

cse547,
DISCRETE MATHEMATICS

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LECTURE 4

CHAPTER 1

PART FIVE: Binary and Relaxed Binary Solutions for Generalized Josephus

Binary Solution

We proved that the **original J-recurrence**:

$$J(1) = 1, \quad J(2n) = 2J(n) - 1, \quad J(2n+1) = 2J(n) + 1 \quad \text{for } n > 1$$

has a beautiful **binary CF solution**

$$J((b_m, b_{m-1}, \dots, b_1, b_0)_2) = (b_{m-1}, b_{m-2}, \dots, b_0, b_m)_2,$$

move b_m !

where $b_m = 1$, as $n = 2^m + 1$

Question: Does the **generalized Josephus GJ** admits a similar solution?

Answer: YES.

Generalized Josephus GF

We **generalized** the function **J** to function $f: N - \{0\} \rightarrow N$ defined as follows

$$f(1) = \alpha$$

$$f(2n) = 2f(n) + \beta, \quad n \geq 1$$

$$f(2n + 1) = 2f(n) + \gamma, \quad n \geq 1$$

Observe that $J = f$ for $\alpha = 1, \beta = -1, \gamma = 1$

We call the function f a **Generalized Josephus GJ**

New Formula for GJ

We re-write the function f as follows

$$f(1) = \alpha;$$

$$f(2n + j) = 2f(n) + \beta_j$$

$$\text{for } j = 0, 1, \quad n \geq 1$$

Assume

$$k = (b_m, b_{m-1}, \dots, b_1, b_0)_2$$

We want to evaluate:

$$f(k) = f((b_m, b_{m-1}, \dots, b_1, b_0)_2)$$

Binary Representation for $k=2n$

Consider case when

$$k = 2n + 0, \quad j = 0.$$

The binary representation of $k = 2n$ is given as:

$$2n = (b_m, b_{m-1}, \dots, b_1, b_0)_2$$

$$2n = 2^m b_m + 2^{m-1} b_{m-1} + \dots + 2b_1 + b_0$$

Binary Representation for $k=2n$

We get $b_m = 1$ and $b_0 = 0$

Hence,

$$n = 2^{m-1}b_m + \dots + b_1$$

$$\mathbf{n} = (\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1)_2$$

Question: What happens when $k = 2n + 1$, $j = 1$?

Binary Representation for $k=2n+1$

Consider case when $k = 2n + j, j = 1$

The binary representation of $k=2n + 1$ is given as:

$$2n + 1 = (b_m, b_{m-1}, \dots, b_1, b_0)_2$$

$$2n + 1 = 2^m b_m + 2^{m-1} b_{m-1} + \dots + 2b_1 + b_0$$

$$b_0 = 1, b_m = 1$$

Binary Representation for $k=2n+1$

We get

$$2n + 1 = 2^m b_m + 2^{m-1} b_{m-1} + \dots + 2b_1 + 1$$

$$2n = 2^m b_m + 2^{m-1} b_{m-1} + \dots + 2b_1$$

$$n = 2^{m-1} b_m + 2^{m-1} b_{m-1} + \dots + b_1$$

$$\mathbf{n} = (\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1)_2$$

Binary Representation

We have **proved** that whether we have a binary representation of $2n = (b_m, b_{m-1}, \dots, b_1, b_0)_2$ or a binary representation of $2n+1 = (b_m, b_{m-1}, \dots, b_1, b_0)_2$, the corresponding representations of n are the same:

$$n = (b_m, b_{m-1}, \dots, b_1)_2$$

Fact

When dealing with **binary representation** we do not need to consider cases of $n \in \text{odd}$ or $n \in \text{even}$ when using our recursive formula

$$f(2n + j) = 2f(n) + \beta_j, \quad j = 0, 1$$

CF in Binary Representation

Here is our recursive formula

$$\text{RF: } f(1) = \alpha, \quad f(2n) = 2f(n) + \beta_0, \quad f(2n + 1) = 2f(n) + \beta_1$$

By the **Fact** evaluate can write RF using n in **binary representation**

$$f((b_m, b_{m-1}, \dots, b_1, b_0)_2) = 2f((b_m, b_{m-1}, \dots, b_1)_2) + \beta_{b_j},$$

where

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} \quad j = 0 \dots m - 1$$

CF in Binary Representation

We evaluate:

$$\begin{aligned} f((b_m, b_{m-1}, \dots, b_1, b_0)_2) &= 2f((b_m, b_{m-1}, \dots, b_1)_2) + \beta_{b_0} \\ &= 2(2f((b_m, b_{m-1}, \dots, b_2)_2) + \beta_{b_1}) + \beta_{b_0} \\ &= 4f((b_m, b_{m-1}, \dots, b_2)_2) + 2\beta_{b_1} + \beta_{b_0} \\ &\vdots \\ &= 2^m f((b_m)_2) + 2^{m-1} \beta_{b_{m-1}} + \dots + 2\beta_{b_1} + \beta_{b_0} \\ &= \mathbf{2^m f((1)_2)} + \mathbf{2^{m-1} \beta_{b_{m-1}}} + \dots + \mathbf{2\beta_{b_1}} + \beta_{b_0} \end{aligned}$$

CF in Binary Representation

We know that $f(1) = \alpha$

So we get (almost) CF formula

$$f((\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) = 2^m \alpha + 2^{m-1} \beta_{\mathbf{b}_{m-1}} + \dots + 2 \beta_{\mathbf{b}_1} + \beta_{\mathbf{b}_0}$$

where

$$\beta_{\mathbf{b}_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} \quad j = 0 \dots m - 1$$

Relaxed Binary CF

We **define** a **relaxed binary** representation as follows

$$2^m \alpha + 2^{m-1} \beta_{b_{m-1}} + \dots + \beta_{b_0} = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2$$

where now β_{b_k} are now **any numbers**, not only 0,1

We write the **relaxed binary CF** as

$$\begin{aligned} \mathbf{f}((\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) &= (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2 \\ \text{"normal"} &= \text{relaxed} \end{aligned}$$

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} \quad j = 0, \dots, m-1$$

Example: Original Josephus

The **GJ** function **f** becomes the **original Josephus** when $\beta_0 = -1, \beta_1 = 1$

Example

Let $n = 100$

Use the **relaxed binary CF** to show that $f(100) = 73 = J(n)$ as we have already evaluated

$$n = (1100100)_2 \\ (b_6 b_5 b_4 b_3 b_2 b_1 b_0)$$

Relaxed coordinates are

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} \quad \text{and hence}$$

$$\beta_{b_j} = \begin{cases} -1 & b_j = 0 \\ 1 & b_j = 1 \end{cases}$$

Example

We have

$$n = (1100100)_2 \\ (b_6 b_5 b_4 b_3 b_2 b_1 b_0)$$

$$\begin{aligned} f((\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) &= (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2 \\ \text{"normal"} &= \text{relaxed} \end{aligned}$$

$$\beta_{b_j} = \begin{cases} -1 & b_j = 0 \\ 1 & b_j = 1 \end{cases}$$

We evaluate

$$\begin{aligned} f(n) &= f((1100100)_2) =^{\text{relax}} (\alpha, \beta_{b_5}, \dots, \beta_{b_0}) \\ &= (1, 1, -1, -1, 1, -1, -1)_2 = 64 + 32 - 16 - 8 + 4 - 2 - 1 = 73 \end{aligned}$$

Cyclic - Shift Property

We **proved** that the original **J-recurrence**:

$$J(1) = 1, \quad J(2n) = 2J(n) - 1, \quad J(2n+1) = 2J(n) + 1 \quad \text{for } n > 1$$

has a beautiful binary **CF** solution, called **cyclic - shift property**, namely

$$J((b_m, b_{m-1}, \dots, b_1, b_0)_2) = (b_{m-1}, b_{m-2}, \dots, b_0, b_m)_2$$

We prove now that the **cyclic - shift property** holds also for the **GF** function f in the case when $\beta_0 = -1, \beta_1 = 1$, i.e.

$$f((b_m, b_{m-1}, \dots, b_1, b_0)_2) = (b_{m-1}, b_{m-2}, \dots, b_0, b_m)_2$$

We know that $b_m = 1$, so we have to prove that:

$$f(1, b_{m-1}, \dots, b_1, b_0)_2 = (b_{m-1}, b_{m-2}, \dots, b_0, 1)_2,$$

for f such that $\beta_0 = -1, \beta_1 = 1$

Cyclic - Shift Property for GJ

We have proved the **relaxed binary CF** solution for **GJ**:

$$CF : f((1, b_{m-1}, \dots, b_1, b_0)_2) = (1, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2$$

where $f(n)$ contains now **1** and **-1** as defined by

$$\beta_{b_j} = \begin{cases} -1 & b_j = 0 \\ 1 & b_j = 1 \end{cases}$$

Example

EXAMPLE

Consider $n = (1, 0, 0, 1, 0, 0, 1)_2$

By CF we have that

$$f((1, 0, 0, 1, 0, 0, 1)_2) = (1, -1, -1, 1, -1, -1, 1)_2$$

General Observation

f transforms a BLOCK of 0's in **normal binary representation** into a BLOCK of -1's in the **relaxed representation**

$$f((1, 0, 0\dots 0)_2) = (1, -1, -1\dots -1)_2$$

ONE BLOCK Transformation

We **prove** now the following relationship between **relaxed** and **normal** representation

ONE BLOCK transformation

$$(1, -1, -1, \dots, -1)_2 = (0, 0, 0, \dots, 0, 1)_2$$

Proof: Let $n = ((-1, -1, \dots, -1)_2$

$$\begin{aligned} n &= (1, -1, -1, \dots, -1)_2 \stackrel{\text{def}}{=} 2^m - 2^{m-1} - 2^{m-2} - \dots - 2^1 - 2^0 \\ &= 2^{m-1} - 2^{m-2} - \dots - 2^1 - 2^0 \\ &= 2^{m-2} - 2^{m-3} - \dots - 2^1 - 2^0 \\ &\quad \vdots \\ &= 2^1 - 2^0 \\ &= 2 - 1 \\ &= 1 = (0, 0, 0, 0, 1)_2 \end{aligned}$$

Many Blocks Transformation

Example for **TWO BLOCKS transformation** plus **binary shift**

$$\begin{aligned} f((1, 0, 0, 1, 1, 0, 0, 1)_2) &= (1, -1, -1, 1, 1, -1, -1, 1)_2 \\ &=^{1bt} (0, 0, 1, 1, 1, -1, -1, 1)_2 \\ &=^{1bt} (0, 0, 1, 1, 0, 0, 1, 1)_2 \\ &= (0, 0, 1, 1, 0, 0, 1, 1)_2 \end{aligned}$$

We know that $f((b_m, \dots, b_1, b_0)_2) = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2$
OBSERVE that each block of binary digits $(1, 0..0)_2$ is **transformed** by f into $(1, -1, \dots)_2$ and multiple applications of **one block transformation** transforms them **back** to $(1, 0..0)_2$, so

$$((\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2 =^{mbt} (b_{m-1}, \dots, b_1, b_0, 1)_2$$

where **mbt** denotes multiple BLOK transformations, and we know that $\alpha = 1$

Cyclic - Shift Property

We now evaluate:

$$\begin{aligned} f((\mathbf{1}, b_{m-1}, \dots, b_1, b_0)_2) &= (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_2 \\ &\stackrel{mbt}{=} (b_{m-1}, \dots, b_1, b_0, \mathbf{1})_2 \end{aligned}$$

This ends the proof of the **Cyclic - Shift Property** for
Generalized Josephus f with $\alpha = 1, \beta_0 = -1, \beta_1 = 1$

Exercise 1

Given

$$f(1) = 5$$

$$f(2n) = 2f(n) - 10$$

$$f(2n + 1) = 2f(n) + 83$$

Exercise 1

Evaluate $f(100)$

Solution: just apply proper formulas!

Exercise 2

Given

$$f(1) = 5$$

$$f(2n) = 3f(n) - 10$$

$$f(2n + 1) = 3f(n) + 83$$

Exercise 2

Evaluate $f(100)$

Observe that now we don't have proper formulas! They work only for base **2**!

Goal **Generalize** f and develop **new formulas** (if possible)

RADIX Representation

We **proved** while solving the Generalized Josephus that

$$\text{RF: } f(1) = \alpha, \quad f(2n + j) = 2f(n) + \beta_j$$

where $j = 0, 1$ and $n \geq 0$

has a **relaxed binary** CF formula

$$\text{CF: } \mathbf{f}((1, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) = (\alpha, \beta_{\mathbf{b}_{m-1}}, \dots, \beta_{\mathbf{b}_0})_2$$

where β_{b_j} are defined by

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} \quad j = 0, \dots, m-1$$

and where the **relaxed binary representation** is defined as

$$(\alpha, \beta_{\mathbf{b}_{m-1}}, \dots, \beta_{\mathbf{b}_0})_2 = 2^m \alpha + 2^{m-1} \beta_{\mathbf{b}_{m-1}} + \dots + \beta_{\mathbf{b}_0}$$

Relaxed Radix Representation

We **generalize GJ** as follows

$$\text{RF: } f(1) = \alpha, \quad f(2n + j) = kf(n) + \beta_j,$$

where $k \geq 2$, $j = 0, 1$ and $n \geq 0$

Exercise: **PROVE** that **RF** has a **relaxed k-representation** closed formula

$$\text{CF: } \mathbf{f}((\mathbf{1}, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) = (\alpha, \beta_{\mathbf{b}_{m-1}}, \dots, \beta_{\mathbf{b}_0})_k$$

where β_{b_j} are defined as before by

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} ; \quad j = 0, \dots, m-1$$

and where we define the **relaxed k-representation** as follows

Relaxed k-Radix Representation

Definition

A relaxed **k**- representation is defined as

$$(\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_k = \alpha k^m + k^{m-1} \beta_{b_{m-1}} + \dots + \beta_{b_0}$$

We repeat the **proof i** directly from there definition following the proof for the case **k = 2**

Proof

$$\begin{aligned} f((b_m, b_{m-1}, \dots, b_1, b_0)_2) &= kf((b_m, b_{m-1}, \dots, b_1)_2) + \beta_{b_0} \\ &= k(kf((b_m, b_{m-1}, \dots, b_2)_2) + \beta_{b_1}) + \beta_{b_0} \\ &= k^2f((b_m, b_{m-1}, \dots, b_2)_2) + k\beta_{b_1} + \beta_{b_0} \\ &= k^3f((b_m, b_{m-1}, \dots, b_3)_2) + k^2\beta_{b_2} + k\beta_{b_1} + \beta_{b_0} \\ &\quad \vdots \\ &= k^mf((b_m)_2) + k^{m-1}\beta_{b_{m-1}} + \dots + k\beta_{b_1} + \beta_{b_0} \\ &= k^m\alpha + k^{m-1}\beta_{b_{m-1}} + \dots + k^2\beta_{b_2} + k\beta_{b_1} + \beta_{b_0} \\ &= (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_1}, \beta_{b_0})_k \end{aligned}$$

$$f((b_m, b_{m-1}, \dots, b_1, b_0)_2) = (\alpha, \beta_{b_{m-1}} \dots \beta_{b_1}, \beta_{b_0})_k$$

base 2

→ base k

Example

Given RF:

$$f(1) = 5$$

$$f(2n) = 6f(n) + 3$$

$$f(2n + 1) = 6f(n) - 10$$

Evaluate: $f(100)$ by the use of the **k-representation** and closed formula

$$CF : \mathbf{f((1, b_{m-1}, \dots, b_1, b_0)_2)} = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_k$$

where β_{b_j} are defined as before by

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} ; \quad j = 0, \dots, m - 1$$

Example Solution

Given

$$f(1) = 5$$

$$f(2n) = 6f(n) + 3$$

$$f(2n + 1) = 6f(n) - 10$$

We evaluate

$$\alpha = 5$$

$$\beta_0 = 3$$

$$\beta_1 = -10$$

Example Solution

Evaluate: $f(100)$

$$\alpha = 5, \beta_0 = 3, \beta_1 = -10, k = 6, n = (1100100)_2 \\ (b_6 b_5 b_4 b_3 b_2 b_1 b_0)$$

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases}, \quad j = 0, \dots, m-1,$$

$$\beta_{b_0} = 3, \beta_{b_1} = 3, \beta_{b_2} = -10, \beta_{b_3} = 3; \beta_{b_4} = 3, \\ \beta_{b_5} = -10, \alpha = 5$$

$$f(100) = f((1100100)_2) = (5, -10, 3, 3, -10, 3, 3)_6$$

More General GJ Function

Further Generalization of GJ

RF:

$$\begin{aligned} f(i) &= \alpha_i, & i &= 1, \dots, d-1 \\ f(dn+j) &= cf(n) + \beta_j, & n &\geq 1, \quad 0 \leq j < d \end{aligned}$$

Exercise

Prove the following closed formula

CF:

$$f((\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_d) = (\alpha_{\mathbf{b}_m}, \beta_{\mathbf{b}_{m-1}} \dots \beta_{\mathbf{b}_1}, \beta_{\mathbf{b}_0})_c$$

Example

$$f(1) = 34$$

$$f(2) = 5$$

$$f(3n) = 10f(n) + 76$$

$$f(3n + 1) = 10f(n) - 2$$

$$f(3n + 2) = 10f(n) + 8$$

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \\ \beta_2 & b_j = 2 \end{cases}, \quad j = 0, \dots, d-1,$$

Example Solution

We evaluate:

$$i = 1, 2$$

$$j = 0, 1, 2$$

$$d = 3$$

$$c = 10$$

$$\alpha_1 = 34$$

$$\alpha_2 = 5$$

$$\beta_0 = 76$$

$$\beta_1 = -2$$

$$\beta_2 = 8$$

Example

Evaluate: $f(19)$

$$19 = (201)_3 = 2 \cdot 3^2 + 0 \cdot 3 + 1$$

$$\alpha_{b_2} = \alpha_2 = 5$$

$$\beta_{b_0} = \beta_0 = 76$$

$$\beta_{b_1} = \beta_1 = -2$$

$$\begin{aligned} f(19) &= f((201)_3) \\ &= (5, 76, -2)_{10} \\ &= 5 \cdot 10^2 + 76 \cdot 10 - 2 \\ &= 500 + 760 - 2 \\ &= 1258 \end{aligned}$$

Short Solution

$$\mathbf{f}((\mathbf{b}_m, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_d) = (\alpha_{\mathbf{b}_m}, \beta_{\mathbf{b}_{m-1}} \dots \beta_{\mathbf{b}_1}, \beta_{\mathbf{b}_0})_c$$

Take

$$19 = (2\ 0\ 1)_3$$

Corresponding solution is

$$(\alpha_2, \beta_0, \beta_1)_{10}$$

we evaluate $\alpha_2 = 5$, $\beta_0 = 76$, $\beta_1 = -2$ and get

Solution:

$$(5, 76, -2)_{10}$$

New Generalization of GJ

New Generalization of GJ

Problem

Use the **repertoire method** to **solve** the following yet **more general four-parameter recurrence RF**

$$h(1) = \alpha;$$

$$h(2n + 0) = 3h(n) + \gamma n + \beta_0;$$

$$h(2n + 1) = 3h(n) + \gamma n + \beta_1, \text{ for all } n \geq 1.$$

Solve means FIND a **closed formula CF** equivalent to **RF**

General Form of CF

Our RF is a FOUR parameters function and it is a **generalization** of the General Josephus GJ function f considered before

So we **guess** that now the **general form** of the CF is also a generalization of the one we already proved for GJ , i.e.

General form of CF is

$$h(n) = \alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n)$$

The **Problem** asks us to use **repertoire method** to prove that CF is **equivalent** to RF

Thinking Time

Solution requires us to develop a system of **8 equations** on $\alpha, \gamma, \beta_0, \beta_1, A(n), B(n), C(n), D(n)$ and accordingly a **4 repertoire functions!**

First : we observe that when $\gamma = 0$, we get that h becomes for Generalized Josephus function f below for $k = 3$:

$$f(1) = \alpha, \quad f(2n + j) = kf(n) + \beta_j,$$

where $k \geq 2, j = 0, 1$ and $n \geq 0$

It seems **worth to examine** the case $\gamma = 0$ **first**

Closed Formula for GJ function f

We **proved** that GJ function f has the **relaxed k -representation** closed formula

$$f((\mathbf{1}, \mathbf{b}_{m-1}, \dots, \mathbf{b}_1, \mathbf{b}_0)_2) = (\alpha, \beta_{\mathbf{b}_{m-1}}, \dots, \beta_{\mathbf{b}_0})_k$$

where β_{b_j} are defined by

$$\beta_{b_j} = \begin{cases} \beta_0 & b_j = 0 \\ \beta_1 & b_j = 1 \end{cases} ; \quad j = 0, \dots, m-1,$$

for the **relaxed k -radix representation** defined as

$$(\alpha, \beta_{\mathbf{b}_{m-1}}, \dots, \beta_{\mathbf{b}_0})_k = \alpha k^m + k^{m-1} \beta_{\mathbf{b}_{m-1}} + \dots + \beta_{\mathbf{b}_0}$$

Special Case of the function h

Consider now a **special case** of the function h for $\gamma = 0$

We know that it now has a **relaxed 3 - representation** closed formula

$$h((1, b_{m-1}, \dots, b_1, b_0)_2) = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_3$$

It means that we get

Fact 0 For any $n = (1, b_{m-1}, \dots, b_1, b_0)_2$,

$$h(n) = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_3$$

Observe that our general form of **CF** in this case becomes

$$h(n) = \alpha A(n) + \beta_0 C(n) + \beta_1 D(n)$$

We must have $h(n) = h(n)$, for all $n, n \in \mathbb{N}$ so from this and **Fact 0** we get the following **Equation 1**

Equation 1

We must have $h(n) = h(n)$, for all $n \in N$

From this and **Fact 0** we get the following

Fact 1 For any $n = (1, b_{m-1}, \dots, b_1, b_0)_2$,

$$\alpha A(n) + \beta_0 C(n) + \beta_1 D(n) = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_3$$

This provides us with the **Equation 1** for finding our general form of **CF**

Next Observation

Observe that $A(n)$ in the Original Josephus was given (and proved to be) by a formula

$$A(n) = 2^k, \text{ for all } n = 2^k + \ell, \ 0 \leq \ell < 2^k$$

We have a **similar solution** for our $A(n)$

Special Case of the function h

We evaluate now few initial values for h in case $\gamma = 0$

$$h(1) = \alpha;$$

$$\begin{aligned} h(2) &= h(2(1) + 0) = 3h(1) + \beta_0 \\ &= 3\alpha + \beta_0; \end{aligned}$$

$$\begin{aligned} h(3) &= h(2(1) + 1) = 3h(1) + \beta_1 \\ &= 3\alpha + \beta_1; \end{aligned}$$

$$\begin{aligned} h(4) &= h(2(2) + 0) = 3h(2) + \beta_0 \\ &= 9\alpha + 4\beta_0; \end{aligned}$$

Equation 2

It is pretty obvious that we do have a similar formula for $A(n)$ as on the Original Josephus **OJ**

We write it as our **Fact 2** and get our

Fact 2

For all $n = 2^k + \ell$, $0 \leq \ell < 2^k$, $n \in N - \{0\}$

$$A(n) = 3^k$$

The proof is almost identical to the one in the **GJ**, we re-write is here for our case as an exercise.

This provides us with the **Equation 2** for finding our general form of **CF**

Reminder

Reminder

We investigate the case when $\gamma = 0$, i.e. now our formulas are

$$\text{RF: } h(1) = \alpha, \quad h(2n + j) = 3h(n) + \beta_j$$

where $j = 0, 1$ and $n \geq 0$ and the closed formula is

$$\text{CF: } h(n) = \alpha A(n) + \beta_0 C(n) + \beta_1 D(n)$$

Proof of the Equation 2

Consider now a further case $\beta_0 = \beta_1 = 0$, and $\alpha = 1$, i.e.

RF : $h(1) = 1$, $h(2n) = 3h(n)$, $h(2n + 1) = 3h(n)$
and CF : $h(n) = A(n)$

We use $h(n) = A(n)$ and re-write RF in terms of $A(n)$

RA : $A(1) = 1$, $A(2n) = 3A(n)$, $A(2n + 1) = 3A(n)$

Fact Closed formula CAR for AR is:

CA : $A(n) = A(2^k + \ell) = 3^k$, $0 \leq \ell < 2^k$

Observe that this **Fact** is equivalent to the following **Fact 2**

Proof of the Fact 2

Fact 2 for all $n = 2^k + \ell$, $0 \leq \ell < 2^k$

$$A(n) = 3^k$$

Proof by induction on k

Base case: $k=0$ i.e. $n=2^0 + \ell$, $0 \leq \ell < 1$, hence $n=1$ and

RA: $A(1) = 1$, and **CA:** $A(1) = 3^0 = 1$, so we have **RA = CA**

Inductive Assumption

$$A(2^{k-1} + \ell) = A(2^{k-1} + \ell) = 3^{k-1}, \text{ for } 0 \leq \ell < 2^{k-1}$$

Inductive Thesis

$$A(2^k + l) = A(2^k + l) = 3^k, \text{ for } 0 \leq l < 2^k$$

Two cases: $n \in \text{even}$, $n \in \text{odd}$

C1: $n \in \text{even}$

$n := 2n$, and we have $2^k + \ell = 2n$ iff $\ell \in \text{even}$

Proof of the Fact 2

We evaluate n as follows

$$2n = 2^k + \ell, \quad n = 2^{k-1} + \frac{\ell}{2}$$

We use n in the **inductive step**

Observe that the **correctness** of using $\frac{\ell}{2}$ follows from that fact that $\ell \in \text{even}$, so $\frac{\ell}{2} \in \mathbb{N}$ and it can be proved formally like on the previous slides

Inductive Proof

$$A(2n) \stackrel{\text{reprn}}{=} A(2^k + \ell) \stackrel{\text{n-eval}}{=} 3A(2^{k-1} + \frac{\ell}{2}) \stackrel{\text{ind}}{=} 3 * 3^{k-1} = 3^k$$

Proof of the Fact 2

C2: $n \in \text{odd}$

$n := 2n+1$, and we have $2^k + \ell = 2n + 1$ iff $\ell \in \text{odd}$

We evaluate n as follows

$$2n + 1 = 2^k + \ell, \quad n = 2^{k-1} + \frac{\ell-1}{2}$$

We use n in the **inductive step**

Observe that the correctness of using $\frac{\ell-1}{2}$ follows from that fact that $\ell \in \text{odd}$, so $\frac{\ell-1}{2} \in \mathbb{N}$

Inductive Proof

$$A(2n + 1) \stackrel{\text{reprn}}{=} A(2^k + \ell) \stackrel{\text{n-eval}}{=} 3A(2^{k-1} + \frac{\ell-1}{2}) \stackrel{\text{ind}}{=} 3 * 3^{k-1} = 3^k$$

It ends the proof of the **Fact 2:** $A(n) = 3^k$

Repertoire Method

We return now to **original functions**:

$$\text{RF: } h(1) = \alpha, h(2n) = 3h(n) + \gamma n + \beta_0,$$

$$h(2n + 1) = 3h(n) + \gamma n + \beta_1,$$

$$\text{CF: } h(n) = \alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n)$$

We have already developed **two equations** (as stated in **Facts 1, 2**) so we need now to consider only **2 repertoire functions** to obtain **4 equations** we need to solve the problem

Repertoire Function 1

Consider a **first repertoire function** : $h(n) = 1$, for all $n \in N - \{0\}$

We put $h(n) = h(n) = 1$, for all $n \in N - \{0\}$

We have $h(1) = 1$, and $h(1) = \alpha$, so we get $\alpha = 1$

We now use $h(n) = h(n) = 1$, for all $n \in N - \{0\}$ and evaluate

$$h(2n) = 3h(n) + \gamma_0 n + \beta_0$$

$$1 = 3 + \gamma_0 n + \beta_0$$

$$0 = 2 + \gamma_0 n + \beta_0$$

$$0 = (2 + \beta_0) + \gamma_0 n$$

$$h(2n + 1) = 3h(n) + \gamma_1 n + \beta_1;$$

$$1 = 3 + \gamma_1 n + \beta_1$$

$$0 = 2 + \gamma_1 n + \beta_1$$

$$0 = (2 + \beta_1) + \gamma_1 n$$

We get $\gamma = 0$, $\beta_0 = \beta_1 = -2$

Solution 1: $\alpha = 1$, $\gamma = 0$, $\beta_0 = \beta_1 = -2$

Equation 3

The general form of **CF** is:

$$h(n) = \alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n)$$

We put $h(n) = \mathbf{h(n)}$ for the **first repertoire function**, i.e. we put $h(n) = \mathbf{h(n) = 1}$, for all $n \in N - \{0\}$, i.e.

$\alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n) = \mathbf{h(n) = 1}$, for all $n \in N - \{0\}$, where $\alpha, \gamma, \beta_0, \beta_1$ already are evaluated in the **Solution 1** as $\alpha = 1, \gamma = 0, \beta_0 = \beta_1 = -2$

We get that **CF** = **RF** if and only if the following holds

Fact 3 For all $n \in N - \{0\}$,

$$A(n) - 2C(n) - 2D(n) = 1$$

This is our **Equation 3**

Repertoire Function 2

Consider a **repertoire function 2**: $h(n) = n$, for all $n \in N - \{0\}$

We put $h(n) = h(n) = n$, for all $n \in N - \{0\}$

$h(1) = \alpha$, $h(1) = 1$ and $h(n) = h(n)$, hence $\alpha = 1$

We now use $h(n) = h(n) = n$, for all $n \in N - \{0\}$ and evaluate

$$h(2n) = 3h(n) + \gamma n + \beta_0$$

$$2n = 3n + \gamma n + \beta_0$$

$$0 = (\gamma + 1)n + \beta_0$$

$$h(2n + 1) = 3h(n) + \gamma n + \beta_1;$$

$$2n + 1 = 3n + \gamma n + \beta_1$$

$$0 = (\gamma + 1)n + (\beta_1 - 1)$$

We get $\gamma = -1$, $\beta_0 = 0$, $\beta_1 = 1$ and

Solution 2: $\alpha = 1$, $\gamma = -1$, $\beta_0 = 0$, $\beta_1 = 1$

Equation 4

$$\text{CF: } h(n) = \alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n)$$

We evaluate **CF** for $h(n) = \mathbf{h(n) = n}$, for all $n \in N - \{0\}$ and for the **Solution 2**: $\alpha = 1, \gamma = -1, \beta_0 = 0, \beta_1 = 1$ and get

CF = **RF** if and only if the following holds

Fact 4 For all $n \in N - \{0\}$

$$A(n) - B(n) + D(n) = n$$

This is our **Equation 4**

Repertoire Method: System of Equations

We obtained the following system of **4 equations** on $A(n)$, $B(n)$, $C(n)$, $D(n)$

1. $\alpha A(n) + \beta_0 C(n) + \beta_1 D(n) = (\alpha, \beta_{b_{m-1}}, \dots, \beta_{b_0})_3$
2. $A(n) = 3^k$
3. $A(n) - 2C(n) - 2D(n) = 1$
4. $A(n) - B(n) + D(n) = n$

We solve it on $A(n)$, $B(n)$, $C(n)$, $D(n)$ and put the solution into $h(n) = \alpha A(n) + \gamma B(n) + \beta_0 C(n) + \beta_1 D(n)$