

CSE509

Computer System Security



2023-02-09

Software Vulnerabilities

Michalis Polychronakis

Stony Brook University

Software Vulnerabilities

Program flaws can turn into exploitable vulnerabilities

Securing user-space applications is equally critical as securing the OS

May run with superuser privileges: system daemons, setuid programs, anything launched by the root account, ...

Non-privileged applications may be a stepping stone to full system compromise → privilege escalation attacks

The OS is software too

Full system compromise may not even be needed (!)

User data is handled by applications → compromising an application may be enough

Browsers, password managers, messaging apps, ...

Types of Software Vulnerabilities

Vast number of different types of programming flaws, weaknesses, omissions, and other oversights

Many different corresponding exploitation techniques

Classifications according to different aspects include:

- Type of bug

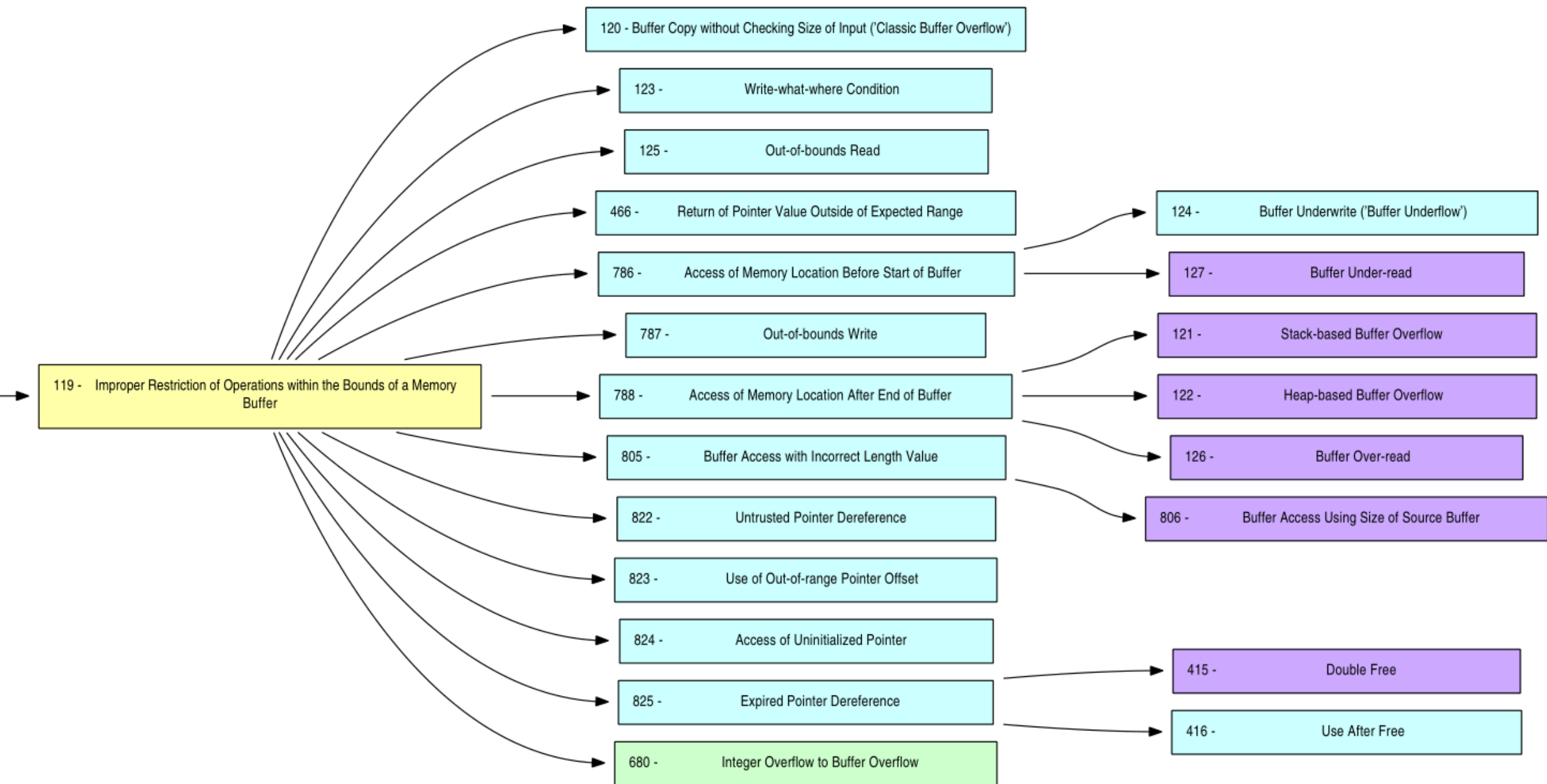
- Exploitation strategy

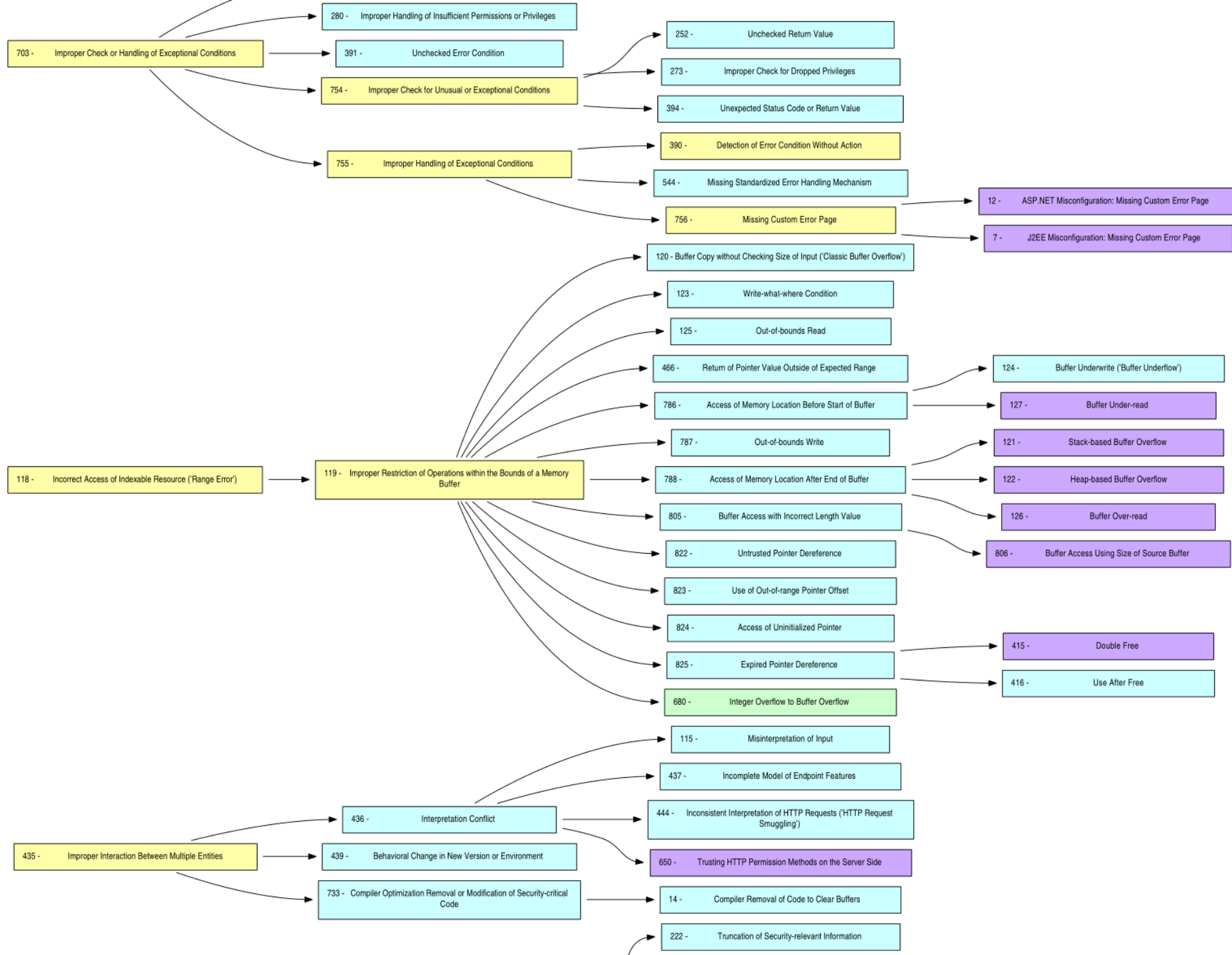
- SDL phase

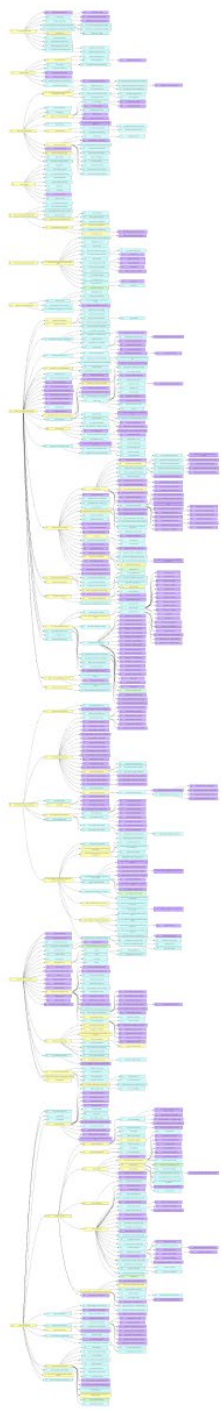
- Programming language

- System layer

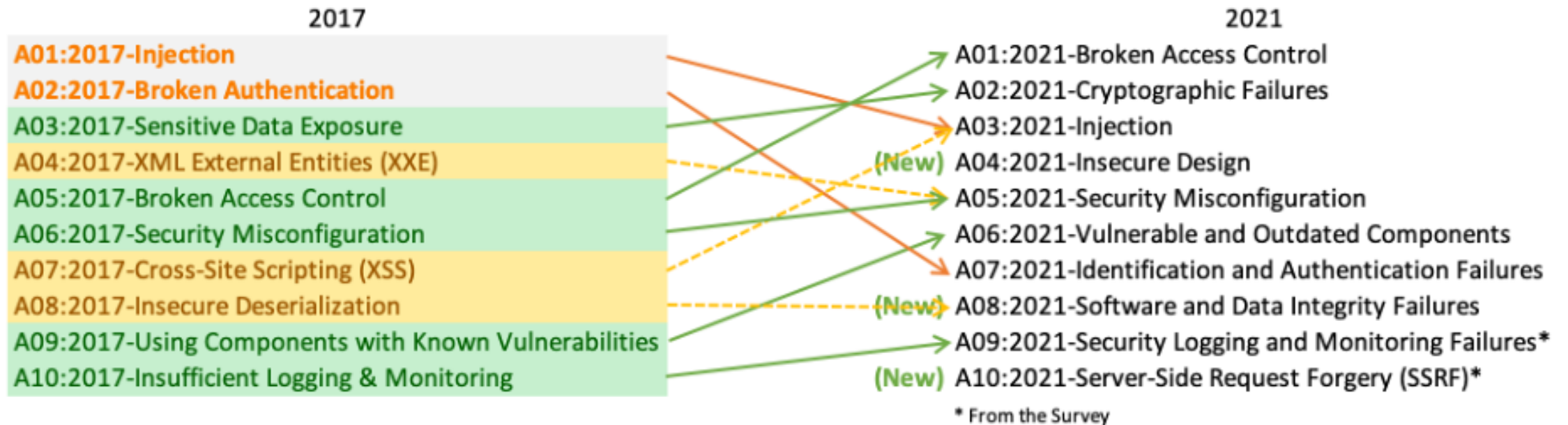
Example: MITRE's Common Weakness Enumeration (CWE) classification







Another example: OWASP Top 10 *“The ten most critical web application security risks”*



Some Basic Types of Software Vulnerabilities

Memory corruption: stack/heap buffer overflow, dangling pointers, ...

Arithmetic errors: arithmetic overflow, signedness, array indexing, ...

Race conditions: synchronization issues, TOCTTOU bugs, ...

Unvalidated input: format strings, SQL injection, command injection, ...

Confused deputy: CSRF, clickjacking, ...

Side channels: timing, power, temperature, ...

Program logic/design/protocol flaws: wrong/missing checks, typos, ...

Arithmetic Overflow

Finite number of bits to represent integers

Let's assume a 32-bit system

Integers are expressed in *two's complement* notation

Signed integers

Positive numbers: `0x00000000` – `0x7fffffff` 0 to $2^{31}-1$

Negative numbers: `0x80000000` – `0xffffffff` $-(2^{31})$ to -1

Unsigned integers

`0x00000000` – `0xffffffff` 0 to $2^{32}-1$

Both can *overflow* or *underflow*

“Only the first five clients can connect”

```
unsigned int connections = 0;
...
/* new connection attempt */
...
connections++;
if (connections < 5) {
    grant_access();
}
else {
    deny_access();
}
```

How can an attacker connect even if there are already five established connections?

“Only the first 5 clients can connect”

```
unsigned int connections = 0;
...
/* new connection attempt */
...
if (connections < 5) {  Upper bound of five connections is enforced
    connections++;
}
if (connections < 5) {
    grant_access();
}
else {
    deny_access();
}
```

Memory-related Errors

Very broad class of *undefined memory access* vulnerabilities

One of the most important and widely exploited types of vulnerabilities

Buffer overflow, null pointer dereference, use after free, uninitialized memory use, ...

In contrast to *memory safe* languages, C and C++ do not safeguard memory against illegal accesses

Under unexpected conditions, attackers may be able to *read* from or *write* to arbitrary memory locations

Lower-level languages → performance

Operating systems, core services, desktop applications, embedded systems, and many other programs are still written in C/C++

Promising recent progress on low-overhead memory-safe languages (e.g., Rust, Go)

Buffer Overflow

C and C++ do not provide any automatic bounds checking capability for allocated chunks of memory

- Arrays: can be indexed past the first (underflow) or last (overflow) item

- Pointers: can point outside the allocated object

Care must be taken when reading/writing user-supplied or user-derived data from/to memory

- More data than expected may be supplied → overflow

- The program should perform explicit bounds checks

An attacker can intentionally overflow the buffer and access out-of-bounds memory

- Modify critical control or program data (overwrite)

- Leak sensitive information (overread)

Simple overflow example: unbounded string copy

```
int main(int argc, char *argv[]) {  
    char buf[16];  
    strcpy(buf, argv[1]);  
    printf("%s\n", buf);  
    return 0;  
}
```

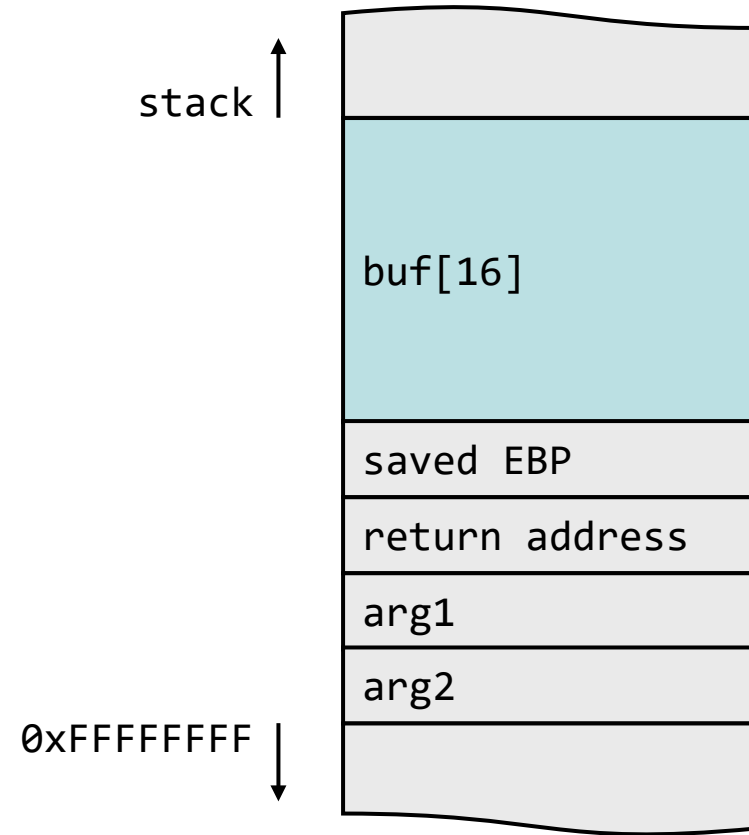
```
$ ./overflow AAAAAAAAAAAAAA
```

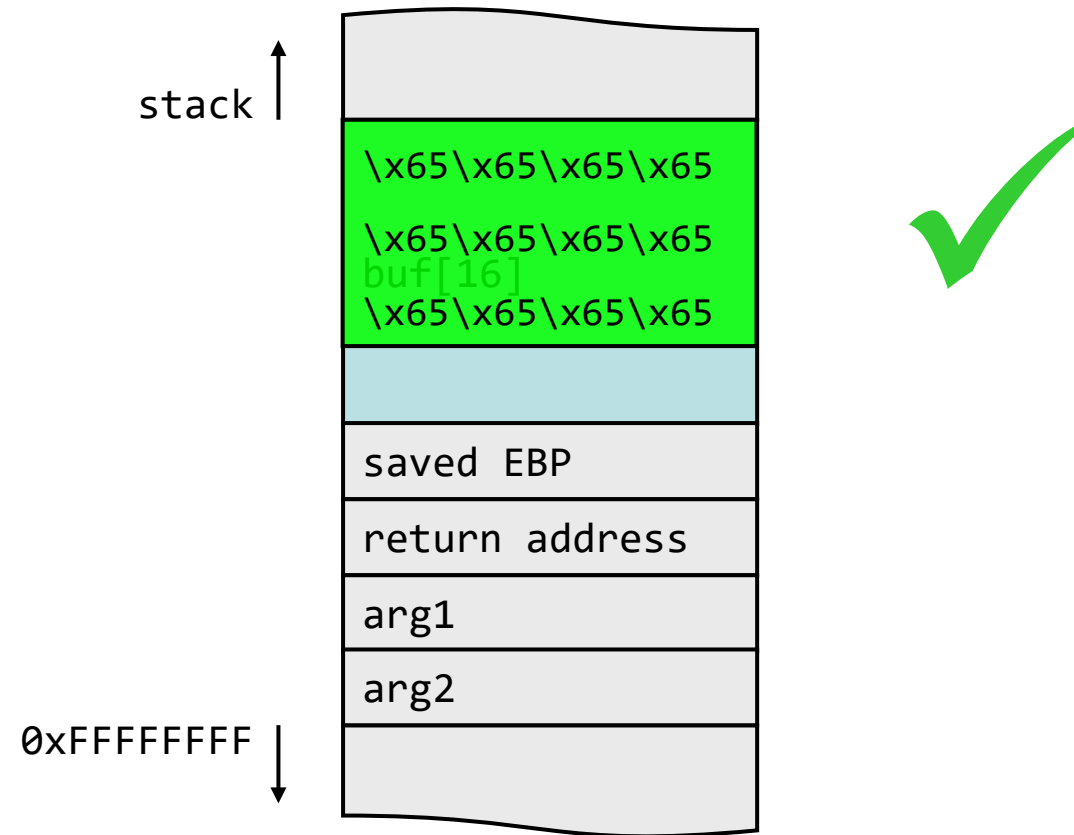
```
AAAAAAAAAAAAAAAAAAAA
```

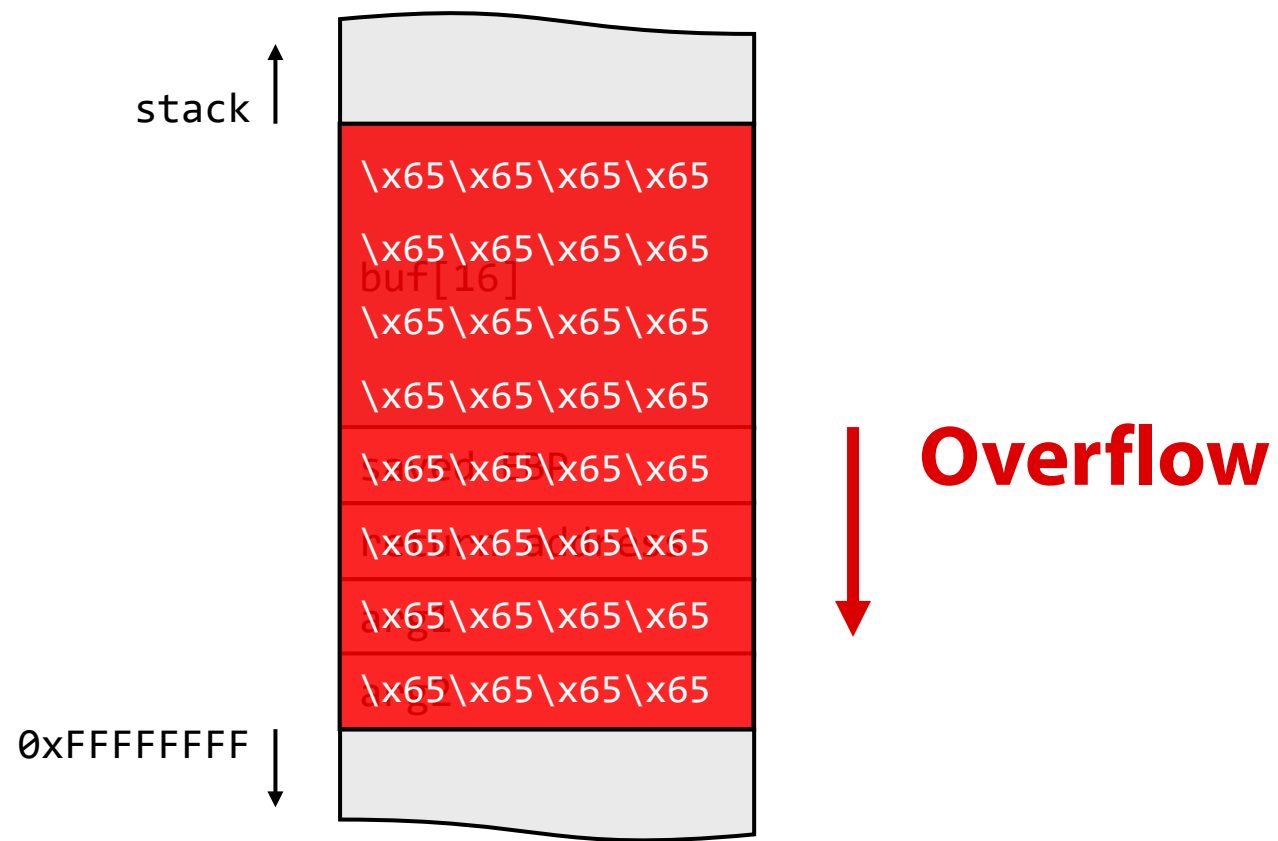
```
$ ./overflow AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
```

```
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
```

```
Segmentation fault (core dumped)
```







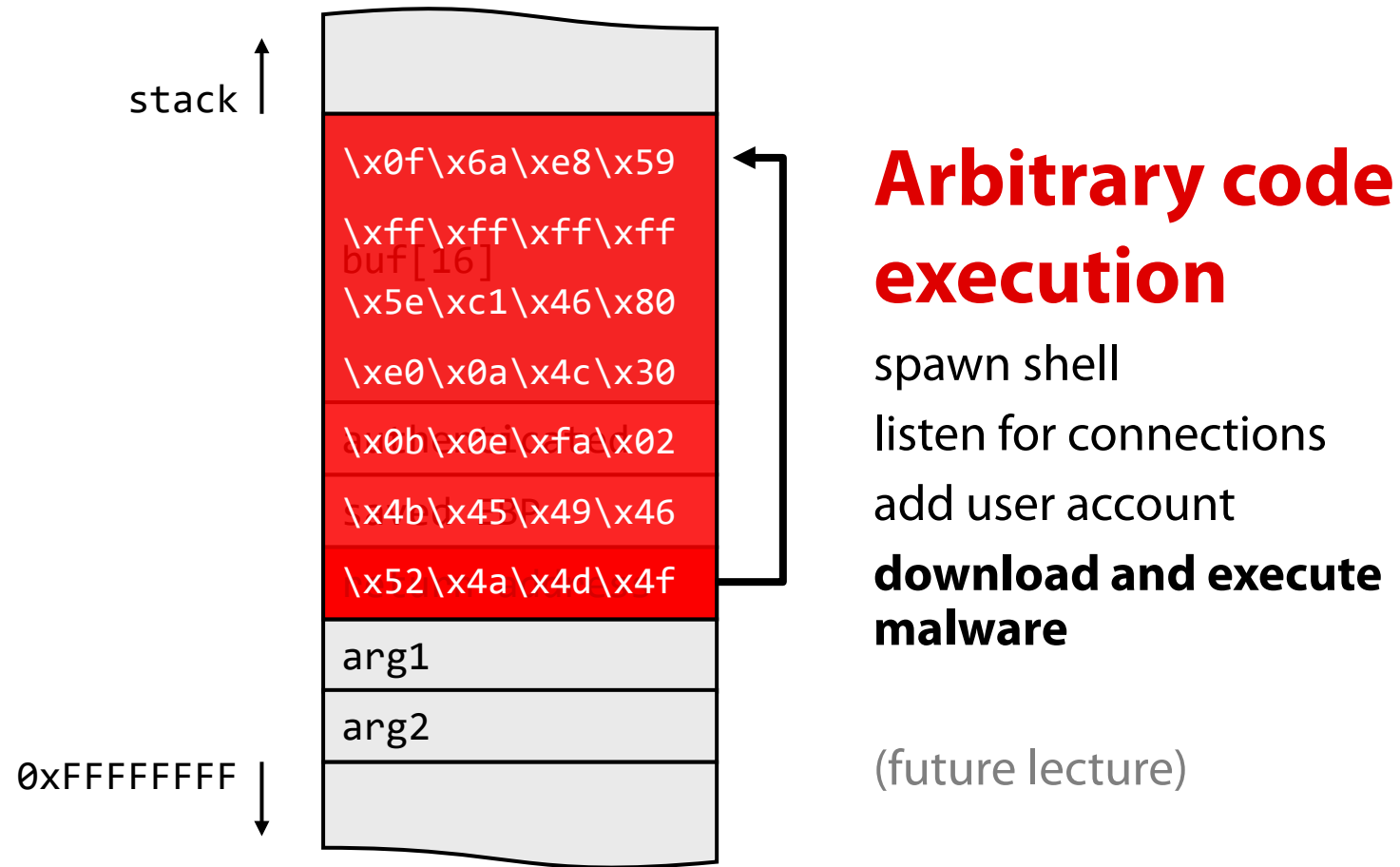
Safer way

```
#define BUF_SIZE 16
```

```
int main(int argc, char *argv[]) {  
    char buf[BUF_SIZE];  
    strncpy(buf, argv[1], BUF_SIZE);  
    printf("%s\n", buf);  
    return 0;  
}
```

← *Explicit bounds check:
strncpy will never write more
than BUF_SIZE bytes into buf*

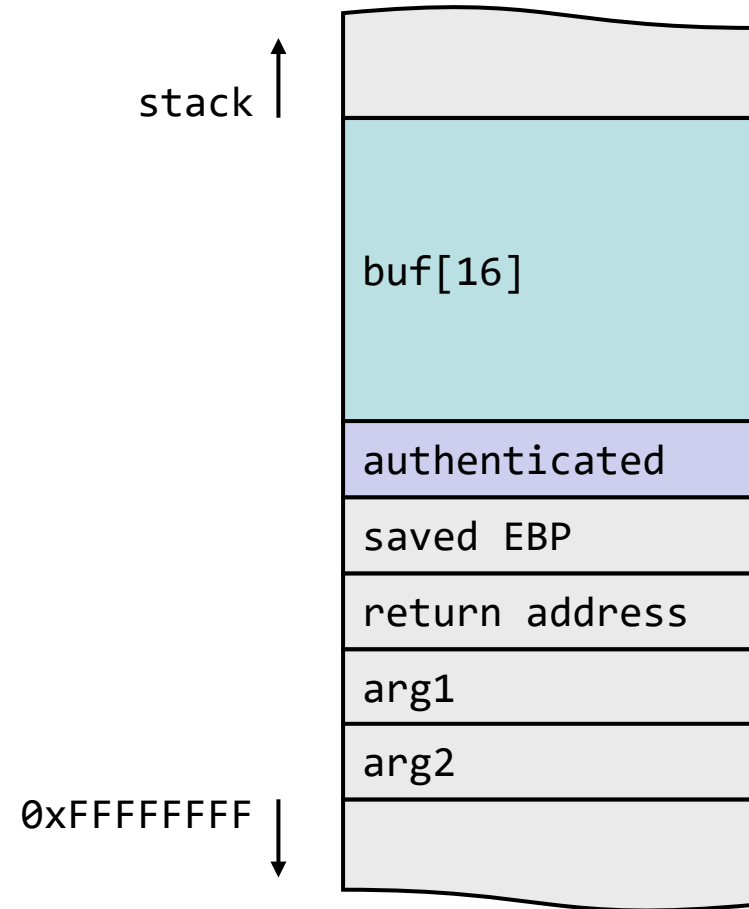
What can the attacker do? **Overwrite control data**

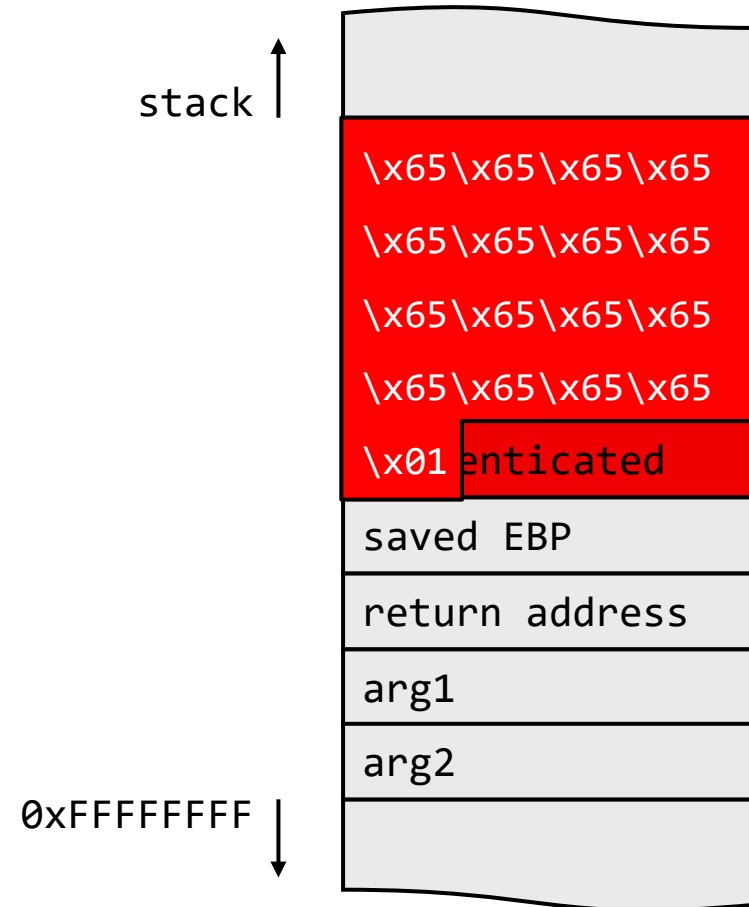


What can the attacker do? ***Overwrite program data***

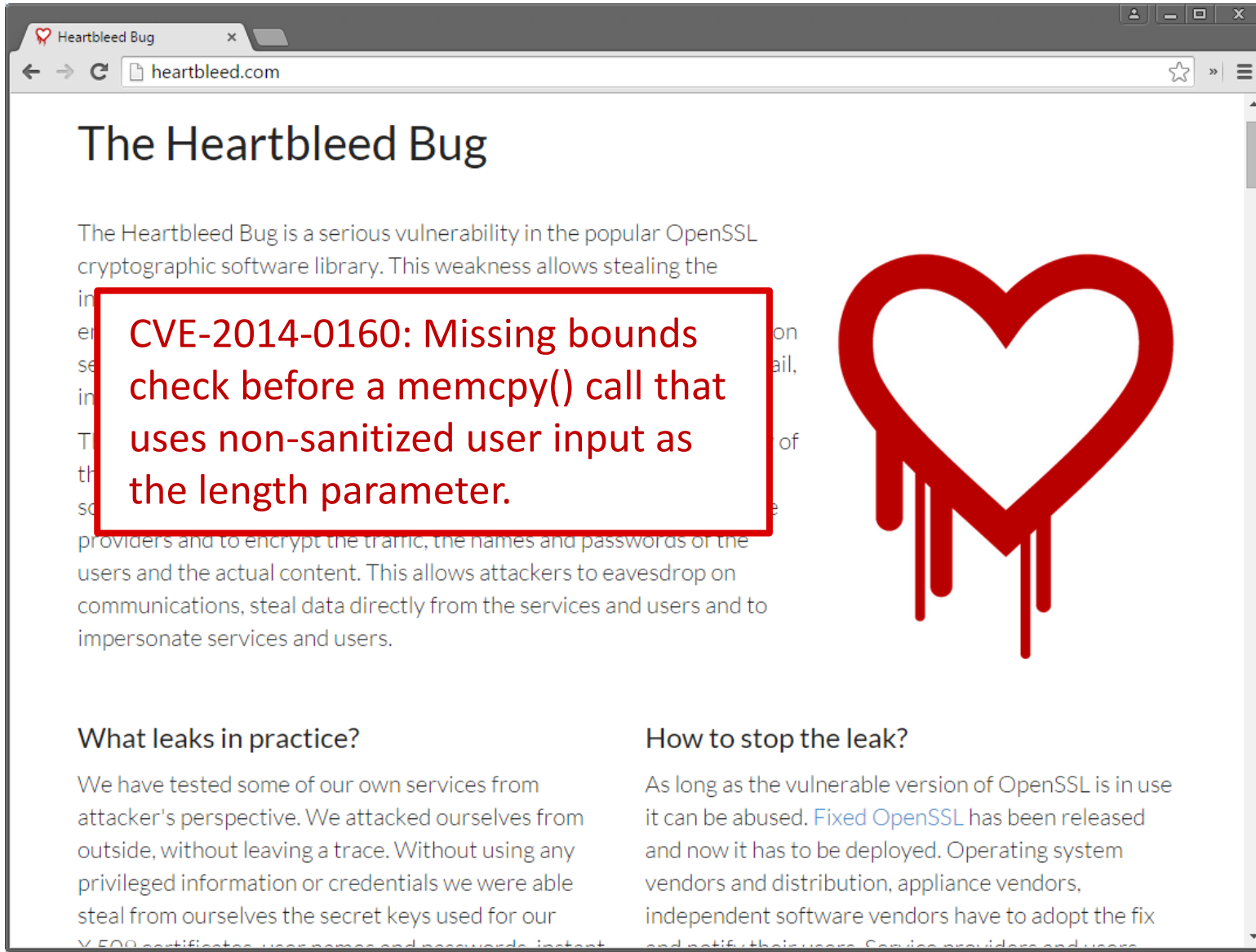
```
int main(int argc, char *argv[]) {
    int authenticated = 0;
    char password[16];
    gets(password);
    if (check_password(password) == TRUE) {
        authenticated = 1;
    }
    return authenticated;
}

$ ./authenticate AAAAAAAAAAAAAAAAAA && echo $?
0
$ ./authenticate AAAAAAAAAAAAAAAAAA && echo $?
65
$ ./authenticate `printf "AAAAAAAAAAAAAAAAAA\x01"` && echo $?
1
```





What can the attacker do? ***Leak data***



The screenshot shows a web browser window with the address bar displaying "heartbleed.com". The page title is "The Heartbleed Bug". The main text describes the vulnerability in OpenSSL. A red rectangular box highlights the following text: "CVE-2014-0160: Missing bounds check before a memcpy() call that uses non-sanitized user input as the length parameter." To the right of the text is a large red heart outline with red liquid dripping down from its base. Below the main text, there are two columns of text: "What leaks in practice?" and "How to stop the leak?".

The Heartbleed Bug

The Heartbleed Bug is a serious vulnerability in the popular OpenSSL cryptographic software library. This weakness allows stealing the information stored in recently exchanged TLS records on a server, including private keys, session cookies, passwords and other personally identifiable information. The weakness is in the heartbeats feature of TLS, which allows a client to request a small amount of data from a server, and the server to respond with the requested data. The vulnerability allows an attacker to request a large amount of data from the server, which will then respond with the requested data. This allows attackers to eavesdrop on communications, steal data directly from the services and users and to impersonate services and users.

CVE-2014-0160: Missing bounds check before a memcpy() call that uses non-sanitized user input as the length parameter.

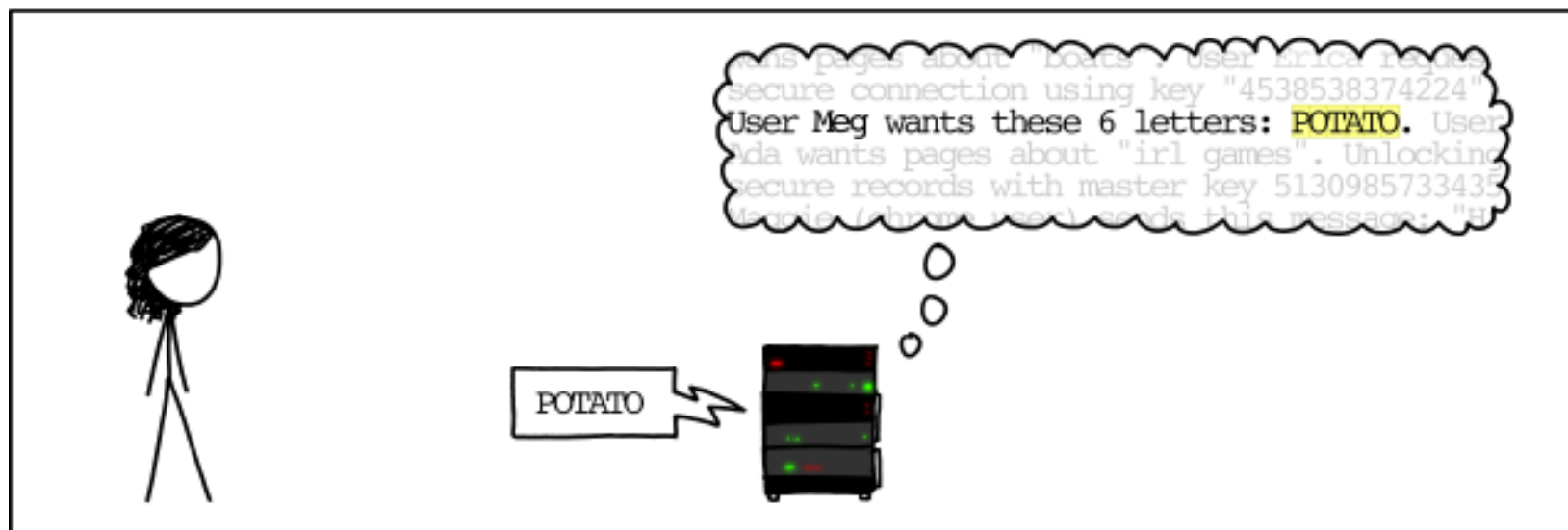
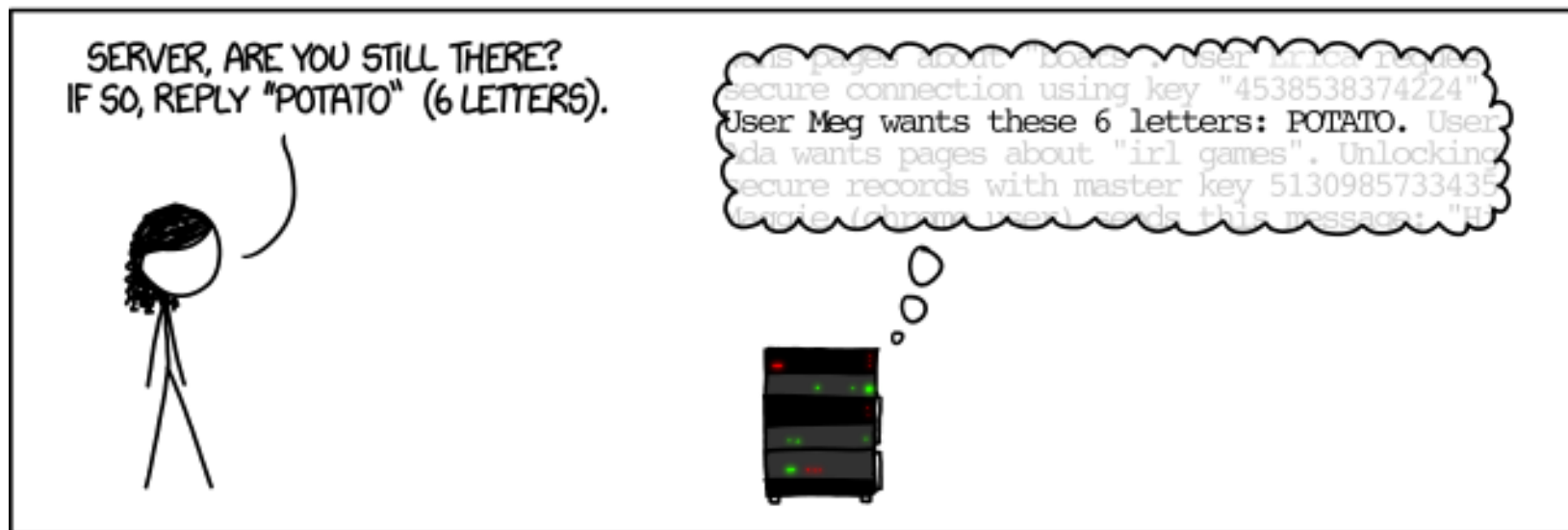
What leaks in practice?

We have tested some of our own services from attacker's perspective. We attacked ourselves from outside, without leaving a trace. Without using any privileged information or credentials we were able to steal from ourselves the secret keys used for our TLS certificates, user names and passwords, instant messages, session cookies, and other personally identifiable information.

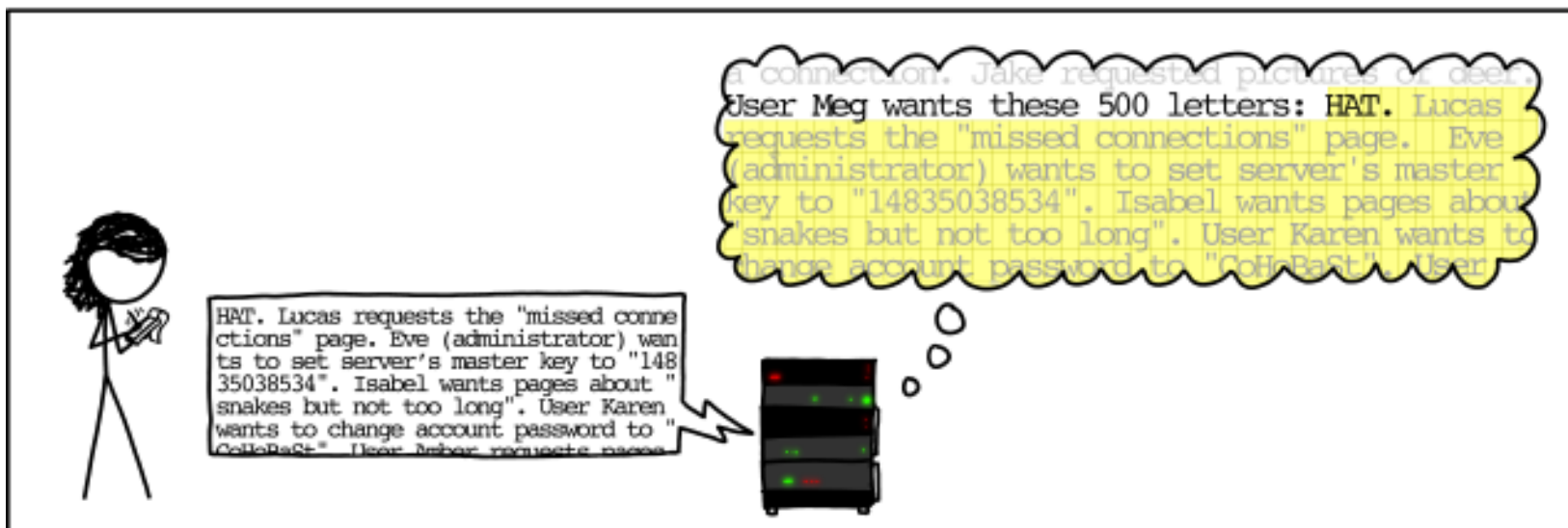
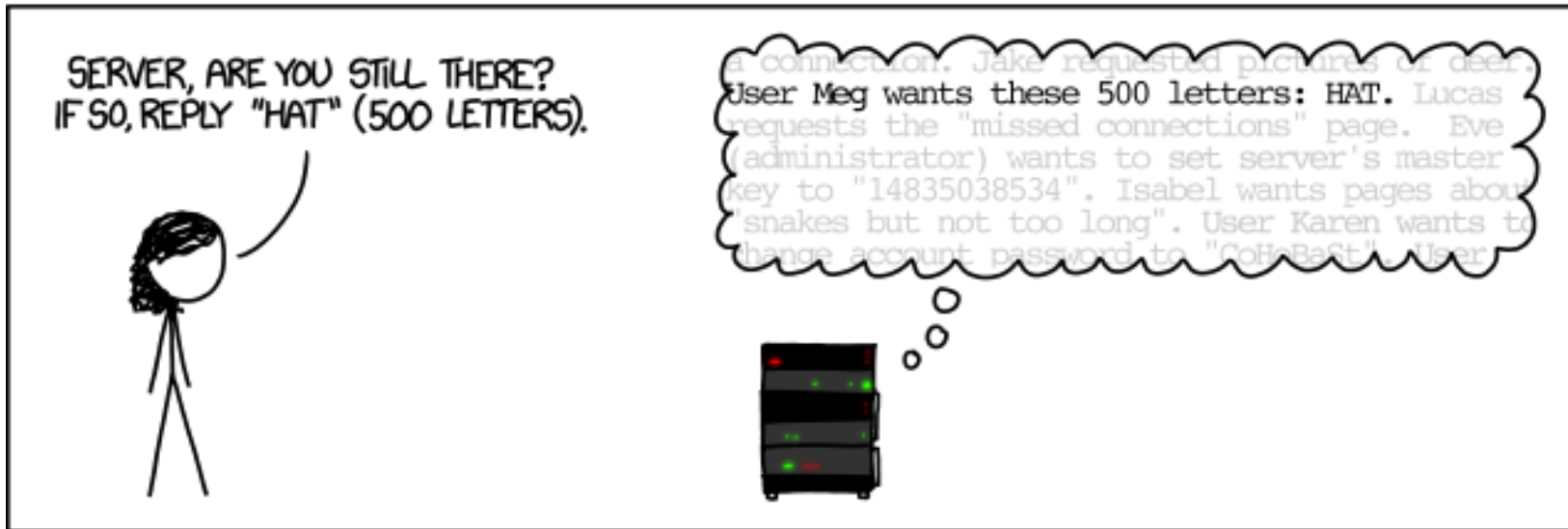
How to stop the leak?

As long as the vulnerable version of OpenSSL is in use it can be abused. [Fixed OpenSSL](#) has been released and now it has to be deployed. Operating system vendors and distribution, appliance vendors, independent software vendors have to adopt the fix and notify their users. Service providers and users should also take action to protect their data.

HOW THE HEARTBLEED BUG WORKS:







Pointer manipulation

```
void vulnerable(void *arg, size_t len) {  
    long val = 0;  
    long *ptr = NULL;  
    char buf[128];  
    ...  
    memcpy(buf, arg, len);  
    *ptr = val;   
    ...  
}
```

“Arbitrary write” capability: The attacker can write controlled data into a controlled location

Common Heap Errors

Initialization mistakes

`malloc()`, `realloc()` and C++ `new` do not initialize the allocated data.

Confusion between scalars and arrays

`new` and `delete` for scalars, `new[]` and `delete[]` for arrays

Failing to check return values

Functions like `malloc()` and `realloc()` return a NULL pointer when out of memory

Writing to already freed memory

Freeing the same memory multiple times

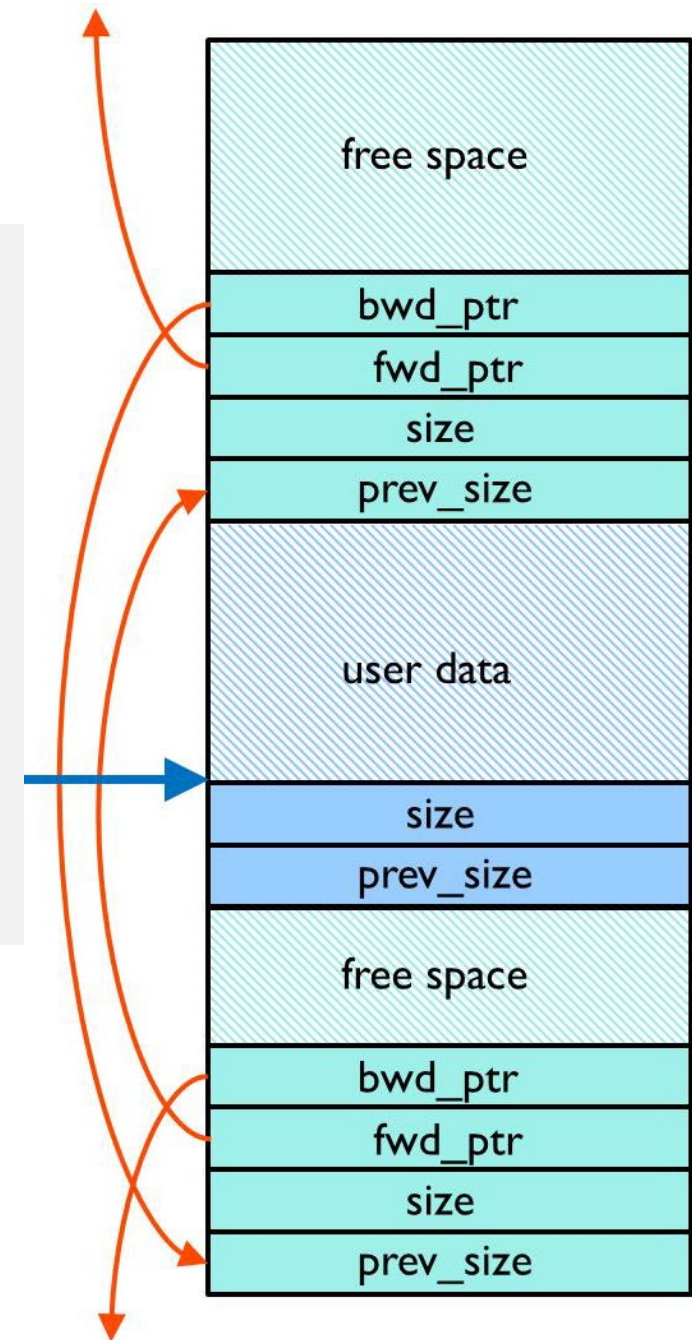
Writing outside of the allocated area

*Corruption of internal memory
allocator data structures or
allocated application data*

Heap-based Overflows

```
int main(int argc, char *argv[]) {  
    char *p, *q;  
    p = malloc(1024);  
    q = malloc(1024);  
    strcpy(p, argv[1]);  
    free(q);  
    free(p);  
    return 0;  
}
```

*Arbitrary write when free() is called
by carefully corrupting heap metadata*



Format String Vulnerabilities

The `printf()` family of functions accept a format string denoting how a variable will be displayed

`printf("%s", str)` → prints `str` variable as string

`printf("%d", num)` → prints `num` as a decimal value

`printf("%x", num)` → prints `num` as a hexadecimal value

Format strings can also *write* to memory

`printf("ABCD%n", &i)` → write the number of bytes output so far to the memory address of the second argument

What if...

The programmer does not supply a format string?

Fewer arguments are passed than the number of format string parameters?

Simple format string error example

```
int main(int argc, char *argv[]) {  
    printf("Input: ");  
    printf(argv[1]);  
    printf("\n");  
}
```

```
$ ./fmt test
```

```
Input: test
```

```
$ ./fmt "%08x %08x %08x %08x"
```

```
input: b773c080 0804846b b7721ff4 08048460
```

```
$ ./fmt $(printf "\x18\xa0\x04\x08")%x%x%x%x%n
```

Safer way

```
int main(int argc, char *argv[]) {  
    printf("Input: ");  
    printf("%s", argv[1]);  
    printf("\n");  
}
```

Other Memory-related Exploitable Conditions

NULL-termination errors

Dangling pointers

NULL pointer dereferences

String truncation

Single-byte overwrite

Off-by-one accesses

Double free

...

Race Conditions

Situations in which the behavior of the program depends on the timing of some event

Critical section

Opens up a window of opportunity for the attacker

Race conditions occur in many different contexts

Multi-threaded programs with different threads operating on the same data

Distributed applications that perform multi-step transactions

Time of check to time of use (TOCTTOU): changes may happen between *checking* a condition and *using* the results of the check

Filesystem race condition example

```
// setuid program
```

```
if (access("file01", W_OK) != 0) { // access() checks the real uid (not eid)
    exit(1);
}
```

← In /etc/passwd file01

```
fd = open("file01", O_WRONLY);
write(fd, buffer, sizeof(buffer)); // write() modifies /etc/passwd
```

iOS 8.1 Hardware-assisted Screenlock Brute-force

Successfully brute-force device PIN *even if “wipe out after 10 failed attempts” is enabled (!)*

Vulnerable code:

1. Display “incorrect pin” message

2. ++attempts;

 ***Power off the device***

Correct code:

1. ++attempts; // gets written to flash memory

2. Display “incorrect pin” message



Program Logic Flaws

Flaws in the design itself, or in the way the design was implemented in the program code

Rather than lower-level bugs (overflows etc.)

Indicative examples

Abuse of functionality (e.g., lock another user out with repeated failed attempts)

Workflow circumvention (e.g., skip payment info and jump straight to delivery info)

Information leakage (e.g., through hidden form fields)

Insufficient/poor validation (e.g., negative number of items)

Diverse and complicated → very difficult to find

GOTO FAIL

iOS 7.0.6 signature verification error

Legitimate-looking TLS certificates with mismatched private keys were unconditionally accepted...

```
if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
    goto fail;
if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
    goto fail;
    goto fail; ← ?!!?!?!?
if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
    goto fail;
...
fail:
    SSLFreeBuffer(&signedHashes);
    SSLFreeBuffer(&hashCtx);
    return err; ← Return with err == 0
```

Check never executed