

Sample Problems and Solutions

Discrete Mathematics

[Propositional Logic]

- [10 points] Give negations of the following statements using formal mathematical language. Your negations must not use conditionals or biconditionals.
 - [1 point] if p , then q
 - [1 point] p only if q
 - [1 point] p if and only if q
 - [1 point] p is necessary for q
 - [1 point] p is sufficient for q
 - [1 point] p is necessary but not sufficient for q
 - [1 point] p is not necessary but sufficient for q
 - [1 point] p is necessary and sufficient for q
 - [2 points] p is neither necessary nor sufficient for q

Solution.

- $p \wedge \sim q$
 - $p \wedge \sim q$
 - $p \oplus q$
(alternate answer) $(p \wedge \sim q) \vee (\sim p \wedge q)$
 - $q \wedge \sim p$
 - $p \wedge \sim q$
 - $(q \wedge \sim p) \vee (\sim p \vee q)$
 - $(\sim q \vee p) \vee (p \wedge \sim q)$
 - $p \oplus q$
(alternate answer) $(q \wedge \sim p) \vee (p \wedge \sim q)$
 - $(\sim q \vee p) \vee (\sim p \vee q) \equiv \mathbf{t}$
- [5 points] Determine if the following deduction rule is valid:
$$\begin{array}{l} p \rightarrow q \\ \sim q \rightarrow \sim r \\ \hline \therefore p \rightarrow r \end{array}$$

Solution.

p	q	r	$p \rightarrow q$	$\sim q$	$\sim r$	$p \rightarrow q$	$\sim q \rightarrow \sim r$	$p \rightarrow r$
T	T	T	T	F	F	T	T	T
T	T	F	T	F	T	T	T	F
T	F	T	F	T	F	F	F	
T	F	F	F	T	T	F	T	
F	T	T	T	F	F	T	T	T
F	T	F	T	F	T	T	T	T
F	F	T	T	T	F	T	F	
F	F	F	T	T	T	T	T	T

The deduction rule is invalid.

3. [5 points] Determine if the following deduction rule is valid:

$$\begin{array}{l}
 p \vee q \\
 \sim p \vee \sim q \\
 \hline
 \end{array}$$

$$\therefore p \oplus q$$

Solution.

p	q	$p \vee q$	$\sim p$	$\sim q$	$\sim p \vee \sim q$	$p \oplus q$
T	T	T	F	F	F	
T	F	T	F	T	T	T
F	T	T	T	F	T	T
F	F	F	T	T	T	

The deduction rule is valid.

4. [5 points] Determine if the following deduction rule is valid:

$$\begin{array}{l}
 p \rightarrow \sim q \\
 p \wedge q \\
 \hline
 \end{array}$$

$$\therefore p \vee q$$

Solution.

p	q	$\sim q$	$p \rightarrow \sim q$	$p \wedge q$	$p \vee q$
T	T	F	F	T	
T	F	T	T	F	
F	T	F	T	F	
F	F	T	T	F	

The deduction rule is valid (vacuous truth).

5. [5 points] Determine the validity of the following statements.

- (a) [1 point] $p \wedge \sim p \equiv \mathbf{t}$
- (b) [1 point] $p \rightarrow p \equiv \mathbf{t}$
- (c) [1 point] $p \oplus q \equiv (\sim p \vee q) \wedge (\sim q \vee p)$
- (d) [1 point] $p \rightarrow (q \rightarrow r) \equiv (p \rightarrow q) \rightarrow r$
- (e) [1 point] $(p \wedge q \wedge \sim r) \vee p \equiv p$

Solution.

- (a) False ($p \wedge \sim p \equiv \mathbf{c}$)
- (b) True
- (c) False ($p \oplus q \equiv (\sim p \wedge q) \vee (\sim q \wedge p)$)
- (d) False (Implication is not associative)
- (e) True

6. [5 points] Determine how many ordered truth assignments satisfy the following statement:

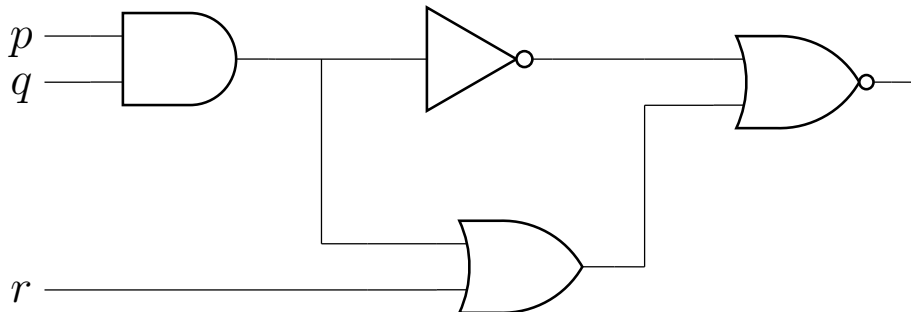
$$\sim (p \wedge q \wedge r) \vee (p \wedge \sim q \wedge r) \vee \sim p$$

Solution.

p	q	r	$p \wedge q \wedge r$	$\sim (p \wedge q \wedge r)$	$p \wedge \sim q \wedge r$	$\sim p$	Final Expression
T	T	T	T	F	F	F	F
T	T	F	F	T	F	F	T
T	F	T	F	T	T	F	T
T	F	F	F	T	F	F	T
F	T	T	F	T	F	T	T
F	T	F	F	T	F	T	T
F	F	T	F	T	F	T	T
F	F	F	F	T	F	T	T

7 truth assignments.

7. [5 points] Given the following logic circuit, construct an Input-Output table.



Solution.

p	q	r	$p \wedge q$	$\sim (p \wedge q)$	$(p \wedge q) \vee r$	$\sim (\sim (p \wedge q) \vee ((p \wedge q) \vee r))$
1	1	1	1	0	1	0
1	1	0	1	0	1	0
1	0	1	0	1	1	0
1	0	0	0	1	0	0
0	1	1	0	1	1	0
0	1	0	0	1	0	0
0	0	1	0	1	1	0
0	0	0	0	1	0	0

8. [5 points] Construct the truth table for the following logical expression:

$$((p \rightarrow q) \wedge r) \leftrightarrow \sim (p \oplus q)$$

Solution.

p	q	r	$p \rightarrow q$	$(p \rightarrow q) \wedge r$	$p \oplus q$	$\sim (p \oplus q)$	$(p \rightarrow q) \wedge r \leftrightarrow \sim (p \oplus q)$
<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>
<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>
<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>F</i>	<i>F</i>
<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>
<i>F</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>
<i>F</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>

9. [5 points] Let us define a new operator \diamond with the following truth table:

p	q	$p \diamond q$
<i>T</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>T</i>
<i>F</i>	<i>T</i>	<i>F</i>
<i>F</i>	<i>F</i>	<i>T</i>

Let us extend another of de Morgan's laws as shown below. Determine whether this property holds for the operator \diamond :

$$p \diamond (q \wedge r) \stackrel{?}{\equiv} (p \diamond q) \vee (p \diamond r)$$

Solution.

p	q	r	$q \wedge r$	$p \diamond q$	$p \diamond r$	$p \diamond (q \wedge r)$	$(p \diamond q) \vee (p \diamond r)$
<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>
<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>F</i>
<i>F</i>	<i>T</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>
<i>F</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>
<i>F</i>	<i>F</i>	<i>F</i>	<i>F</i>	<i>T</i>	<i>T</i>	<i>T</i>	<i>T</i>

Yes, the property holds.