i.e., nearly linear speedup) provided $\frac{T_1}{T_\infty} \gg p$.

Proof:

Given, $\frac{T_1}{T_\infty} \gg p \Rightarrow \frac{T_1}{p} \gg T_\infty$

$\therefore$ From Graham-Brent Theorem:

$$T_p \leq \frac{T_1}{p} + T_\infty \approx \frac{T_1}{p}$$

$$\Rightarrow \frac{T_1}{T_p} \approx p \Rightarrow S_p \approx p$$