

CSE/MAT371 QUIZ 1 SOLUTIONS Fall 2018

QUESTION 1

Write the following natural language statement:

From the fact that each natural number is greater than zero we deduce that it is not possible that Anne is a boy or, if it is possible that Anne is not a boy, then it is necessary that it is not true that each natural number is greater than zero

in the following two ways

1. As a formula $A_1 \in \mathcal{F}_1$ of a language $\mathcal{L}_{\{\neg, \Box, \Diamond, \cup, \Rightarrow\}}$, 2. As a formula $A_2 \in \mathcal{F}_2$ of a language $\mathcal{L}_{\{\neg, \cup, \Rightarrow\}}$

Solution of 1.

Propositional Variables: a, b

a denotes statement: *each natural number is greater than zero*,

b denotes a statement: *Anne is a boy*

Formula is $(a \Rightarrow (\neg \Diamond b \cup (\Diamond \neg b \Rightarrow \Box \neg a)))$

Solution of 2.

Propositional Variables: a, b, c, d

a denotes statement: *each natural number is greater than zero*,

b denotes a statement: *it is not possible that Anne is a boy*,

c denotes a statement: *it is possible that Anne is not a boy*, and

d denotes a statement: *it is necessary that it is not true that each natural number is greater than zero*

Formula is $(a \Rightarrow (\neg b \cup (c \Rightarrow d)))$

QUESTION 2

Circle formulas that are propositional/ predicate **tautologies**

1. $\mathcal{S}_1 = \{ (A \Rightarrow (A \cup \neg B)), ((a \Rightarrow b) \cap (a \Rightarrow c)) \Rightarrow (a \Rightarrow b), ((a \Rightarrow b) \cup (a \cap \neg b)), (A \cup (A \Rightarrow B)), (a \cup \neg b) \}$

Solution

$\neq (a \cup \neg b)$, all other formulas are tautologies

2. $\mathcal{S}_2 = \{ (\forall x A(x) \Rightarrow \exists x A(x)), (\forall x (\neg P(x, y) \cup P(x, y)) \Rightarrow \exists x P(x, y)), ((\exists x A(x) \cap \exists x B(x)) \Rightarrow \exists x (A(x) \cap B(x))) \}$

Solution

$\models (\forall x A(x) \Rightarrow \exists x A(x))$, all other formulas are NOT tautologies

QUESTION 3

Here is a mathematical statement **S**:

For all real numbers x the following holds: If $x < 0$, then there is a natural number n , such that $x + n < 0$.

1. Re-write **S** as a symbolic mathematical statement **SF** that only uses mathematical and logical symbols.
2. Translate the symbolic statement **SF** into to a corresponding formula with **restricted quantifiers**
3. Translate your **restricted domain** quantifiers logical formula into a correct formula A of the predicate language \mathcal{L}

Solution

The statement **S** is:

"For all real numbers x the following holds: If $x < 0$, then there is a natural number n , such that $x + n < 0$."

S becomes a symbolic mathematical statement **SF** $\forall_{x \in R}(x < 0 \Rightarrow \exists_{n \in N} x + n < 0)$

We write $R(x)$ for $x \in R$, $N(y)$ for $n \in N$, a constant c for the number 0

We use $L \in \mathbf{P}$ to denote the relation $<$, we use $f \in \mathbf{F}$ to denote the function $+$.

The statement $x < 0$ becomes an **atomic formula** $L(x, c)$.

The statement $x + n < 0$ becomes an **atomic formula** $L(f(x, y), c)$.

The symbolic mathematical statement **SF** $\forall_{x \in R}(x < 0 \Rightarrow \exists_{n \in N} x + n < 0)$

becomes a **restricted quantifiers** formula $\forall_{R(x)}(L(x, c) \Rightarrow \exists_{N(y)}L(f(x, y), c))$.

We apply now the **transformation rules** and get a corresponding formula $A \in \mathcal{F}$:

$\forall x(N(x) \Rightarrow (L(x, c) \Rightarrow \exists y(N(y) \cap L(f(x, y), c))))$.

QUESTION 4

Given a formula $A : \forall x \exists y P(f(x, y), c)$ of the predicate language \mathcal{L} , and

two **model structures** $\mathbf{M}_1 = (Z, I_1)$, $\mathbf{M}_2 = (N, I_2)$ with the interpretations defined as follows.

$P_{I_1} := , \quad f_{I_1} : +, \quad c_{I_1} : 0$ and $P_{I_2} := >, \quad f_{I_2} : \cdot, \quad c_{I_2} : 0$.

1. Show that $\mathbf{M}_1 \models A$

Solution

$A_{I_1} : \forall_{x \in Z} \exists_{y \in Z} x + y = 0$ is a **true** statement;

For each $x \in Z$ exists $y = -x$ and $-x \in Z$ and $x - x = 0$.

2. Show that $\mathbf{M}_2 \not\models A$.

Solution

$A_{I_2} : \forall_{x \in N} \exists_{y \in N} x \cdot y > 0$ is a **false** statement for $x = 0$.

QUESTION 5

Given a formula $A = (\neg a \Rightarrow (\neg b \cup (b \Rightarrow \neg c)))$

and its **restricted model** $v_A : \{a, b, c\} \rightarrow \{T, F\}$, $v_A(a) = T$, $v_A(b) = T$, $v_A(c) = F$

Extend v_A to the set of all propositional variables VAR to obtain 2 different, non restricted **models** for A

Solution

Model w_1 is a function

$w_1 : VAR \rightarrow \{T, F\}$ such that

$w_1(a) = v_A(a) = T$, $w_1(b) = v_A(b) = T$,

$w_1(c) = v_A(c) = F$, and $w_1(x) = T$, for all $x \in VAR - \{a, b, c\}$

Model w_2 is defined by a formula

$w_2(a) = v_A(a) = T$, $w_2(b) = v_A(b) = T$,

$w_2(c) = v_A(c) = F$, and $w_2(x) = F$, for all $x \in VAR - \{a, b, c\}$

REMINDER: we define H semantics operations \cup and \cap as $x \cup y = \max\{x, y\}$, $x \cap y = \min\{x, y\}$

the implication and negation are defined as

$$x \Rightarrow y = \begin{cases} T & \text{if } x \leq y \\ y & \text{otherwise} \end{cases}$$
$$\neg x = x \Rightarrow F$$

QUESTION 6

We know that $v : VAR \rightarrow \{F, \perp, T\}$ is such that $v^*((a \cap \neg b) \Rightarrow (a \Rightarrow c)) = \perp$ under H semantics.

Evaluate $v^*((b \Rightarrow a) \Rightarrow b)$. You can use SHORTHAND notation.

Solution

The Truth Tables for Implication and Negation are:

H-Implication

\Rightarrow	F	\perp	T
F	T	T	T
\perp	F	T	T
T	F	\perp	T

H Negation

\neg	F	\perp	T
	T	F	F

We KNOW that

$v^*((a \cap \neg b) \Rightarrow (a \Rightarrow c)) = \perp$ under **H** semantics if and only if (we use shorthand notation)

iff $(a \cap \neg b) = T$ and $(a \Rightarrow c) = \perp$

iff $a = T, b = F,$ and $(T \Rightarrow c) = \perp$

iff $a = T, b = F, c = \perp$

Hence we evaluate

$v^*((b \Rightarrow a) \Rightarrow b) = (F \Rightarrow T) \Rightarrow F = T \Rightarrow F = F$