

**CSE371 PRACTICE MIDTERM 1 Fall 2017**  
**(10 extra pts)**

**NAME**

**ID:**

**MY POINTS ARE:**

TAKE test as a practice - and **correct it yourself** to see how much points you would get.

Solutions to all problems and questions are somewhere on our webpage! so you CAN correct yourself- but do it **ONLY AFTER you complete it all by yourself.**

This is the **goal** of the PRACTICE TEST!

**PLEASE SUBMIT your SOLUTIONS that have been CORRECTED BY YOU - Write corrections in RED. You WILL GET 10 points for THAT!** even if all problems you solved were first wrong- and then CORRECTED!

**Write a sum of POINTS** you give yourself for your solutions -after you check your answers for corrections.

The **real midterm will have less problems**; I will make sure you will be able to complete it within 1 hour and 15 minutes.

**BRING YOUR solved-corrected TEST to class on Monday, October 23**

**PART ONE: DEFINITIONS (10pts)**

**All Definitions** are for language  $\mathcal{L} = \mathcal{L}_{\{\neg, \Rightarrow, \cup, \cap\}}$  and **classical semantics**

Write carefully the following DEFINITIONS

**D1.** Extentional Connectives

**D2.** Given the truth assignment  $v : VAR \rightarrow \{T, F\}$ . Write the definition of its **extension**  $v^*$  to the set  $\mathcal{F}$  of all formulas of  $\mathcal{L}$

**D3.** Restricted MODEL for a given formula  $A \in \mathcal{F}$

**D4.** Proof System S

**D5.** Formal proof from  $\Gamma$  in a system S

**D6.** Sound rule of inference in a system S

**D7.** Sound proof system S

**D8.** Soundness and Completeness Theorem for S (classical semantics)

**D9.** A non-empty set  $\mathcal{G} \subseteq \mathcal{F}$  **consistent** (classical semantics)

**D10.** A non-empty set  $\mathcal{G} \subseteq \mathcal{F}$  **inconsistent** (classical semantics)

## PART TWO: PROBLEMS (65pts)

**Problem 1 (5pts), all other problems (10pts)**

### Problem 1

Given a mathematical statement **S** written with logical symbols

$$(\exists x \in N \ x \leq 5 \ \cap \ \forall y \in Z \ y = 0)$$

1. Translate **S** it into a proper logical formula that **uses** the restricted domain quantifiers.
2. Translate your restricted quantifiers formula into a correct formula **without** restricted domain quantifiers.

Write a **short** solution.

### Problem 2

Given a formula  $A : \forall x \exists y P(f(x, y), c)$  of the predicate language  $\mathcal{L}$ , and

two **model structures**  $\mathbf{M}_1 = (Z, I_1)$ ,  $\mathbf{M}_2 = (N, I_2)$  with the interpretations defined as follows.

$$P_{I_1} : =, \quad f_{I_1} : +, \quad c_{I_1} : 0 \quad \text{and} \quad P_{I_2} : >, \quad f_{I_2} : \cdot, \quad c_{I_2} : 0.$$

1. Show that  $\mathbf{M}_1 \models A$

2. Show that  $\mathbf{M}_2 \not\models A$

### Problem 3

$S$  is the following proof system:

$$S = ( \mathcal{L}_{\{\Rightarrow, \cup, \neg\}}, \mathcal{F}, LA = \{(A \Rightarrow (A \cup B))\} \ (r1), (r2) )$$

**Rules** of inference:

$$(r1) \frac{A ; B}{(A \cup \neg B)}, \quad (r2) \frac{A ; (A \cup B)}{B}$$

1. Verify whether  $S$  is sound/not sound under classical semantics.
  
2. Verify whether  $S$  is sound/not sound under K semantics.
  
3. Find a formal proof of  $\neg(A \Rightarrow (A \cup B))$  in  $S$ , i. e. show that  $\vdash_S \neg(A \Rightarrow (A \cup B))$
  
4. Does above point **3.** prove that  $\models \neg(A \Rightarrow (A \cup B))$ ?

**Problem 4**

1. Given a formula  $A = ((a \cap \neg c) \Rightarrow (\neg a \cup b))$  of a language  $\mathcal{L}_{\{\neg, \cap, \cup, \Rightarrow\}}$ .

**Find** a formula  $B$  of a language  $\mathcal{L}_{\{\neg, \Rightarrow\}}$ , such that  $A \equiv B$ . **List** all proper logical equivalences used at at each step.

2. Prove that  $\mathcal{L}_{\{\neg, \cap, \cup, \Rightarrow\}} \equiv \mathcal{L}_{\{\neg, \Rightarrow\}}$ .

### Problem 5

Consider the Hilbert system  $H_1 = (\mathcal{L}_{\{\Rightarrow\}}, \mathcal{F}, \{A1, A2\}, (MP) \frac{A; (A \Rightarrow B)}{B})$  where

$A1; (A \Rightarrow (B \Rightarrow A)), \quad A2 : ((A \Rightarrow (B \Rightarrow C)) \Rightarrow ((A \Rightarrow B) \Rightarrow (A \Rightarrow C)))$  and A, B are any formulas from  $\mathcal{F}$ .

Use **Deduction Theorem** to prove  $(A \Rightarrow B), (B \Rightarrow C) \vdash_{H_1} (A \Rightarrow C)$ .

Write comments how each step was obtained.

### Problem 6

**Complete** the steps  $B_1, \dots, B_5$  of the formal proof in  $H_2$  of  $(B \Rightarrow \neg\neg B)$  by writing all details for each step of the proof.

**You can use** the following already proved facts:

**F1**  $(A \Rightarrow B), (B \Rightarrow C) \vdash_{H_2} (A \Rightarrow C)$

**F2**  $\vdash_{H_2} (\neg\neg B \Rightarrow B)$

Here are the steps

$B_1 = ((\neg\neg\neg B \Rightarrow \neg B) \Rightarrow ((\neg\neg\neg B \Rightarrow B) \Rightarrow \neg\neg B))$

$B_2 = (\neg\neg\neg B \Rightarrow \neg B)$

$$B_3 = ((\neg\neg\neg B \Rightarrow B) \Rightarrow \neg\neg B)$$

$$B_4 = (B \Rightarrow (\neg\neg\neg B \Rightarrow B))$$

$$B_5 = (B \Rightarrow \neg\neg B)$$

**Problem 7** We define, for  $A, b_1, b_2, \dots, b_n$  and truth assignment  $v$  a corresponding formulas  $A'$ ,  $B_1, B_2, \dots, B_n$  as follows:

$$A' = \begin{cases} A & \text{if } v^*(A) = T \\ \neg A & \text{if } v^*(A) = F \end{cases} \quad B_i = \begin{cases} b_i & \text{if } v(b_i) = T \\ \neg b_i & \text{if } v(b_i) = F \end{cases}$$

We proved the following **Main Lemma**: For any formula  $A = A(b_1, b_2, \dots, b_n)$  and any truth assignment  $v$ ,

if  $A', B_1, B_2, \dots, B_n$  are corresponding formulas defined above, then  $B_1, B_2, \dots, B_n \vdash A'$ .

Let  $A$  be a formula  $((\neg a \Rightarrow \neg b) \Rightarrow (b \Rightarrow a))$ , and let  $v$  be such that  $v(a) = T$ ,  $v(b) = F$ .

Write what **Main Lemma** asserts for the formula  $A$ .