

Exploratory Visual Analysis of Anomalous Runtime Behavior in Streaming High Performance Computing Applications

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Abstract. The online analysis of abnormal runtime behavior is essential in streaming scientific workflows for performance tuning. Integration of anomaly detection and visualization is necessary to support efficient and scalable human-centered analysis, such as verification of candidate anomalies utilizing domain knowledge. In this work, we propose a visual analytics system for the online performance analysis of scientific workflows toward the exascale scenario. Our approach adopts the call stack tree representation, which encodes the structural and temporal information of the function executions. Based on the call stack tree features (e.g., execution time of the root function or vector representation of the tree structure), we employ online anomaly detection approaches to identify the candidate anomalous function executions. We also devise a set of visualization tools for the verification and exploration in a level-of-detailed manner. The general information such as the distribution of the execution times are provided in an overview visualization. The detailed structure (e.g., function invocation relations) and the temporal information (e.g., message communication) of the execution call stack of interest are also visualized. The usability and efficiency of our methods are verified in the NWChem use case.

Keywords: Anomaly Detection · High Performance Computing · Streaming Analysis · Trace Events · Visual Analytics.

1 Introduction

Performance analysis is a critical task for the diagnosis of parallel High Performance Computing (HPC) applications. In particular, domain scientists are typically interested in the online detection of abnormal runtime behaviors in the execution of HPC applications. Since the supercomputer resources in use are limited and costly, the timely identification of the causes of adverse performance events (e.g., abnormal communication latencies) is essential. We have been working with a group of chemists who use an HPC cluster to solve complex molecular equations. The development of the system subject of this paper was driven by their need to monitor and identify computation latencies at runtime.

Anomalous function executions are usually identified by examining the detailed trace events collected in the HPC cluster. A trace is a sequence of events collected by program instrumentation in a computing core during application execution, such as function entry, function exit, and message passing. Fig. 1 (a) shows an example sequence of trace events inside one execution of the `main` function in an HPC core. It represents the call stack information (Fig. 1 (b)) inside the execution of the root function. Domain scientists typically detect anomalous function executions based on the extracted temporal information [18] [2] (e.g., execution time and exit timestamps) or structural information [23] (e.g., call relations from parent function to children) from the trace events. While

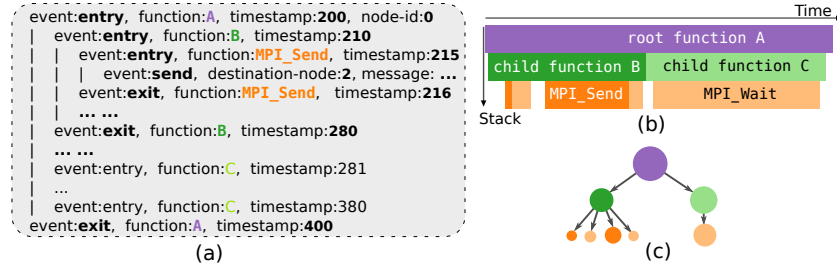


Fig. 1. (a) Example trace events during an execution of `main` in a compute core. (b) The call stack of the execution reconstructed from the trace events. (c) The call stack represented as a Call Stack Tree (CSTree), which is a directed tree with vertex weights.

existing detection approaches achieve good performance and scalability in offline analysis, detecting abnormal runtime behavior in online analysis remains challenging. One reason for this is that the complexities of most feature extraction and anomaly detection algorithms are too high to support online training. It would be prohibitively slow to use a standard algorithm like Local Outlier Factor (LOF) [3] to update the learning result every time a new trace event is generated. Conventional anomaly detection models do not support continuous updates, rather they need to be re-trained in each time window. Furthermore, in order for human experts to stay aware and gain insights into the performance data an interactive visual interface will be most appropriate. However, offline visual analysis tools are not designed to provide visualizations sufficiently responsive in a streaming data environment. For example, standard visualization algorithms such as Multidimensional Scaling (MDS) [12] or t-Distributed Stochastic Neighbor Embedding (t-SNE) [13] are not fast enough to plot and visualize tens of thousands of points at the rate of streaming data.

We recently devised a visual analysis approach for the detection of abnormal HPC runtime behavior [23]. This tool, however, was designed for offline analysis and cannot be used for online analysis of streaming data. In the current work we build on this approach to create a new method for streaming trace event data. Similar to our earlier work we make use of the call stack tree (CSTree) representation we devised there (Fig. 1 (c)), but we now use it to encode the structural as well as the temporal information inside a function execution via

a directed tree. Abnormal behavior of function execution can then be detected by identifying abnormal trees in a call stack forest. We also modify our original tree feature extraction process to now support online tree representation and online anomaly detection. We propose a set of new online visualizations to assist the scientist in the understanding and exploration of the function executions. In particular, we visualize the feature vectors and the learned anomaly labels of the function executions in an overview projection. The structural and temporal information of the CSTree are visualized in a structure and timeline visualization, respectively, capable of dealing with large amounts of data in real time.

The remainder of our paper is structured as follows. Section 2 reviews related work. Section 3 defines the problem and gives an overview of our approach. Section 4 introduces our online algorithms for vector representation and anomaly detection of CSTrees. Section 5 describes our visual analytics approach for anomalous CSTree exploration. A case study is used to validate our approach in Section 6. Section 7 ends with conclusions and future work.

2 Related Work

Domain scientists usually employ instrumentation and measurement tools [1] [5] [10] [16] to generate the performance data, such as trace events, and so on. Many techniques have been proposed for the evaluation and diagnosis of these data [4]. In this paper, we focus on the most important two tasks in the diagnosis of the HPC application performance: detection of abnormal runtime behavior and visualization of performance.

2.1 Abnormal Runtime Behavior Detection

For anomalous runtime behavior detection, most existing approaches identify the candidates based on the temporal information [7]. For example, a large function execution has a very high probability of being abnormal since it has a large negative impact on the downgraded performance. However, diagnosis without the execution context makes it difficult to determine the source of a major latency. For example, a function execution may spend most of the time in the children functions, but the delay can also be caused by the HPC node it communicates with. Furthermore, semantics of the executions are also necessary for the diagnosis. For example, the initialization function usually takes more time to set up the calculation or communication. However, a long initialization executions cannot be identified as abnormal as they are running with expected behavior.

In contrast, the Call Stack Tree representation (CSTree) [23] can encode the temporal information as well as the context structure obtained from the call stacks to identify the potential anomalous executions. The embedding vectors are then generated from the CSTrees. Each vector encodes the structural information from the call stack of one function execution. A conventional anomaly detection algorithm such as the One-Class Support Vector Machine (OCSVM) [15] is then able to take the vectors and identify the anomalies.

However, both the tree representation and the anomaly detection have high complexities for online analysis. To deal with that, in this paper, we modify the training strategy of these learning algorithms for the analysis of streaming data.

2.2 Performance Visualization

Performance visualization makes problem detection and diagnosis of HPC applications [7] [22] more transparent to domain experts. Most visualization approaches support the comprehension of different levels and aspects of the performance data, including the trace timeline, call stack structures, and the messages.

For the trace events, a common practice is to visualize the events along a time axis, as is done in Vampir [9] and Jumpshot [24]. Most existing temporal visualizations provide level-of-detail explorations. The overview of the temporal trend or statistics is usually shown first. Then for the time window of interest, users can zoom in to see more detailed events [21]. Other temporal visualizations are also capable of presenting the relationships between threads, such as SyncTrace [8].

Call path visualization (e.g., the call relationship between parent and children functions) is critical for understanding the behavior of the runtime execution. Existing approaches employ a directed tree or graph to present the structure in a call stack, such as Vampir [9]. These visualizations usually use the visual properties of the tree to encode detailed information from the call path. For example, CSTree [23] utilizes the color and size of a tree node to encode the type and the execution duration of a function.

Communication delay is usually the main reason for application latency, and so visualizing the message passing activities is necessary for diagnosis. A straightforward approach to encode the message passing is to draw a directed line between the sending and receiving functions, which is adopted in the Jumpshot [24] implementation. Since the messages can also be regarded as directed edges, the communication between threads or processes can also be summarized in terms of an adjacency matrix [9].

One major issue with the existing visualization approaches is their limited capability for online performance evaluation. Some offline analysis tools focus on the complete event or message passing structure which is available only when the communication is finished. In addition, most visualization paradigms, such as MDS, t-SNE, etc. are too complex for streaming data. They use global optimization to produce a consistent 2D layout that preserves the relations the data have in high-dimensional space. In order to handle online applications, we need to update this global optimization which can be costly. While methods for progressive t-SNE have emerged [14], they require users to make trade-offs in speed and accuracy. On the other hand, we also need to deal with online data reduction and sampling to prevent overdraw. In our approach, we adjust the common projection, timeline, and tree visualizations for streaming to facilitate incremental updates and data visualizations.

3 Problem Formulation and Approach Overview

We focus on the following problem: Given a set of functions of interest (FOIs) $\{A\}$ and all of their invoked executions in an HPC cluster, determine which executions are associated with anomalous runtime behavior. An FOI A is usually a key function for computation or communication, which is specified by the domain experts. An anomalous runtime behavior is then in most cases indicated by the temporal or structural features. For example, a deadlock will cause large execution time and an infinite loop will generate unexpected call path structure.

3.1 Call Stack Tree Representation and Problem Formulation

As mentioned in Section 2, the Call Stack Tree (CSTree) (Fig. 1 (c)) representation provides a comprehensive way to encode the execution of an FOI A since it takes advantage of the execution’s context information. Each execution of A is converted to a CSTree T where a vertex in T is a function invoked in the call stack and a directed edge represents the call from the parent to the child function. All executions of A collectively give rise to a forest $\mathcal{T} = \{T_i\}$ where each tree T represents a single execution of A . The runtime behavior can be observed directly from the features of a CSTree. For example, a very large vertex in a tree can represent a delayed function execution, while a large set of children vertices indicates that the same child function is invoked many times by the parent function in a loop.

Given this representation the detection of anomalous behavior is then formulated as the problem of finding anomalous tree structures in the call stack forest. Our visual interface exposes the CSTrees whose structure are most different, which are the candidate anomalies.

3.2 Approach Overview and System Architecture

Based on the CSTree representation, our online visual analytics approach for detecting anomalous CSTrees uses the following four steps with the architecture³ in Fig. 2:

Step 1 Data processing: The trace events (Fig. 2 (a)) generated during the execution of an HPC application are paired and ordered by their type (i.e., message or function event) in the data analysis server (Fig. 2 (b)). Given a set of functions of interest $\{A\}$, the call stack trees rooted at A can also be generated and inserted into the call stack forest $\mathcal{T} = \{T_i\}$.

Step 2 Tree feature extraction and anomaly detection: Each new CSTree in the forest is converted into a feature vector, which can consist of the temporal features of the root functions (e.g., total execution duration). Other options for feature extraction such as the graph kernel are also provided. Then anomaly detection is performed based on these features.

³ Please visit <https://github.com/CODARcode/ChimbukoVisualization> to see the project page and the source code.

Step 3 Overview visualization: An overview projection is calculated by the visualization server (Fig. 2 (c)). The visualization platform is the web browser (Fig. 2 (d)) showing the feature vector distribution and learned labels of CSTrees.

Step 4 Detailed visual exploration: The detailed visualizations enable the user to identify the true anomalies. By specifying the candidate anomalous CSTrees in the overview, the user can view their structures and message communications, which are requested from the visualization server.

The step 3 and 4 can be performed repeatedly for better insights into the anomaly detection results.

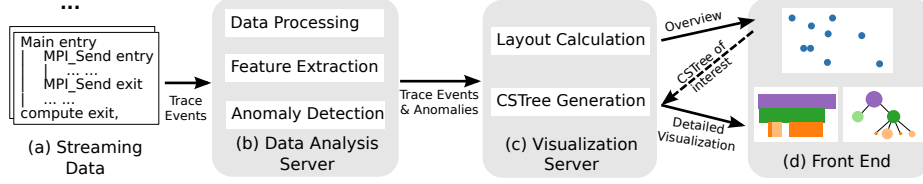


Fig. 2. Our system architecture for online anomalous execution detection. (a) The streaming trace events. (b) Feature vectors of the new CSTrees are generated in the data analysis server. Then the anomaly labels (normal or abnormal) of the CSTrees are learned online based on their feature vectors. (c) The visualization server calculates the layout of the overview and generates the CSTree of interest for detailed exploration. (d) On the browser end, the user is allowed to interact with the overview of the tree distribution in a 2D view. Using the candidates of interest in the overview, further investigation are supported by detailed visualizations.

4 Online Anomaly Detection of Call Stack Trees

Our approach begins with updating the call stack forest in the data server. For a FOI *A*, a new CSTree is generated and inserted into the call stack forest when an exit event of *A* is received in the trace stream. Since the user will focus on the recent data, the CSTrees older than a given threshold (e.g., 1 hour) will be removed from the forest. Then the feature extraction and anomaly detection are performed for the updated call stack forest in the data analysis server.

4.1 Feature Vector Representation

One of the major challenge of our approach is how to extract the temporal and structural features from the CSTree for anomaly detection. We provide different options, including time-based and tree-based feature extraction.

Time-Based Representation A straightforward way to represent runtime behavior is to utilize the information from the root function of the CSTrees. The domain expert is able to customize the specific set of features to be extracted, such as the execution duration, message frequency, and the exit time of the execution. With this pre-designed feature set, a vector of each CSTree is constructed and used as the input for anomaly detection.

Tree-Based Representation Thus far there is no contextual information in the above temporal feature construction, however, structural representation is critical for the diagnosis of the runtime behavior. To provide an option for structural feature extraction, we follow the Graph Kernel [23] approach for the vector representation, which uses an analogy to document analysis. A tree is analogous to a document while the subtrees are analogous to the words in a document. The duration of each subtree in a CSTree is synonymous to “word frequency” in a document. The subtrees (Fig.3 (b)) are extracted using Weisfeiler-Lehman Graph Kernels [17], and then the CSTree (Fig.3 (a)) is represented as a bag-of-subtrees (Fig.3 (c)).

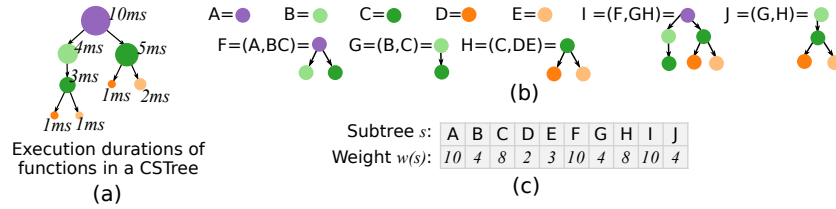


Fig. 3. (a) An example CSTrees with the function execution time of the vertices. (b) An example of the result substructure generated from a CSTree using Weisfeiler-Lehman Graph Kernels. (c) The bag-of-subtree representation of the CSTree.

4.2 Online Anomaly Detection

The extracted feature vectors of the CSTrees are used as the input of the anomaly detection. Conventional algorithms such as Local Outlier Factor can be too complex for online training. For better online performance in the presence of a very large amount of streaming data we chose a fairly straightforward statistical approach. Using a Gaussian distribution to model the feature vector distribution of the CSTrees, we calculate the standard deviation of the feature vectors in the call stack forest. We then label all CSTrees at 3σ (i.e., 0.003%) of the data distribution as anomaly candidates. In our own experiments we found this to be a good threshold, but the confidence level can be user-adjusted based on the estimated percentage of anomalies in the dataset.

5 Visual Exploration of Anomalous Call Stack Trees

It is necessary that human experts can verify the learning results to make sure the identified candidates are true anomalies. In addition, exploration of the temporal and structural patterns of the candidates also helps the user understand why they are potentially anomalous. In this section, we describe our level-of-detail visualization system designed for the exploration and verification of the CSTrees in different granularities.

5.1 Overview Visualization

The overview visualization is calculated in the visualization server and displayed in a browser-based interface. The overview shows the general distribution of the feature vectors of CSTrees. As noted above, the most common approach available for this purpose is a low-dimensional embedding method such as MDS. However, after some experimentation we determined, as we also mentioned above, that these methods are not sufficiently fast and scalable for large streaming datasets.

To deal with this problem, we resorted to a standard bivariate projective scatterplot approach. In this visualization the x and y positions are calculated by the basic attributes of the CSTrees specified the user. For example, in Fig. 4 (a), the CSTrees are visualized as points in the scatter plot. They are projected by the execution time and entry time of their root functions. The color of each points represent the FOI type. The points highlighted with thick borders are candidate anomalies which are detected in the previous step.

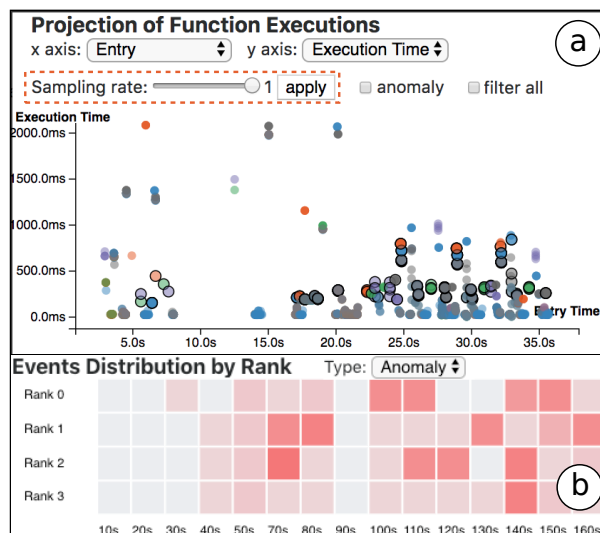


Fig. 4. (a) In the NWChem application, we project the CSTrees by execution time and entry time of the root functions. The user can specify other axis-encoding schema for projection. Filtering and negative down sampling are supported to reduce visual clutter. (b) The heatmap visualizes the temporal distribution of all the history data.

Display scalability can still be a problem when the data volume is massive. After discussing this with a domain expert we decided to adopt a negative down sampling strategy to reduce the data that are displayed. Since users are typically mainly interested in the candidate anomalies, they can set a down sampling rate (see the dashed box in Fig. 4 (a)) to keep only a portion of the normal points in the projection. At the extreme, when the down sampling rate is set to 0, the

projection will only visualize the candidate anomalies. In addition, users can also filter the displayed data by FOI type. This helps users to put more focus on a subset of more important FOIs.

In an online analysis, as opposed to a postmortem analysis, users will usually want to focus on the latest trace data. Hence, the visualization server will only maintain a subset of the data, namely those collected within a given time period (e.g., 1 hour). To keep the display fresh, we set up a fixed time window (e.g., 60 seconds) and update the scatter plot whenever the latest CSTrees appear. On the other hand, even in online analysis, keeping a historical perspective is desirable. We provide this view by visualizing the distribution of all history data in form of a heatmap (see Fig. 4 (b)). In this map the user is free to specify the color encoding, such as for highlighting the number of anomalies or trace events.

Further, the user can specify the points of interest in the scatter plot, which will send requests to the visualization server for more detail on these selections. The structure and timeline visualizations will then provide the corresponding CSTree structures and events sequences for further exploration.

5.2 Structure Visualization

The details of the selected CSTree are visualized in the Structure Visualization, as shown in Fig. 5 (a). The tree vertex size and color represent the execution and function name of the corresponding tree node. The directed edge shows the call relationship from a parent to a child function. Force-directed Layout [11] is employed to calculate the position of the tree vertices. For clearer representation, we set a maximum depth to only keep the important parent functions which have major influences on the tree. This reduces the visual clutter, which is helpful especially for CSTrees with a large number of descendant functions. On the other hand, in order to provide complete anomaly information, the abnormal substructures are preserved in the visualization even if they are beyond the maximum depth limit.

5.3 Timeline Visualization

While the structure visualization is effective, the message communication between different HPC cores is not visualized. To deal with this problem, we propose the Timeline Visualization (Fig. 5 (b)), which basically follows the visual design of Vampir [9] and Jumpshot [24]. Our visualization shows the event sequence as well as the message passing in a stack timeline. The x and y axes of the timeline encode the timestamp and the growing direction of the call stack. Since communication is one of the main reasons for HPC application delays, messages are also visualized, via lines between different cores, as indicated in Fig. 5 (b). Users can zoom in and out of the x-axis (see red box in Fig. 5 (b)) and so explore the detailed call stack in different time ranges, as shown in Fig. 5 (c).

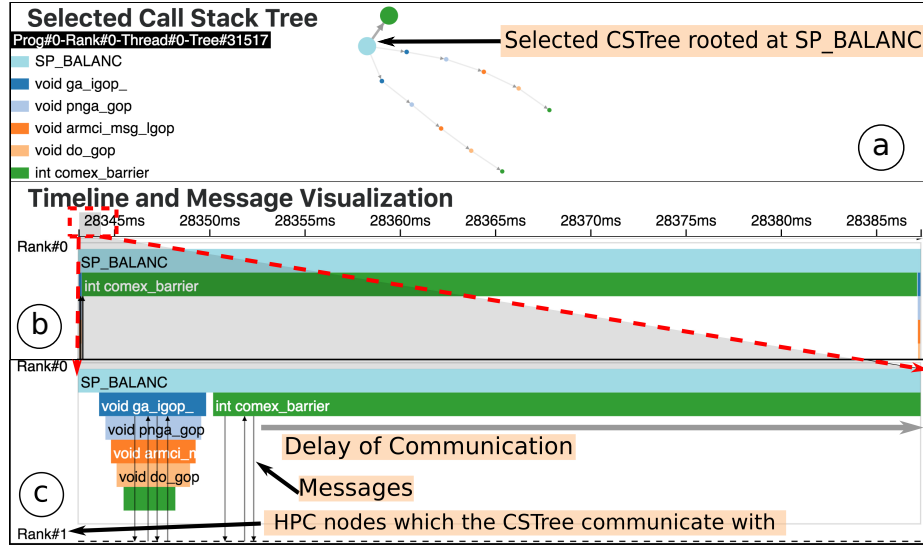


Fig. 5. (a) The call stack tree structure visualization only shows the tree at a limited depth or abnormal substructures. (b) The timeline follows the same visualization as in the offline analysis. (c) The timeline can be zoomed in for detailed exploration.

6 Case Study

We conducted a case study with a scientist from at Brookhaven National Laboratory. He was working on the NWChem [19] application, which is a parallel computational chemistry application in the HPC cluster. In this study we focused on the molecular dynamics functionality.

6.1 Experiment Settings and Online Performance Evaluation

Our participant is a chemist and not an expert in visualization or machine learning and sought to use our system to find anomalous runtime behaviors in the function executions. In the onset, we had a number of thorough discussions with this domain scientist to learn about functions of interest and possible anomalous behavior patterns. We also held a short training session to introduce our visualization system.

We applied the temporal feature extraction described in Section 4.1 since it was faster for large scale datasets. We also employed the anomaly detection by standard deviation described in Section 4.2. After studying the dataset in our interface, the participant set the significance level defining a candidate anomaly to be roughly 3σ (i.e., 0.003%) to reduce false positives.

The streaming trace events of NWChem were generated by SOSFlow [20] and ADIOS [6] in real time. We tested different NWChem application settings with different molecular system sizes (small and large in the first columns of

Table 1) as well as the scalability in different HPC settings (the second columns of Table 1).

The details for the six datasets are shown in Table 1. We found that our feature extraction and anomaly detection algorithms did not cause significant delays for the streaming analysis. The throughput of our system was acceptable according to our user.

NWChem Setting	Number of HPC cores	Number of trace events	Number of trace events per second	Number of anomalies	Throughput (MBps)
small	2	393,542	58.5k	350	5.2
small	4	1,201,533	128.1k	445	11.5
small	8	4,024,651	224.2k	2235	20.5
large	2	784,122	52.9k	818	4.7
large	4	2,386,634	101.0k	1121	9.1
large	8	7,972,872	172.9k	3683	15.8

Table 1. Summary of experiment datasets and the throughput.

6.2 Case 1: Delay of Communication

During the online analysis, our participant first examined the overview distribution of the CSTrees. He was interested in the `SP_BALANC` function, which was designed to redistribute the work over the processors to minimize the time spend waiting in communication functions. He suggested that it was a critical factor to the overall performance of the code. In the scatter plot, he noticed that most of the points in the projection were normal. To put more focus on the potential abnormal CSTrees, he reduced the negative sampling rate (Fig. 4 (a)). He noticed an abnormal point, which was execution #31517 of the Thread #0 in Rank #0 of the program #0, as shown in the upper left of Fig. 5 (a).

From the Structure Visualization, he found that there was a big green function of `comex_barrier` which spent the execution majority of the time in `SP_BALANC`. `comex_barrier` was a function responsible for the communication between different HPC cores. He made an assumption that the barrier function was the major reason for the latency of this execution.

To learn about the temporal pattern and message communication, our participant examined the Timeline Visualization. He found that barrier function invoked some communications, as shown in the message passing visualizations with Rank #1 in Fig. 4 (c). After comparing with other regular executions of `SP_BALANC`, he concluded that this execution waited for a long time for the response of Rank #1. As a result, the communication delay to other computing cores made this execution a candidate anomaly. He concluded that our system helped him understand the one of the reasons for the performance fluctuation of `SP_BALANC`, which provided insights of how to optimize the source code to improve the overall program performance.

6.3 Case 2: Delay of Computation

Our participant continued to explore other functions of interest. In the scatter plot, he only visualized `CF_CENMAS` function, which computed the center of mass coordinates of individual molecules in the NWChem simulation. He located another candidate anomaly and showed its detailed structure and timeline. From the structure visualization (Fig. 6 (a)) he learned that it was abnormal since the root node in the CSTree was very huge. He zoomed into the timeline (Fig. 6 (b)) and found that the all of the children functions were executed as expected. However, they were invoked after waiting a long time for `CF_CENMAS`.

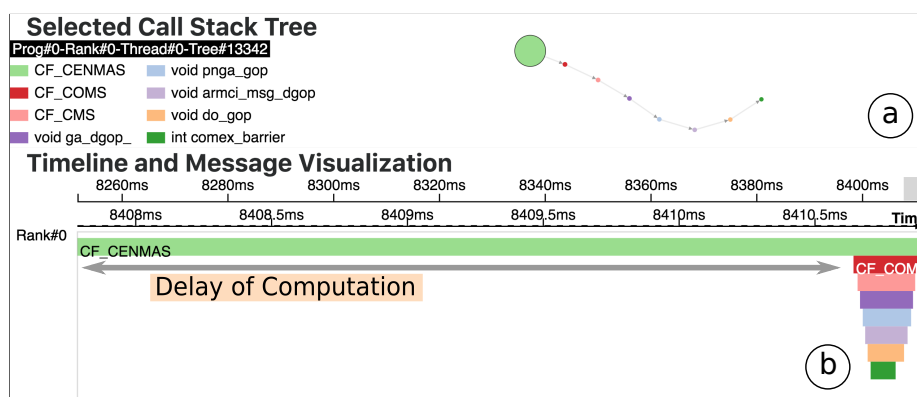


Fig. 6. (a) The CSTree structure of the `CF_CENMAS` function. (b) The user zoomed in the timeline to see the detailed temporal sequence at the beginning of the execution, as shown in the gray range in the time axis.

He expressed that this could happen when the computation in `CF_CENMAS` took a long time. The descendant functions (e.g., `comex_barrier`) had to wait for the computation to finish to communicate with other HPC cores. Our expert noted that the visual exploration our system provided was a critically important supplement to the learned labels (i.e., normal or abnormal) of the automatic algorithm. He stated that our system provided the much needed comprehensive analysis support for understanding and diagnosing the runtime behavior.

6.4 Feedback and Discussion

We conducted an interview session where we asked our participant to evaluate the learning cost and usability of our system. The participant mentioned that the visualizations were easy to understand since they are also commonly used by the existing performance analysis tools [7] [4]. He also compared our system with the commonly available tools in his community. He indicated that with current tools such as Jumpshot [24] he would only be able to learn about the executions after a

lengthy session with the system. He would first manually locate the time window of an anomaly candidate and then look at a detailed view on the respective call paths and messages. Conversely, he said that our interface was much more direct, enabling him to quickly discover the anomaly and succinctly explain it by using the three linked views.

7 Conclusion and Future Work

We described a visual analytics approach for the online detection of anomalous function executions in HPC clusters and their visual exploration. Our approach is based on the CSTree representation. It provides effective anomaly detection and visualization tools that address challenges with streaming performance evaluation for parallel computation at scale. We demonstrated our approach with a real world NWChem application.

In the case study, we learned that our bag-of-subtree vector can be too sparse since the subtree corpus will be massive. To cope with this problem, a Stack2vec embedding [23] can be a viable option to generate the embedding vectors from the sparse bag-of-subtree vector. We also plan to integrate machine log analysis and source code examination into our system which will provide more insights into the execution scheduling and code design optimization.

Acknowledgments

This research was partially supported by NSF grant IIS 1527200, BNL LDRD grant 16-041 and 18-009, ECP CODAR project 17-SC-20-SC, and the MSIP (Ministry of Science, ICT and Future Planning), Korea, under “IT Consilience Creative Program (ITCCP)” supervised by NIPA.

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