CSE581 Computer Science Fundamentals: Theory

Professor Anita Wasilewska

LOGIC LECTURE 0

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GENERAL INFORMATION

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Course Book B1

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Logics for Computer Science: Classical and Non-Classical

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You can get the book in Hard cover, or in Electronic form Springer also has an option of providing you with chapters of your choice

https://www.springer.com/us/book/9783319925905

Book B1

I wrote **Book B1** with students on my mind so that they can read and learn by themselves, **before** coming to class For sure, it is also essential to study after the class

We start our Part 1: LOGIC Lectures from Chapter 2 It is an intuitive introduction to the Classical Propositional and Predicate Logic appropriate for somebody without much of previous knowledge of Symbolic Logic

Chapter 3 provides a generalization of classical propositional part of **Chapter 2** Formal Propositional Languages and Extensional Semantics: Classical and Many Valued

First Task when one builds a symbolic logic, or foundations of mathematics, or foundations of computer science, is to define formally a proper symbolic language

We distinguish and define two kind of languages: **propositional** and **predicate**

They are also called also zero and first order languages, respectively

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Second Task is to define formally what does it mean that formulas of a symbolic language are considered to be true, and always true i.e. we have to define a notion of a tautology

It means that we **define** what is called a **semantics** for a given **language**

The same languages can have different semantics

For example, the languages for classical and intuitionistic logics can be the same, but their the semantics are different

Third Tasks is to define a syntactical notion of a proof in a proof system based on a given language

It allows us to find out what can, or cannot be **proved** if certain axioms and rules of inference are assumed

This part of syntax is also called a proof theory

Fourth Task is to investigate the relationship between a **syntactical** notion of a **proof system** based on a given language and a **semantics** for that language

It means we establish **formal** relationship between the **syntax** and **semantics** for a given **language**

This **relationship** is established by providing answers to the following **two questions**

Fourth Task is to pose and answer the following questions

Q1: Is everything one **proves** in a given proof system **tautology** under a given semantics?

The positive answer to the question **Q1** is called **Soundness Theorem** for a given proof system and a given semantics proof system

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Such proof system is called a sound proof system

Logic Soundness Theorem

We write the Soundness Theorem symbolically as follows

Soundness Theorem (with respect to a semantics **M**) Let **S** be a proof system and **A** any formula of its language, then the following holds

IF $\vdash_S A$ THEN $\models_M A$

Logic Main Questions

Q2: Is it also possible to guarantee a **provability** in a **sound** proof system of everything we know to be a **tautology** under a given semantics?

The positive answer to the question **Q2** is called **Completeness Theorem** for a proof system under a given semantics

Such proof system is called **complete proof system** with respect to the given semantics

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Completenss Theorem

We write the Completeness Theorem symbolically as follows

CompletenessTheorem (with respect to a semantics **M**) Let **S** be a proof system and **A** any formula of its language, then the following holds

 $\vdash_{S} A$ if and only if $\models_{M} A$

Fifth Task is to develop proof systems in which a process of finding proofs can be carried fully automatically These are **automated** theorem proving systems The book presents various Gentzen Type **automated** theorem proving systems and discusses various methods of proving the **Completeness Theorem** for them

In our **Part 1: LOGIC** Lectures we cover, as an example, a theorem proving system RS and and give a constructive and very easy proof of the **Completeness Theorem**

Main Goals

The first set of Main Goals of the BOOK is to formally define and develop the above FIVE TASKS in case of Classical Propositional and Predicate Logic

The second set of **Main Goals** is to develop and discuss the FIVE TASKS for some **Non-Classical** Propositional Logics, namely for some extensional Many Valued logics, for the Intuitionistic logic, and Modal S4, S5 logics

Main Goals of the Book

The third set of Main Goals of the book is to formally define and develop the notion of a formal theory based on a given proof system for a first order logic It discusses notions of a model of a theory, its semantical and syntactical consistency and completeness The book presents some Formal Theories based on classical predicate logic

In particular it presents the Peano Arithmetic of

Natural Numbers PA and discusses and proves the

Gödel Incompleteness Theorems