

CSE 541 - Logic in Computer Science

Sample Natural Deduction Proofs

Exercise 1.2.1(i). Proof of $(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \wedge q \rightarrow r$, first in sequent notation

1	$p \wedge q \vdash p \wedge q$	axiom
2	$(p \rightarrow r) \wedge (q \rightarrow r) \vdash (p \rightarrow r) \wedge (q \rightarrow r)$	axiom
3	$p \wedge q \vdash p$	$\wedge e_1$ 1
4	$(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \rightarrow r$	$\wedge e_1$ 2
5	$(p \rightarrow r) \wedge (q \rightarrow r), p \wedge q \vdash r$	$\rightarrow e$ 3,4
6	$(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \wedge q \rightarrow r$	$\rightarrow i$ 5

and then using boxes

1	$(p \rightarrow r) \wedge (q \rightarrow r)$	premise
2	$p \wedge q$	assumption
3	p	$\wedge e_1$ 2
4	$p \rightarrow r$	$\wedge e_1$ 1
5	r	$\rightarrow e$ 3,4
6	$p \wedge q \rightarrow r$	$\rightarrow i$ 2-5

Exercise 1.2.1(j). Proof of $q \rightarrow r \vdash (p \rightarrow q) \rightarrow (p \rightarrow r)$:

1	$q \rightarrow r$	premise
2	$p \rightarrow q$	assumption
3	p	assumption
4	q	$\rightarrow e$ 3,2
5	r	$\rightarrow e$ 4,1
6	$p \rightarrow r$	$\rightarrow i$ 3-5
7	$(p \rightarrow q) \rightarrow (p \rightarrow r)$	$\rightarrow i$ 2-6

Exercise 1.2.1(q). Proof of $\vdash q \rightarrow (p \rightarrow (p \rightarrow (q \rightarrow p)))$:

1	q	assumption
2	p	assumption
3	p	assumption
4	q	assumption
5	p	copy 3
6	$q \rightarrow p$	$\rightarrow i$ 4-5
7	$p \rightarrow (q \rightarrow p)$	$\rightarrow i$ 3-6
8	$p \rightarrow (p \rightarrow (q \rightarrow p))$	$\rightarrow i$ 2-7
9	$q \rightarrow (p \rightarrow (p \rightarrow (q \rightarrow p)))$	$\rightarrow i$ 1-8

Exercise 1.2.2(e). Proof of $p \rightarrow q \vee r, \neg q, \neg r \vdash \neg p$ without using Modus Tollens:

1	$p \rightarrow q \vee r$	premise
2	$\neg q$	premise
3	$\neg r$	premise
4	p	assumption
5	$q \vee r$	$\rightarrow$ e 4,1
6	q	assumption
7	$\perp$	$\neg$ e 6,2
8	r	assumption
9	$\perp$	$\neg$ e 8,3
10	$\perp$	$\vee$ e 5,6-7,8-9
11	$\neg p$	$\neg$ i 4-10

Exercise 1.2.3(b). Proof of $\neg p \vdash p \rightarrow q$, first in sequent notation

1	$p \vdash p$	axiom
2	$\neg p \vdash \neg p$	axiom
3	$p, \neg p \vdash \perp$	$\neg$ e 1,2
4	$p, \neg p \vdash q$	$\perp$ e 3
5	$\neg p \vdash p \rightarrow q$	$\rightarrow$ i 4

and then using boxes

1	$\neg p$	axiom
2	p	assumption
3	$\perp$	$\neg$ e 1,2
4	q	$\perp$ e 3
5	$p \rightarrow q$	$\rightarrow$ i 2-4

Exercise 1.2.3(e). Proof of $\neg(p \rightarrow q) \vdash q \rightarrow p$:

1	$\neg(p \rightarrow q) \vdash \neg(p \rightarrow q)$	axiom
2	$p, q \vdash q$	axiom
3	$q \vdash (p \rightarrow q)$	$\rightarrow$ i 2
4	$\neg(p \rightarrow q), q \vdash \perp$	$\neg$ e 1,3
5	$\neg(p \rightarrow q), q \vdash p$	$\perp$ e 4
6	$\neg(p \rightarrow q) \vdash q \rightarrow p$	$\rightarrow$ i 5

Exercise 1.2.3(o). Proof of $\neg(\neg p \vee \neg q) \vdash p \wedge q$:

1	$\neg(\neg p \vee \neg q)$	premise
2	$\neg p$	assumption
3	$\neg p \vee \neg q$	$\vee i_1$ 2
4	$\perp$	$\neg e$ 1,3
5	p	PBC 2-4
6	$\neg q$	assumption
7	$\neg p \vee \neg q$	$\vee i_2$ 6
8	$\perp$	$\neg e$ 1,7
9	q	PBC 6-8
9	$p \wedge q$	$\wedge i$ 5,9

Exercise 1.2.3(q). Proof of $\vdash (p \rightarrow q) \vee (q \rightarrow r)$:

1	$q \vee \neg q$	LEM
2	q	assumption
3	p	assumption
4	q	copy 2
5	$p \rightarrow q$	$\rightarrow i$ 3-4
6	$(p \rightarrow q) \vee (q \rightarrow r)$	$\vee i$ 5
7	$\neg q$	assumption
8	q	assumption
9	$\perp$	$\neg e$ 7,8
10	r	$\perp e$ 9
11	$q \rightarrow r$	$\rightarrow i$ 8-10
12	$(p \rightarrow q) \vee (q \rightarrow r)$	$\vee i$ 11
13	$(p \rightarrow q) \vee (q \rightarrow r)$	$\vee e$ 1,2-6,7-12

Exercise 1.2.3(u). Proof of $(s \rightarrow p) \vee (t \rightarrow q) \vdash (s \rightarrow q) \vee (t \rightarrow p)$.

1	$\vdash p \vee \neg p$	LEM
2	$p, t \vdash p$	axiom
3	$p \vdash t \rightarrow p$	$\rightarrow$ i 2
4	$p \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ i 3
5	$\neg p \vdash \neg p$	axiom
6	$s \rightarrow p \vdash s \rightarrow p$	axiom
7	$s \rightarrow p, \neg p \vdash \neg s$	MT 6,5
8	$s \vdash s$	axiom
9	$s \rightarrow p, \neg p, s \vdash \perp$	$\neg$ e 7,8
10	$s \rightarrow p, \neg p, s \vdash q$	$\perp$ e 9
11	$s \rightarrow p, \neg p \vdash s \rightarrow q$	$\rightarrow$ i 10
12	$s \rightarrow p, \neg p \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ i 11
13	$s \rightarrow p \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ e 1,4,12
14	$\vdash q \vee \neg q$	LEM
15	$q, s \vdash q$	axiom
16	$q \vdash s \rightarrow q$	$\rightarrow$ i 15
17	$q \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ i 16
18	$\neg q \vdash \neg q$	axiom
19	$t \rightarrow q \vdash t \rightarrow q$	axiom
20	$t \rightarrow q, \neg q \vdash \neg t$	MT 19,18
21	$t \vdash t$	axiom
22	$t \rightarrow q, \neg q, t \vdash \perp$	$\neg$ e 20,21
23	$t \rightarrow q, \neg q, t \vdash q$	$\perp$ e 22
24	$t \rightarrow q, \neg q \vdash s \rightarrow q$	$\rightarrow$ i 23
25	$t \rightarrow q, \neg q \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ i 24
26	$t \rightarrow q \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ e 14,17,25
27	$(s \rightarrow p) \vee (t \rightarrow q) \vdash (s \rightarrow p) \vee (t \rightarrow q)$	axiom
28	$(s \rightarrow p) \vee (t \rightarrow q) \vdash (s \rightarrow q) \vee (t \rightarrow p)$	$\vee$ e 27,13,26

Note that lines 1–13 and 14–26 represent similar subproofs.

Exercise 1.2.5(a). Proof of $\vdash ((p \rightarrow q) \rightarrow q) \rightarrow ((q \rightarrow p) \rightarrow p)$ in sequent style notation,

1	$p, \neg p \vdash q$	lemma, cf. 1.2.3(b)
2	$\neg p \vdash p \rightarrow q$	$\rightarrow$ i 1
3	$(p \rightarrow q) \rightarrow q \vdash (p \rightarrow q) \rightarrow q$	axiom
4	$(p \rightarrow q) \rightarrow q, \neg p \vdash q$	$\rightarrow$ e 3,2
5	$q \rightarrow p \vdash q \rightarrow p$	axiom
6	$\neg p \vdash \neg p$	axiom
7	$q \rightarrow p, \neg p \vdash \neg q$	MT 6,5
8	$(p \rightarrow q) \rightarrow q, q \rightarrow p, \neg p \vdash \perp$	$\neg$ e 7,4
9	$(p \rightarrow q) \rightarrow q, q \rightarrow p \vdash p$	PBC 8
10	$(p \rightarrow q) \rightarrow q \vdash (q \rightarrow p) \rightarrow p$	$\rightarrow$ i 9
11	$\vdash ((p \rightarrow q) \rightarrow q) \rightarrow ((q \rightarrow p) \rightarrow p)$	$\rightarrow$ i 10

and a different proof of $(p \rightarrow q) \rightarrow q \vdash (q \rightarrow p) \rightarrow p$ using boxes,

1	$(p \rightarrow q) \rightarrow q$	premise
2	$q \rightarrow p$	assumption
3	$(p \rightarrow q) \vee \neg(p \rightarrow q)$	LEM
4	$p \rightarrow q$	assumption
5	q	$\rightarrow$ e 1,4
6	p	$\rightarrow$ e 2,5
7	$\neg(p \rightarrow q)$	assumption
8	$\neg p$	assumption
9	p	assumption
10	$\perp$	$\neg$ e 8,9
11	q	$\perp$ e 10
12	$p \rightarrow q$	$\rightarrow$ i 9-11
13	$\perp$	$\neg$ e 4,12
14	p	$\neg$ i 5-13
15	p	$\vee$ e 4-6,7-14
16	$(q \rightarrow p) \rightarrow p$	$\rightarrow$ i 2-15

(courtesy of Shang Yang, Spring 2007).