

CSE508

Network Security



2021-03-18

Authentication

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Authentication

The process of verifying someone's identity or role

User, device, service, request, ...

What is identity?

Which characteristics uniquely identify an entity?

Authentication is a critical service

Enables communicating parties to verify the identity of their peers

Many other security mechanisms rely on it

Two main types

Human to computer

Computer to computer

SOLAR LOGIN

Stony Brook ID#

Password

Sign In

✔ This system is online.

[SOLAR Account & Password Help](#)

Credentials

Evidence used to prove an identity

User Authentication: credentials supplied by a person

Something you know

Something you have

Something you are

Computer authentication: crypto, location

Computers (in contrast to humans) can “remember” large secrets (keys) and perform complex cryptographic operations

Location: evidence that an entity is at a specific place (IP, subnet, switch port, ...)

Authentication can be delegated

The verifying entity accepts that a trusted third party has already established authentication

Something You Know: Password-based Authentication

Passwords, passphrases, pins, key-phrases, access codes, ...

Good passwords are easy to remember and hard to guess

Easy to remember → easy to guess

Hard to guess → hard to remember

Bad ideas: date of birth, SSN, zip code, favorite team name, ...

Password space (bits) depends on:

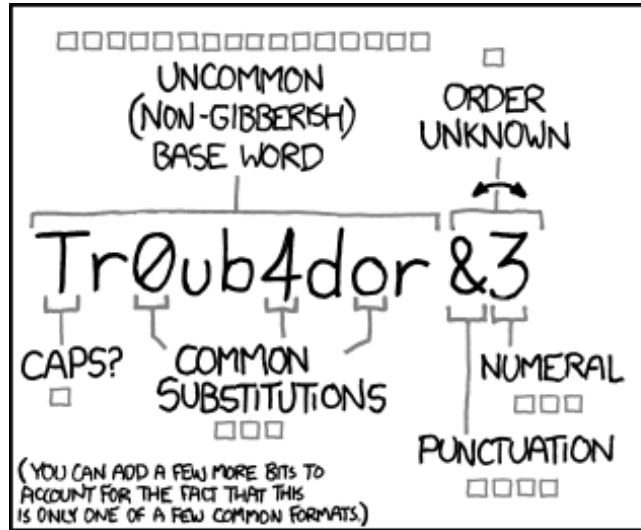
Password length

Character set

Better way to think about strong passwords

Long passphrases

Can be combined with custom variations, symbols, numbers, capitalization, ...



~28 BITS OF ENTROPY

$2^{28} = 3 \text{ DAYS AT } 1000 \text{ GUESSES/SEC}$

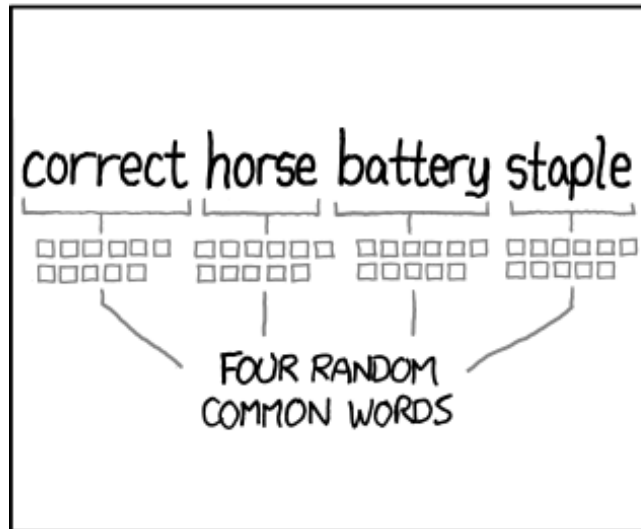
(PLAUSIBLE ATTACK ON A WEAK REMOTE WEB SERVICE. YES, CRACKING A STOLEN HASH IS FASTER, BUT IT'S NOT WHAT THE AVERAGE USER SHOULD WORRY ABOUT.)

DIFFICULTY TO GUESS: **EASY**

WAS IT TROMBONE? NO, TROUBADOR. AND ONE OF THE 0s WAS A ZERO?

AND THERE WAS SOME SYMBOL...

DIFFICULTY TO REMEMBER: **HARD**



~44 BITS OF ENTROPY

$2^{44} = 550 \text{ YEARS AT } 1000 \text{ GUESSES/SEC}$

DIFFICULTY TO GUESS: **HARD**

THAT'S A BATTERY STAPLE.

CORRECT!

DIFFICULTY TO REMEMBER: YOU'VE ALREADY MEMORIZED IT

THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.

Password Policies (often have the opposite effect)

Password rules (often miss the point)

“At least one special character,” “Minimum/Maximum length of 8/12 characters,” “Must contain at least one number,” “Must contain at least one capital letter”

Make passwords hard to remember! → encourage password reuse

Better: encourage long passphrases, evaluate strength on-the-fly

Periodic password changing

“You haven’t changed your password in the last 90 days”

Probably too late anyway if password has already been stolen

Makes remembering passwords harder → more password resets

Hinders the use of password managers (!)

What users do: password1 → password2 → password3 → ...

Attacking Passwords

Offline cracking }
Online guessing } *Brute force attacks*

Eavesdropping

Capturing

Password Storage

Storing passwords as plaintext is disastrous

Better way: store a cryptographic hash of the password

Even better: store the hash of a “salted” version of the password

Defend against *dictionary attacks*: prevent precomputation of hash values (wordlists of popular passwords, rainbow tables, ...)

Even if two users happen to have the same password, their hash values will be different → need to be cracked separately

Salting *does not* make brute-force guessing a given password harder!

Username	Salt	Password hash
Bobbie	4238	h(4238, \$uperman)
Tony	2918	h(2918, 63%TaeFF)
Mitsos	6902	h(6902, zour1da)
Mark	1694	h(1694, Rockybrook#1)

Password databases are still getting leaked...

Password Cracking

Exhaustive search → infeasible for large password spaces

Dictionary attacks (words, real user passwords from previous leaks, ...)

Variations, common patterns, structure rules

Prepend/append symbols/numbers/dates, weird capitalization, l33tspeak, visually similar characters, intended misspellings, ...

Target-specific information

DOB, family names, favorite team, pets, hobbies, anniversaries, language, slang, ...

Easy to acquire from social networking services and other public sites


Particularly effective against “security questions”

Advanced techniques

Probabilistic context-free grammars, Markov models, ...

example_hashes [hashcat] x

Secure | https://hashcat.net/wiki/doku.php?id=example_hashes



hashcat
advanced password recovery

hashcat Forums Wiki Tools Events

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Example hashes

If you get a "line length exception" error in hashcat, it is often because the hash mode that you have requested does not match the hash. To verify, you can test your commands against example hashes.

Unless otherwise noted, the password for all example hashes is **hashcat**.

Table of Contents

- Example hashes
- Generic hash types
- Specific hash types
- Legacy hash types

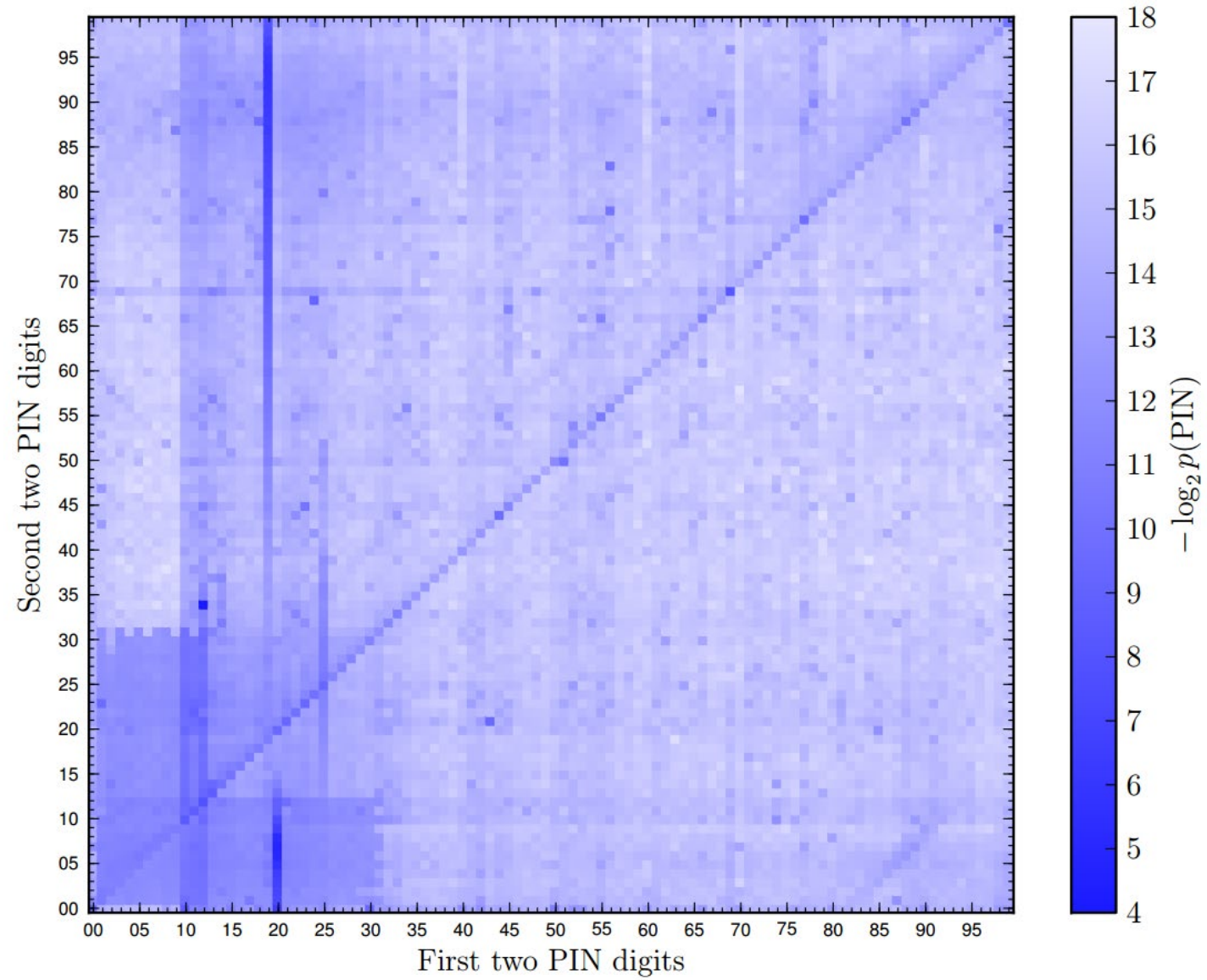
Generic hash types

Hash-Mode	Hash-Name	Example
0	MD5	8743b52063cd84097a65d1633f5c74f5
10	md5(\$pass.\$salt)	01dfae6e5d4d90d9892622325959afbe:7050461
20	md5(\$salt.\$pass)	f0fda58630310a6dd91a7d8f0a4ceda2:4225637426
30	md5(utf16le(\$pass).\$salt)	b31d032cfdcf47a399990a71e43c5d2a:144816
40	md5(\$salt.utf16le(\$pass))	d63d0e21fdc05f618d55ef306c54af82:13288442151473
50	HMAC-MD5 (key = \$pass)	fc741db0a2968c39d9c2a5cc75b05370:1234
60	HMAC-MD5 (key = \$salt)	bfd280436f45fa38eaacac3b00518f29:1234
100	SHA1	b89eaac7e61417341b710b727768294d0e6a277b
110	sha1(\$pass.\$salt)	2fc5a684737ce1bf7b3b239df432416e0dd07357:2014
120	sha1(\$salt.\$pass)	cac35ec206d868b7d7cb0b55f31d9425b075082b:5363620024
130	sha1(utf16le(\$pass).\$salt)	c57f6ac1b71f45a07dbd91a59fa47c23abcd87c2:631225
140	sha1(\$salt.utf16le(\$pass))	5db61e4cd8776c7969cfd62456da639a4c87683a:8763434884872
150	HMAC-SHA1 (key = \$pass)	c898896f3f70f61bc3fb19bef222aa860e5ea717:1234
160	HMAC-SHA1 (key = \$salt)	d89c92b4400b15c39e462a8caa939ab40c3aeaaa:1234
200	MySQL323	7196759210defdc0
300	MySQL4.1/MySQL5	fcf7c1b8749cf99d88e5f34271d636178fb5d130

50 Most-used (Worse) Passwords

123456	1234567	123	ashley	evite
123456789	qwerty	omgpop	987654321	123abc
picture1	abc123	123321	unknown	123qwe
password	Million2	654321	zxcvbnm	sunshine
12345678	000000	qwertyuiop	112233	121212
111111	1234	qwer123456	chatbooks	dragon
123123	iloveyou	123456a	20100728	1q2w3e4r
12345	aaron431	a123456	123123123	5201314
1234567890	password1	666666	princess	159753
senha	qqww1122	asdfghjkl	jacket025	0123456789

Distribution of 4-digit sequences within RockYou passwords



Wordlists

ce#ebc.dk	4637324	gea8mw4yz	fujinshan	masich	gothpunksk8er	20081010
goddess5	bugger825	kukumbike	counter	pengaiwei	rftaeo48	leelou44
20071002	marmaris	260888	N8mr0n	coalesce	8d7R0K	8UfjeGb0
271075711	jinjin111	jordi10	520057	56402768	5172032	200358808
zs3cu7za	170383gp	lexusis	adc123	thesis	aics07	dellede
scoopn	3484427	kj011a039	bmaster	aabbcc894	34mariah	liang123.
frygas1411	fl33321	c84bwlr	qbjh04zg	marion&maxime	dongqinwei	captainettekt
SL123456s1	zwqrfg	priyanka05	ueldaa79	614850	samarica	kwiki-mart
12345687ee123	67070857	loveneverdies	EMANUELLI	ydz220105	cap1014	mdovydas
xuexi2010	432106969	u8Aqobj576	yanjing	584521584521	0167387943	tigmys2001
daigoro	6856	FGYfgy77	assynt	txudecp	AE86Trueno	denial
12345614	704870704870	659397	62157173	84410545	19700913	678ad5251
DICK4080	pv041886	327296	0704224950753	pietro.chiara	mcsuap	woaiwuai
567891234	20060814	74748585	6903293	jman1514	bu56mpbu	1591591591212
tilg80	512881535	19720919	axaaxa	heryarma	danbee	hNbDGN
6z08c861	milanimilani	050769585	hilall	39joinmam	passw<>	cardcap
:zark:	472619	nicopa	30091983	timelapse	money521	13985039393
ravishsneha	dbyxw888	2232566	2510618981	mwinkar	conan83	001104
150571611369	85717221	bearss	soukuokpan	251422	nxfjpl	desare11
661189	cc841215	n0tpublic	tosecondlife	willrock	rateg143	412724198
passme	ariana19321	isitreal00	p4os8m6q	YHrtfgDK	kojyihen	nibh1kab
trolovinasveta	bbbnnn	ashraf19760	015614117	xys96exq	058336257	asferg
abdukhaleque	ang34hehiu	48144	acw71790	mercadotecnia	sarah4444	hqb555
007816	wj112358	22471015	lsyljm2	8s5sBEx7	7363437	xgames7
xLDSX	Brenda85	antyzhou115	2xgialdl	0125040344	freindship	muckerlee
Florida2011	786525pb	0167005246	gaybar9	margitka	JytmvW0848	choqui67
037037	shi461988	ec13kag	88203009	omaopa	sb inbau	12130911
WestC0untry	pingu	226226226226	MKltyh87	dfTi6nh	30907891	lierwei120
hitsugaiya	yeybozip	6767537/33	quiggle	1314520521	0515043111	skytdvn
955998126	71477nak	mimilebrock	2063775206	pixma760	1973@ati	milena1995
3n3rmax	stokurew	gueis8850	fr3iH3it	pearpear	wlxgjf	kambala11

LEAKED LISTS

Complete left lists from public leaks

ID	Name	Last Update	Num of Hashes	Progress	Left Hashes	Found
6505	H4v3 1 b33n pwn3d (SHA1)	02.10.2017 - 02:03:24	320'294'464	319'837'535 (99.86%)	Get	Get
5638	P4y4sUGym (MD5)	02.10.2017 - 02:04:19	241'266	221'152 (91.66%)	Get	Get
4920	L1nk3d1n (SHA1)	02.10.2017 - 03:24:58	61'829'262	60'147'825 (97.28%)	Get	Get
3282	4mzr3v13w7r4d3r.c0m (MYSQL5)	02.10.2017 - 03:25:32	41'823	39'166 (93.65%)	Get	Get
3186	X5pl17 (SHA1)	02.10.2017 - 03:32:38	2'227'254	2'162'101 (97.07%)	Get	Get
2499	Hashkiller 32-hex left total	02.10.2017 - 11:48:14	9'976'651	1'723'709 (17.28%)	Get	Get
2498	Hashkiller 40-hex left total	02.10.2017 - 13:22:34	1'739'204	350'788 (20.17%)	Get	Get
1619	4m4t3urc0mmuni7y.c0m	02.10.2017 - 13:33:26	197'302	57'407 (29.1%)	Get	Get
1535	b73r.c0m (MD5)	02.10.2017 - 13:34:43	63'070	32'543 (51.6%)	Get	Get
1427	4v17r0n.fr	02.10.2017 - 13:34:43	2'405	2'334 (97.05%)	Get	Get
1366	v0d4f0n3 (MD5(\$pass."s+(_a*)")	02.10.2017 - 13:34:44	322	307 (95.34%)	Get	Get
1314	1141407_5_07 (MD5)	02.10.2017 - 13:34:44	176	99'459 (57%)	Get	Get

518

pwned websites

10,624,652,379

pwned accounts








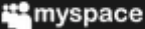


113,998

pastes









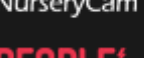
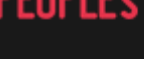
199,730,234

paste accounts

Largest breaches

	772,904,991	Collection #1 accounts
	763,117,241	Verifications.io accounts
	711,477,622	Onliner Spambot accounts
	622,161,052	Data Enrichment Exposure From PDL Customer accounts
	593,427,119	Exploit.In accounts
	457,962,538	Anti Public Combo List accounts
	393,430,309	River City Media Spam List accounts
	359,420,698	MySpace accounts
	268,765,495	Wattpad accounts
	234,842,089	NetEase accounts

Recently added breaches

	11,788	WeLeakInfo accounts
	465,141	Liker accounts
	637,279	Travel Oklahoma accounts
	66,521	Gab accounts
	1,834,006	Oxfam accounts
	1,921,722	Ticketcounter accounts
	20,339,937	SuperVPN & GeckoVPN accounts
	645,786	Filmai.in accounts
	10,585	NurseryCam accounts
	358,822	People's Energy accounts

Password Hashing Functions

Hash functions are very fast to evaluate → facilitate fast password cracking

Solution: slow down the guessing process (password “stretching”)

Benefit: cracking becomes very inefficient (e.g., 10-100ms per check)

Drawback: increased cost for the server if it must authenticate many users

Make heavy use of available resources

Fast enough computation to validate honest users, but render password guessing infeasible

Adaptable: flexible cost (time/memory complexity) parameters

Bcrypt [Provos and Mazières, 1999]

Cost-parameterized, modified version of the Blowfish encryption algorithm

Tunable cost parameter (exponential number of loop iterations)

Alternatives: Scrypt (memory-hard), PBKDF2 (PKCS standard)

Online Guessing

Similar strategy to offline guessing, but rate-limited

Connect, try a few passwords, get disconnected, repeat...

Prerequisite: know a valid user name

Credential stuffing: try username + password combinations from previous breaches

Many failed attempts can lead to a system reaction

Introduce delay before accepting future attempts (exponential backoff)

Shut off completely (e.g., ATM capturing/disabling the card after 3 tries)

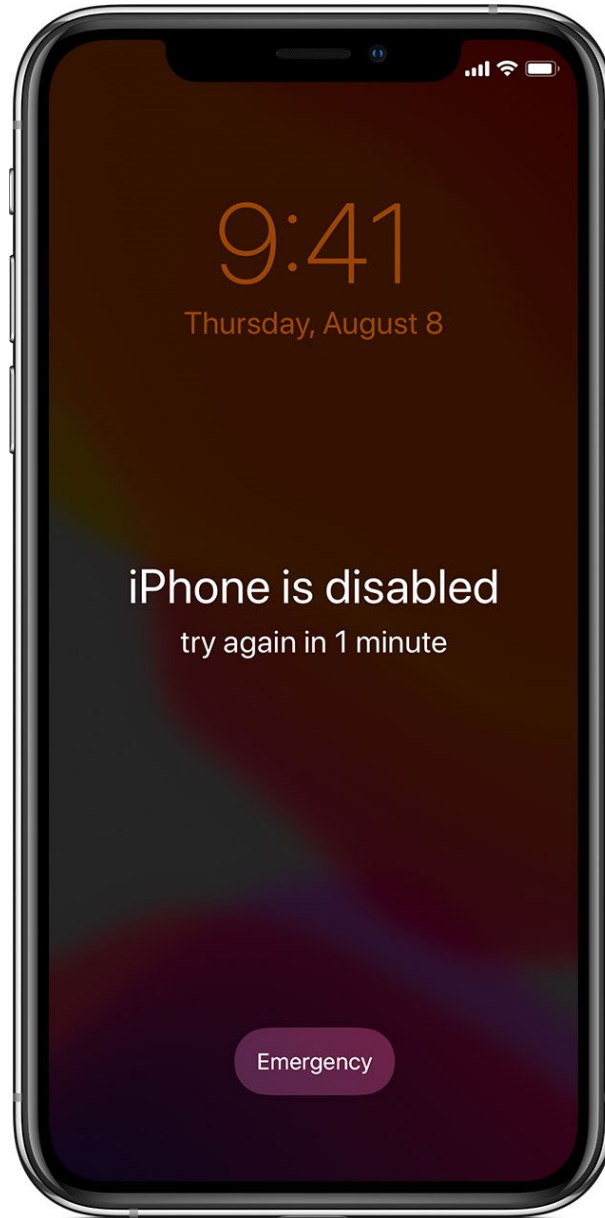
Ask user to solve a CAPTCHA

Very common against publicly accessible SSH, VPN, RDP, and other servers

Main reason people move sshd to a non-default port

Fail2Ban: block IP after many failed attempts → attackers may now be able to lock you out

Better: disable password authentication and use a key pair → cumbersome if having to log in from many/others' computers



LOGIN: mitch
PASSWORD: FooBar!-7
SUCCESSFUL LOGIN

(a)

LOGIN: carol
INVALID LOGIN NAME
LOGIN:

(b)

LOGIN: carol
PASSWORD: Idunno
INVALID LOGIN
LOGIN:

(c)

(a) Successful login

(b) Login rejected after name is entered

(c) Login rejected after name and password are typed → less information makes guessing harder

Default Router Passwords x

www.routerpasswords.com

Home | Add Password | About

RouterPasswords.com

Welcome to the internet's largest and most updated default router passwords database,

Select Router Manufacturer:

CISCO

Try the default first...

Find Password

Manufacturer	Model	Protocol	Username	Password
CISCO	CACHE ENGINE	CONSOLE	admin	diamond
CISCO	CONFIGMAKER		cmaker	cmaker
CISCO	CNR Rev. ALL	CNR GUI	admin	changeme
CISCO	NETRANGER/SECURE IDS	MULTI	netrangr	attack
CISCO	BBSM Rev. 5.0 AND 5.1	TELNET OR NAMED PIPES	bbsd-client	changeme2
CISCO	BBSD MSDE CLIENT Rev. 5.0 AND 5.1	TELNET OR NAMED PIPES	bbsd-client	NULL

Eavesdropping and Replay

Physical world

- Watch user type password (shoulder surfing)

- Cameras (e.g., ATM skimmers)

- Lift fingerprints (e.g., Apple Touch ID)

- Post-it notes, notebooks, ...

Network makes things easier

- Sniffing (LAN, WiFi, ...)

- Man-in-the-Middle attacks

Defenses

- Encryption

- One-time password schemes

Kerberos



Long-lived vs. session keys

Use long-lived key for authentication and negotiating session keys

Use "fresh," ephemeral session keys (prevent replay, cryptanalysis, old compromised keys) for encrypted communication, MACs, ...

Kerberos: most widely used (non-web) single sign-on system

Originally developed at MIT, now used in Unix, Windows, ...

Authenticate users to services: using their password as the initial key, without having to retype it for every interaction

A Key Distribution Center (KDC) acts as a trusted third party for key distribution

Online authentication: variant of Needham-Schroeder protocol

Assumes a non-trusted network: prevents eavesdropping

Assumes that the Kerberos server and user workstations are secure...

Use cases: workstation login, remote share access, printers, ...

Password Capture

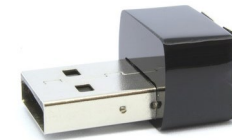
Hardware bugs/keyloggers

Software keyloggers/malware

Cameras

Phishing

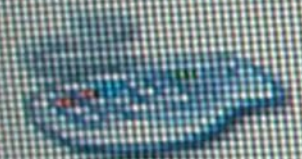
Social engineering





Microsoft
Windows
Professional

Copyright © 1985-2001
Microsoft Corporation



Press Ctrl-Alt-Delete to begin.

Requiring this key combination at startup helps keep
computer secure. For more information, click Help.



(a)



(b)

(a) Correct login screen

(b) Phony login screen

Something You Have: Authentication Tokens

One-time passcode tokens

Time-based or counter-based

Various other authentication tokens

Store certificates, encryption keys, challenge–response, ...

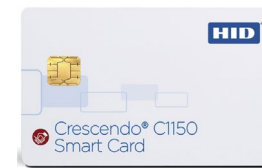
Smartcards (contact or contactless)

Identification, authentication, data storage, limited processing

Magnetic stripe cards, EMV (chip-n-pin credit cards), SIM cards, RFID tags, ...

USB/NFC tokens, mobile phones, watches, ...

Can be used as authentication devices



Multi-factor Authentication

Must provide several separate credentials of different types

Most common: *two-factor authentication (2FA)*

Example: Password + hardware token/SMS message/authenticator app, ...

Example: ATM card + PIN

Motivation: a captured/cracked password is not enough to compromise a victim's account → **not always true**

Man-in-the-Middle: set up fake banking website, relay password to real website, let the user deal with the second factor...

Man-in-the-Browser: hijack/manipulate an established session after authentication has completed (banking Trojans)

Dual infection: compromise both PC and mobile device

More importantly: the most commonly used 2nd factor (SMS) is the least secure

SMS Is Not a Secure 2nd Factor

(but still better than no 2nd factor)

Social engineering

Call victim's mobile operator and hijack the phone number

SIM swap, message/call forwarding, ...

Message interception

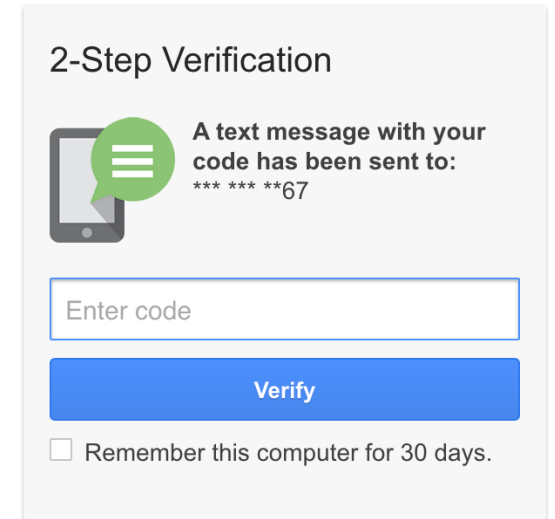
Rogue cell towers: IMSI catchers, StingRays,...

Some phones even display text messages on the lock screen (!)

SS7 attacks

The protocol used for inter-provider signaling is severely outdated and vulnerable

Allows attackers to spoof change requests to users' phone numbers and intercept calls or text messages



home

Scams

'Sim swap' gives fraudsters access-all-areas via your mobile phone

There's a new, little-known scam designed to empty your bank account, as one Vodafone customer found to her cost



1908 15

Anna Tims

Saturday 26 September 2015 02.00 EDT



Most popular in US



Las Vegas shooting: death toll rises to 58 as police name suspect - latest updates



Confusion follows reports of Tom Petty death after heart attack



Las Vegas gunman may have used special device to fire faster, expert says



A Hacker Got All My Texts for \$16

A gaping flaw in SMS lets hackers take over phone numbers in minutes by simply paying a company to reroute text messages.



By [Joseph Cox](#)

March 15, 2021, 1:10pm



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I hadn't been SIM swapped, where hackers trick or bribe telecom employees to port a target's phone number to their own SIM card. Instead, the hacker used a service by a company called Sakari, which helps businesses do SMS marketing and mass messaging, to reroute my messages to him. This

overlooked attack vector shows not only how unregulated commercial SMS tools are but also how there are gaping holes in our telecommunications infrastructure, with a hacker sometimes just having to pinky swear they

Better Alternative: Authenticator App

Six/eight digit code provided after successful password validation

Time-based one-time password (TOTP)

Code computed from a shared secret key and the current time (using HMAC)

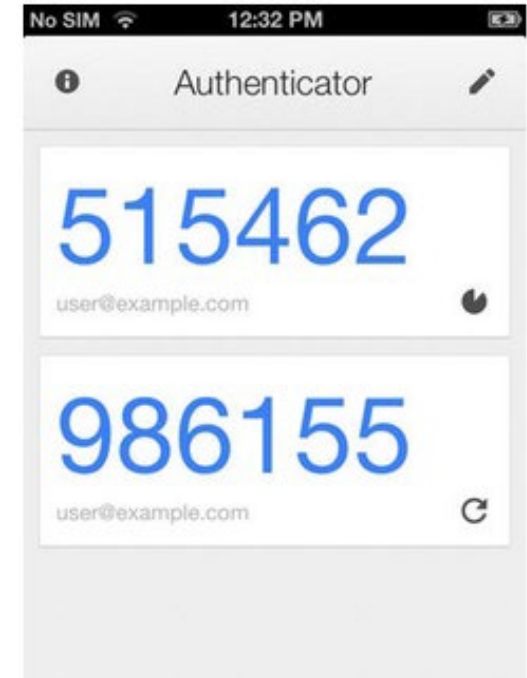
The key is negotiated during registration

Requires “rough” client–server synchronization

Code constantly changes in 30-second intervals

Phishing is still possible!

The attacker just needs to proxy the captured credentials in real time (rather than collecting them for later use)



Evilginx2 <https://github.com/kgretzky/evilginx2>

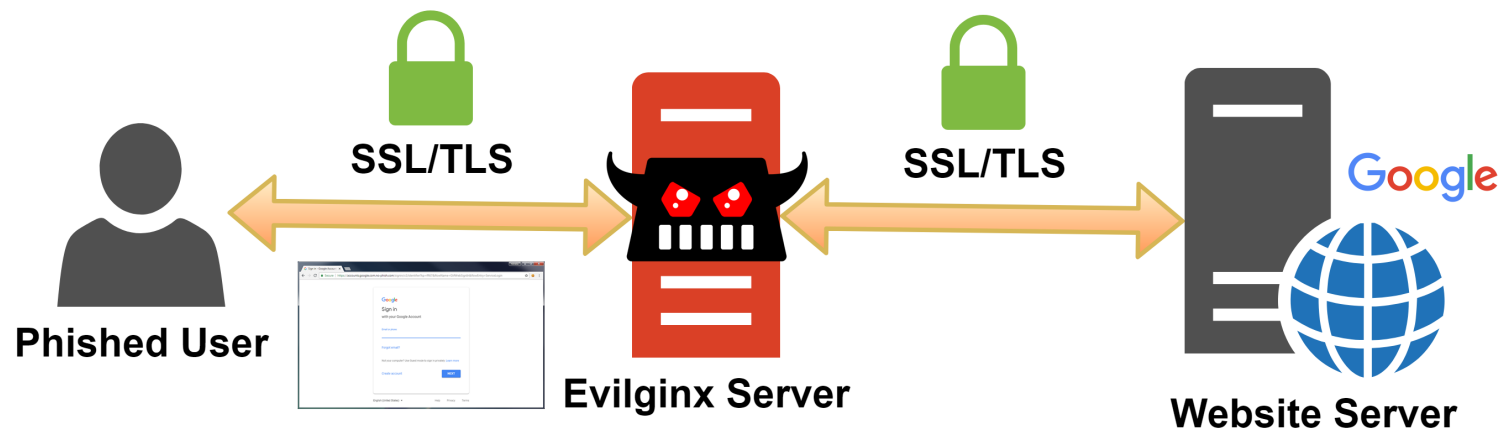
Man-in-the-middle attack framework for phishing login credentials along with session cookies

Bypasses 2-factor authentication

No need for HTML templates: just a web proxy

Victim's traffic is forwarded to the real website

TLS termination at the proxy (e.g., using a LetsEncrypt certificate)





Google

Sign in

with your Google Account

Email or phone

[Forgot email?](#)

Not your computer? Use Guest mode to sign in privately. [Learn more](#)

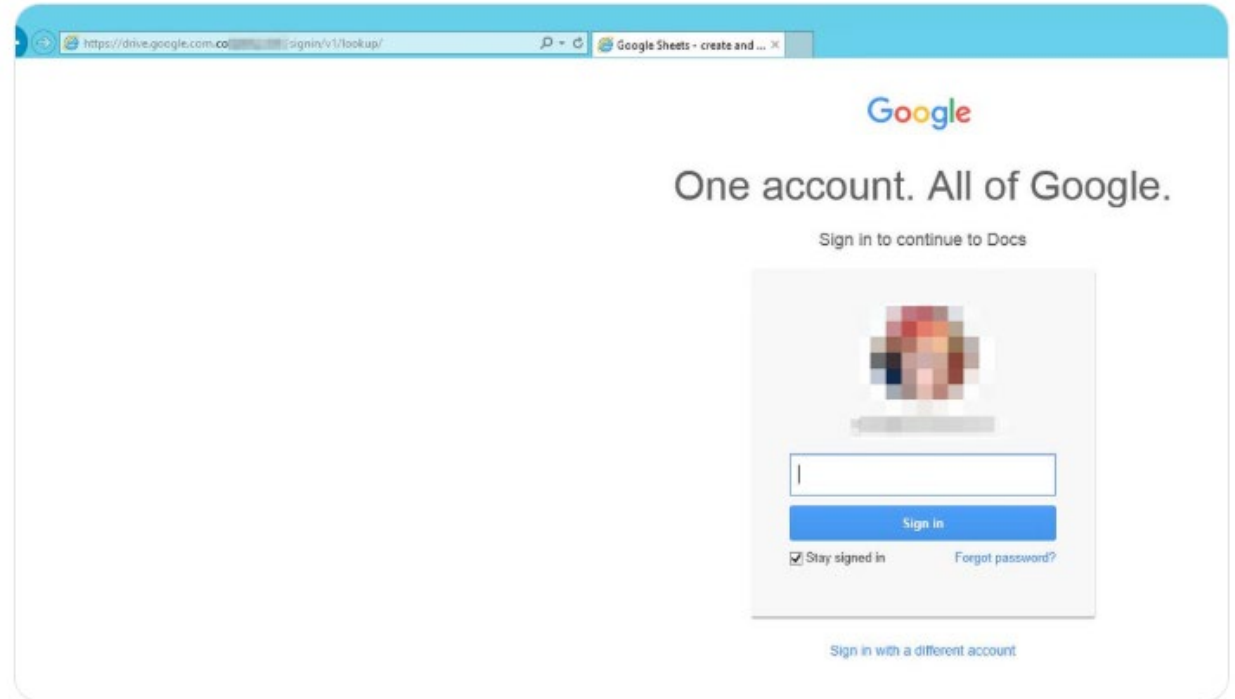
[Create account](#) [NEXT](#)



Justin Warner
@sixdub

Follow

I love digging through certificate transparency logs. Today, I saw a fake Google Drive landing page freshly registered with Let's Encrypt. It had a hardcoded picture/email of presumably the target. These can be a wealth of info that I recommend folks checking out.



5:21 PM - 22 Jul 2018

Evilginx2's Tokenized phishing URLs

Scanners look into public certificate transparency logs for newly registered domains

"For some phishing pages, it took usually one hour for the hostname to become banned and blacklisted by popular anti-spam filters"

Solution: create unique phishing URLs

Response to scanner: benign page

<https://totally.not.fake.linkedin.foo.com/auth/signin>

Response to victim: malicious page

https://totally.not.fake.linkedin.foo.com/auth/signin?tk=secret_token

Additional countermeasure: temporarily hide the phishing page

While submitting it to bit.ly, sending it through email, appearing on CT log, ...

Modlishka <https://github.com/drk1wi/Modlishka>

Phishing reverse proxy

Support for the majority of 2FA authentication schemes

No website templates

User credential harvesting (with context based on URL parameter passed identifiers)

Web panel with a summary of collected credentials and user session impersonation

```
>>>> "Modlishka" Piotr Duszynski @drk1wi - Reverse Proxy started <<<<

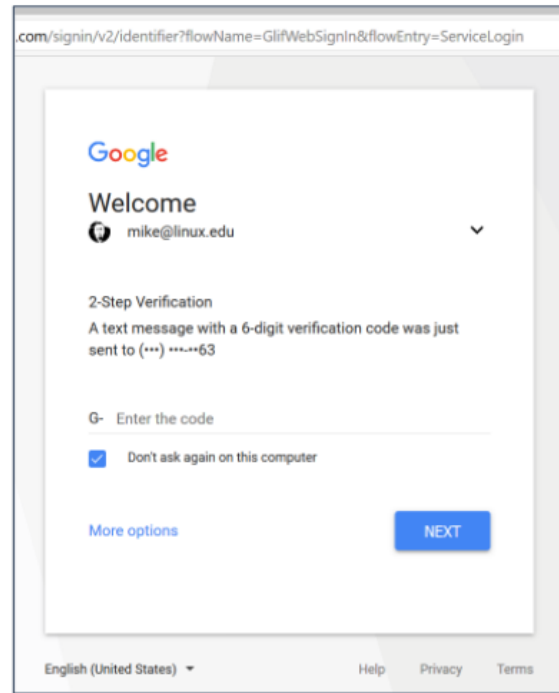
[127.0.0.1:443]
[127.0.0.1:443]

Listening on: [127.0.0.1:443]
Proxying [phishing.com.dev:443] via --> [https://google.com]
[Sat Dec 22 14:02:41 2018] INF Username collected ID:[42bc12cf-eea6-4cc1-acc9-84fe10b81f4c] username: phishingng
[Sat Dec 22 14:02:47 2018] INF Credentials collected ID:[42bc12cf-eea6-4cc1-acc9-84fe10b81f4c] username: phishingng password: supersecretpass
[Sat Dec 22 14:03:23 2018] INF [P] Tracking victim via initial parameter 9a0d22a9-19be-4c13-bc61-ff1dae2d7170
[Sat Dec 22 14:03:46 2018] INF Credentials collected ID:[9a0d22a9-19be-4c13-bc61-ff1dae2d7170] username: testuser password: yetanothersecretpass
```

CredSniper <https://github.com/ustayready/CredSniper>

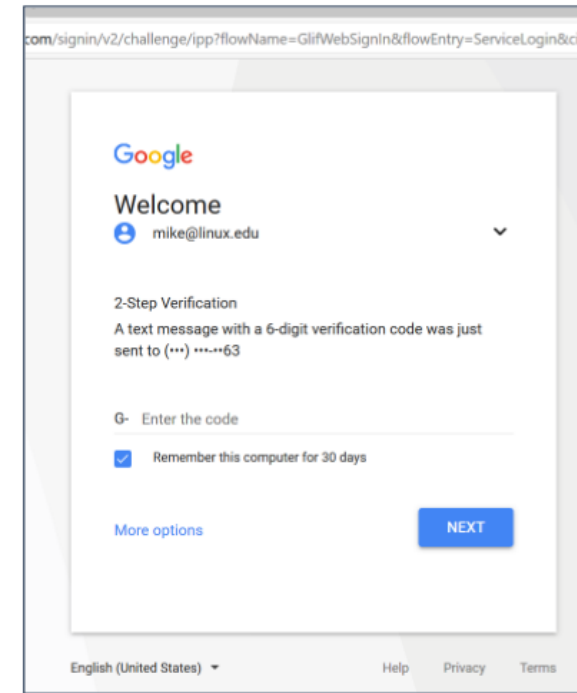
Exact login form clones for realistic phishing

Supports TLS via Let's Encrypt, and phishing 2FA tokens



Fake

**Real
Or
Fake?**



Real

Even Better Alternative: U2F Tokens

Universal Second Factor (U2F)

FIDO (Fast IDentity Online) alliance: Google, Yubico, ...

Supported by all popular browsers and many online services



A different key pair is generated for each origin during registration

Origin = <protocol, hostname, port>

Private key stored re-generated on device

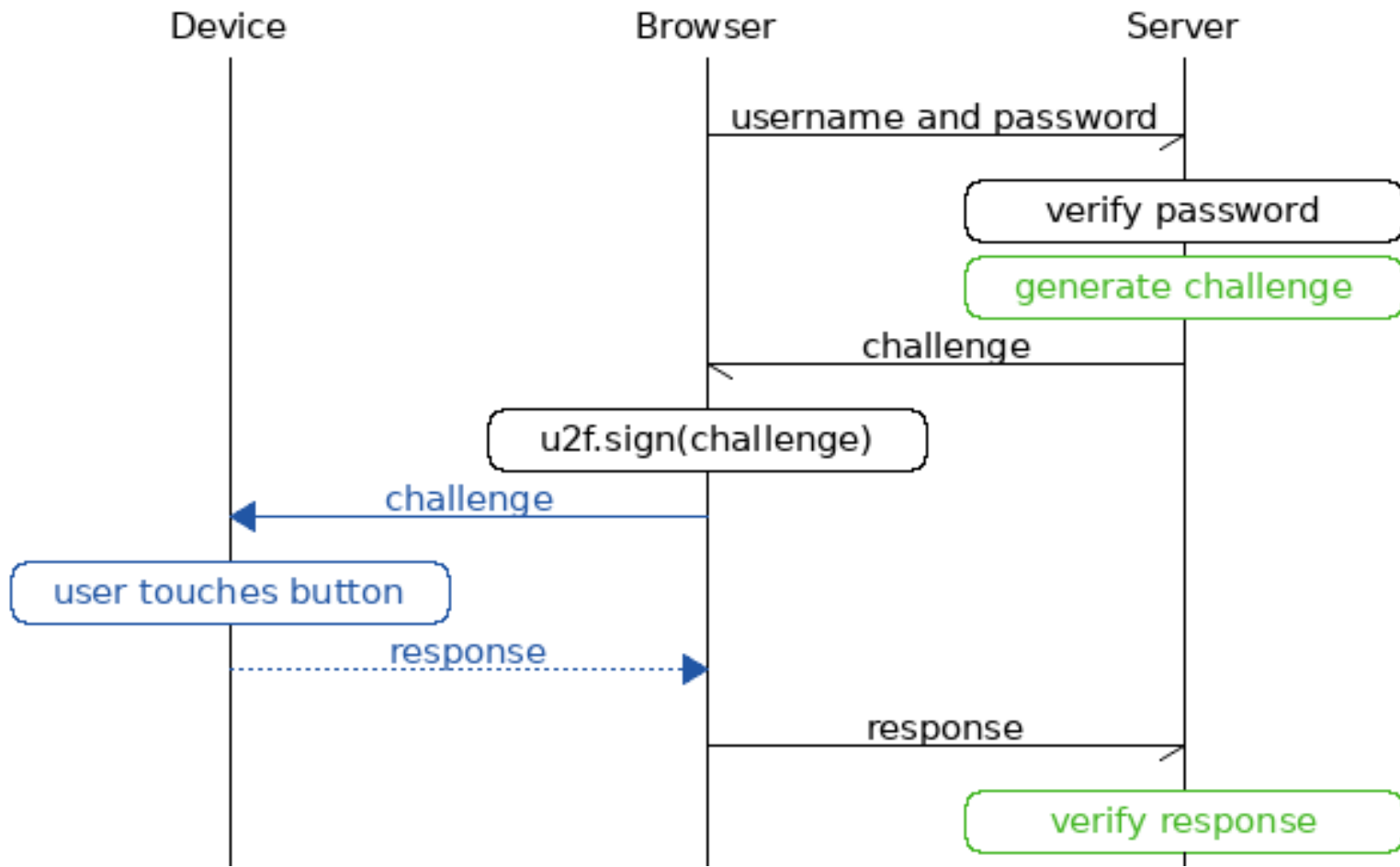
Public key sent to server

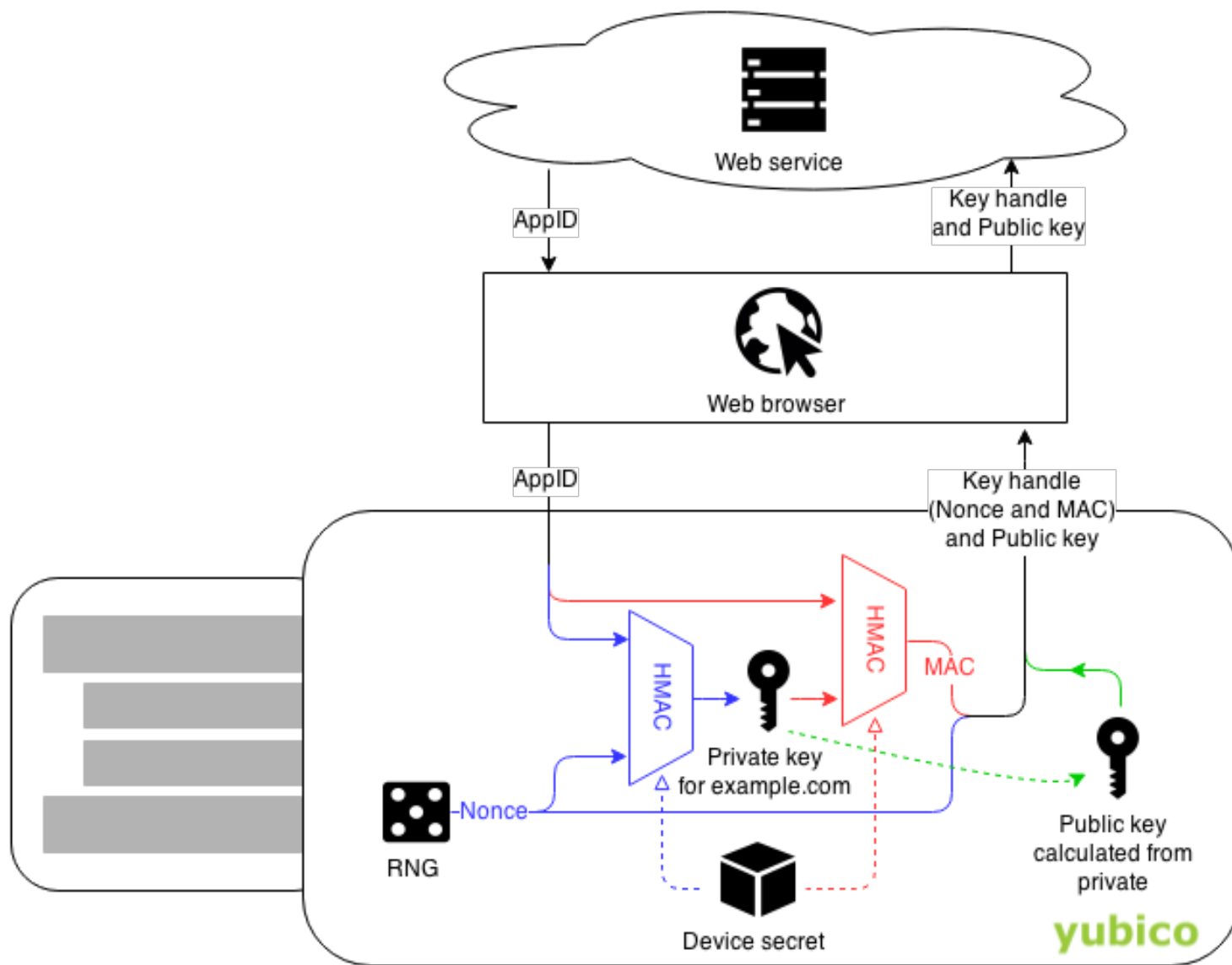
Additions to the authentication flow:

Origin (URI): *prevents phishing*

TLS Channel ID (optional): *prevents MitM*







U2F tokens

Benefits

Easy: just tap the button (no typing)

Works out of the box (no drivers to install)

USB, NFC, Bluetooth communication

No shared secret between client and server

Origin checking → effective against phishing!



Drawbacks

Can be lost → a fallback is needed (second U2F token, Authenticator App, ...)

Cumbersome: have to pull keychain out and plug token in (or have an always pugged-in token, in which case though it can be stolen along with the device)

Cost (\$10–\$70)

2FA Recap – *What threats does it prevent?*

SMS: useful against two main threats

Credential stuffing (people tend to reuse passwords across different services)

Leaked passwords (post-it, hardware keyloggers, cameras, shoulder surfing, ...)

Introduces new security/privacy issues: SIM swapping, SMS forwarding, SMS spam...

Authenticator Apps/Push Auth: much better alternative than SMS

Protects against the same threats without relying on phone numbers

U2F: additional protection against phishing

Modern phishing toolkits bypass SMS/Authenticator/Push 2FA through MitM

Humans fall for typosquatting, but U2F's origin check doesn't

None of the above protect against session hijacking and Man-in-the-Browser

Game over anyway if the host is compromised after the user has successfully logged in

Password Managers

Have become indispensable

- Encourage the use of complex/non-memorable passwords

- Obviate the need for password reuse: unique passwords per site/service

Protection against phishing: auto-fill won't work for incorrect domains

- As long as users don't copy/paste passwords out of the password manager (!)

Various options: third-party applications, OS-level, in-Browser

Password synchronization across devices

- Can the service provider access all my passwords or not?

- Preferable option: passwords should be encrypted with master password never visible to the cloud service

WebAuthn

W3C Web Authentication standard (FIDO2): Successor of FIDO U2F

Use cases

Low friction and phishing-resistant 2FA (in conjunction with a password)

Passwordless, biometrics-based re-authorization

2FA *without* a password (passwordless logins)

Authenticators: devices that can generate private/public key pairs and gather consent (simple tap, fingerprint read, ...)

Built-in: fingerprint readers, cameras, ...

External: USB, BLE, NFC, ...



Single Sign-on/Social Login

Pros

Convenience: fewer passwords to remember

Easier development: outsource user registration/management

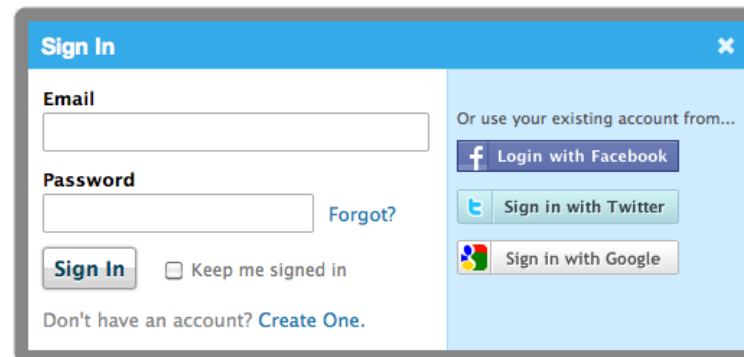
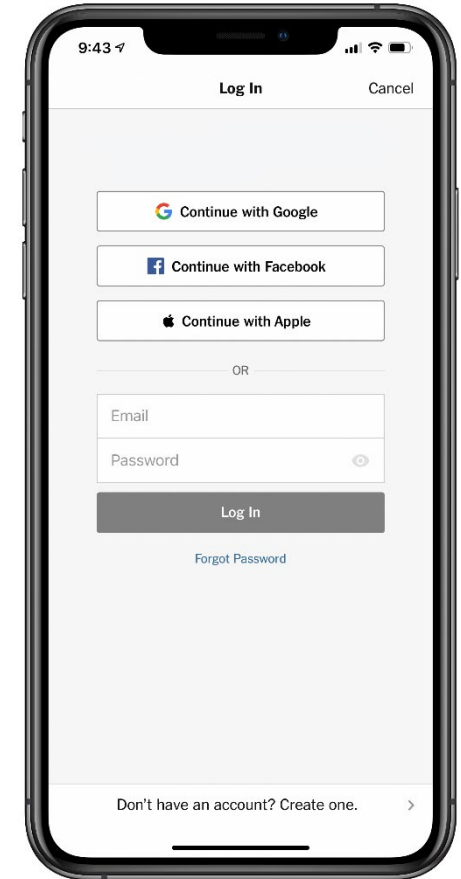
Rich experience through social features

Cons

Same credentials for multiple sites: single point of failure

Access to user's profile

User tracking



Biometrics

Fingerprint reader



Face recognition

Depth sensing, infrared cameras, ...

“liveness” detection (pulse, thermal) to foil simple picture attack

Retina/iris scanner

Voice recognition

...

Related concept: continuous authentication

Keystroke timing, usage patterns, ...

Crypto-based Authentication

Rely on a cryptographic key to prove a user's identity

User performs a requested cryptographic operation on a value (challenge) that the verifier supplies

Usually based on knowledge of a key (secret key or private key)

Can use symmetric (e.g., Kerberos) or public key (e.g., U2F) schemes

How can we trust a key? Why is it authentic?

Need to establish a level of trust

Different approaches: **TOFU, PKI, Web of Trust**

Emerging approach: PKI based on ~~blockchain~~ distributed ledger

Trust on First Use (aka Key Continuity)

Use case: SSH

Performs *mutual authentication*

Server *always* authenticates the client

password, key pair, ...

Client almost always authenticates the server – *except the first time!*

First connection: server presents its public key

No other option for the user but to accept it: MitM opportunity

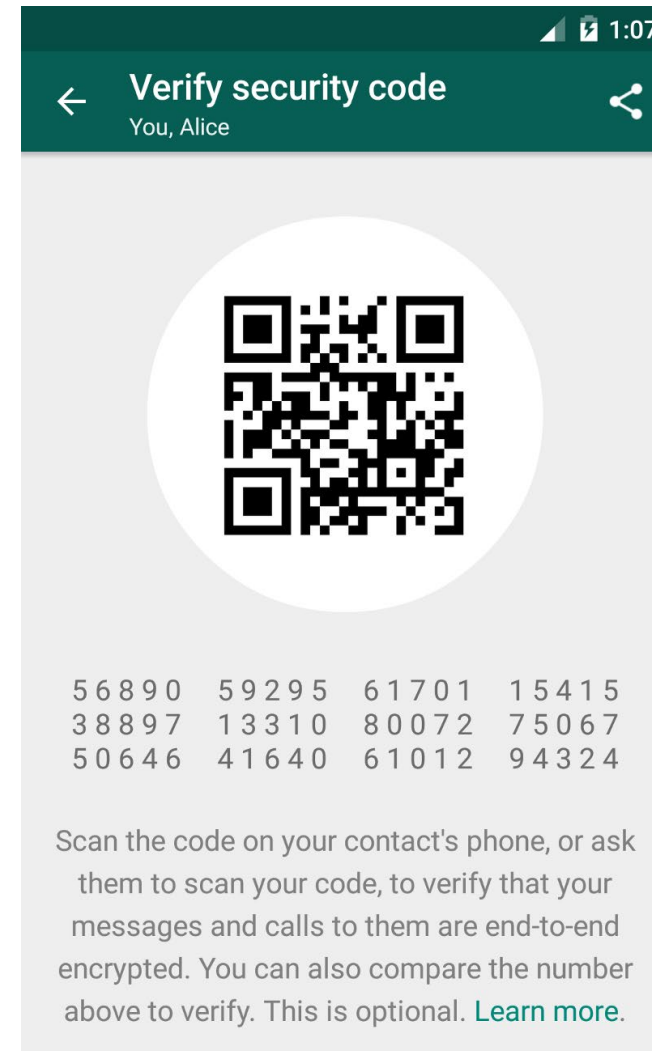
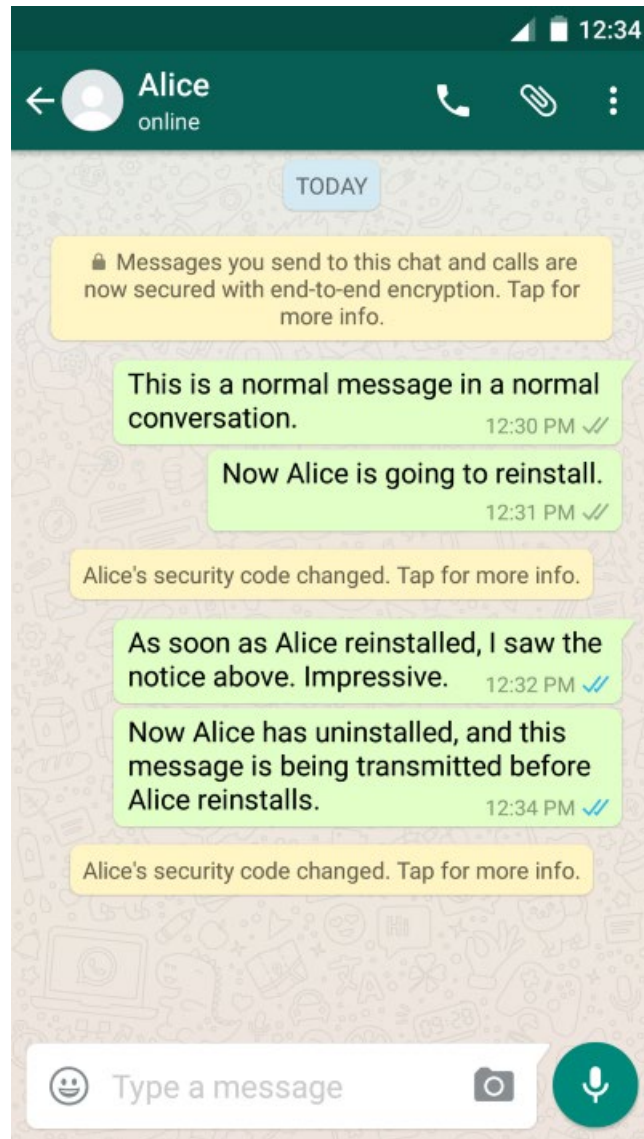
Subsequent connections: client remembers server's key, and triggers an alert on key mismatch

Pragmatic solution, but shifts the burden to users

Users must determine the validity of the presented key

Accepting a key change without verifying the new key offers no protection against MitM (unfortunately, that's what most users do)


```
@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@
@   WARNING: REMOTE HOST IDENTIFICATION HAS CHANGED!   @
@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@
IT IS POSSIBLE THAT SOMEONE IS DOING SOMETHING NASTY!
Someone could be eavesdropping on you right now (man-in-the-middle attack)!
It is also possible that the RSA host key has just been changed.
The fingerprint for the RSA key sent by the remote host is
df:c8:52:aa:cd:e3:da:8c:ec:50:46:db:4d:21:d9:c7.
Please contact your system administrator.
Add correct host key in /root/.ssh/known_hosts to get rid of this message.
Offending key in /root/.ssh/known_hosts:1
RSA host key for 192.168.2.5 has changed and you have requested strict checking.
Host key verification failed.
```



SCAN CODE

Certificates

How can we distribute “trusted” public keys?

Public directory → risk of forgery and tampering

More practical solution: “certified” public keys

A certificate is a digitally signed message that contains an identity and a public key

Makes an association between a user/entity and a private key

Valid until a certain period

Why trust a certificate?

Because it is signed by an “authority”

Requiring a signature by a third party prevents straightforward tampering



Public Key Infrastructures (PKI)

Facilitate the authentication and distribution of public keys with the respective identities of entities

People, organizations, devices, applications, ...

Set of roles, policies, hardware, software, and procedures to create, manage, distribute, use, store, and revoke digital certificates and manage public key encryption

An issuer signs certificates for subjects

Trust anchor

Methods of certification

Certificate authorities (hierarchical structure – root of trust)

Web of trust (decentralized, peer-to-peer structure)

Certificate Authorities

Trusted third-parties responsible for certifying public keys

- Most CAs are tree-structured

- Single point of failure: CAs can be compromised!

Why should we trust an authority?

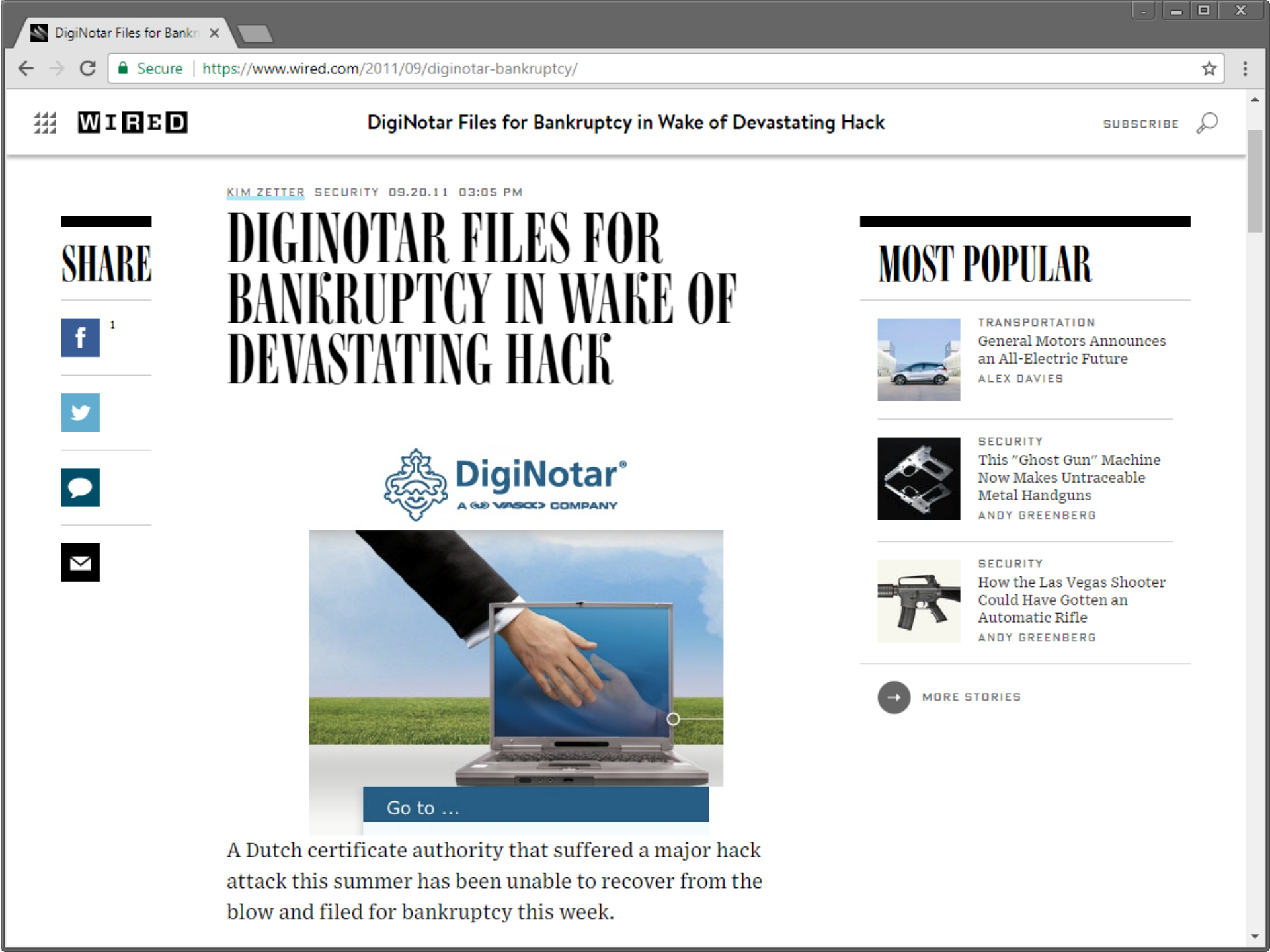
- How do we know the public key of the Certificate Authority (CA)?

CA's public key (trust anchor) must somehow be provided out of band

- Trust has to start somewhere

Operating systems and browsers are pre-configured with tens/hundreds of trusted root certificates

- A public key for any website in the world will be accepted without warning if it has been certified by any of these CAs (more on that in the TLS lecture)



SHARE



1



KIM ZETTER SECURITY 09.20.11 03:05 PM

DIGINOTAR FILES FOR BANKRUPTCY IN WAKE OF DEVASTATING HACK



Go to ...

A Dutch certificate authority that suffered a major hack attack this summer has been unable to recover from the blow and filed for bankruptcy this week.

MOST POPULAR



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ALEX DAVIES



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ANDY GREENBERG



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ANDY GREENBERG

→ MORE STORIES

Web of Trust (mainly used in PGP – more in the email lecture)

Entirely decentralized authentication

- No single point of failure

- No need to buy certs from CAs

Users sign other users' keys

- Only if they deem them trustworthy

- Certificate signings can form an arbitrarily complex graph

- Users can verify path to as many trust anchors as they wish

Drawbacks

- Hard to use, requires in-person verification: key signing parties!

- Hard to know what trust level to assign transitively

WoT Alternative: Online Social "Tracking"

The screenshot shows a web browser window with the address bar displaying `https://keybase.io/mikepo`. The page features the Keybase logo and navigation links for 'Join', 'Login', and user actions. The profile section includes a circular profile picture of Michalis Polychronakis, his name, and the URL `keybase.io/mikepo`. Below this, there are links for his public key (`8EBD 8F30 8899 8AFF`), a Twitter handle (`polychronakis`), and a GitHub handle (`polychronakis`). A green notification box states: "mikepo has an invitation available. If you know mikepo, you can ask them for an invitation to Keybase." Below the profile are two buttons: 'Encrypt' and 'Verify'. The main content area shows a code block with instructions for joining and logging in from the command line. On the right, there are two columns: 'Tracking (6)' and 'Trackers (6)', each listing users with their profile pictures and usernames: `hargikas`, `mstamat`, and `gianluca_string` in the Tracking column; and `hargikas`, `kontaxis`, and `mstamat` in the Trackers column.

Michalis Polychronakis (m x)

`https://keybase.io/mikepo`

Keybase

Join Login

`keybase.io/mikepo`

`8EBD 8F30 8899 8AFF`

`polychronakis` tweet

`polychronakis` gist

✓ mikepo has an invitation available
If you know mikepo, you can ask them for an invitation to Keybase.

Encrypt Verify

mikepo from the [command line](#)

```
# first
keybase join # if you're new, or
keybase login # if you're not.

# then
keybase push # if you already have a public key, or
keybase gen # if this is all new to you
```

Tracking (6)

hargikas

mstamat

gianluca_string

Trackers (6)

hargikas

kontaxis

mstamat

Keybase.io

In essence, a directory associating public keys with names

Identity established through *public signatures*

Identity proofs: *"I am Joe on Keybase and MrJoe on Twitter"*

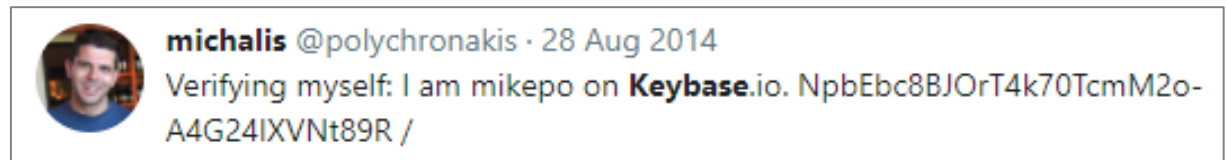
Follower statements: *"I am Joe on Keybase and I just looked at Chris's identity"*

Key ownership: *"I am Joe on Keybase and here's my public key"*

Revocations: *"I take back what I said earlier"*

Keybase identity = sum of public identities

Twitter, Facebook, Github, Reddit,
domain ownership, ...



An attacker has to compromise all connected identities

The more connected identities, the harder to impersonate a user

Best Practices

Use long passphrases instead of passwords

Never reuse the same password on different services

Use two-factor authentication when available

Avoid SMS if possible! Use an authenticator app or U2F instead

Remove phone number from account after authenticator/U2F setup

Store your backup codes in a safe location

Use a password manager

Pick non-memorable passwords and avoid copy/pasting them

Password auto-fill helps against phishing (auto-fill won't work if the domain is wrong)

Not only for passwords! Also for "security" questions

Use SSH keys instead of passwords