The Logic of Quantified Statements

CSE 215: Foundations of Computer Science

Stony Brook University

http://www.cs.stonybrook.edu/~liu/cse215

The Logic of Quantified Statements

All men are mortal.

Socrates is a man.

∴ Socrates is mortal.

- Propositional calculus: analysis of ordinary compound statements
- Predicate calculus: symbolic analysis of predicates and **quantified statements ($\forall x$, $\exists x$)**
  - *P* is a *predicate symbol*

    P stands for "*is a student at SBU*"

    P(x) stands for "*x is a student at SBU*"
  - *x* is a *predicate variable* **argument of the predicate**

Predicates and Quantified Statements

- A **predicate** is a sentence that contains a finite number of variables and becomes a **statement** when specific values are substituted for the variables.

- The **domain** of a predicate variable is the set of all values that may be substituted in place of the variable.

- Example:

  $P(x)$ is the predicate "$x^2 > x$", the domain of x is the set **R** of all real numbers

  $P(2): 2^2 > 2.$ True

  $P(1/2): (1/2)^2 > 1/2.$ False

Truth Set of a Predicate

- If $P(x)$ is a predicate and x has domain D, the truth set of $P(x)$, $\{x \in D \mid P(x)\}$, is the set of all elements of D that make $P(x)$ true when they are substituted for x.

- Example:

  $Q(n)$ is the predicate for "n is a factor of 8."

  If the domain of n is the set **Z** of all integers,

  the truth set is $\{1, 2, 4, 8, -1, -2, -4, -8\}$

The Universal Quantifier: ∀

- Quantifiers are words that refer to quantities ("some" or "all") and tell for how many elements a given predicate is true.

- **Universal quantifier:** ∀ "for all"

- Example:

  "All human beings are mortal"

  ∀ human beings x, x is mortal.

  If H is the set of all human beings

  $\forall x \in H$, x is mortal

Universal statements

- A **universal statement** is a statement of the form

  "$\forall x \in D, Q(x)$" where $Q(x)$ is a predicate and D is the domain of x.

  - $\forall x \in D, Q(x)$ is true if, and only if, $Q(x)$ is true for every x in D
  - $\forall x \in D, Q(x)$ is false if, and only if, $Q(x)$ is false for at least one x in D (the value for x is a **counterexample**)

- Example:

  $\forall x \in D, x^2 \geq x$ for $D = \{1, 2, 3, 4, 5\}$

  $1^2 \geq 1, \quad 2^2 \geq 2, \quad 3^2 \geq 3, \quad 4^2 \geq 4, \quad 5^2 \geq 5$

  - Hence "$\forall x \in D, x^2 \geq x$" is true.

The Existential Quantifier: ∃

- **Existential quantifier:** ∃ "there exists"

- Example:

  "There is a student in CSE 215"

  ∃ a person *p* such that *p* is a student in CSE 215

  More formally:

  $\exists p \in P$ such that *p* is a student in CSE 215

  where *P* is the set of all people

The Existential statement

- An **existential statement** is a statement of the form

  "$\exists x \in D$ such that $Q(x)$" where $Q(x)$ is a predicate and D the domain of x

  - $\exists x \in D$ s.t. $Q(x)$ is true if, and only if, $Q(x)$ is true for at least one x in D
  - $\exists x \in D$ s.t. $Q(x)$ is false if, and only if, $Q(x)$ is false for all x in D

- Example:
  - $\exists m \in \boldsymbol{Z}$ s.t. $m^2 = m$

  $1^2 = 1$ True

- Notation: such that = s.t.

Formal versus Informal Language

- Translating from formal to informal languages
- Translating from informal to formal Language
- Look at Exercise 1 answers

Will look again next time.

More extra-credit homework will be give on Friday.

Universal Conditional Statements

- **Universal conditional statement:**

  $\forall x$, *if P(x) then Q(x)* $\qquad \forall x,\ P(x) \longrightarrow Q(x)$

- Example:

  If a real number is greater than 2 then its square is greater than 4.

  $\forall x \in \boldsymbol{R}$, *if* $x > 2$ *then* $x^2 > 4$

Equivalent Forms of Universal and Existential Statements

- $\forall x \in U$, *if P(x) then Q(x)* can be rewritten in the form

  $\forall x \in D$, *Q(x)* by narrowing *U* to be the domain *D* consisting of all values of variable *x* that make *P(x)* true.

  - Example: $\forall x$, *if x is a square then x is a rectangle*

    $\forall$ *squares x, x is a rectangle.*

- $\exists x$ *s.t. P(x) and Q(x)* can be rewritten in the form

  $\exists x \in D$ *s.t. Q(x)* where *D* consists of

  all values of variable *x* that make *P(x)* true.

Implicit Quantification

- $P(x) \Rightarrow Q(x)$

  means that every element in the truth set of $P(x)$ is in the truth set of $Q(x)$, or, equivalently,

  $\forall x,\ P(x) \longrightarrow Q(x)$

- $P(x) \Leftrightarrow Q(x)$

  means that $P(x)$ and $Q(x)$ have identical truth sets, or, equivalently,

  $\forall x,\ P(x) \leftrightarrow Q(x)$

Negations of Quantified Statements

- Negation of a Universal Statement:

  The negation of a statement of the form $\forall x \in D, Q(x)$ is logically equivalent to a statement of the form $\exists x \in D, \sim Q(x)$, that is,

  $$\sim(\forall x \in D, Q(x)) \equiv \exists x \in D, \sim Q(x)$$

- Example:
  - "All mathematicians wear glasses"
  - Its negation is: "There is at least one mathematician who does not wear glasses"
  - Its negation is NOT "No mathematicians wear glasses"

Negations of Quantified Statements

- Negation of an Existential Statement

  The negation of a statement of the form $\exists x \in D, Q(x)$ is logically equivalent to a statement of the form $\forall x \in D, \sim Q(x)$, that is,

$$\sim(\exists x \in D, Q(x)) \equiv \forall x \in D, \sim Q(x)$$

- Example:
  - "Some snowflakes are the same."
  - Its negation is:

  "No snowflakes are the same" ≡ "All snowflakes are different."

- Better example: "Some students finished HW1"

  negation: "All students did not finish HW1"

Negations of Quantified Statements

- More Examples:
  - ~(∀ primes *p*, *p* is odd) ≡ ∃ a prime *p* s.t. *p* is **not** odd
  - ~(∃ a triangle *T* s.t. the sum of the angles of *T* equals 200°)
    ≡ ∀ triangles *T*, the sum of the angles of *T* **does not** equal 200°
  - ~(∀ politicians *x*, *x* is **not** honest)
    ≡ ∃ a politician *x* s.t. *x* is honest **(by double negation)**
  - ~(∀ computer programs *p*, *p* is finite)
    ≡ ∃ a computer program *p* that is not finite
  - ~(∃ a computer hacker *c*, *c* is over 40)
    ≡ ∀ computer hacker *c*, *c* is 40 or under
  - ~(∃ an integer *n* between 1 and 37 s.t. 1,357 is divisible by *n*)
    ≡ ∀ integers *n* between 1 and 37, 1,357 is not divisible by *n*

Negations of Universal Conditional Statements

- $\sim(\forall x, P(x) \rightarrow Q(x)) \equiv \exists x \text{ s.t. } P(x) \wedge \sim Q(x)$
- Proof:

  $\sim(\forall x, P(x) \rightarrow Q(x)) \equiv \exists x \text{ s.t. } \sim(P(x) \rightarrow Q(x))$

  $\sim(P(x) \rightarrow Q(x)) \equiv P(x) \wedge \sim Q(x)$

- Examples:
  - $\sim$($\forall$ people p, if p is blond then p has blue eyes)
    $\equiv \exists$ a person p s.t. p is blond and p does not have blue eyes
  - $\sim$(If a computer program has more than 100,000 lines, then it contains a bug) $\equiv$ There is at least one computer program that has more than 100,000 lines and does not contain a bug

The Relation among ∀, ∃, ∧, and ∨

- $D = \{x_1, x_2, \ldots, x_n\}$ and $\forall x \in D, Q(x)$

  $\equiv$

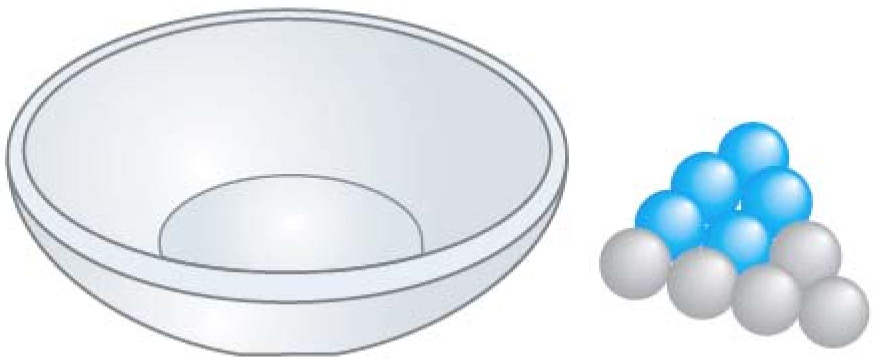
  $Q(x_1) \wedge Q(x_2) \wedge \cdots \wedge Q(x_n)$

- $D = \{x_1, x_2, \ldots, x_n\}$ and $\exists x \in D$ s.t. $Q(x)$

  $\equiv$

  $Q(x_1) \vee Q(x_2) \vee \cdots \vee Q(x_n)$

Vacuous Truth of Universal Statements



All the balls in the bowl are blue True

$\forall x$ in D, if $P(x)$ then $Q(x)$ is *vacuously true* or *true by default* if, and only if, $P(x)$ is false for every x in D

Variants of Universal Conditional Statements

- Universal conditional stmt: $\forall x \in D$, if $P(x)$ then $Q(x)$
- **Contrapositive:** $\forall x \in D$, if $\sim Q(x)$ then $\sim P(x)$

  $\forall x \in D$, if $P(x)$ then $Q(x)$ $\equiv$ $\forall x \in D$, if $\sim Q(x)$ then $\sim P(x)$

  Proof: for any x in D by the logical equivalence between statement and its contrapositive
- **Converse:** $\forall x \in D$, if $Q(x)$ then $P(x)$.
- **Inverse:** $\forall x \in D$, if $\sim P(x)$ then $\sim Q(x)$.
- Example: $\forall x \in R$, if $x > 2$ then $x^2 > 4$

  Contrapositive: $\forall x \in R$, if $x^2 \leq 4$ then $x \leq 2$

  Converse: $\forall x \in R$, if $x^2 > 4$ then $x > 2$

  Inverse: $\forall x \in R$, if $x \leq 2$ then $x^2 \leq 4$

Necessary and Sufficient Conditions

- Necessary condition:

  "$\forall$x, r (x) is a **necessary condition** for s(x)" means

  "$\forall$x, if ~r (x) then ~s(x)" ≡ "$\forall$x, if s(x) then r(x)"

  (by contrapositive)

- Sufficient condition:

  "$\forall$x, r (x) is a **sufficient condition** for s(x)" means

  "$\forall$x, if r (x) then s(x)"

Necessary and Sufficient Conditions

- Examples:
  - Squareness is a **sufficient condition** for rectangularity;

    Formal statement: ∀x, if x is a square, then x is a rectangle

  - Being at least 35 years old is a **necessary condition** for being President of the United States

    ∀ people x, if x is younger than 35, then x cannot be President of the United States ≡

    ∀ people x, if x is President of the United States then x is at least 35 years old (by contrapositive)

Only If

- Only If:

  "$\forall$x, r(x) **only if** s(x)" means

  "$\forall$x, if ~s(x) then ~r (x)" ≡ "$\forall$x, if r(x) then s(x)."

- Example:

  A product of two numbers is 0 only if one of the numbers is 0.

  If neither of two numbers is 0, then their product is not 0 ≡

  If a product of two numbers is 0, then one of the numbers is 0

  (by contrapositive)

Statements with Multiple Quantifiers

- Example:

  "There is a person supervising every detail of the production process"

  - What is the meaning?

  "There is one single person who supervises all the details of the production process"?

  OR

  "For any particular production detail, there is a person who supervises that detail, but there might be different supervisors for different details"?

- NATURAL LANGUAGE IS AMBIGUOUS

  LOGIC IS CLEAR

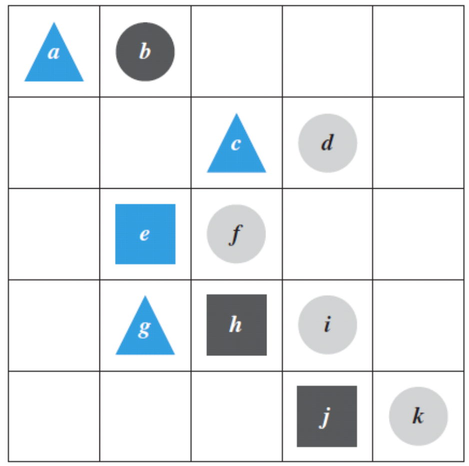
Statements with Multiple Quantifiers

- Quantifiers are performed in the order in which the quantifiers occur:
- Example:

  ∀x in set D, ∃y in set E s.t. x and y satisfy property P(x, y)

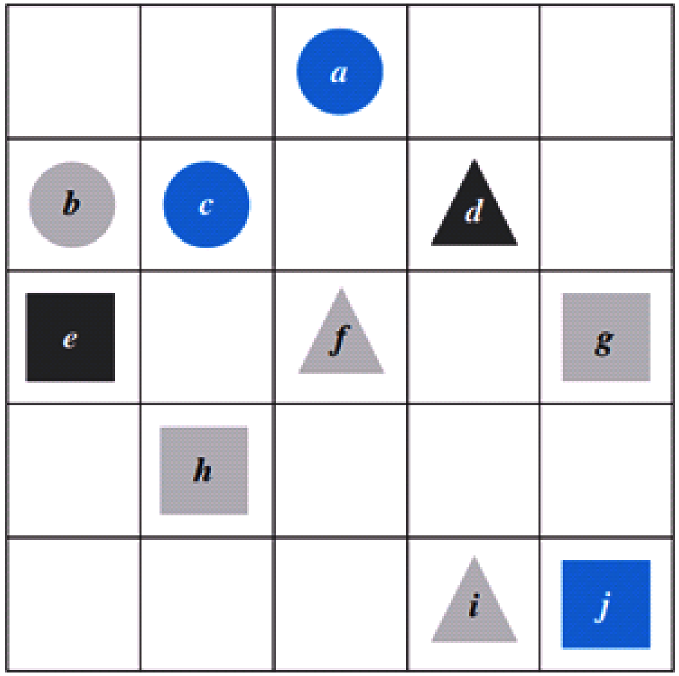
Tarski's World

- Blocks of various sizes, shapes, and colors located on a grid



- $\forall t$, Triangle(t) $\longrightarrow$ Blue(t) True
- $\forall x$, Blue(x) $\longrightarrow$ Triangle(x). False
- $\exists y$ s.t. Square(y) $\wedge$ RightOf(d, y). True
- $\exists z$ s.t. Square(z) $\wedge$ Gray(z). False

Statements with Multiple Quantifiers in Tarski's World



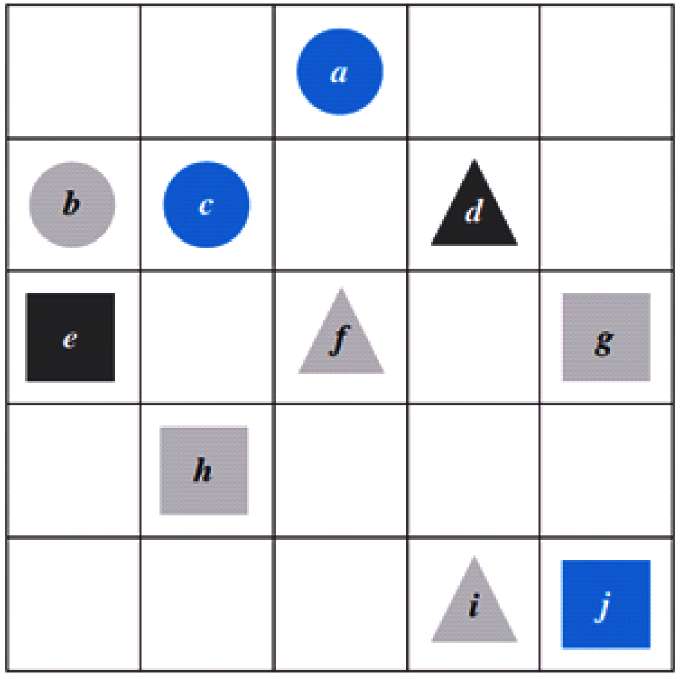
∀∃

- For all triangles x, there is a square y s.t. x and y have the same color

True

Given $x =$	**choose** $y =$	**and check that** y **is the same color as** x**.**
d	e	yes ✓
f or i	h or g	yes ✓

Statements with Multiple Quantifiers in Tarski's World



$\exists \forall$

- There is a triangle x s.t. for all circles y, x is to the right of y

True

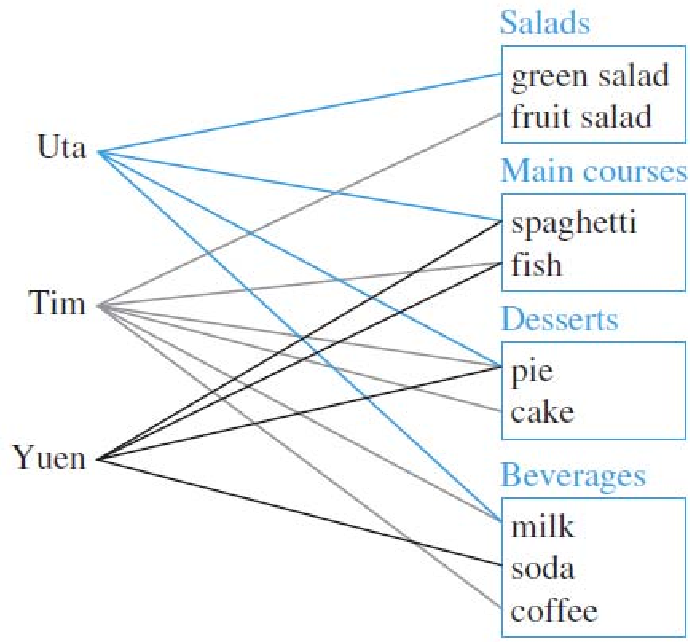
Choose $x =$	Then, given $y =$	check that x is to the right of y.
d or i	a	yes ✓
	b	yes ✓
	c	yes ✓

Interpreting Statements with Two Different Quantifiers

- ∀x in D, ∃y in E s.t P(x, y)
  - for whatever element x in D, you must find an element y in E that "works" for that particular x

- ∃x in D s.t. ∀y in E, P(x, y)
  - find one particular x in D that "works" no matter what y in E anyone might choose

Interpreting Statements with Multiple Quantifiers



- ∃ an item I s.t. ∀ students S, S chose I . True
- ∃ a student S s.t. ∀ stations Z, ∃ an item I in Z s.t. S chose I. True
- ∀ students S and ∀ stations Z, ∃ an item I in Z s.t. S chose I . False

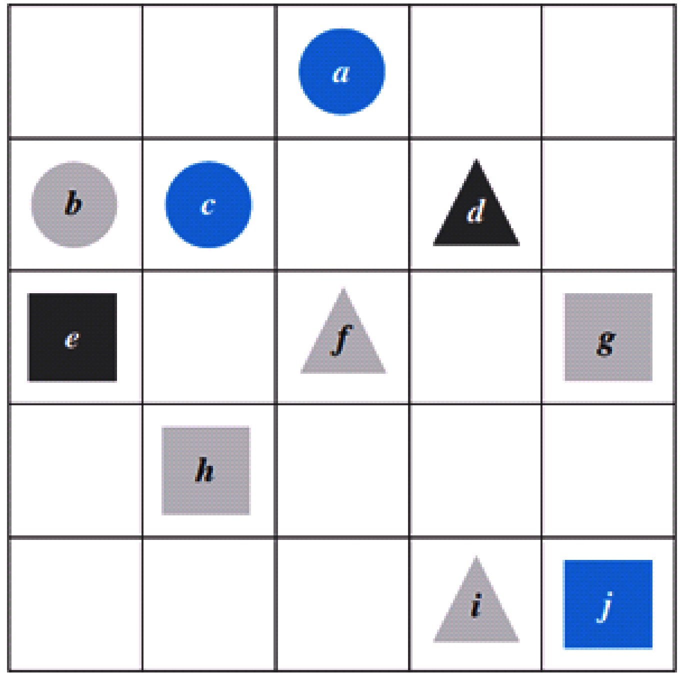
Negations of Multiply-Quantified Statements

- Apply negation to quantified statements from left to right:

$$\sim(\forall x \text{ in } D, \exists y \text{ in } E \text{ s.t. } P(x, y))$$
$$\equiv \exists x \text{ in } D \text{ s.t. } \sim(\exists y \text{ in } E \text{ s.t. } P(x, y))$$
$$\equiv \exists x \text{ in } D \text{ s.t. } \forall y \text{ in } E, \sim P(x, y)$$

$$\sim(\exists x \text{ in } D \text{ s.t. } \forall y \text{ in } E, P(x, y))$$
$$\equiv \forall x \text{ in } D, \sim(\forall y \text{ in } E, P(x, y))$$
$$\equiv \forall x \text{ in } D, \exists y \text{ in } E \text{ s.t. } \sim P(x, y)$$

Negating Statements in Tarski's World



- For all squares x, there is a circle y s.t. x and y have the same color

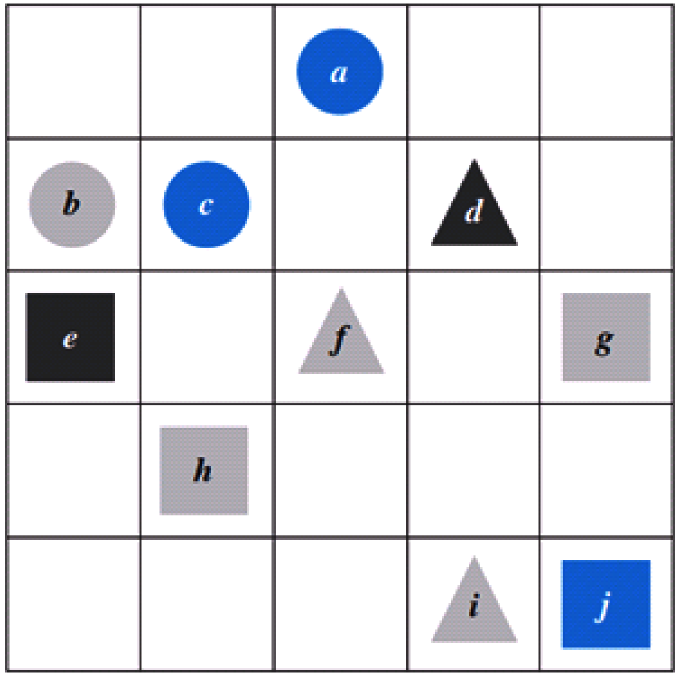
Negation:

∃ a square x s.t. ~(∃ a circle y s.t. x and y have the same color)

≡ ∃ a square x s.t. ∀ circles y, x and y do not have the same color True

Square e is black and no circle is black.

Negating Statements in Tarski's World



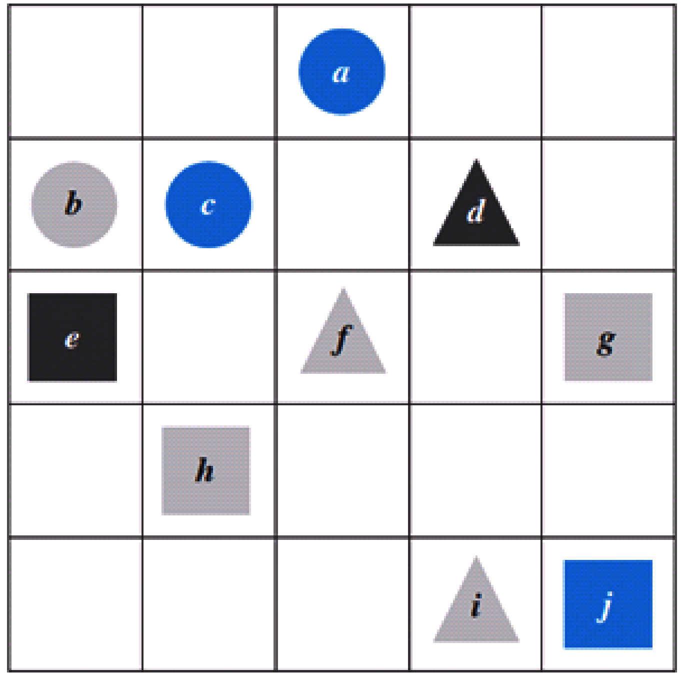
- There is a triangle x s.t. for all squares y, x is to the right of y

Negation:

∀ triangles x, ~(∀ squares y, x is to the right of y)

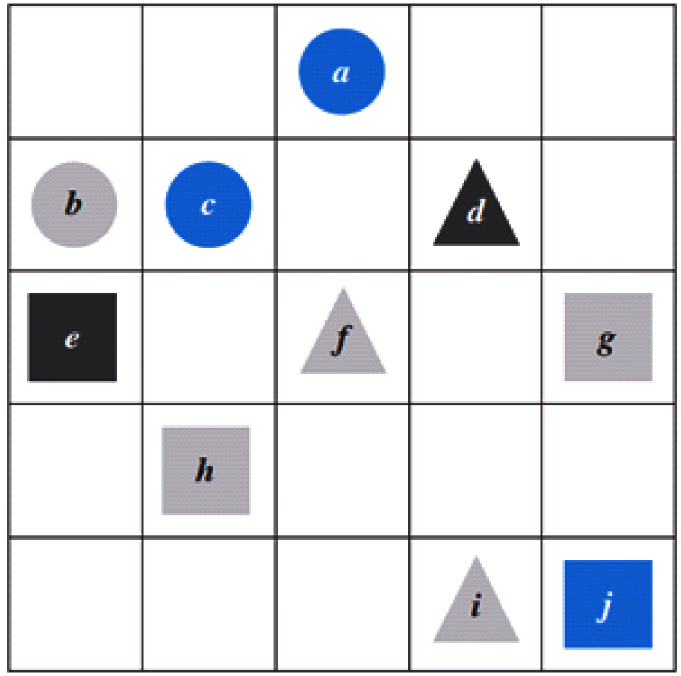
≡ ∀ triangles x, ∃ a square y s.t. x is not to the right of y   True

Quantifier Order in Tarski's World



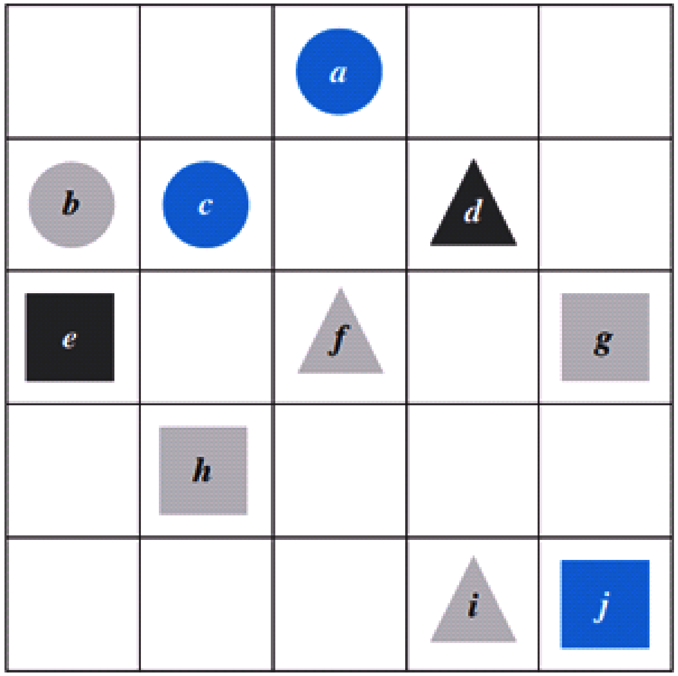
- For every square x there is a triangle y s.t. x and y have different colors
- There exists a triangle y s.t. for every square x, x and y have different colors

Quantifier Order in Tarski's World



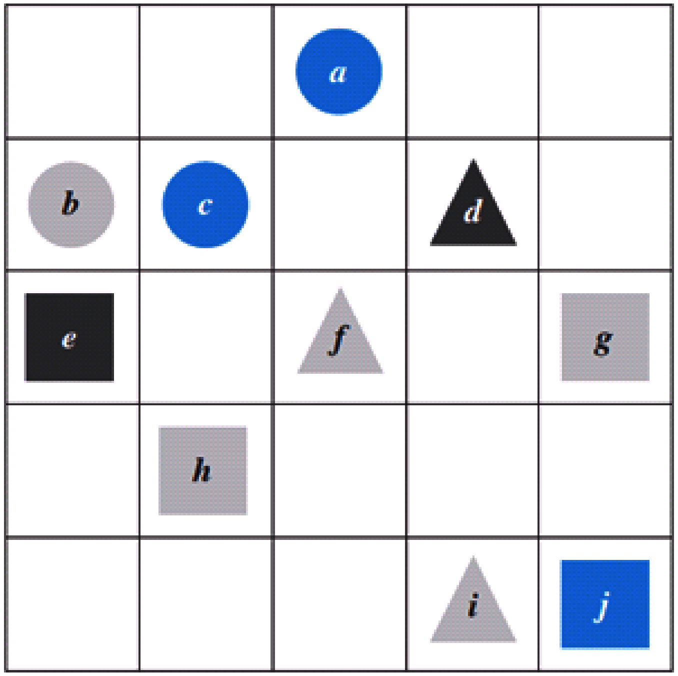
- For every square x there is a triangle y s.t. x and y have different colors
  True
- There exists a triangle y s.t. for every square x, x and y have different colors
  False

Formalizing Statements in Tarski's World



- Triangle(x) means "x is a triangle"
- Circle(x) means "x is a circle"
- Square(x) means "x is a square"
- Blue(x) means "x is blue"
- Gray(x) means "x is gray"
- Black(x) means "x is black"
- RightOf(x, y) means "x is to the right of y"
- Above(x, y) means "x is above y"
- SameColorAs(x, y) means "x has the same color as y"
- x = y denotes the predicate "x is equal to y"

Formalizing Statements in Tarski's World



- For all circles x, x is above f

  $\forall x$, Circle(x) $\rightarrow$ Above(x, f)

- Negation:

  $\sim(\forall x$, Circle(x) $\rightarrow$ Above(x, f))

  $\equiv \exists x$, $\sim$(Circle(x) $\rightarrow$ Above(x, f))

  $\equiv \exists x$, Circle(x) $\wedge$ $\sim$Above(x, f)

Formalizing Statements in Tarski's World

- There is a square x s.t. x is black

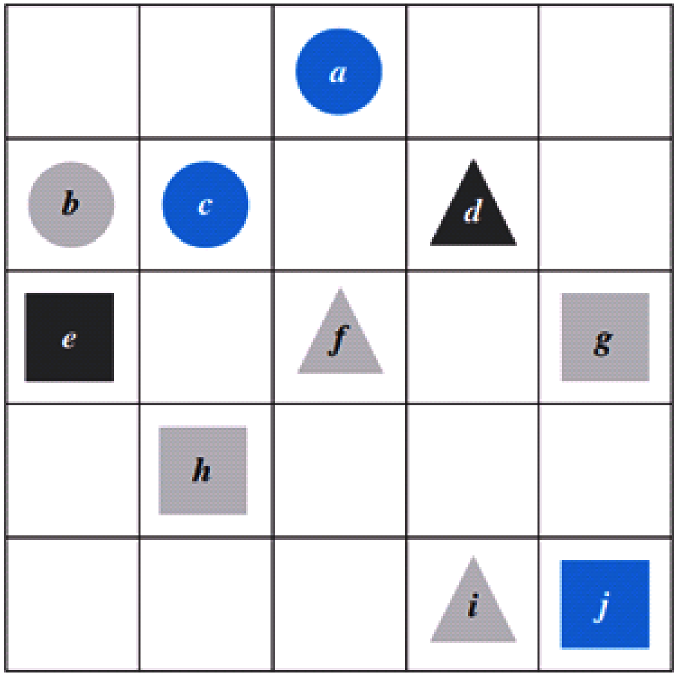
  $\exists x$, Square(x) $\wedge$ Black(x)

- Negation:

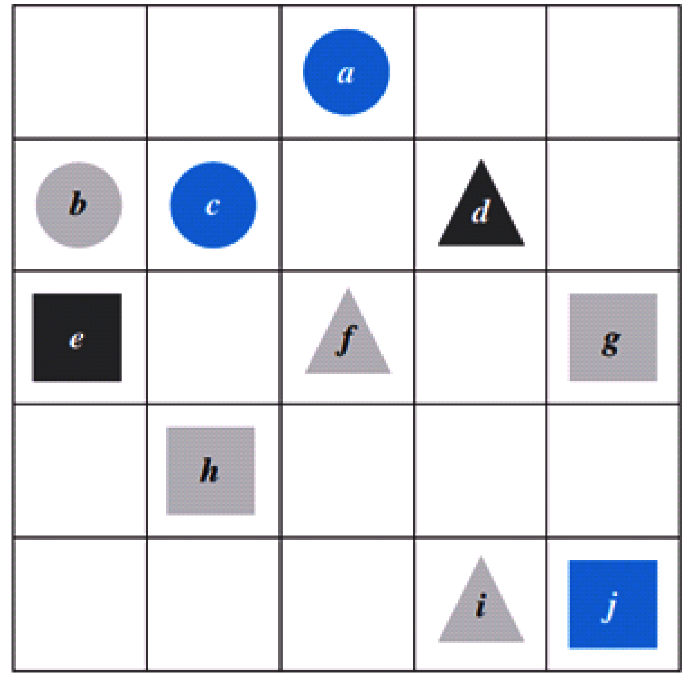
  $\sim(\exists x$, Square(x) $\wedge$ Black(x))

  $\equiv \forall x \sim$(Square(x) $\wedge$ Black(x))

  $\equiv \forall x$, $\sim$Square(x) $\vee$ $\sim$Black(x)



Formalizing Statements in Tarski's World



- For all circles x, there is a square y s.t. x and y have the same color

  $\forall x, \text{Circle}(x) \rightarrow \exists y, \text{Square}(y) \wedge \text{SameColor}(x, y)$

- Negation:

  $\sim(\forall x, \text{Circle}(x) \rightarrow \exists y, \text{Square}(y) \wedge \text{SameColor}(x, y))$

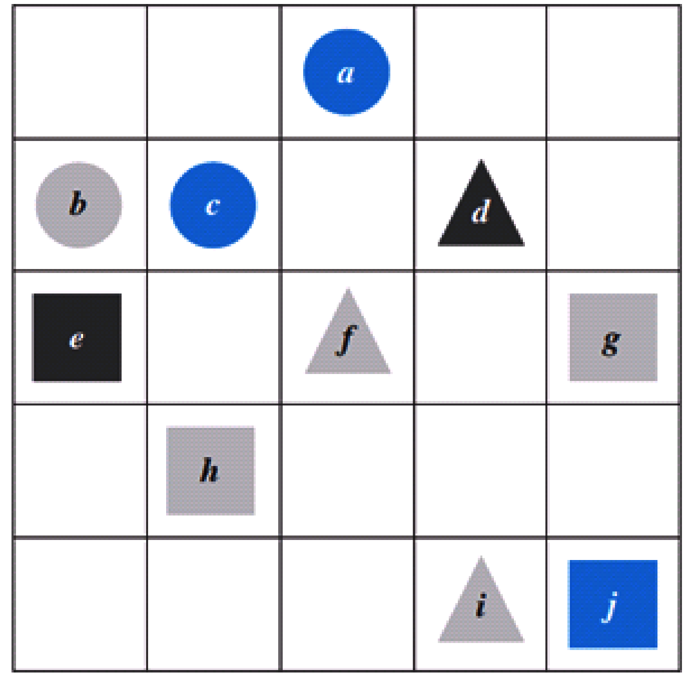
  $\equiv \exists x, \sim(\text{Circle}(x) \rightarrow \exists y, \text{Square}(y) \wedge \text{SameColor}(x, y))$

  $\equiv \exists x, \text{Circle}(x) \wedge \sim(\exists y, \text{Square}(y) \wedge \text{SameColor}(x, y))$

  $\equiv \exists x, \text{Circle}(x) \wedge \forall y, \sim(\text{Square}(y) \wedge \text{SameColor}(x, y))$

  $\equiv \exists x, \text{Circle}(x) \wedge \forall y, \sim\text{Square}(y) \vee \sim\text{SameColor}(x, y)$

Formalizing Statements in Tarski's World



- There is a square x s.t. for all triangles y, x is to right of y

  $\exists x, \text{Square}(x) \wedge \forall y, \text{Triangle}(y) \rightarrow \text{RightOf}(x, y)$

- Negation:

  $\sim(\exists x, \text{Square}(x) \wedge \forall y, \text{Triangle}(y) \rightarrow \text{RightOf}(x, y))$

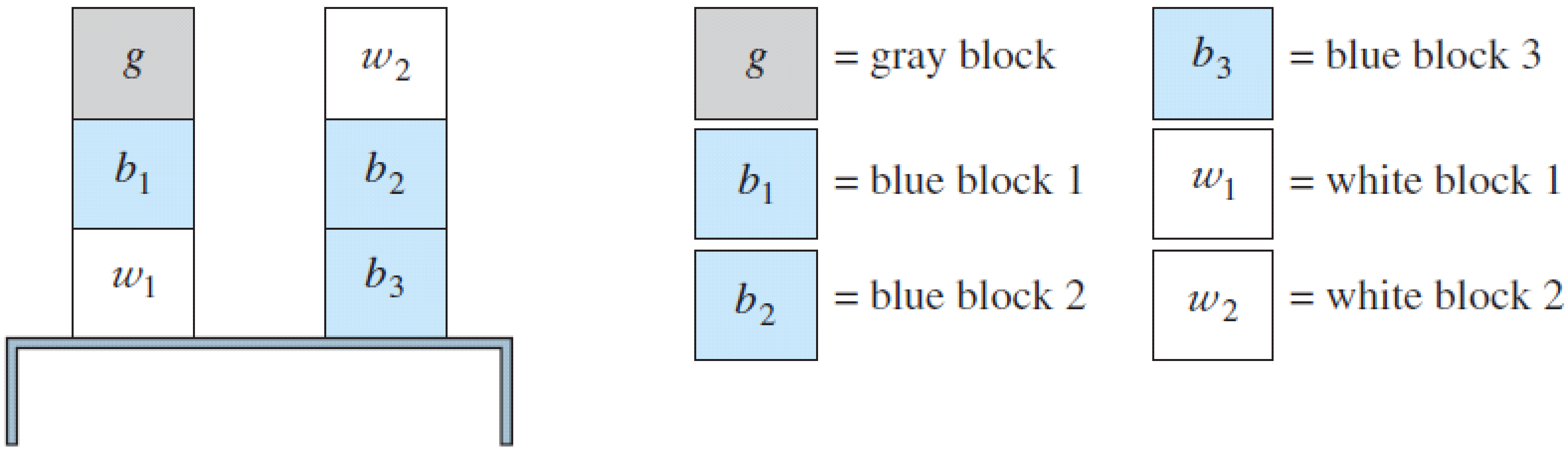
  $\equiv \forall x \sim(\text{Square}(x) \wedge \forall y, \text{Triangle}(x) \rightarrow \text{RightOf}(x, y))$

  $\equiv \forall x, \sim\text{Square}(x) \vee \sim(\forall y, \text{Triangle}(y) \rightarrow \text{RightOf}(x, y))$

  $\equiv \forall x, \sim\text{Square}(x) \vee \exists y, \sim(\text{Triangle}(y) \rightarrow \text{RightOf}(x, y))$

  $\equiv \forall x, \sim\text{Square}(x) \vee \exists y, \text{Triangle}(y) \wedge \sim\text{RightOf}(x, y)$

Prolog (Programming in logic)



- Prolog statements:

color(g, gray). color(b_1, blue). color(b_2, blue). color(b_3, blue).

color(w_1, white). color(w_2, white).

isabove(g, b_1). isabove(b_1,w_1). isabove(w_2, b_2). isabove(b_2, b_3).

isabove(X, Z) :- isabove(X,Y), isabove(Y, Z).

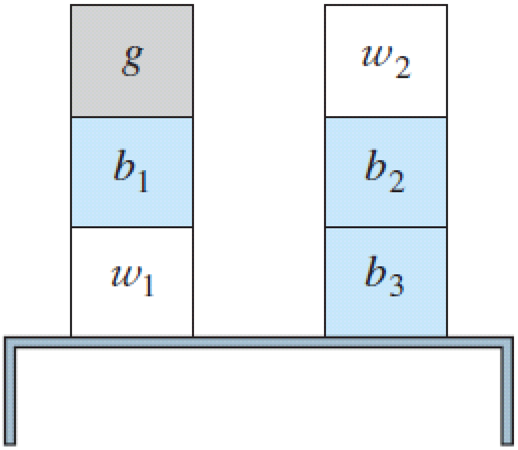
?- color(b_1, blue).

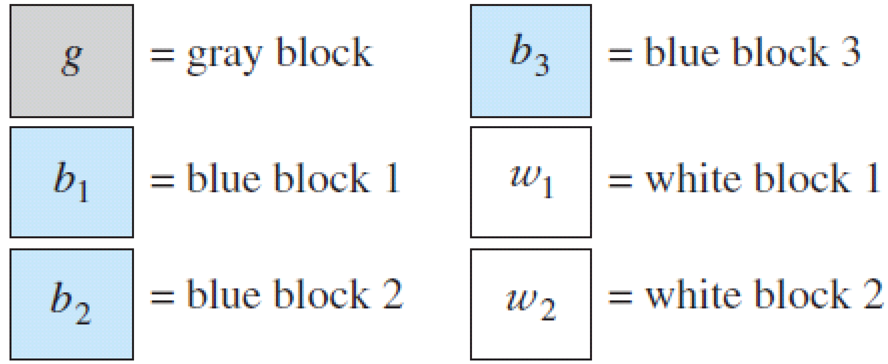
True

?- isabove(X,w_1).

X=b_1; X=g

Prolog (Programming in logic)





?- isabove(b_1,w_1).

True

?- color(w_1, X) .

X = white

?- color(X, blue).

X = b1; X = b2; X = b3.

Logic and practice of programming

- CS: Logic/math versus practice of programming
- Practice: Python has become the most widely used language, by both the least experienced and the most experienced. C…
- Python's inspiration and predecessor was ABC --- for beginners; Python was also influenced by C --- system programming.
- ABC was inspired by SETL--- based on mathematical theory of sets.
- Logic and practice should be together --- much easier and simpler, providing much more assurance, and fun! …DistAlgo, Alda

Precise quantifiers each and some

- Universal quantification

```
∀ x ∈ S  ,   P(x)
∀ x ∈ S  |   P(x)
∀ x ∈ S  :   P(x)
...
```

```
each x in S  has  P(x)
each(x in S, has= P(x))

all( P(x)  for x in S )
forall x in S  |  P(x)
```

- Existential quantification

```
∃ x ∈ S s.t. P(x)     textbook
∃ x ∈ S  ,   P(x)     others
...        |
...
```

```
some x in S  has  P(x)       ABC,ideal
some(x in S, has= P(x))      da in py

any( P(x)  for x in S )      py
exists x in S  |  P(x)       setl
```

Arguments with Quantified Statements

- Universal instantiation: if some property is true of everything in a set, then it is true of any particular thing in the set.

- Example:

  All men are mortal.

  Socrates is a man.

  ∴ Socrates is mortal.

Universal Modus Ponens

Formal Version	*Informal Version*
∀x, if P(x) then Q(x).	If x makes P(x) true, then x makes Q(x) true.
P(a) for a particular a.	a makes P(x) true.
∴ Q(a).	∴ a makes Q(x) true.

- Example:

∀x, if E(x) then S(x).	If an integer is even, then its square is even.
E(k) for a particular k.	k is a particular integer that is even.
∴ S(k).	∴ k^2 is even.

Universal Modus Tollens

Formal Version	*Informal Version*
∀x, if P(x) then Q(x).	If x makes P(x) true, then x makes Q(x) true.
~Q(a), for a particular a.	a does not make Q(x) true.
∴ ~P(a).	∴ a does not make P(x) true.

- Example:

∀x, if H(x) then M(x)	All human beings are mortal.
~M(Z)	Zeus is not mortal.
∴ ~H(Z).	∴ Zeus is not human.

Validity of Arguments with Quantified Statements

- An argument form is **valid**, if and only if, for any particular predicates substituted for the predicate symbols in the premises **if the resulting premise statements are all true, then the conclusion is also true**

- Using diagrams to test for validity:

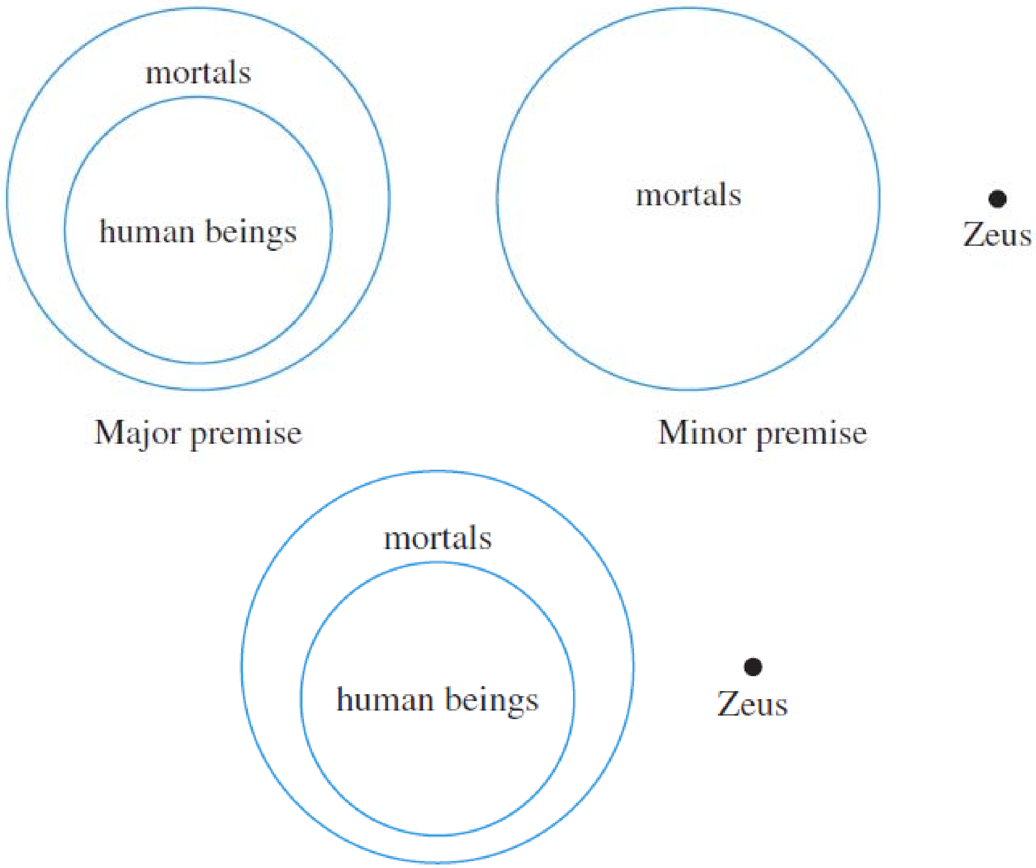
$$\forall \textit{ integers } n, n \textit{ is a rational number}$$

Using Diagrams to Test for Validity

All human beings are mortal.

Zeus is not mortal.

∴ Zeus is not a human being.

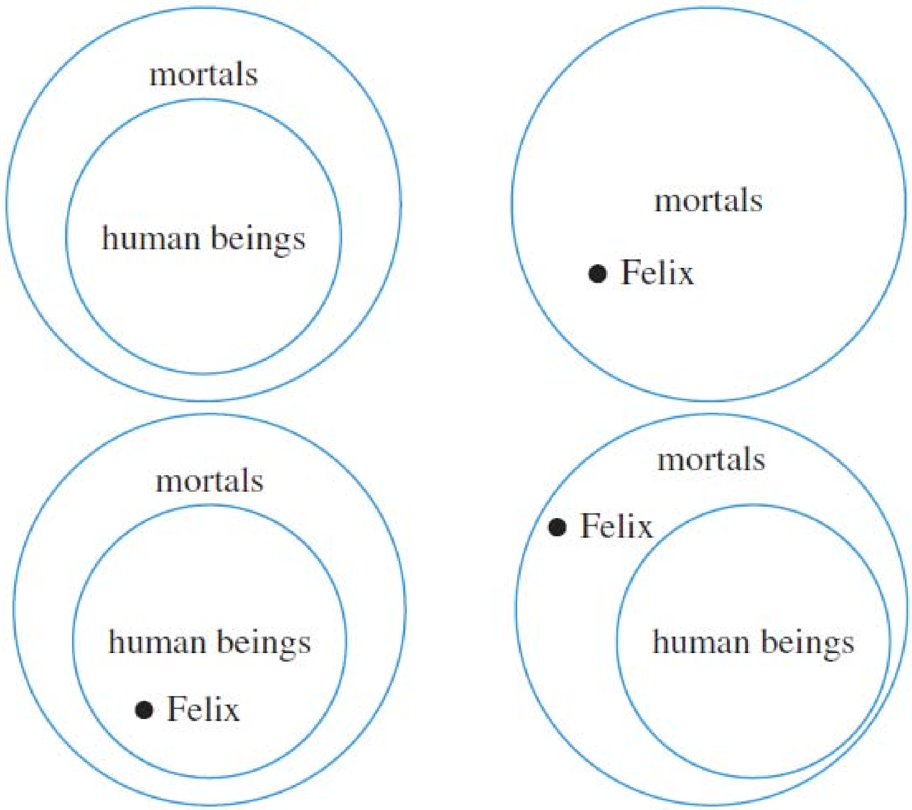


Using Diagrams to Show Invalidity

All human beings are mortal.

Felix is mortal.

∴ Felix is a human being.



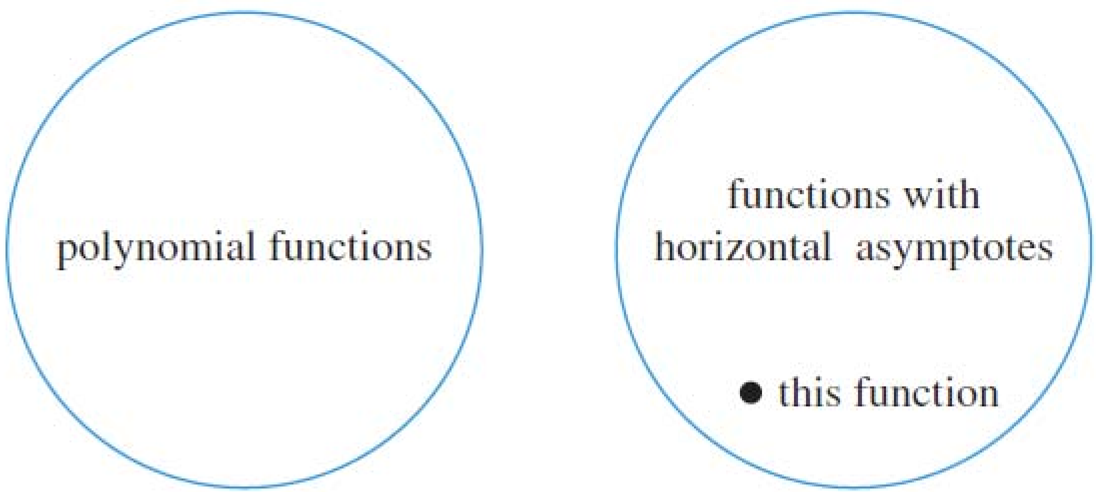
Using Diagrams to Test for Validity

- Universal modus tollens example:

  No polynomial functions have horizontal asymptotes.

  This function has a horizontal asymptote.

  ∴ This function is not a polynomial function.



Universal Transitivity

Formal Version	*Informal Version*
$\forall x, P(x) \rightarrow Q(x).$	Any x that makes P(x) true makes Q(x) true.
$\forall x, Q(x) \rightarrow R(x).$	Any x that makes Q(x) true makes R(x) true.
$\therefore \forall x, P(x) \rightarrow R(x).$	∴ Any x that makes P(x) true makes R(x) true.

- Example from Tarski's World:

∀x, if x is a triangle, then x is blue.

∀x, if x is blue, then x is to the right of all the squares.

∴ ∀x, if x is a triangle, then x is to the right of all the squares.

Converse Error (Quantified Form)

Formal Version	*Informal Version*
$\forall x$, if P(x) then Q(x).	If x makes P(x) true, then x makes Q(x) true.
Q(a) for a particular a.	a makes Q(x) true.
∴ P(a).	∴ a makes P(x) true.

invalid conclusion

Inverse Error (Quantified Form)

Formal Version	*Informal Version*
∀x, if P(x) then Q(x).	If x makes P(x) true, then x makes Q(x) true.
∼P(a), for a particular a.	a does not make P(x) true.
∴ ∼Q(a).	∴ a does not make Q(x) true.

invalid conclusion