CSE303 Q2 SOLUTIONS

PART 1   YES/NO QUESTIONS

Circle the correct answer. Write SHORT justification.

1. The set $F$ of final states of any non-deterministic finite automaton is always non-empty
   \textbf{Justify}: The definition says that $F$ is a finite set, i.e. can be empty, hence for some $M$, $L(M) = \emptyset$ \hspace{1em} \textbf{n}

2. Given an automaton $M = (K, \Sigma, \delta, s, F)$, a binary relation $\vdash_M \subseteq (K \times \Sigma^*) \times (K \times \Sigma^*)$ is a transition relation iff the following condition holds
   \textbf{Justify}: Proper condition is: \hspace{1em} \textbf{y}

3. If $M = (K, \Sigma, \Delta, s, F)$ is a non-deterministic as defined in the book, then $M$ is also non-deterministic, as defined in the lecture.
   \textbf{Justify}: $\Sigma \cup \{e\} \subseteq \Sigma^*$ \hspace{1em} \textbf{y}

PART 2  Very Short Questions

Q1: $M_1$ has components: $K = \{q_0, q_1, q_2\}$, $s = q_0$, $\Sigma = \{a, b\}$, $F = \{q_1\}$
   \[ \delta = \{(q_0, a, q_1), (q_1, a, q_1), (q_0, b, q_2)\} \]

Solution

1. $M_1$ is non-deterministic; $\delta$ is not a function with the domain $K \times \Sigma$. It can be completed to a function by adding some trap states. But the trap states information was not stated in the problem - so $M_1$ is NDFA
2. $L(M_1) = aa^*$

Q2: $M_2$ has $K = \{q_0, q_1, q_2\}$, $\Sigma = \{a, b\}$, $F = \{q_1, q_2\}$
   \[ \Delta = \{(q_0, a, q_1), (q_0, b, q_2), (q_1, a, q_1), (q_1, b, q_2), (q_2, ab, q_2)\} \]

Solution

$M_2$ is \textbf{not} DFA and is \textbf{not} NDFA because $M_2$ is \textbf{NOT} an automaton It does not have the INITIAL state.

$L(M_2)$ does not exist - as languages are defined for automatas

PART 3 PROBLEMS

QUESTION 1

Components of an automaton $M$ are:
   \[ K = \{q_0, q_1, q_2, q_3\}, \quad s = q_0, \quad \Sigma = \{a, b\}, \quad F = \{q_1, q_3\} \quad \text{and} \quad \delta = \{(q_0, a, q_1), (q_1, b, q_2), (q_2, a, q_2), (q_2, b, q_3)\} \]

\textbf{Complete} the components to a full definition of a deterministic $M$ by adding trap state(s)

\textbf{Write} a regular expression defining $L(M)$

Solution

\textbf{YOU ONLY DRAW the diagram} - and compare with the listed components below

\textbf{You do NOT write components}
We add a new state $q_4$ and extend $\delta$ to function $\delta_1$ such that

$$\delta_1 : (K \cup \{q_4\}) \times \Sigma \rightarrow K \cup \{q_4\}$$

as follows

$$\delta_1 = \delta \cup \{(q_0, b, q_4), (q_1, a, q_4), (q_3, a, q_4), (q_3, b, q_4)\}$$

$$L(M) = a \cup aba^*b$$

**QUESTION 2**

Let $M = (K, \Sigma, s, \Delta, F)$ for $K = \{q_0, q_1, q_2\}$, $s = q_0$, $\Sigma = \{a, b, c\}$, $F = \{q_1, q_2\}$ and

$$\Delta = \{(q_0, abc, q_0), (q_0, ab, q_1), (q_0, b, q_2)\}$$

**Draw the diagram** of an automaton $M'$ such that $M' \equiv M$ and $M'$ is defined by the BOOK definition.

**Solution**

**ONLY DRAW the diagram** - and compare with the listed BELOW components

**DO NOT write components**

We apply the "stretching" technique to $M$ and the new $M'$ is as follows.

$$M' = (K \cup \{p_1, p_2, p_3\}, \Sigma, s = q_0, \Delta', F' = F)$$

$$\Delta' = \{(q_0, b, q_2), (q_0, a, p_1), (p_1, b, p_2), (p_2, c, q_0), (q_0, b, q_2), (q_0, a, p_3), (p_3, b, q_1)\}$$