

# Database and Distributed Computing Foundations of Blockchains

Sujaya Maiyya, Victor Zakhary, Mohammad Javad Amiri, Divyakant Agrawal, Amr El Abbadi

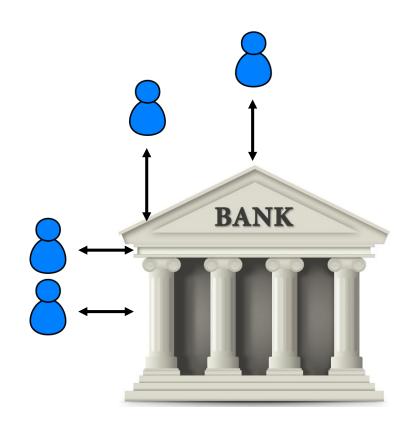
{sujaya-maiyya, victorzakhary, amiri, agrawal, amr}@cs.ucsb.edu

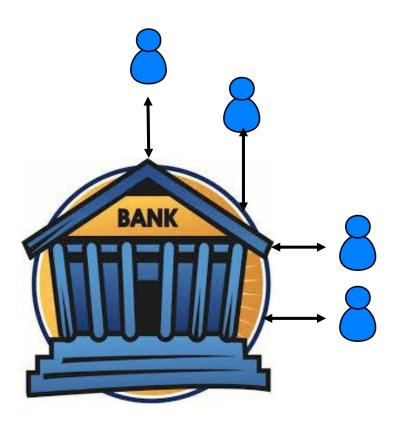


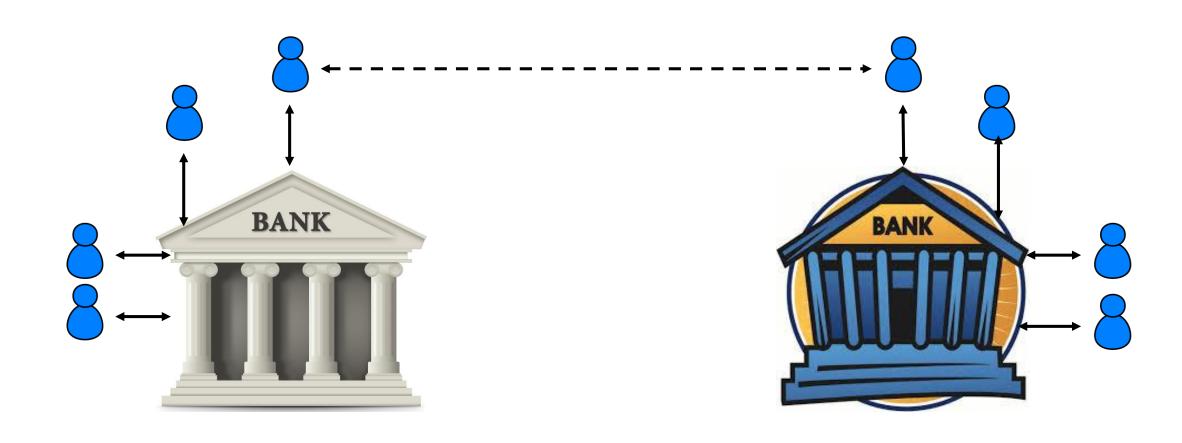


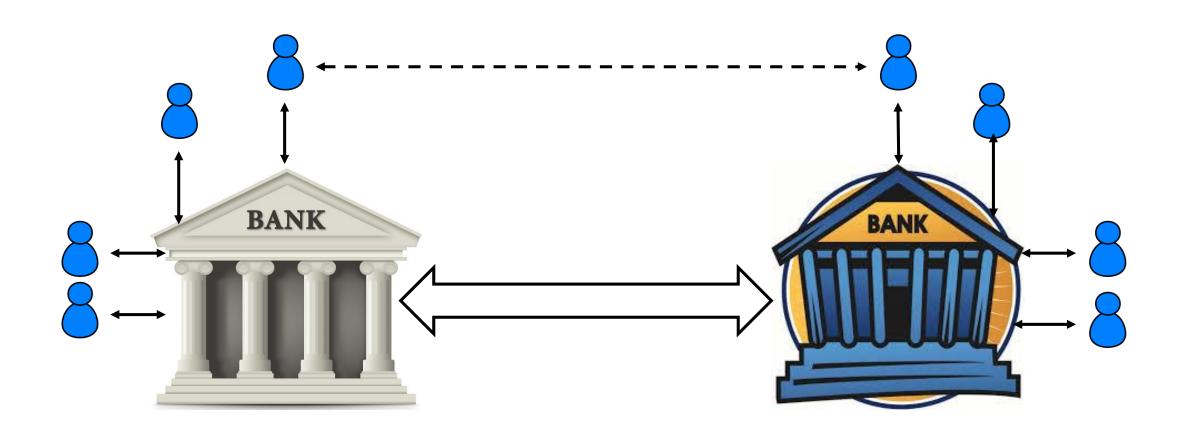


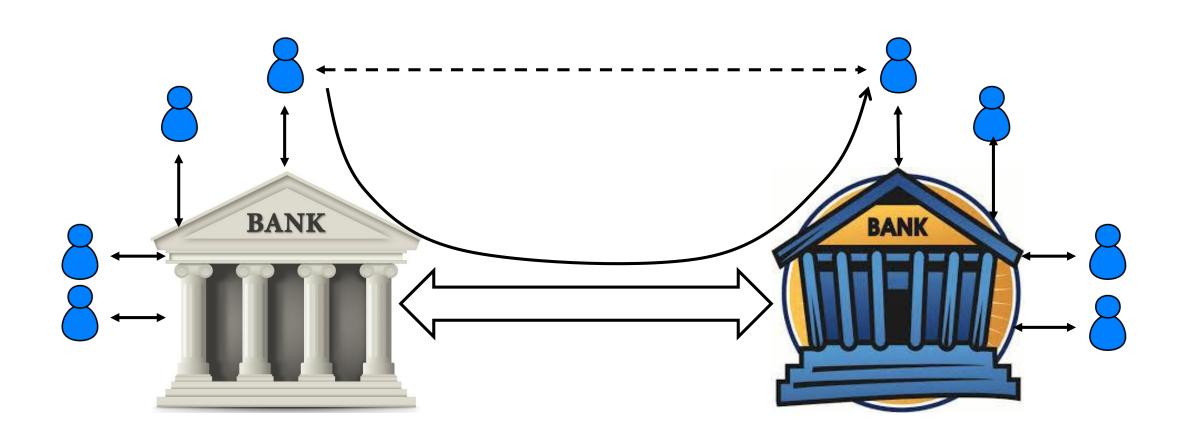












## Traditional Banking Systems

From Database and Distributed Computing Perspective



- From Database and Distributed Computing Perspective
- Identities and Signatures





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  - You are your signature [ID, username and password]





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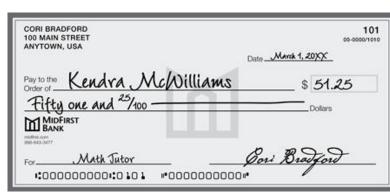




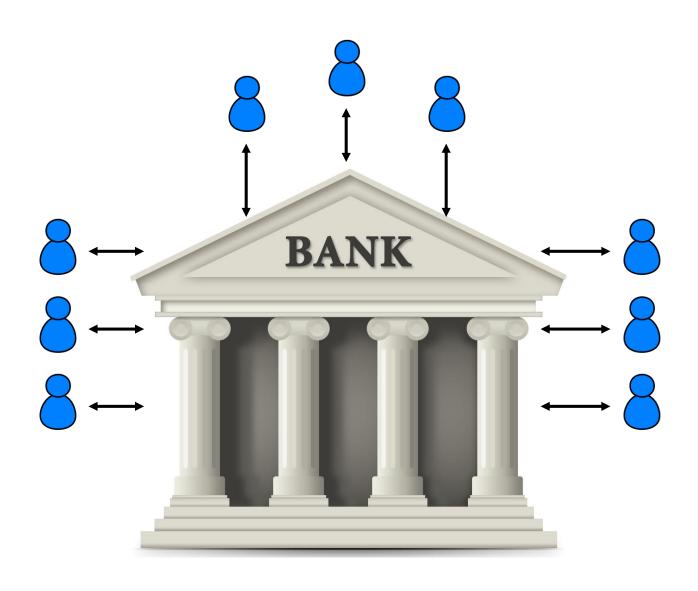
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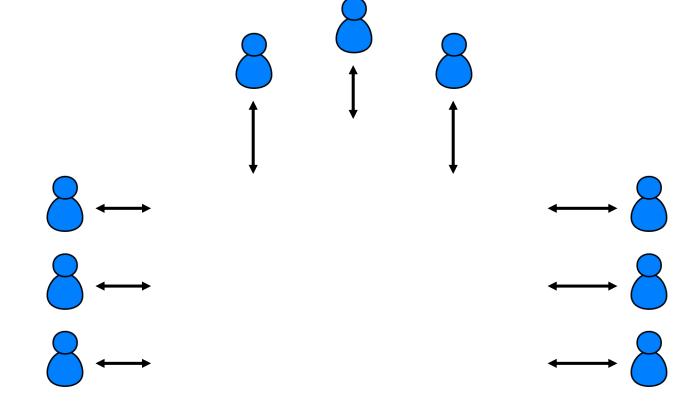
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    - Log is immutable and tamper-free (end-users trust this)



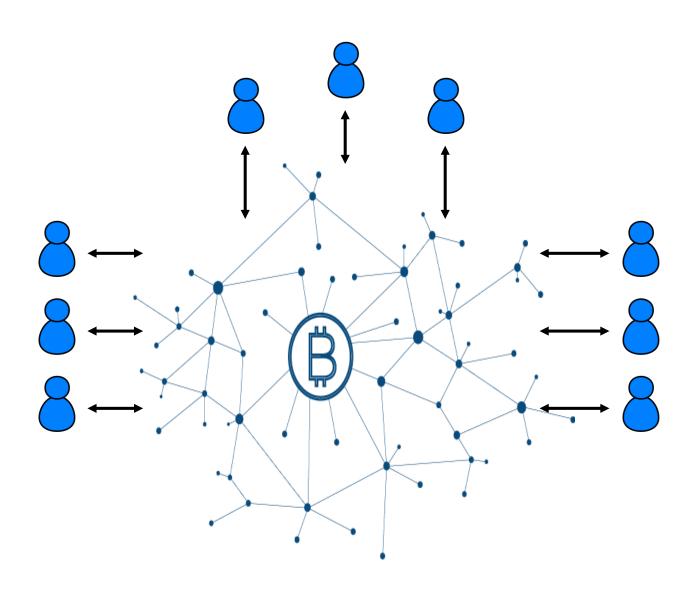
### Bitcoin



# Bitcoin



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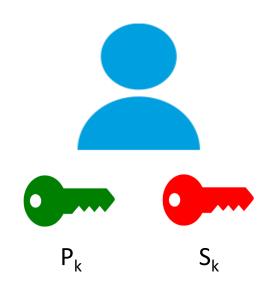
### Bitcoin: A Peer-to-Peer Electronic Cash System

- From Database and Distributed Computing Perspective
- Identities and Signatures
  - Public/Private key pair
- Ledger
  - The balance of each identity (saved in the blockchain)
- Transactions
  - Move bitcoins from one identity to another
  - Concurrency control to serialize transactions (Mining and PoW)
  - Typically backed by a transactions log (blockchain)
    - Log is persistent (replicated across the network nodes)
    - Log is immutable and tamper-free (PoW and Hash pointers)

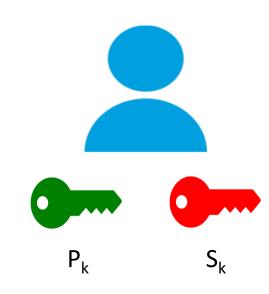


# Digital Signatures

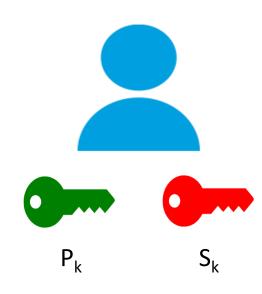
•  $P_k$ ,  $S_k \leftarrow Keygen(keysize)$ 



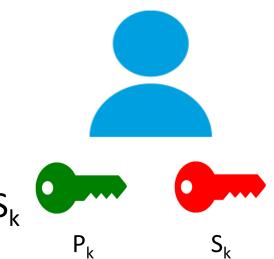
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- S<sub>k</sub> is private



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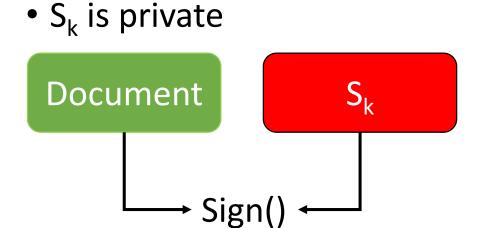
• S<sub>k</sub> is private

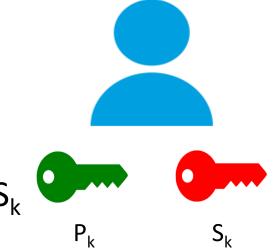
Document

 $S_k$ 

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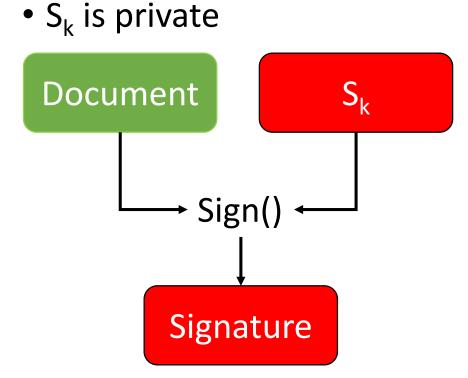


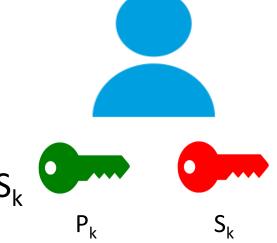


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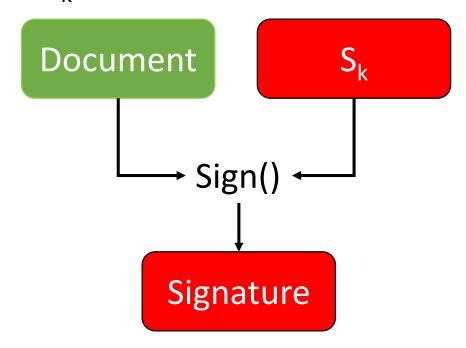


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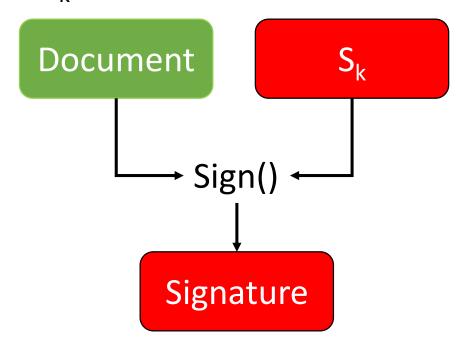
 $P_k$ 

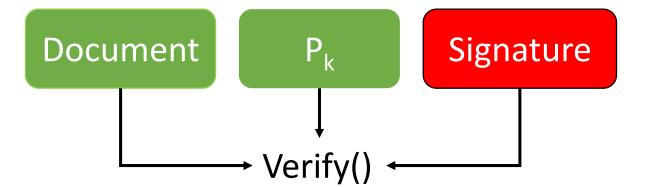
Signature

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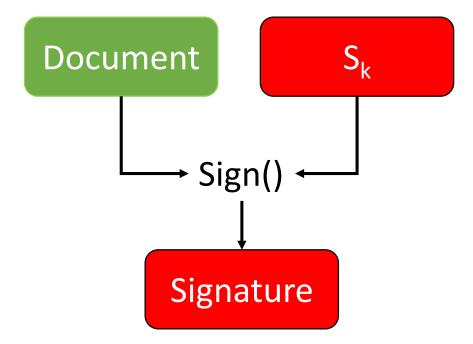


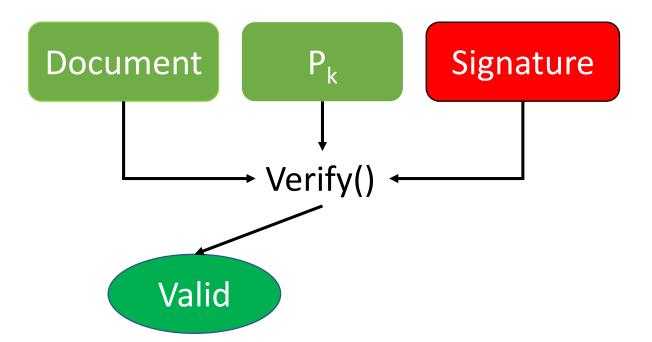
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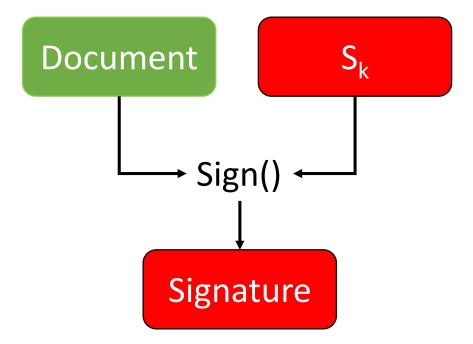


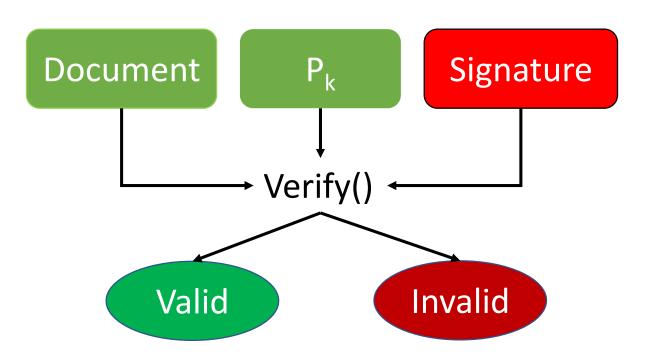
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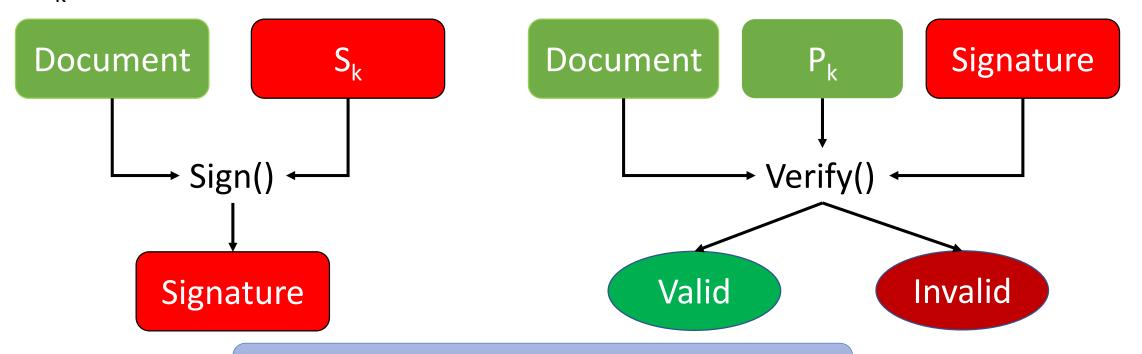


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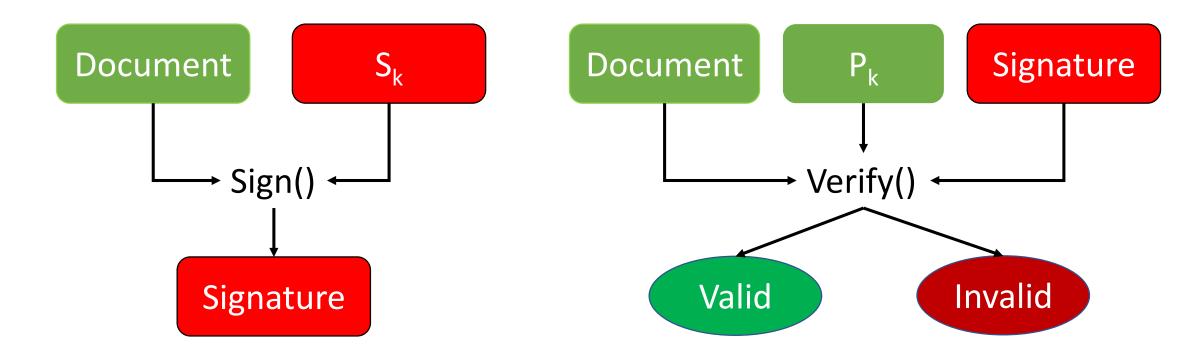
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**Used for Authentication not privacy** 

## Digital Signatures

- Unique to the signed document
- Mathematically hard to forge
- Mathematically easy to verify

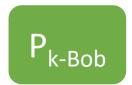


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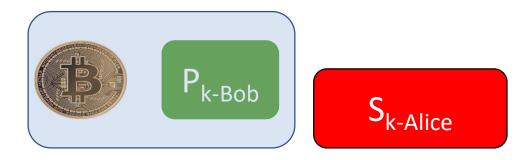




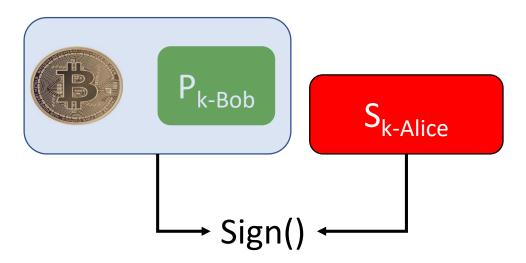
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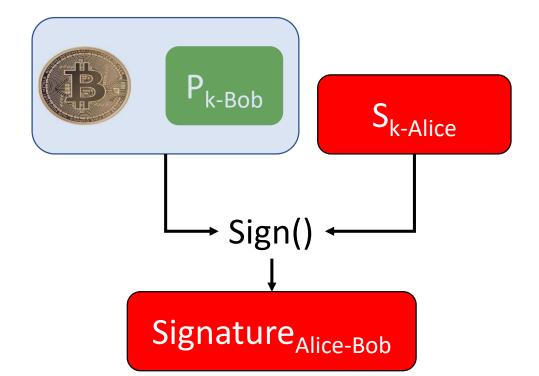
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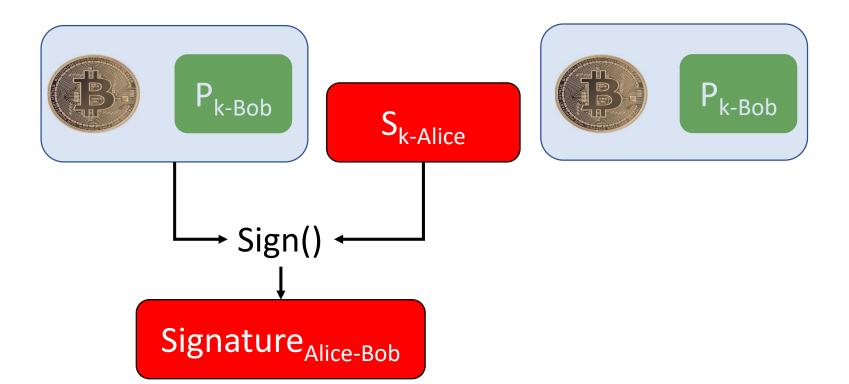
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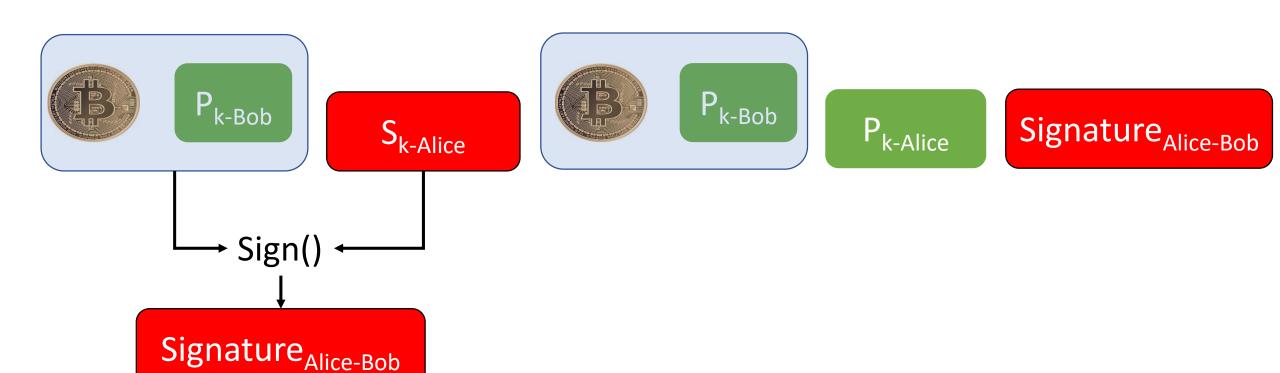
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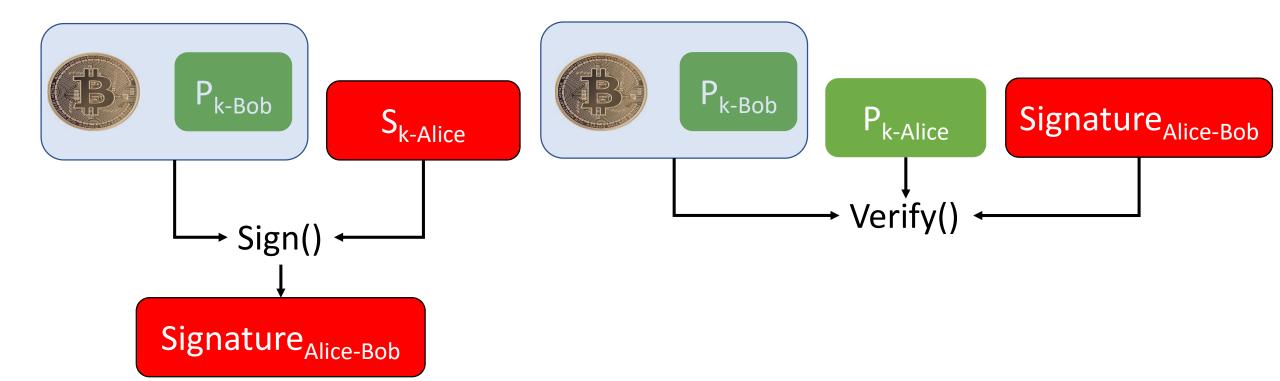
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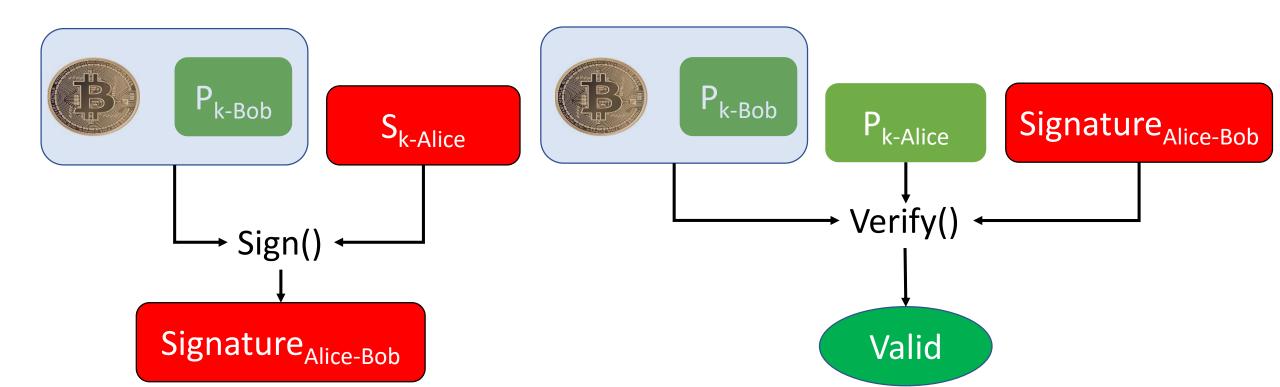
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## Digital Signatures and Bitcoin



Now what if Bob wants to move his coins to Diana

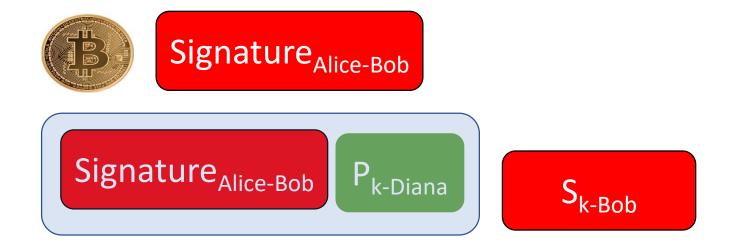


Signature<sub>Alice-Bob</sub>

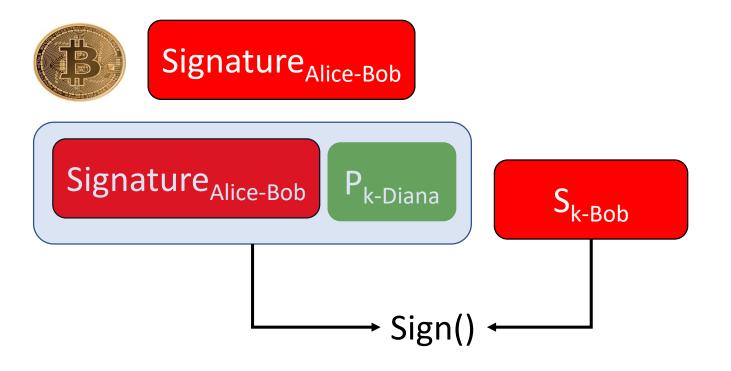




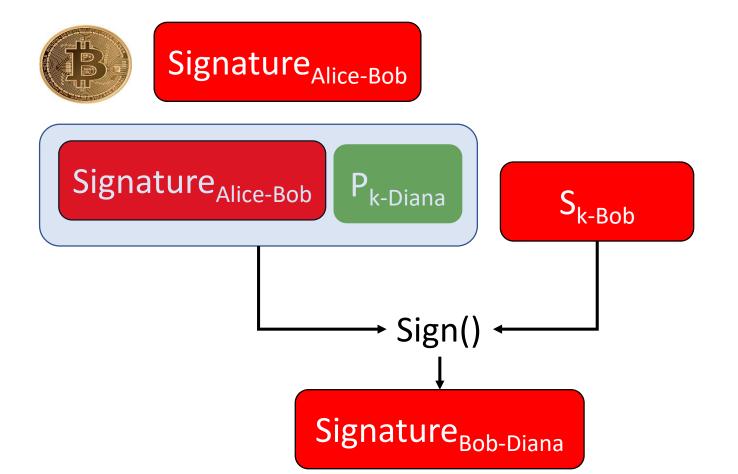












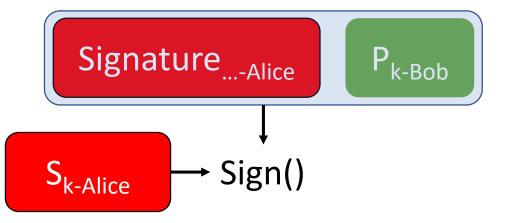
## A Bitcoin Big Picture

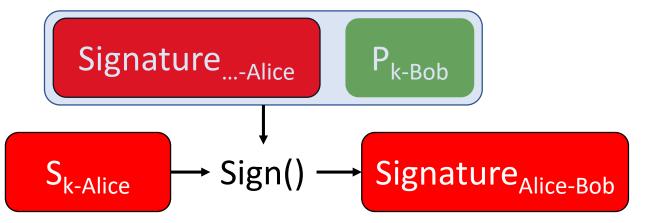
Signature<sub>...-Alice</sub>

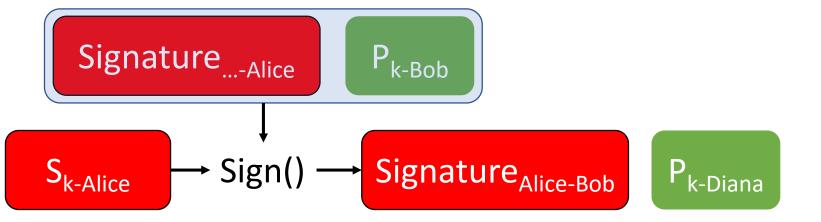
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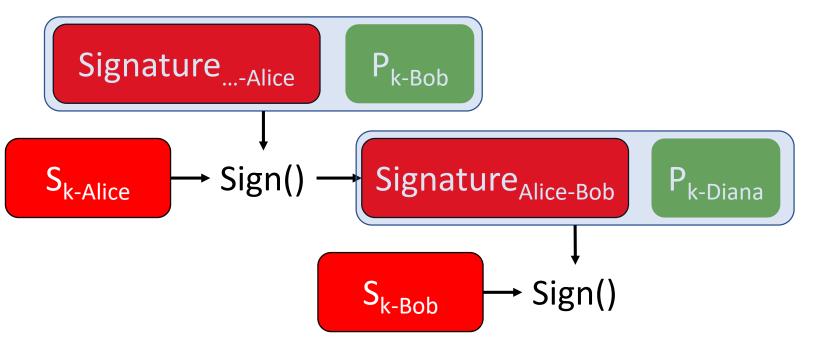
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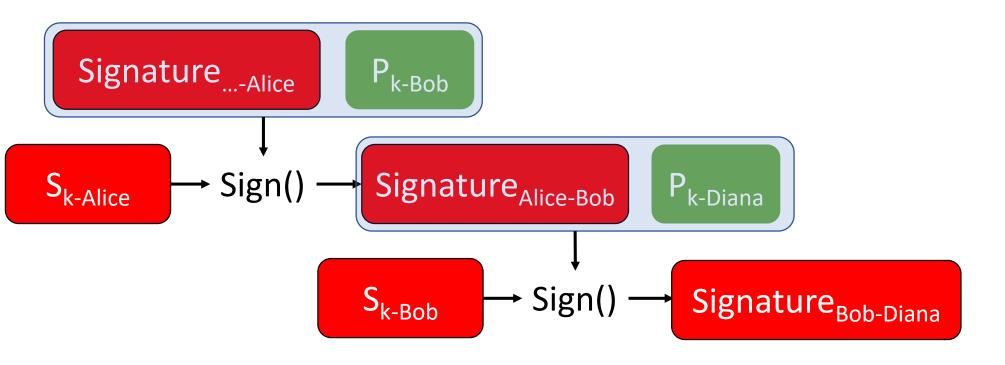


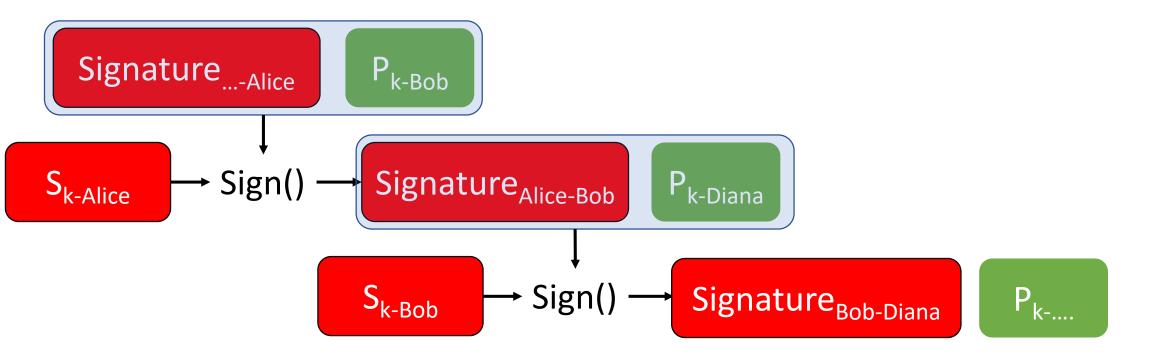


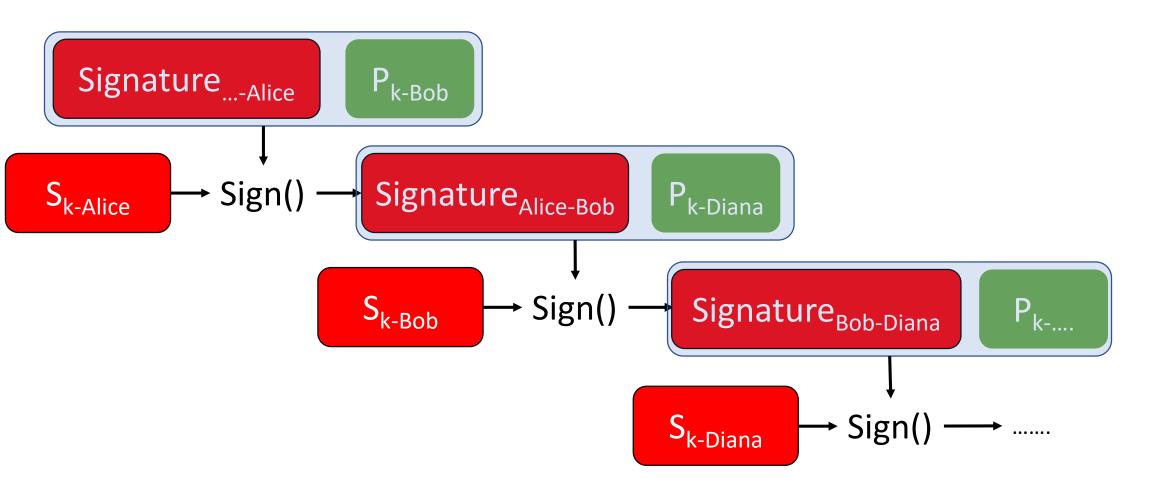


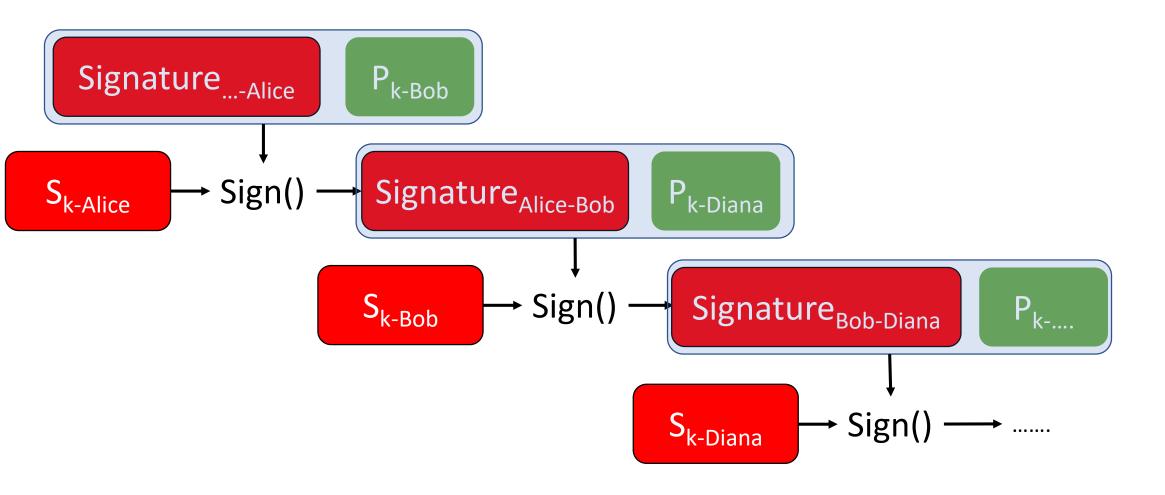


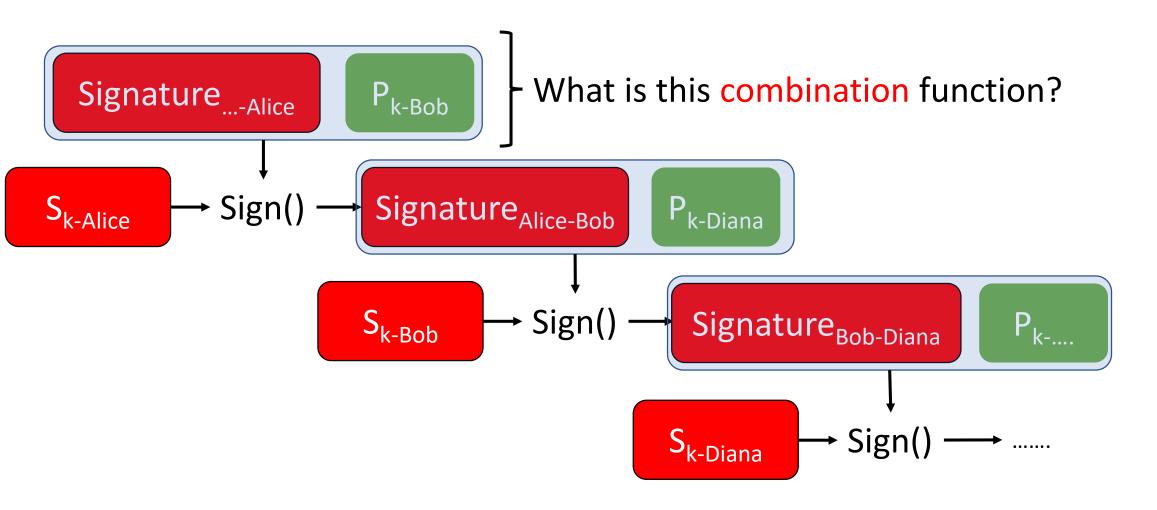


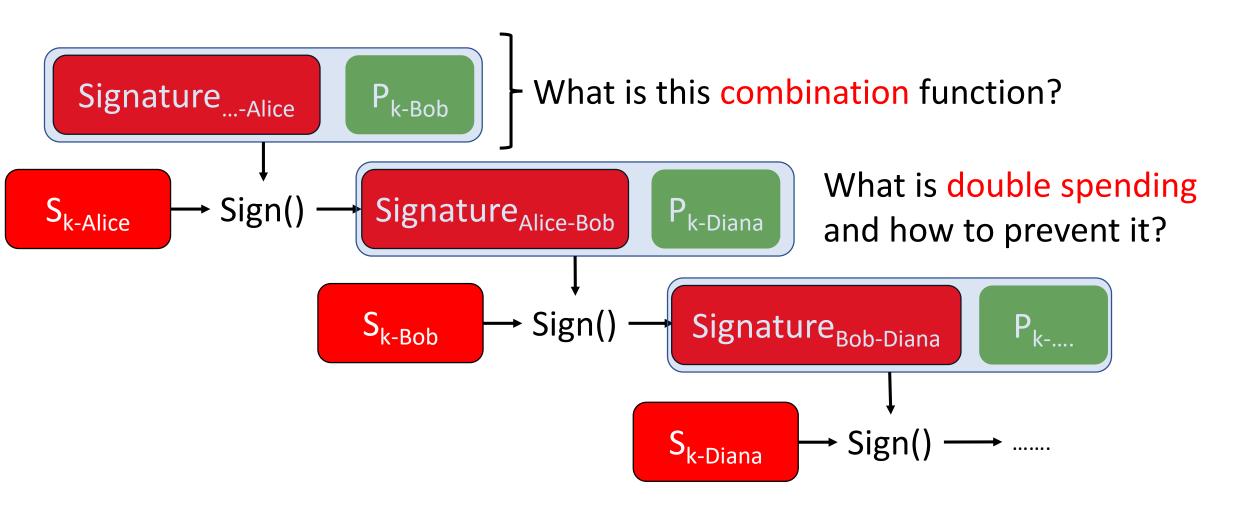


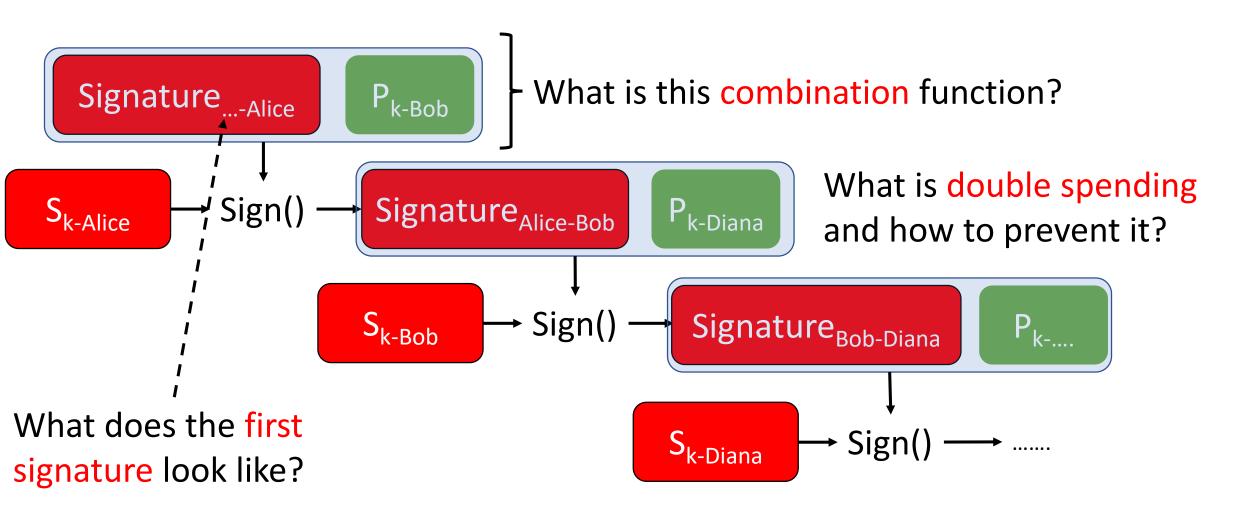












Hashing H(x)

Signature<sub>Alice-Bob</sub>

 $P_{k-Diana}$ 

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Signatures and public keys are combined using Hashing

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- Takes any string x of any length as input
- **Fixed** output size (e.g., 256 bits)



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- Signatures and public keys are combined using Hashing
- Takes any string x of any length as input
- **Fixed** output size (e.g., 256 bits)
- Efficiently computable.
- Satisfies:
  - Collision Free: no two x, y s.t. H(x) = H(y)
    - Message digest.
  - Hiding: Given H(x) infeasible to find x (one-way hash function)
    - Commitment: commit to a value and reveal later
  - Puzzle Friendly: Given a random puzzle ID and a target set Y it is hard to find x such that: H(ID | x) ε Y







```
Signature<sub>Alice-Bob</sub> P<sub>k-Diana</sub>
```

```
SHA256( Signature<sub>Alice-Bob</sub> | P<sub>k-Diana</sub> ) = 256-bit (32-byte) unique string
```













SHA256(abc) = ba7816bf8f01cfea414140de5dae2223b00361a396177a9cb410ff61f20015ad







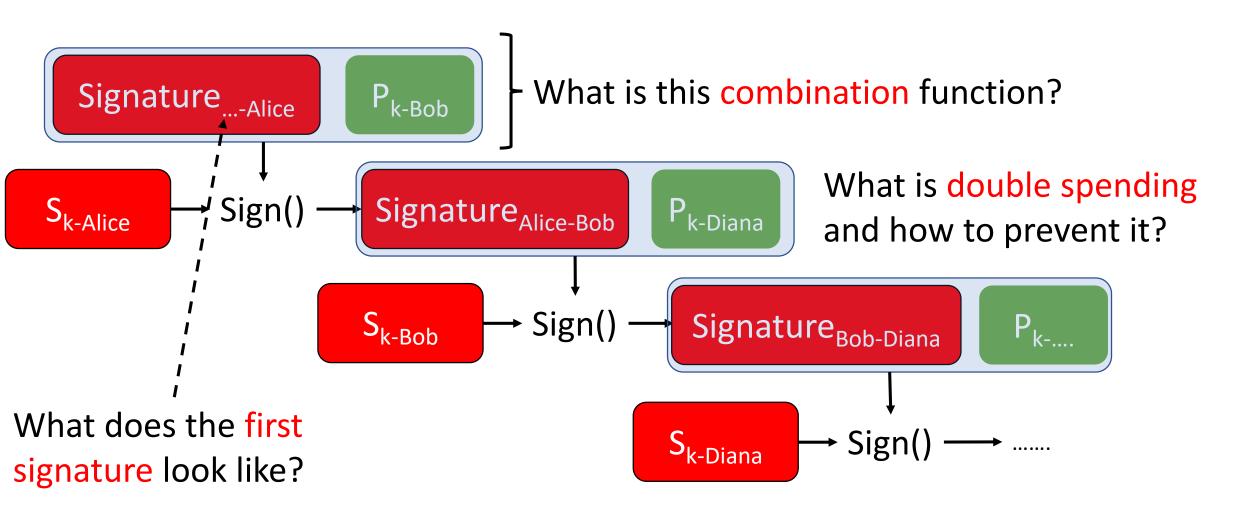
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ba7816bf8f01cfea414140de5dae2223b00361a396177a9cb410ff61f20015ad

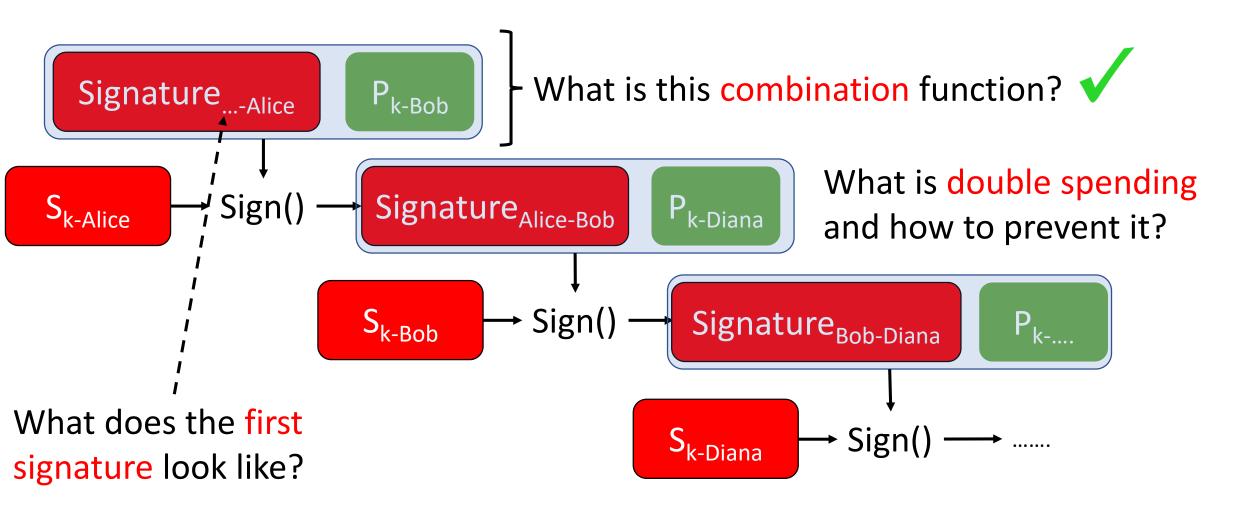
SHA256(abC) =

0a2432a1e349d8fdb9bfca91bba9e9f2836990fe937193d84deef26c6f3b8f76

#### What About's?



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- Impossible to do in physical cash
- Prevented in traditional banking systems through concurrency control

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Signature<sub>Alice-Bob</sub>



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 $P_{k ext{-Diana}}$ 

Signature Alice-Bob

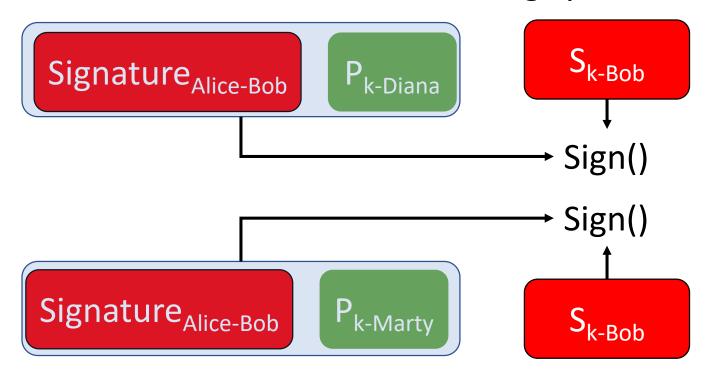


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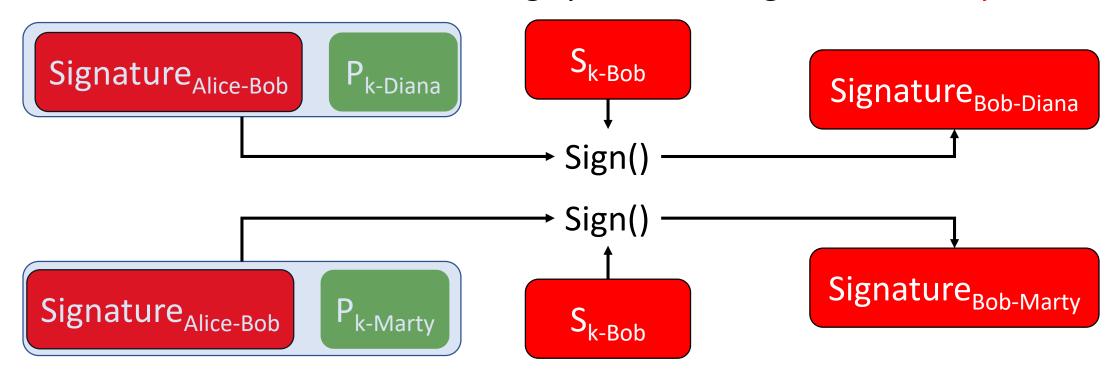




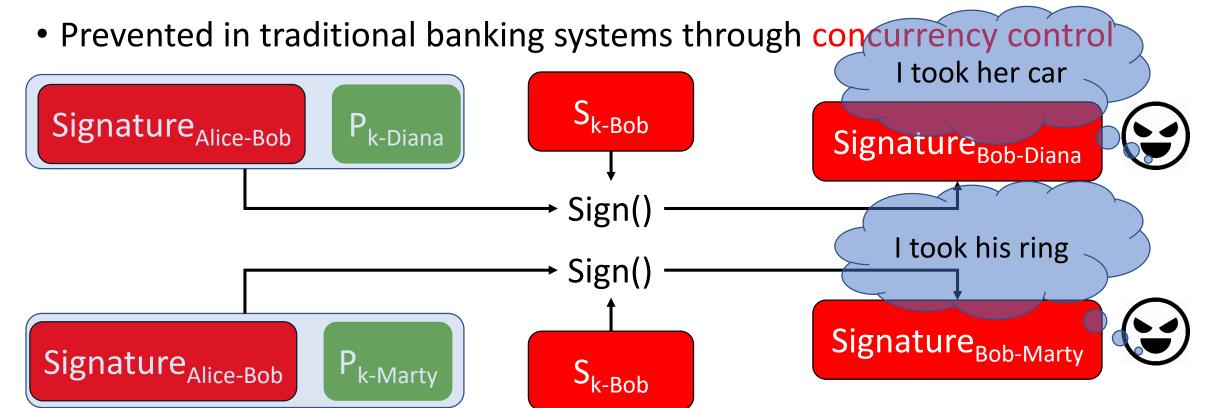
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# Double Spending Prevention

Centralized

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  - Transactions on coins go through a trusted 3<sup>rd</sup> party (Trent)





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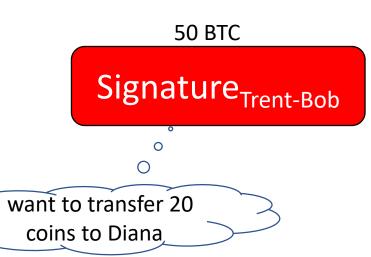
50 BTC

Signature<sub>Trent-Bob</sub>



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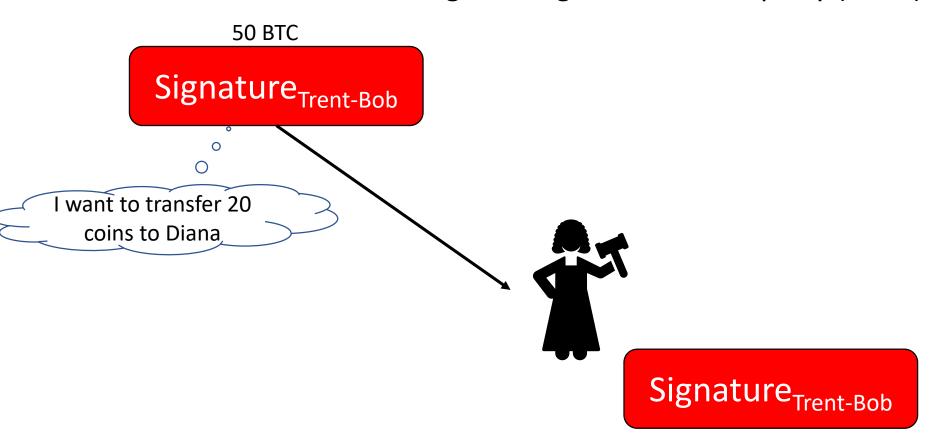






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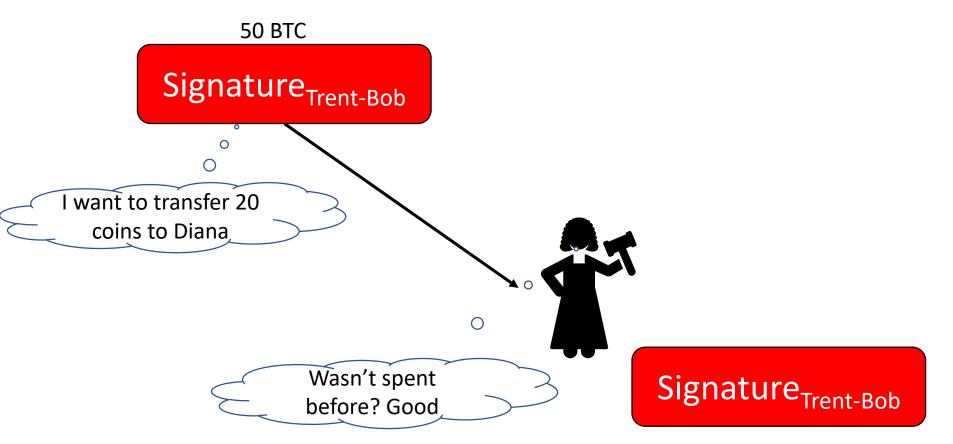






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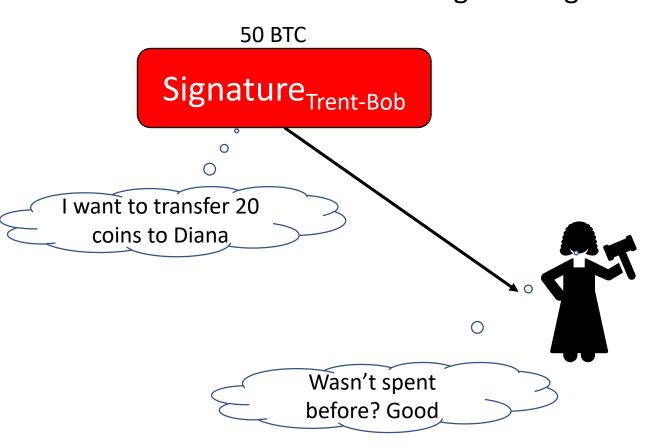






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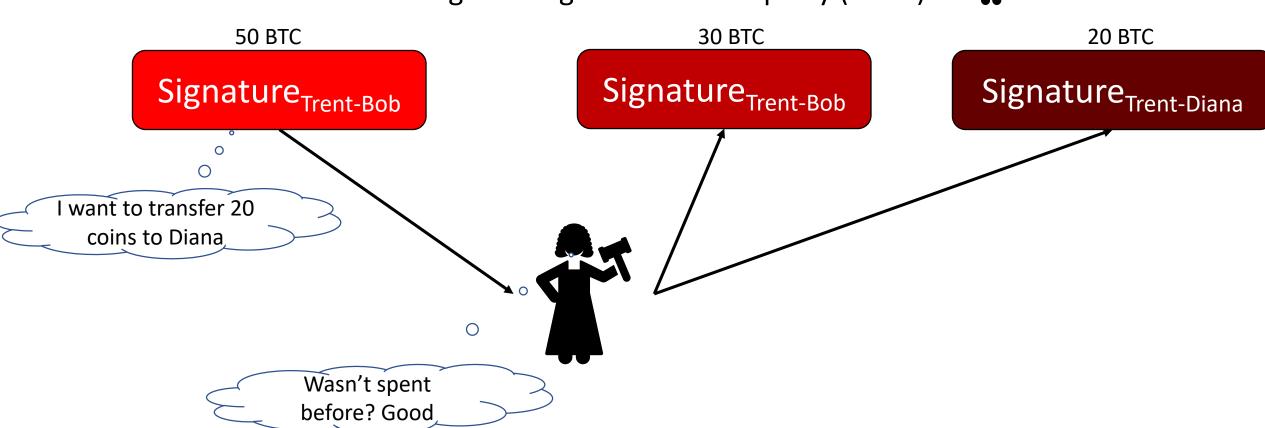


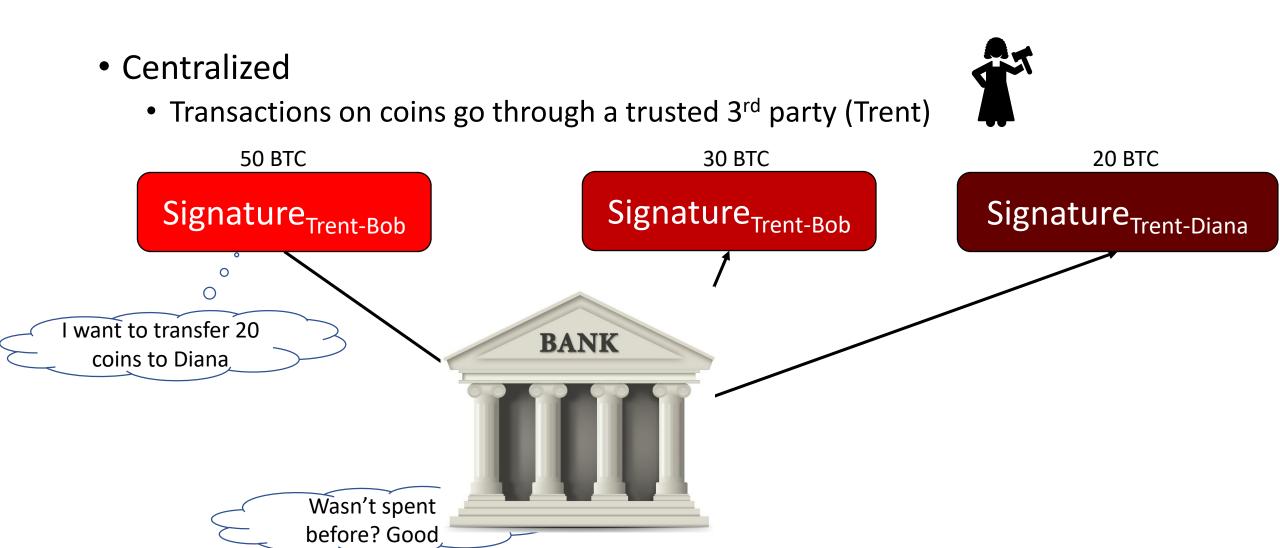
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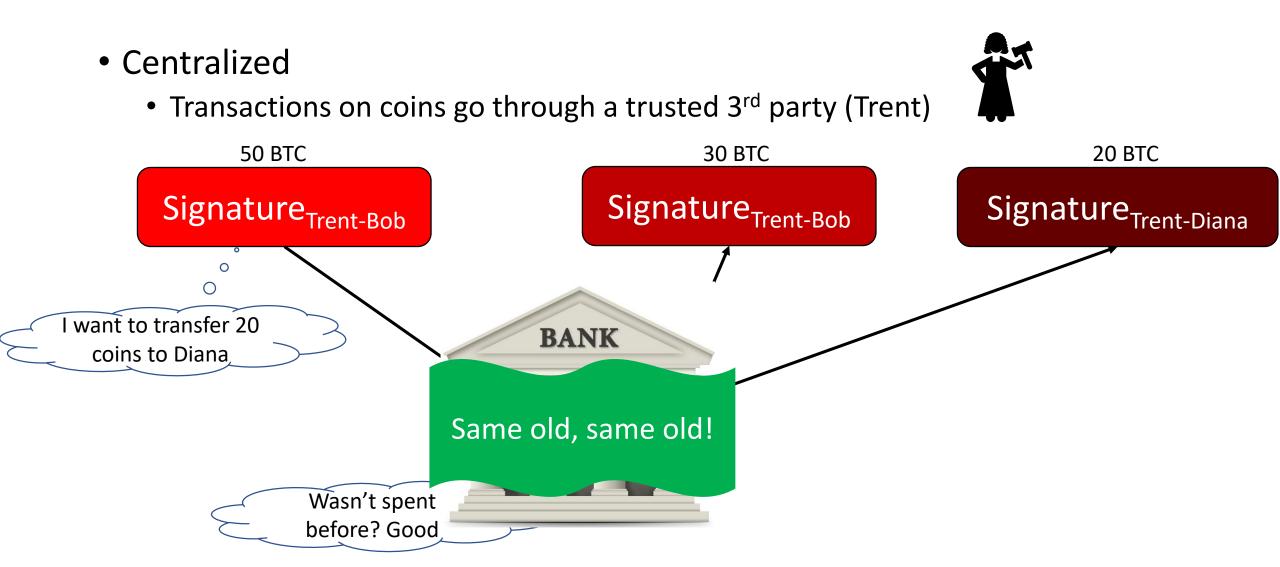


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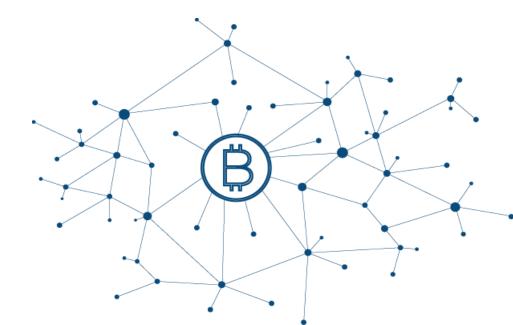




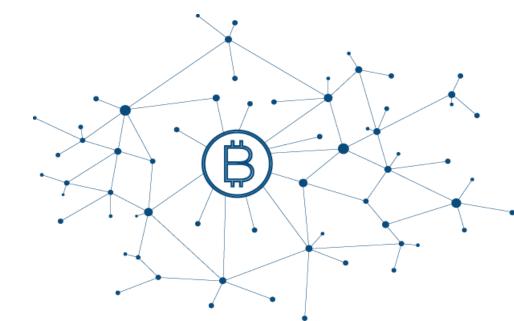




Decentralized

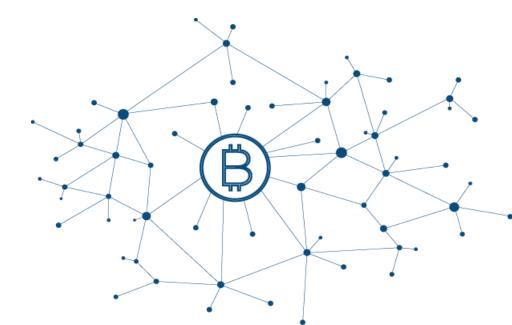


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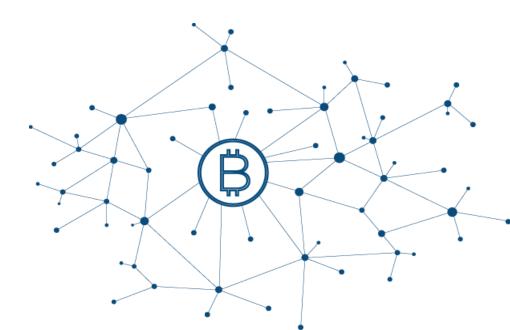




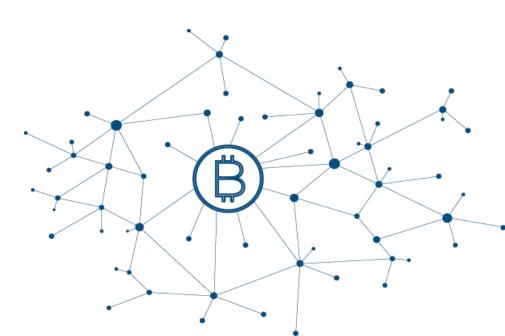
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  - Network nodes work to agree on transactions order
    - Serializing transactions on every coin prevents double spending



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  - A network of nodes maintains a ledger
  - Network nodes work to agree on transactions order
    - Serializing transactions on every coin prevents double spending
  - What is the ledger?



- Decentralized
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What is the Ledger?

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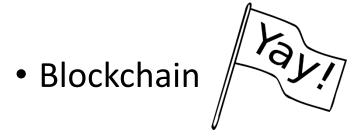
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What is the Ledger?

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# What is the Ledger?



Transactions are grouped into blocks

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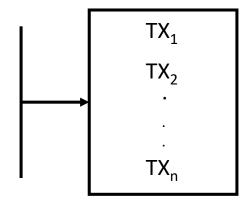
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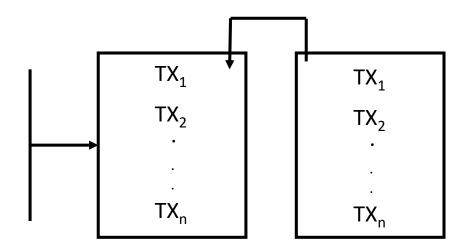


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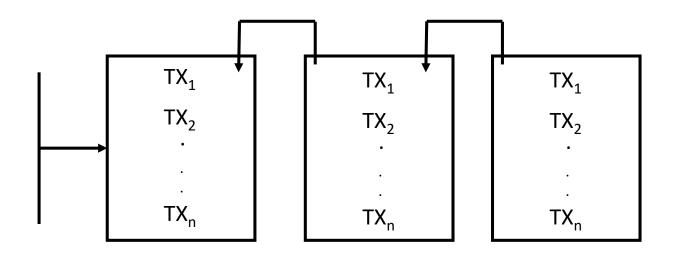


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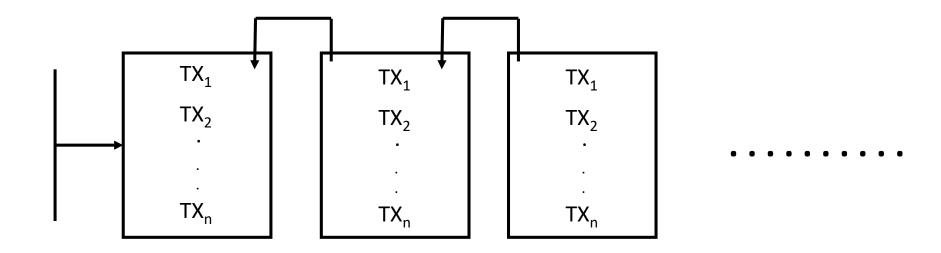


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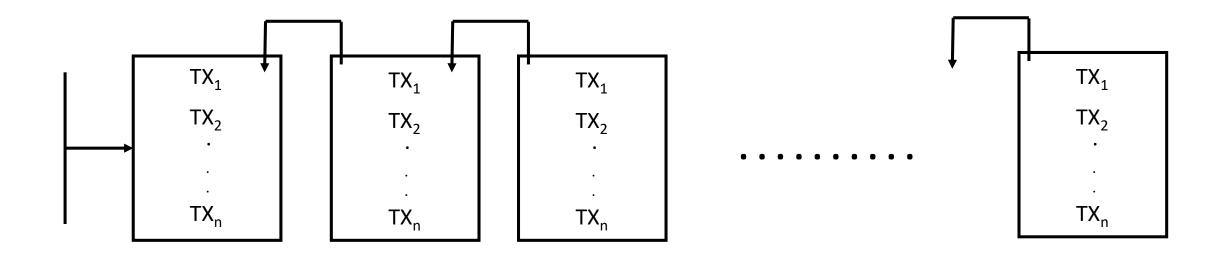
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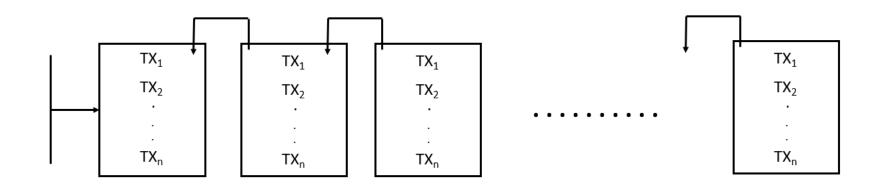


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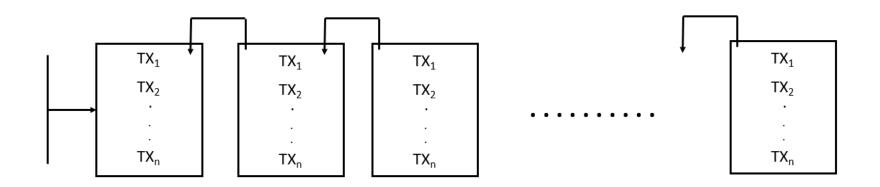
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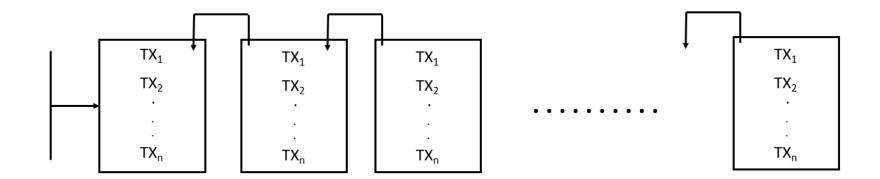


## The Ledger's What About's?

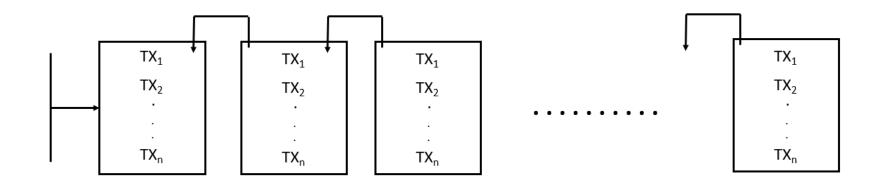
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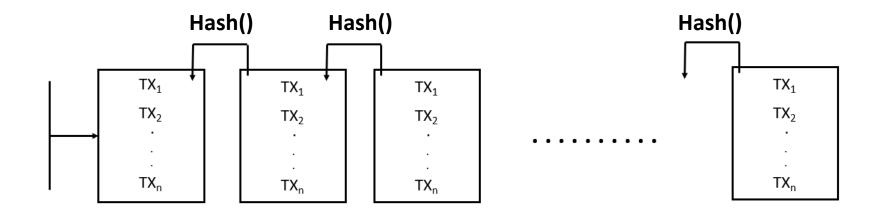
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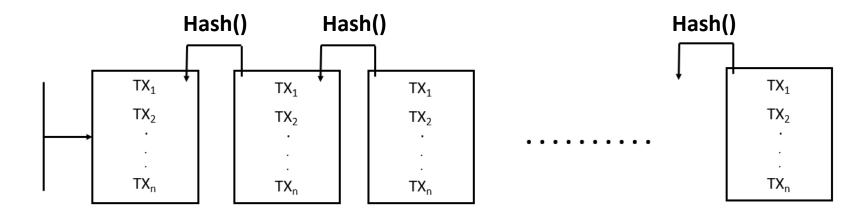


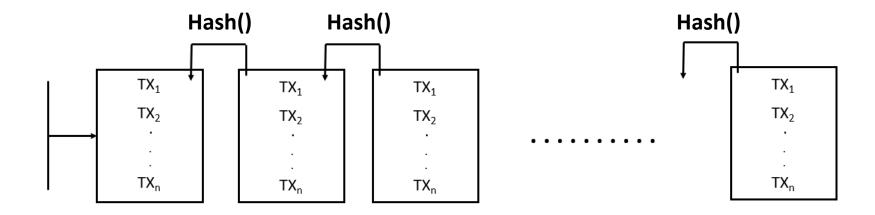
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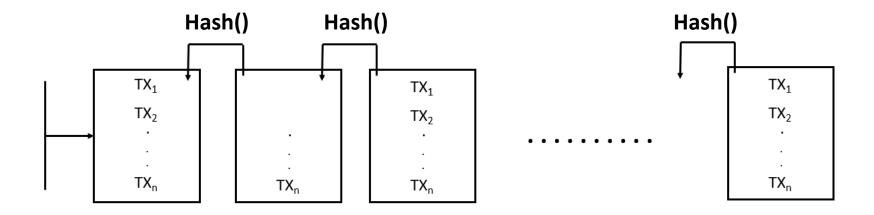


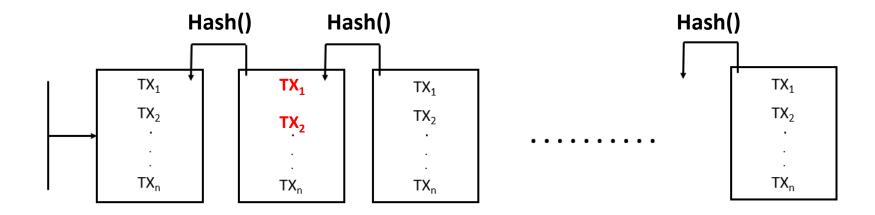


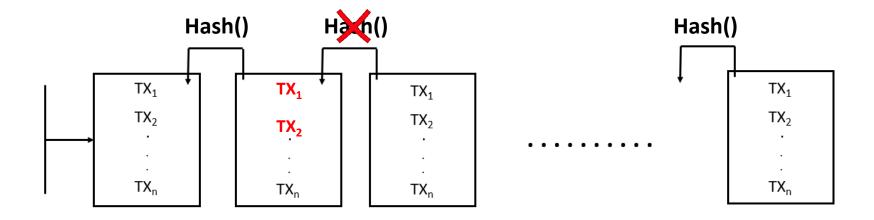
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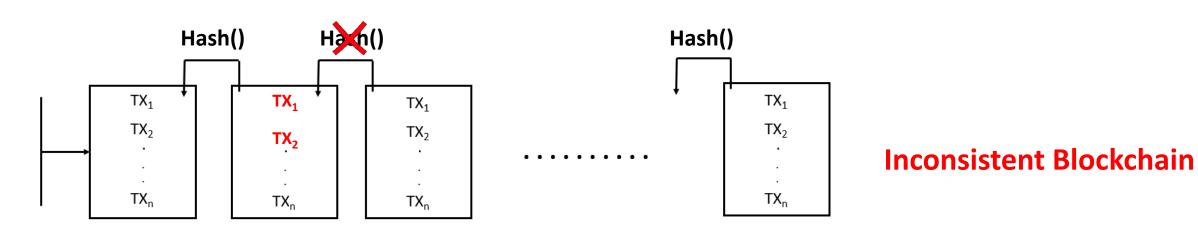




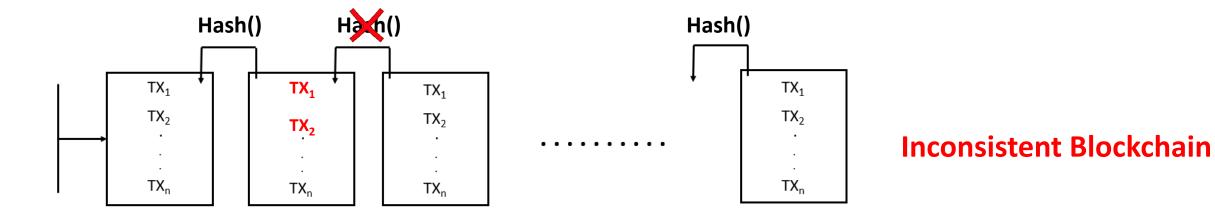






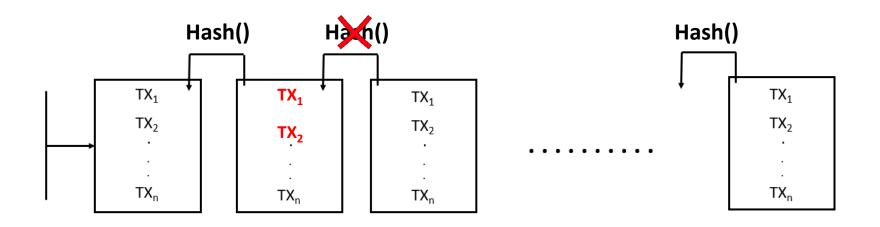


# Tampering with the Ledger



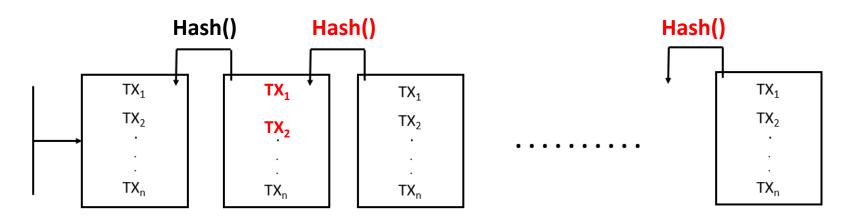
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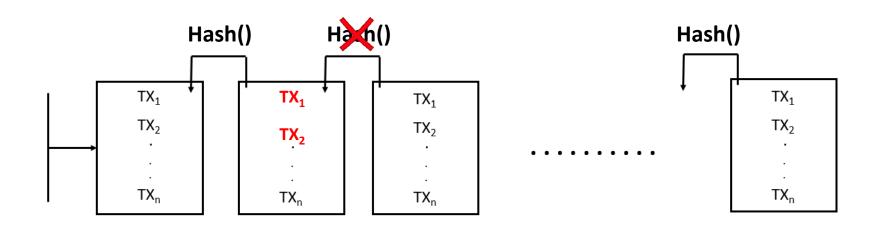


**Inconsistent Blockchain** 

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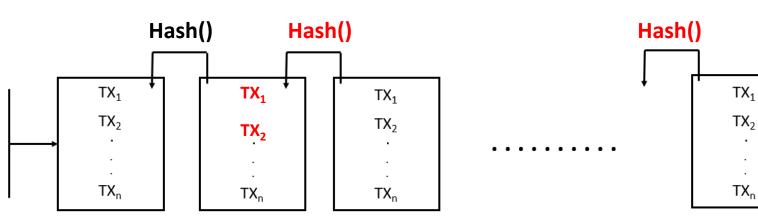


## Tampering with the Ledger



**Inconsistent Blockchain** 

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TX<sub>1</sub>
TX<sub>2</sub>
Consistent Blockchain

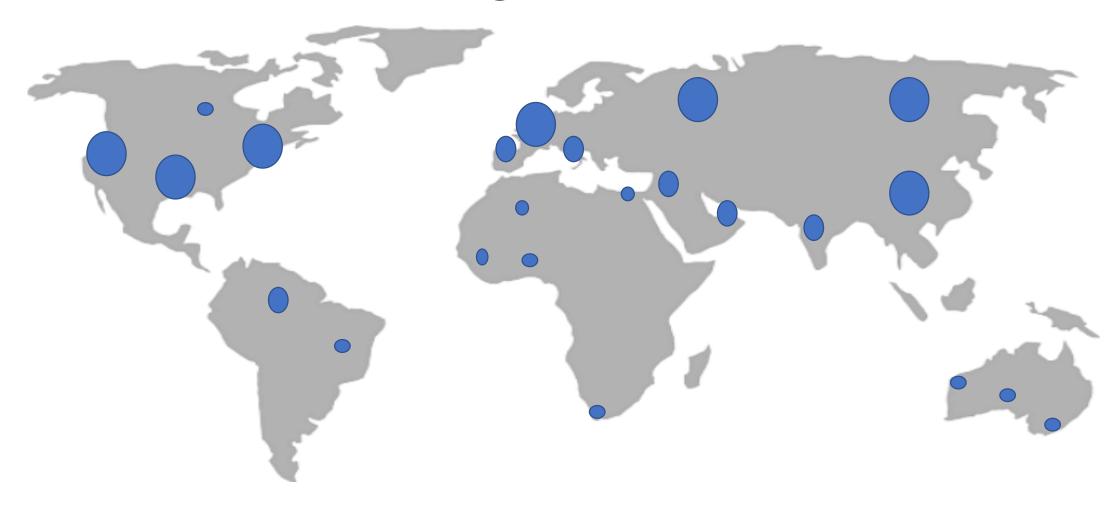
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  - 2. Replacing a consistent blockchain with another tampered consistent block chain should be **made very hard**, How?

# Network Nodes Big Picture



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## Making Progress

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  - **Synchronous** systems:
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Α \_\_\_\_\_

Majority

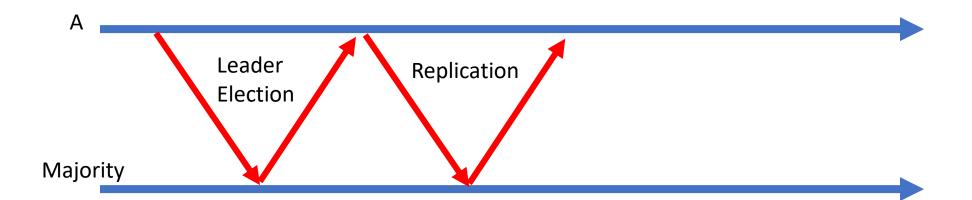
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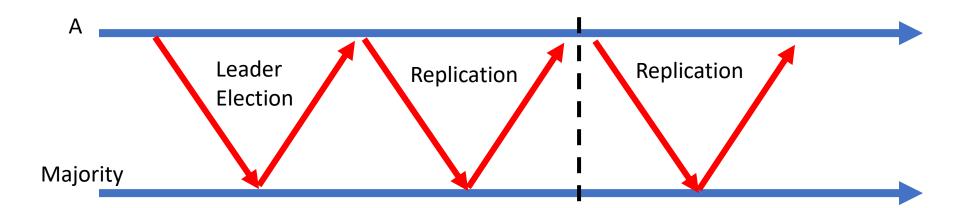
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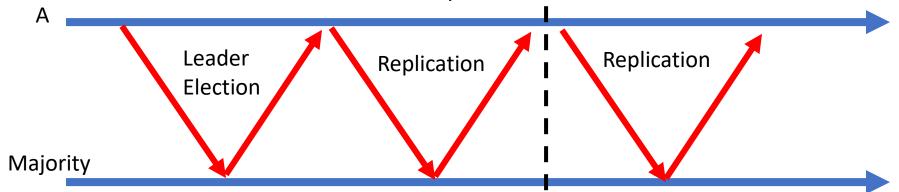
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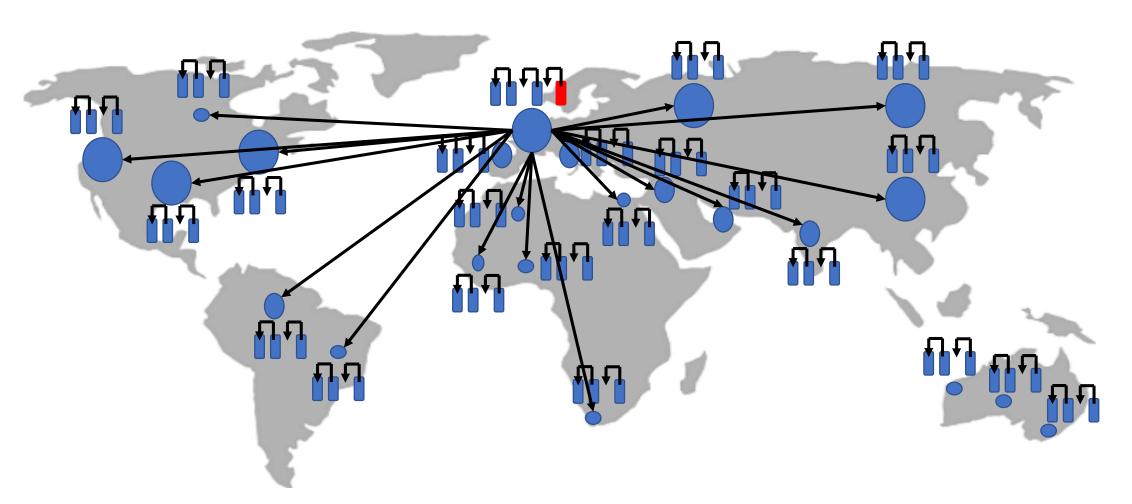
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## Paxos Consensus

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- Also, Paxos has high network overhead

# Practical Byzantine Fault Tolerance (PBFT)

 Goal: Implement a deterministic replication service with arbitrary malicious faults in an asynchronous environment

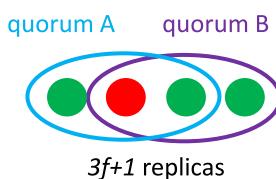
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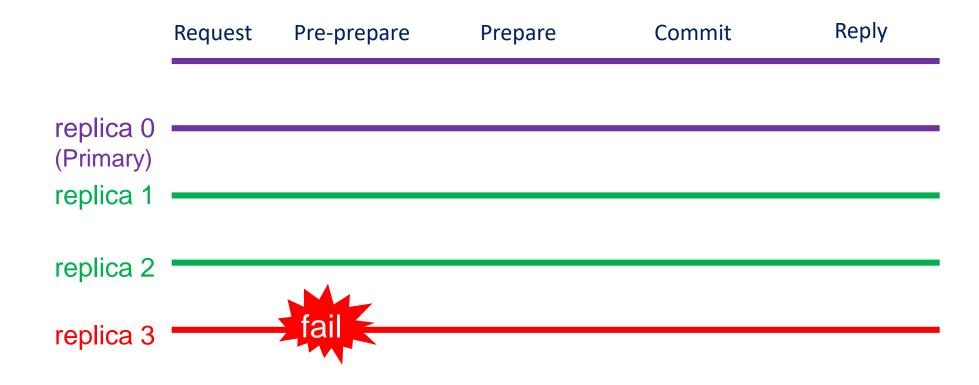
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- Assumptions:
  - 3f+1 replicas to tolerate f Byzantine faults (optimal)
    - quorums have at least 2f+1 replicas
    - quorums intersect in f+1, hence have at least one correct replica
  - Strong cryptography
  - Only for liveness: eventual time bounds



# Algorithm

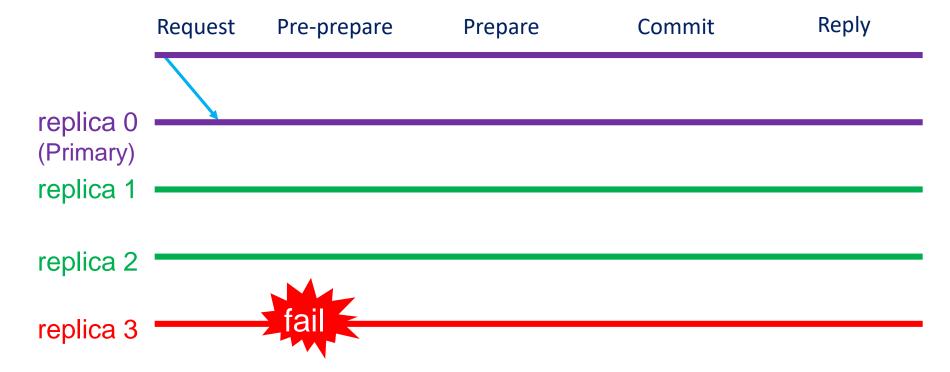
The algorithm has three main phases: (1) *pre-prepare* picks order of requests (2) *prepare* ensures order within views, (3) *commit* ensures order across views



# Algorithm

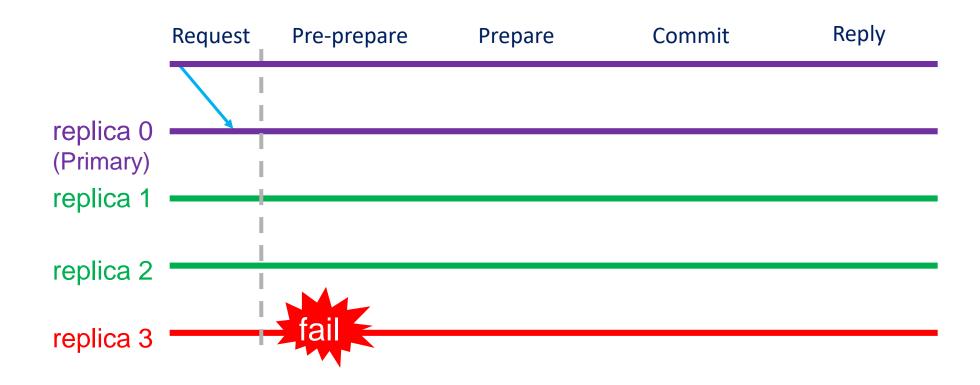
The algorithm has three main phases: (1) *pre-prepare* picks order of requests (2) *prepare* ensures order within views, (3) *commit* ensures order across views

(1) A client sends a request for a service to the primary



# Algorithm

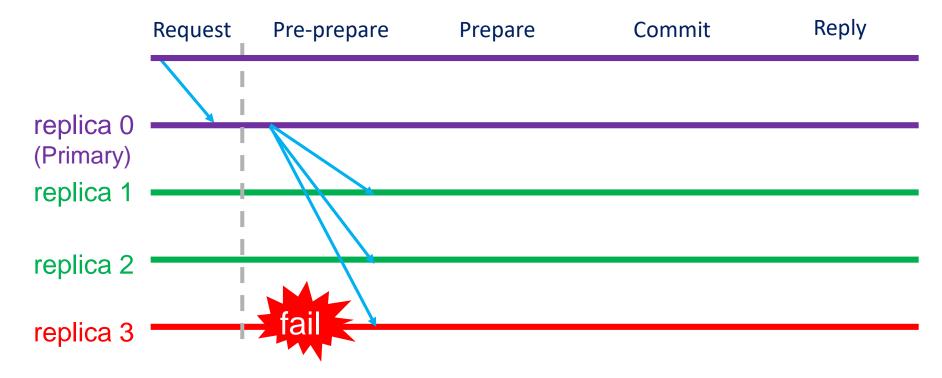
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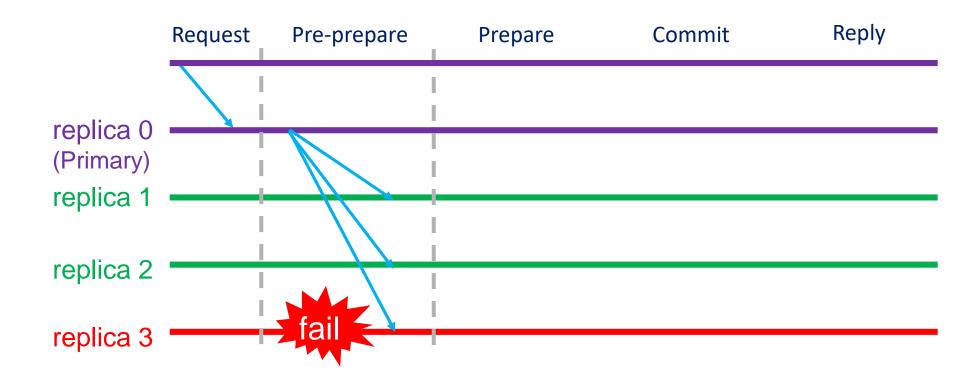
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(2) The primary multicasts the request to the backups



# Algorithm

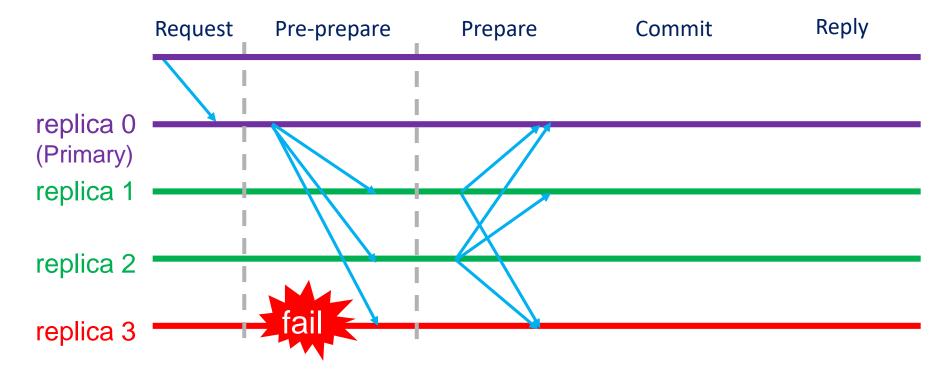
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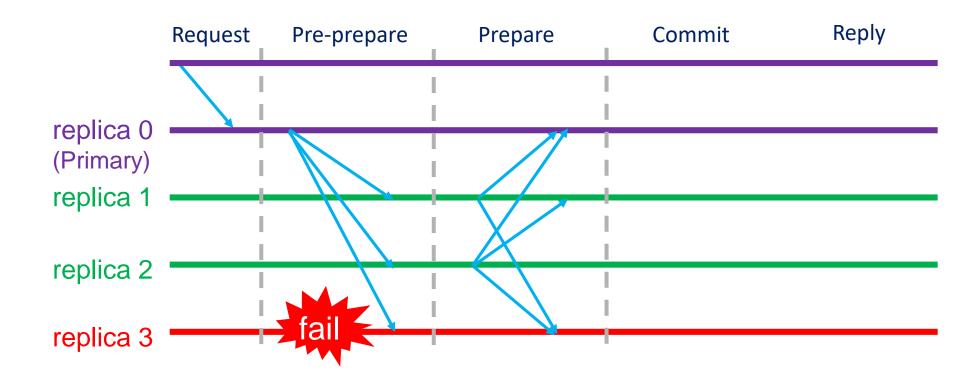
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(3) Backups multicast PREPARE message



## Algorithm

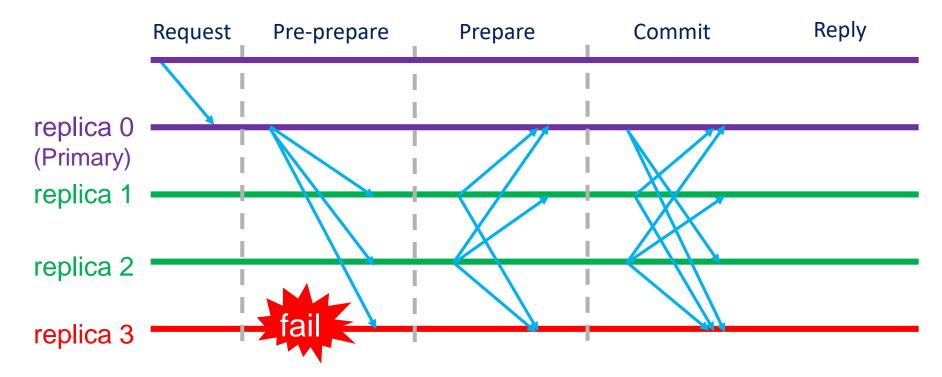
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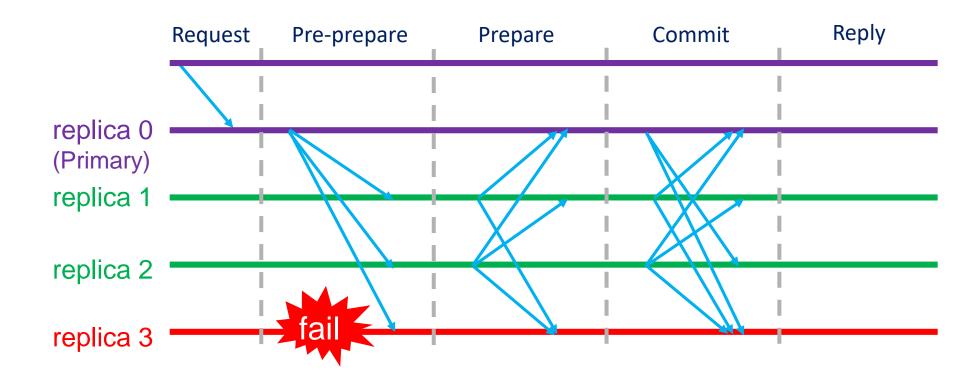
The algorithm has three main phases: (1) *pre-prepare* picks order of requests (2) *prepare* ensures order within views, (3) *commit* ensures order across views

(4) If a replica receives at least 2f matching PREPARE message, multicasts a COMMIT message



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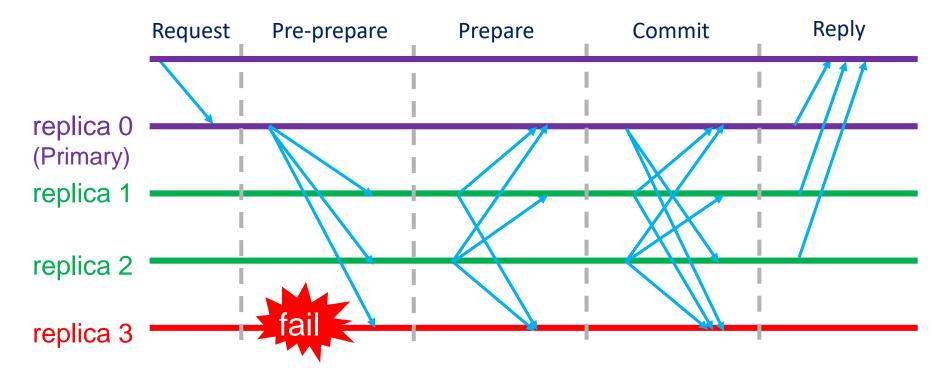
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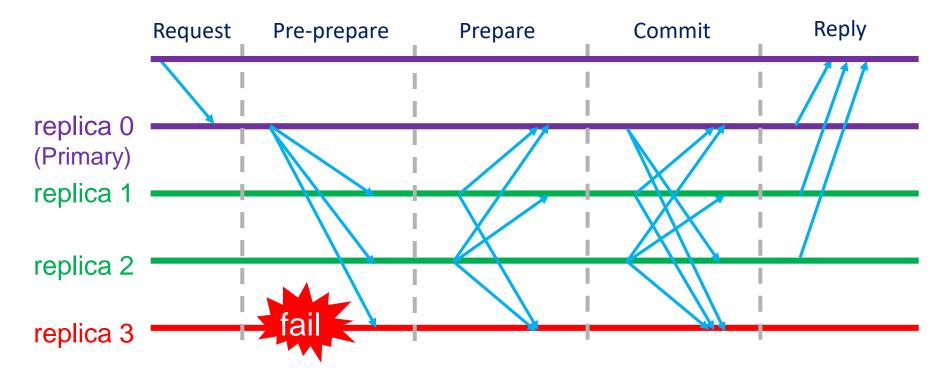
(5) If a replica receives at least 2f COMMIT messages, reply the result to the client



## Algorithm

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(6) The client waits for f+1 replies from different replicas with the same result

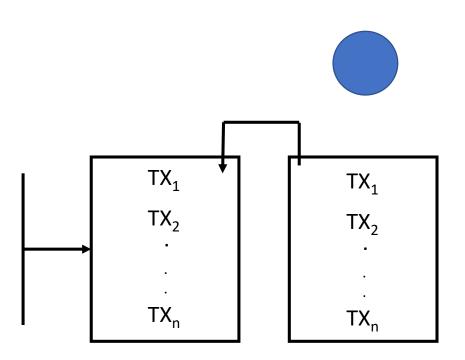


### **PBFT Consensus**

- Tolerates Byzantine (Malicious) failures
  - To make progress, at least 2/3 of the participants should be correct
  - Progress is not guaranteed (FLP impossibility)
- However, PBFT is Permissioned
  - All participants should be known a priori
- Also, PBFT has high network overhead O(N<sup>2</sup>) [number of messages]
  - Every node multi-casts their responses to every other node

### Nakamoto's Consensus

- Intuitively, network nodes race to solve a puzzle
- This puzzle is computationally expensive
- Once a network node finds (mines) a solution:
  - It adds its block of transactions to the blockchain
  - It multi-casts the solution to other network nodes
  - Other network nodes accept and verify the solution

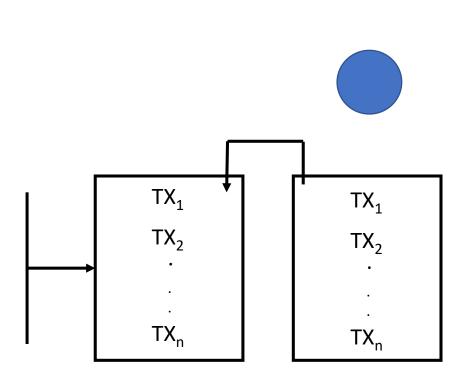


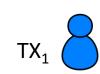






# Mining Details



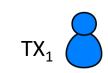




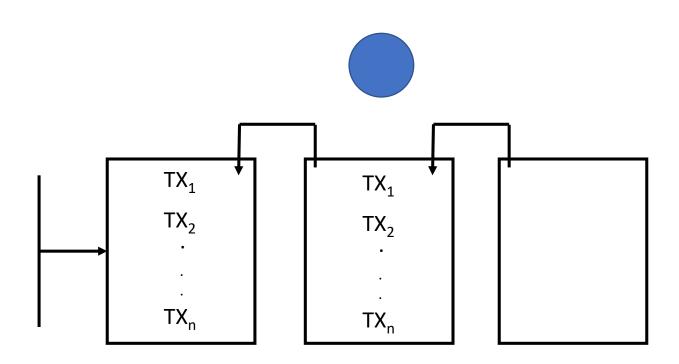
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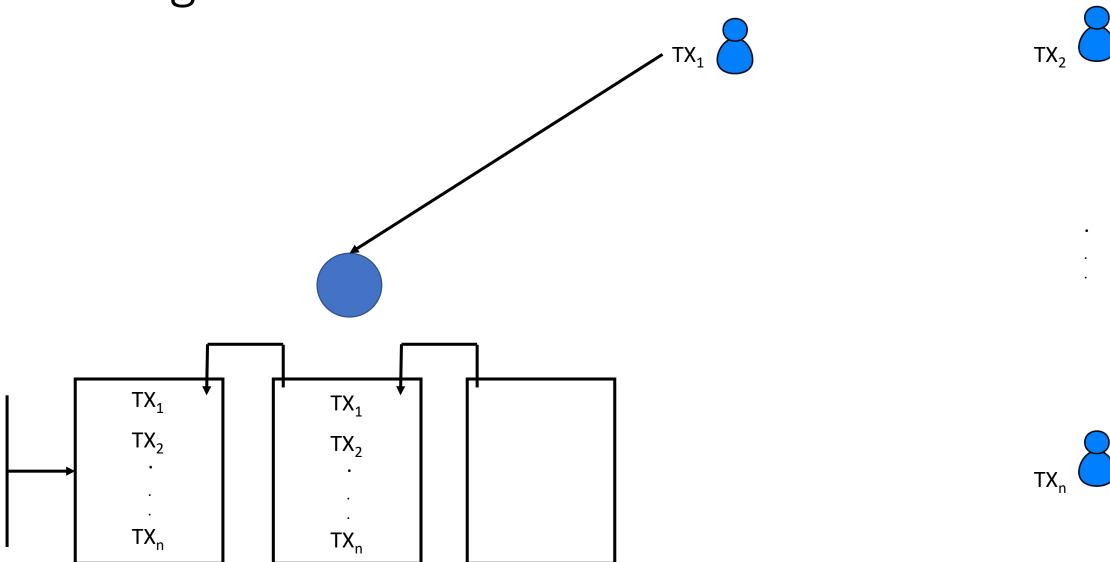


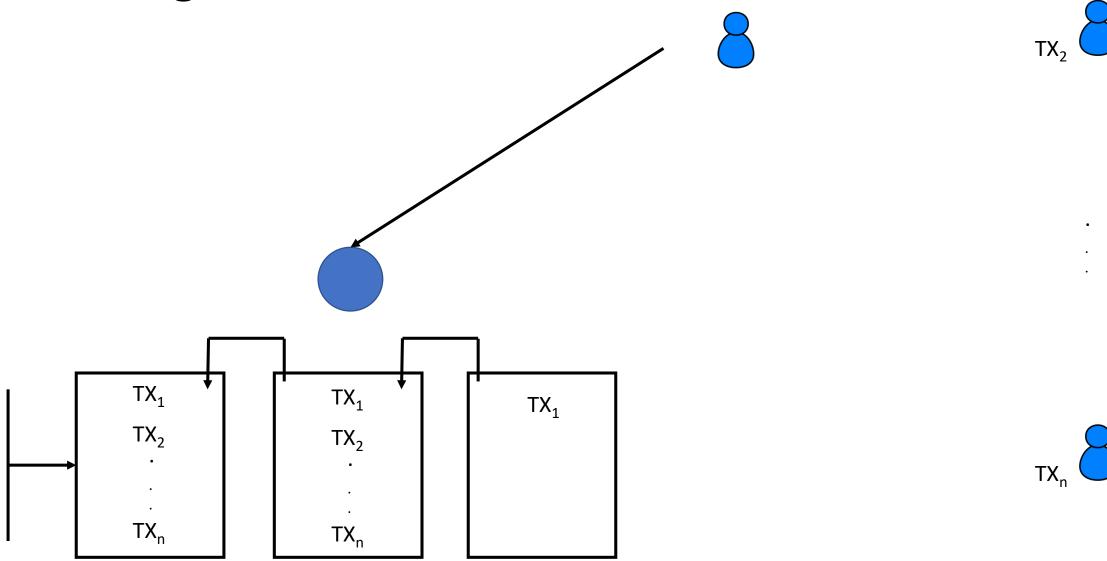


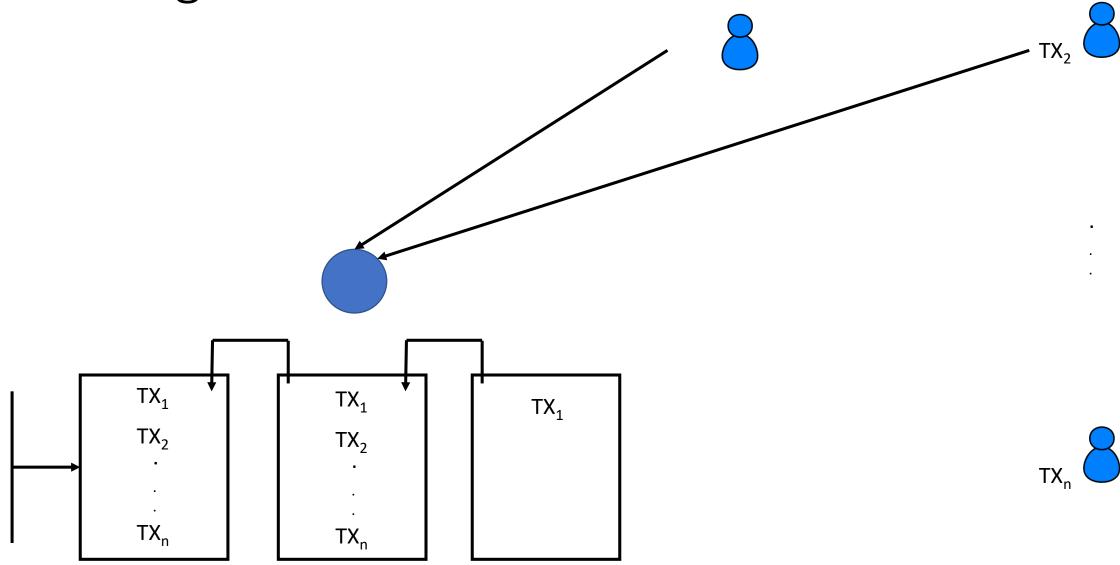


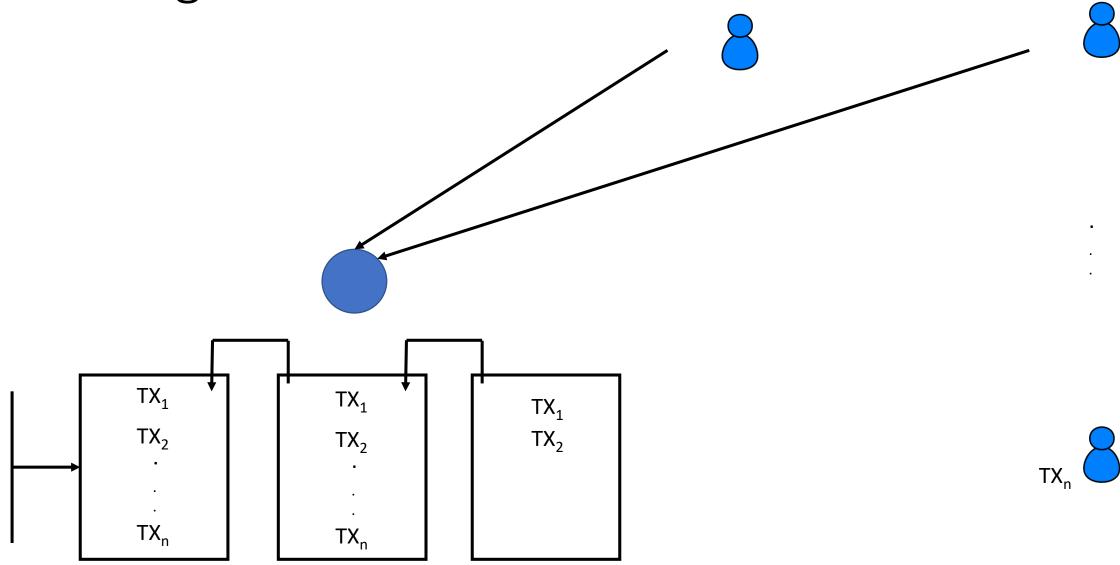
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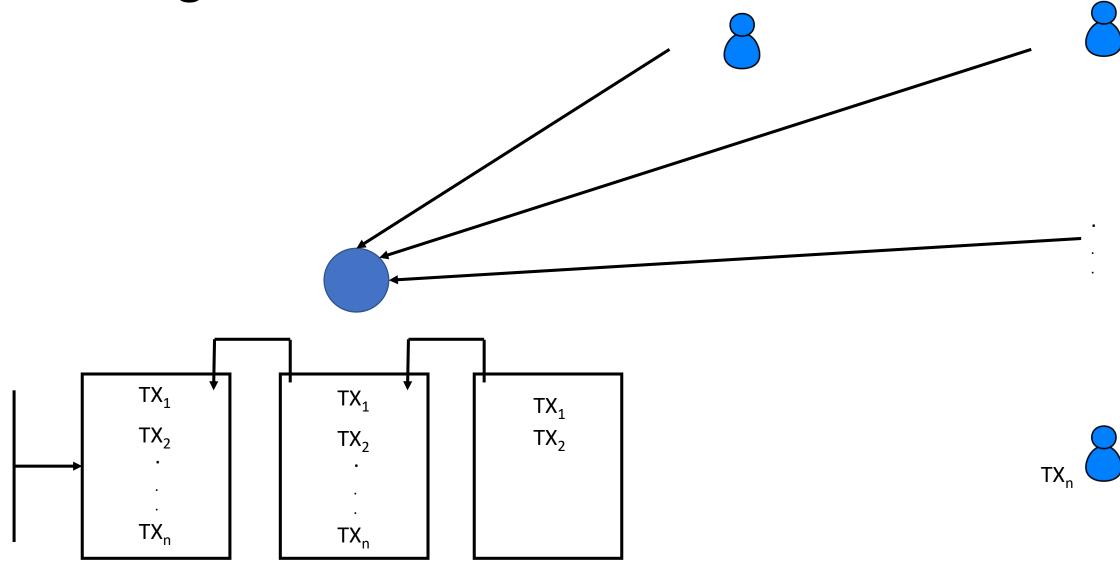


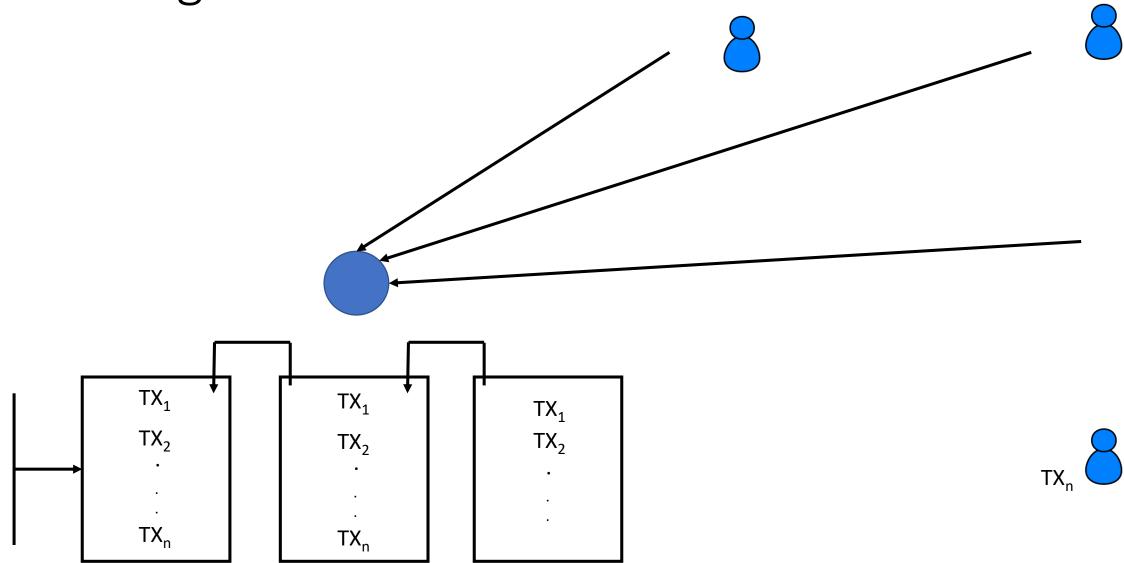


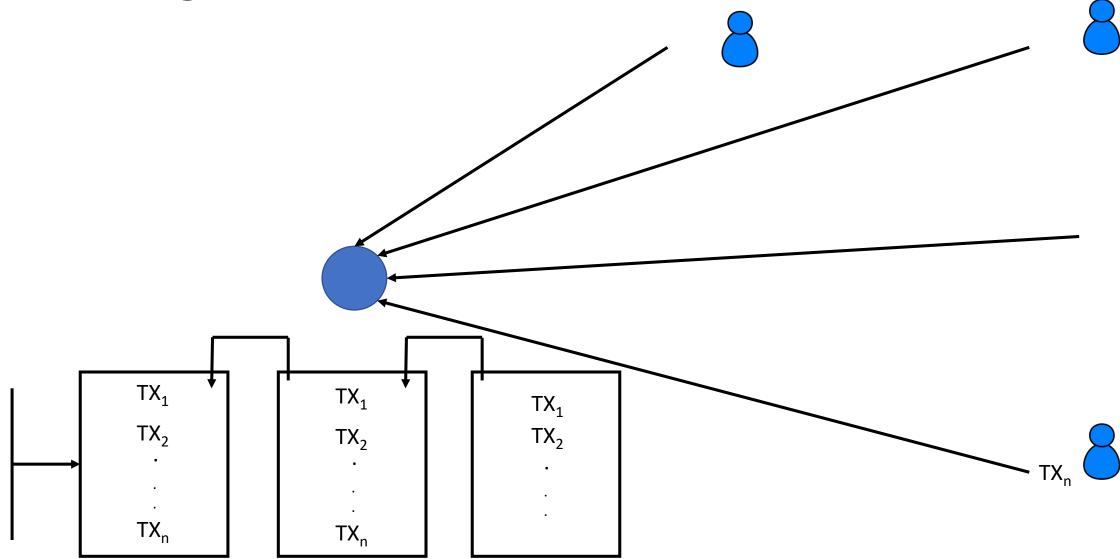


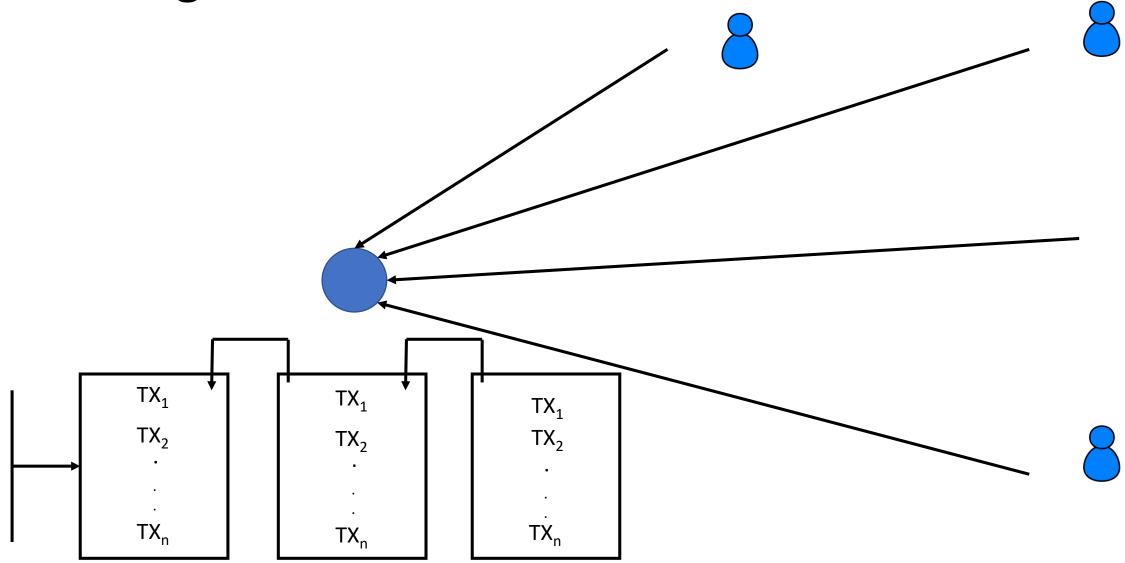


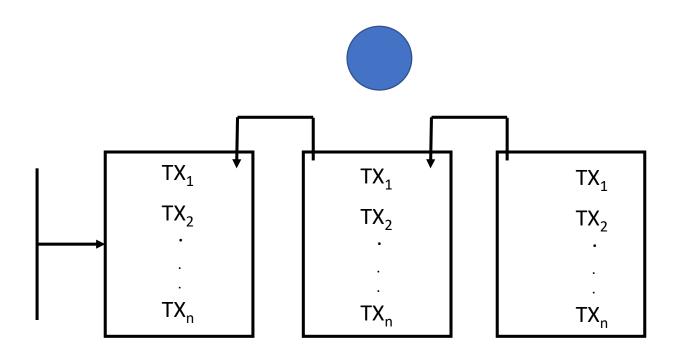


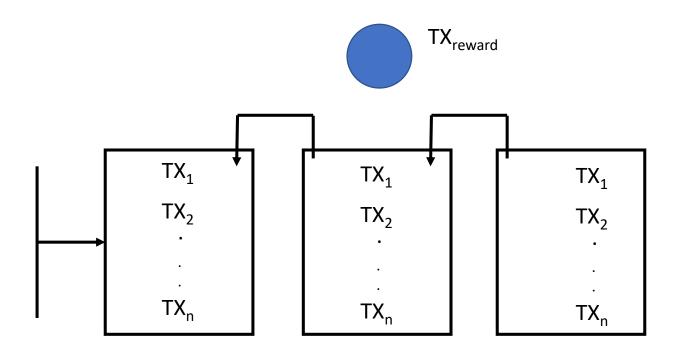


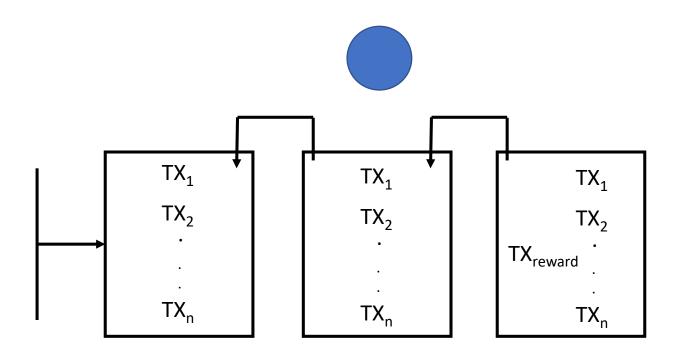


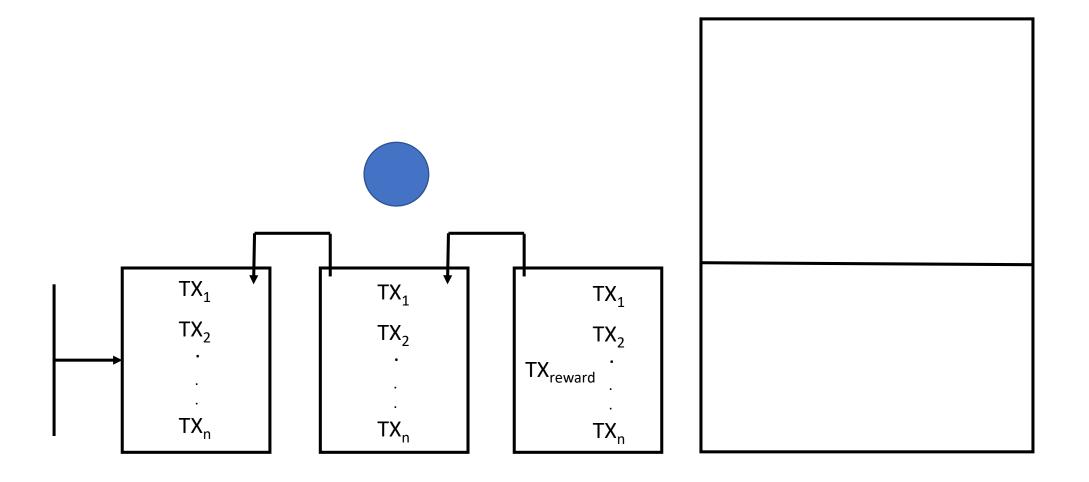


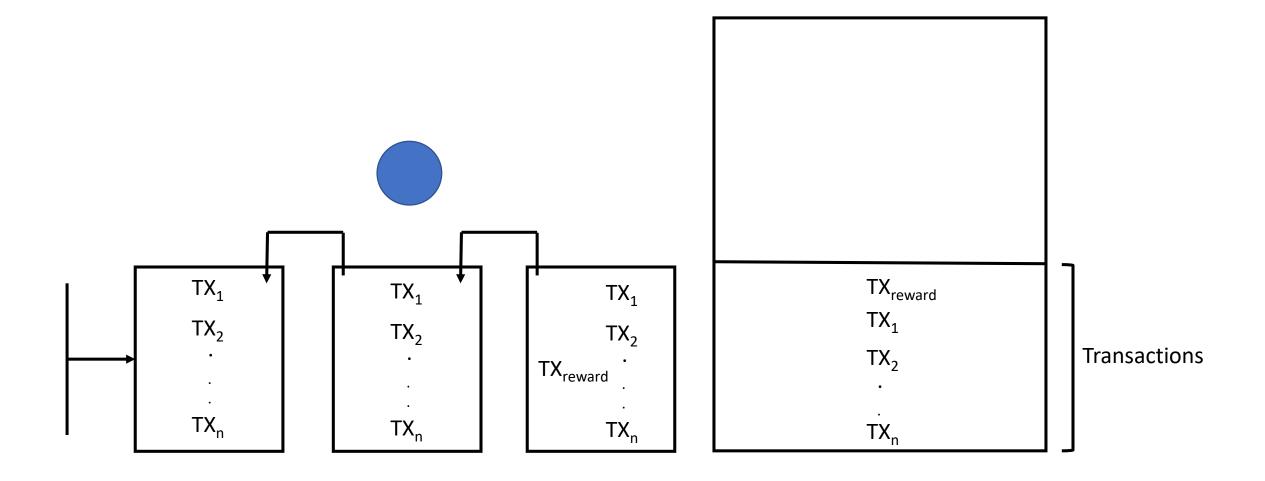


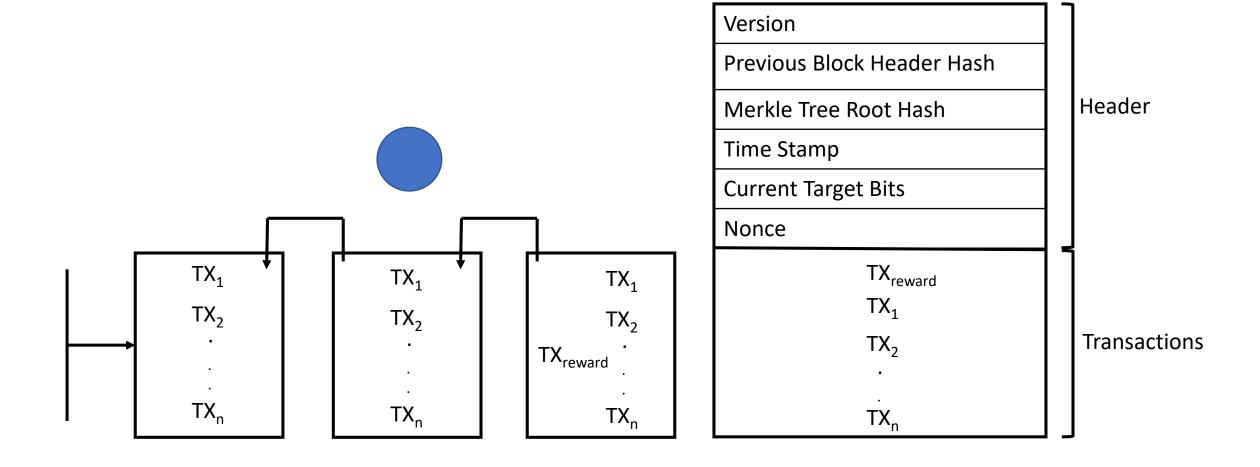


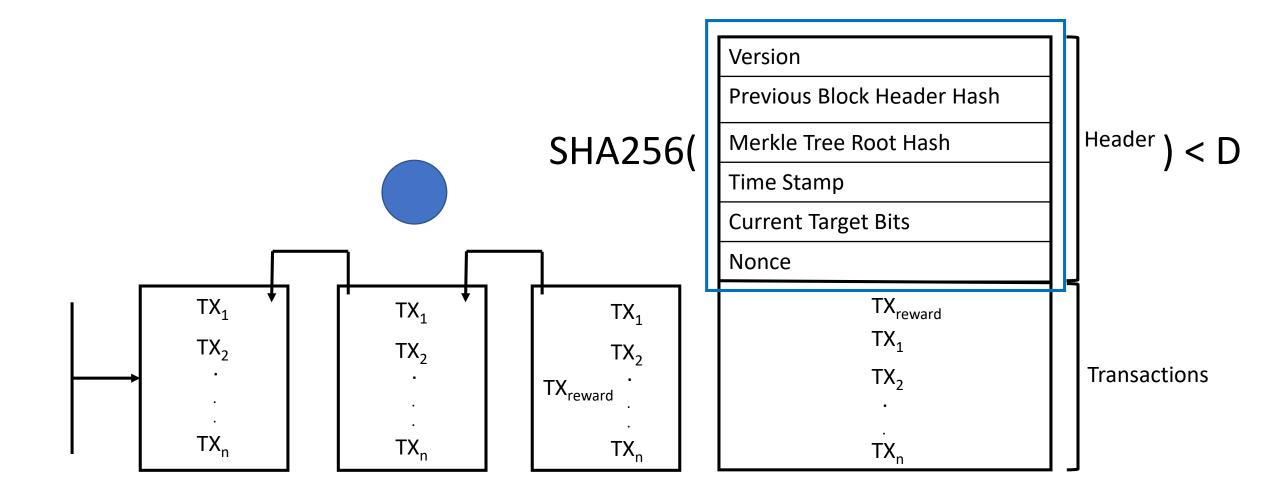


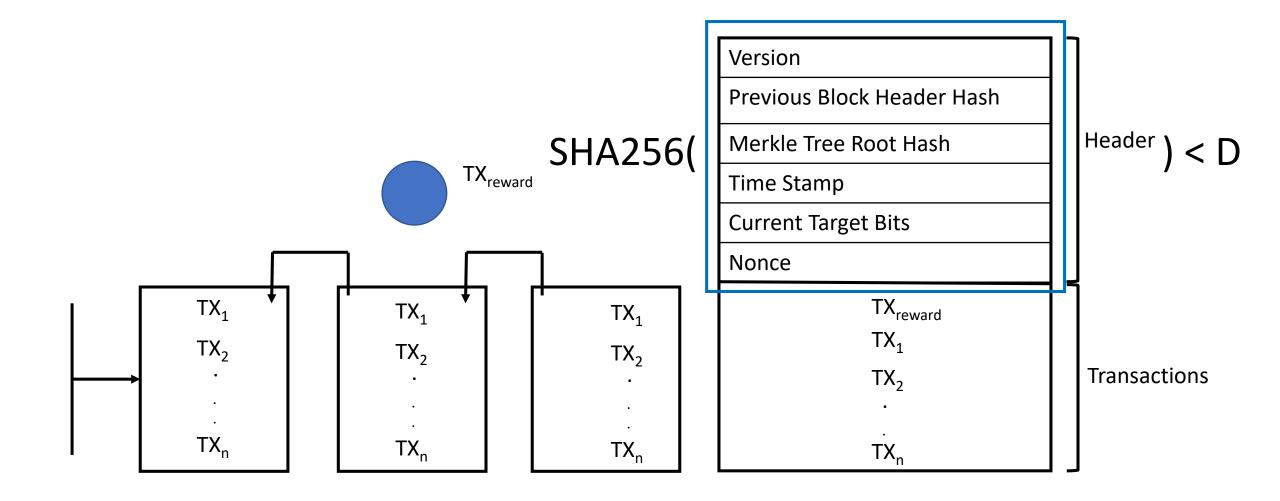


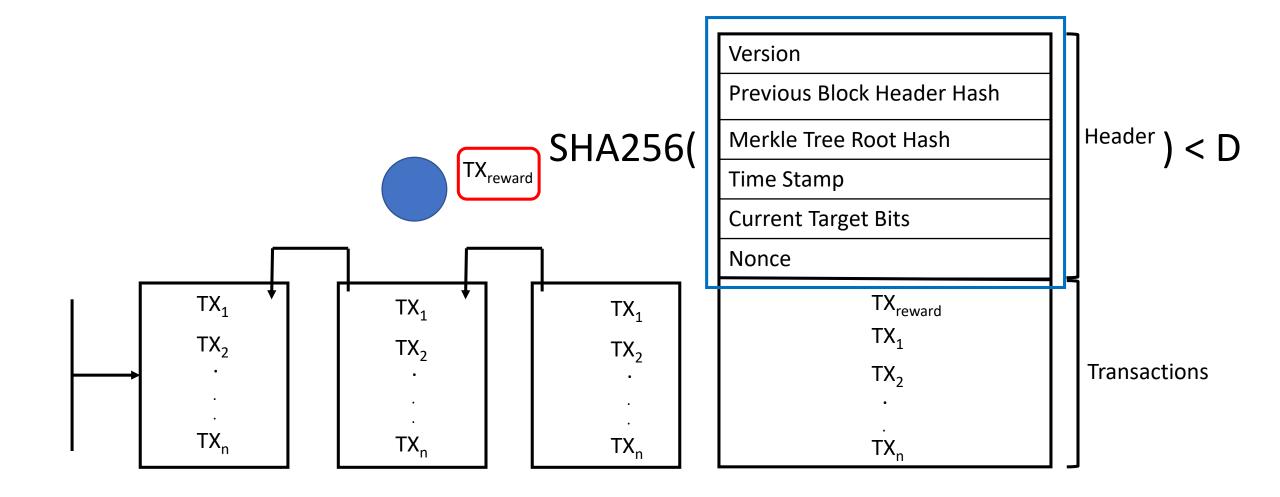




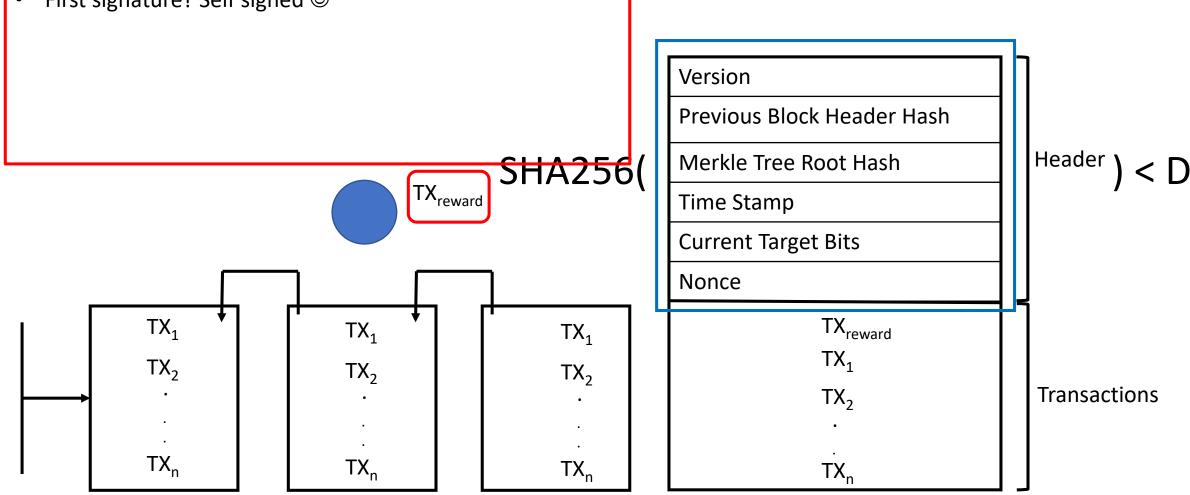




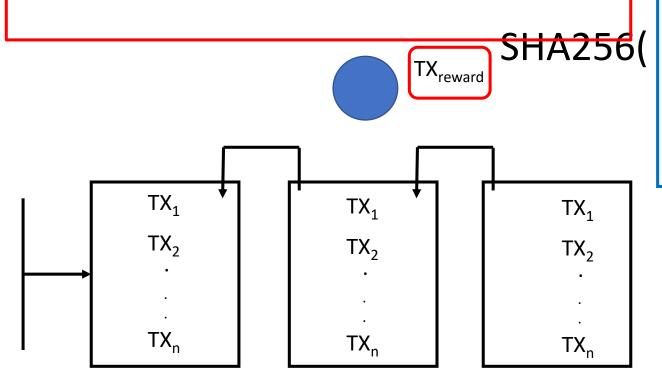


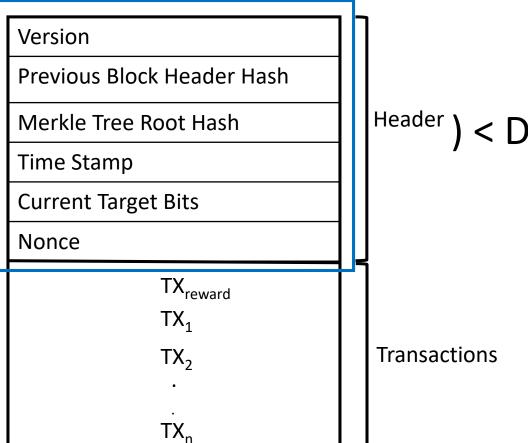


- TX<sub>reward</sub> is self signed (also called coinbase transaction)
- First signature? Self signed ©

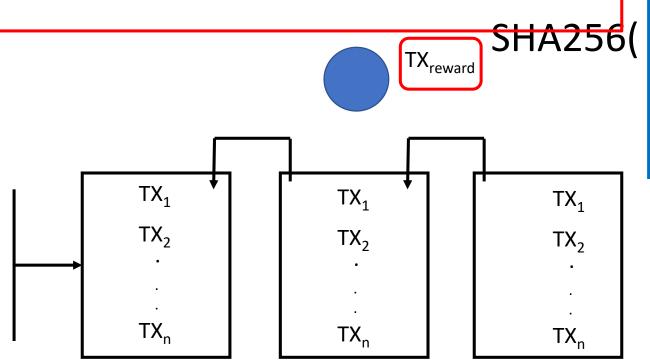


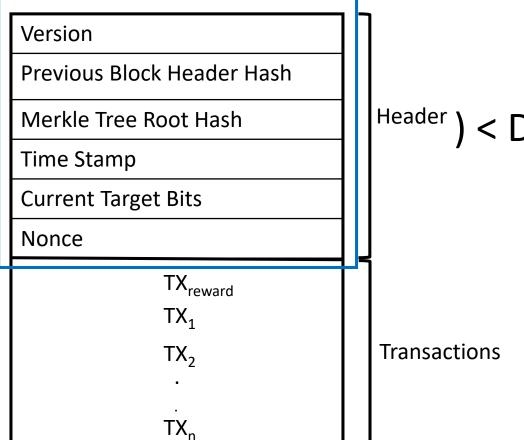
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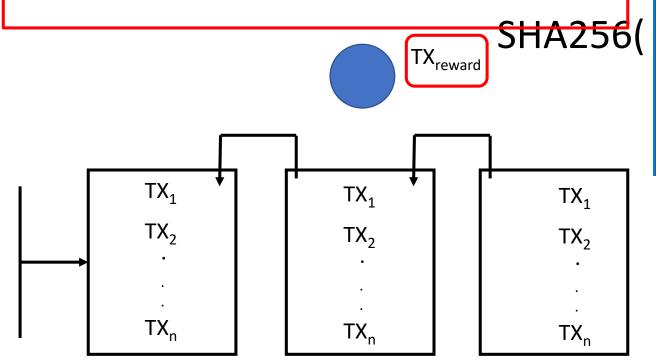


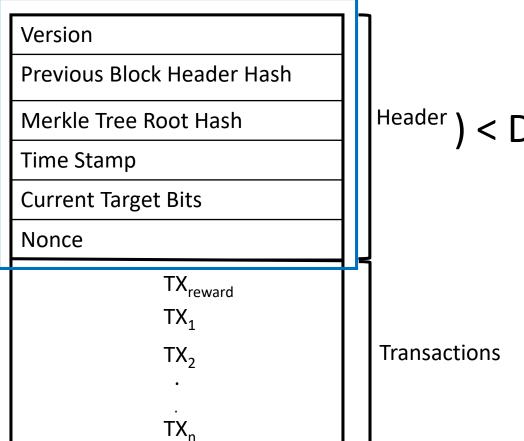
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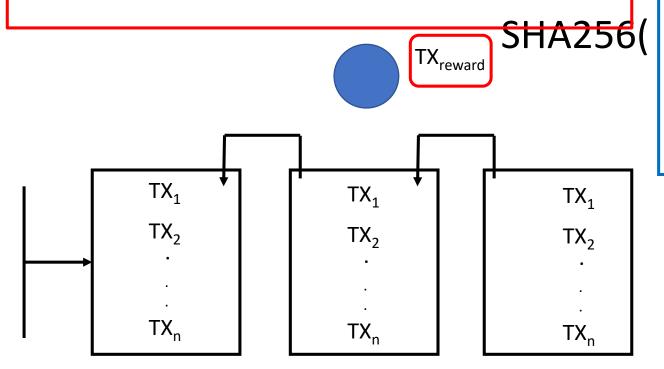


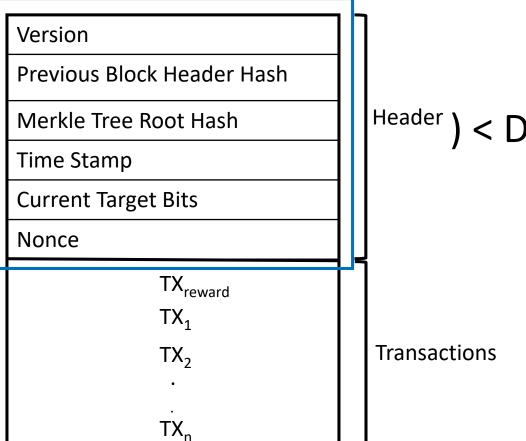
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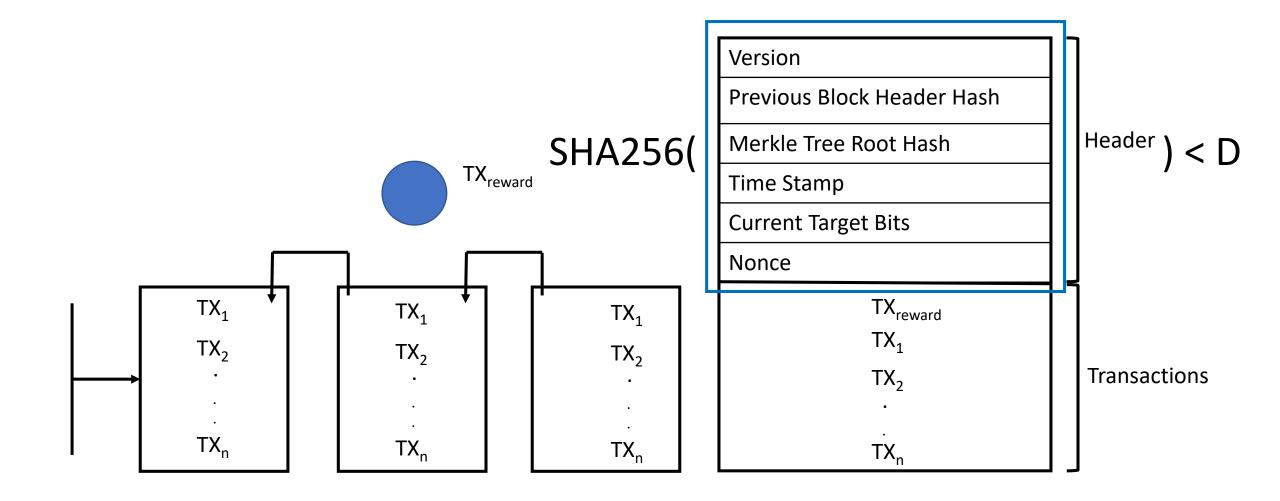


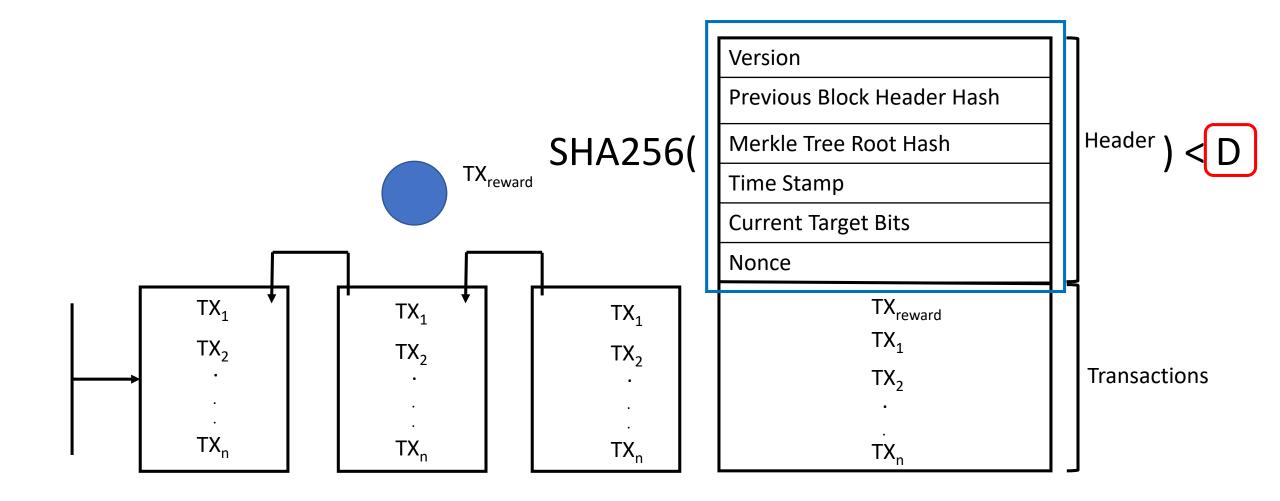


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- Currently, it's 12.5 Bitcoins per block
- Incentives network nodes to mine



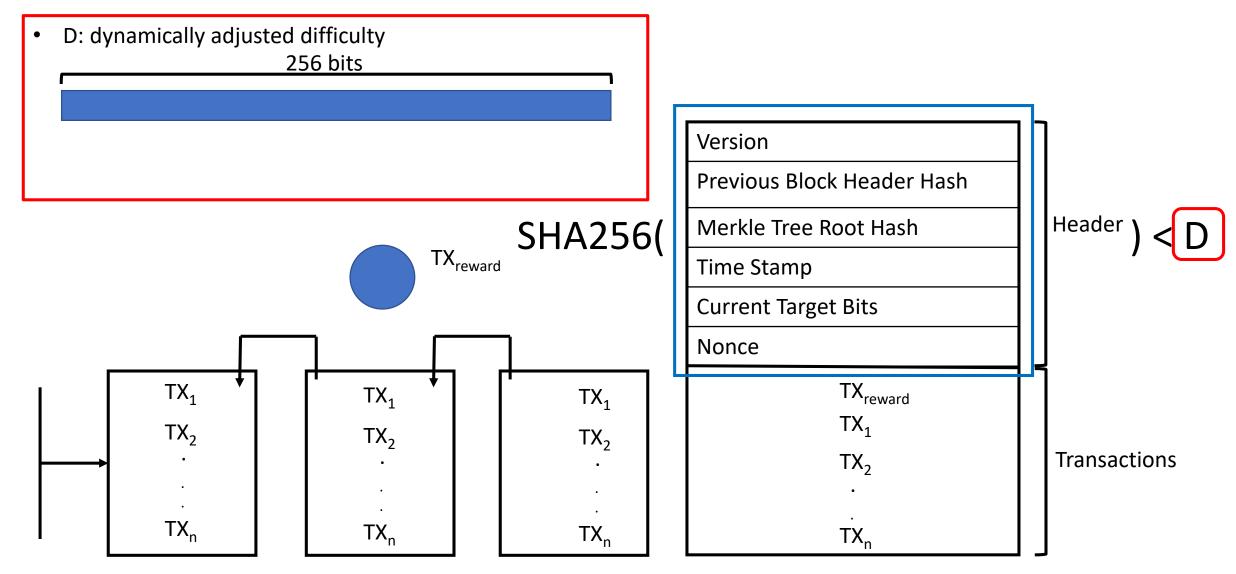


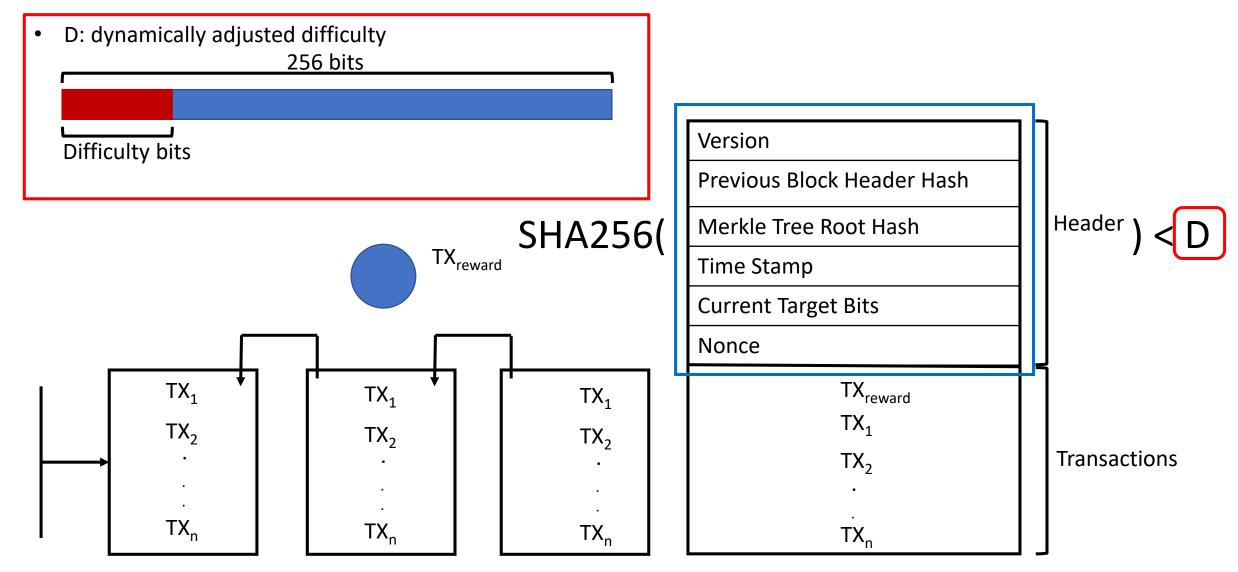


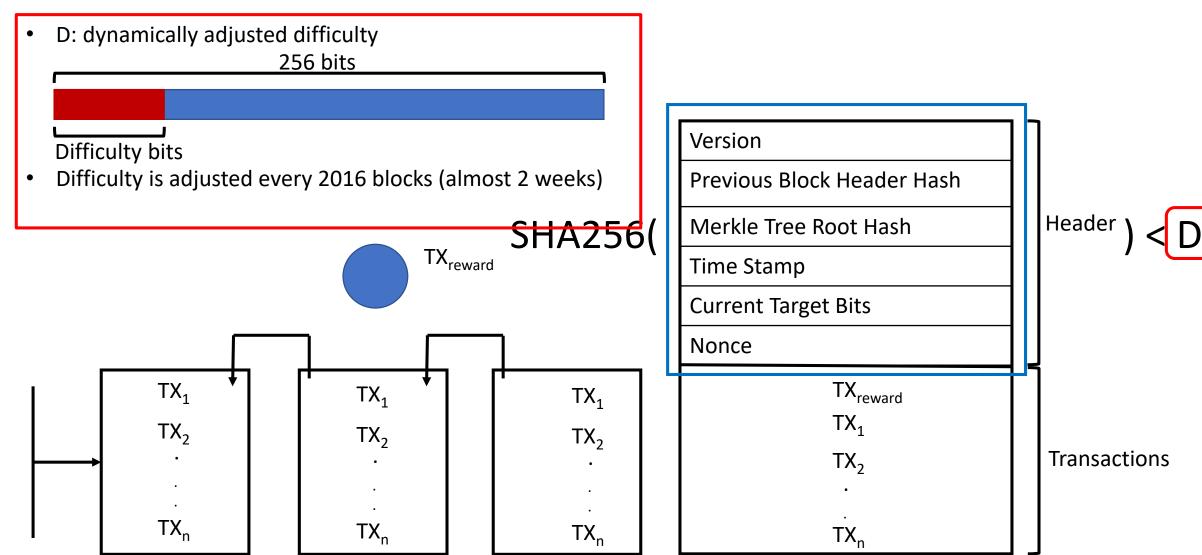


## Mining Details

D: dynamically adjusted difficulty Version Previous Block Header Hash Header Merkle Tree Root Hash SHA256(  $\mathsf{TX}_{\mathsf{reward}}$ Time Stamp **Current Target Bits** Nonce  $\mathsf{TX}_{\mathsf{reward}}$  $TX_1$  $\mathsf{TX}_1$  $\mathsf{TX}_1$  $\mathsf{TX}_1$  $TX_2$  $TX_2$  $TX_2$ **Transactions**  $TX_2$  $TX_n$  $TX_n$  $TX_n$  $TX_n$ 







## Difficulty

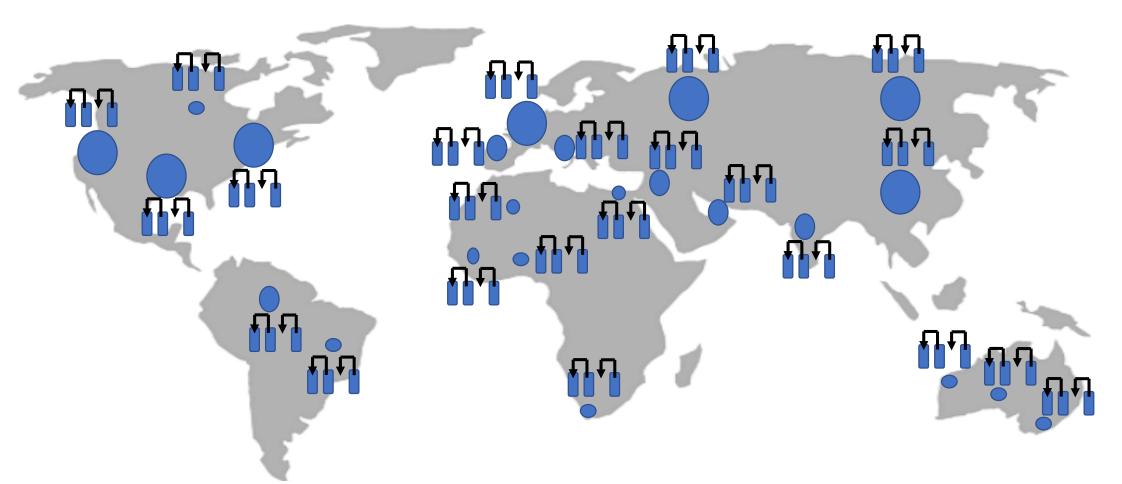
Adjust difficulty every 2016 blocks

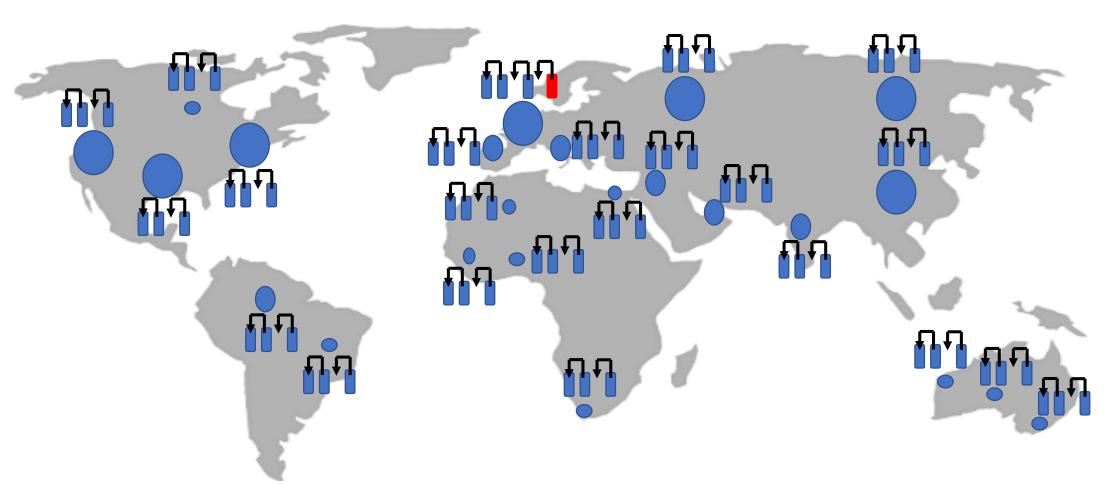
- Adjust difficulty every 2016 blocks
- Expected 20160 mins to mine (10 mins per block)

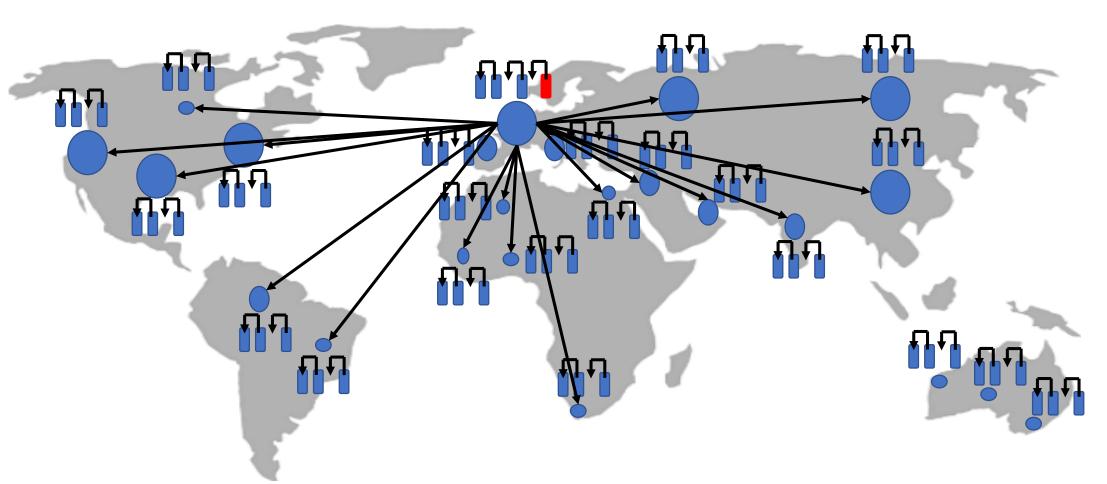
- Adjust difficulty every 2016 blocks
- Expected 20160 mins to mine (10 mins per block)
- Actual time = timestamp of block 2016 time stamp of block 1

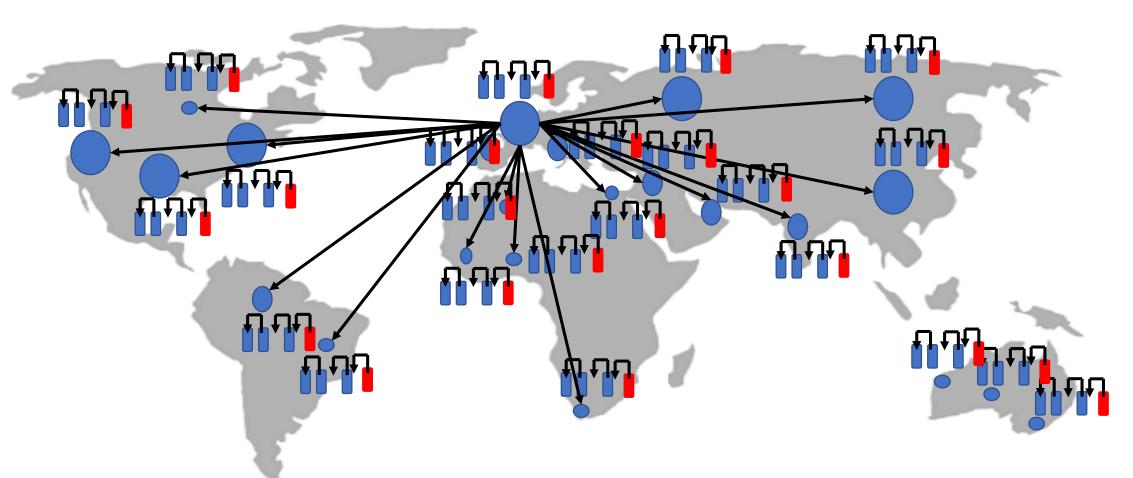
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- Expected 20160 mins to mine (10 mins per block)
- Actual time = timestamp of block 2016 time stamp of block 1
- New\_difficulty = old\_difficulty \* expected/actual
- Difficulty decreases if actual > expected, otherwise, increases









### Mining Details

Find a nonce that results in SHA256(block) < Difficulty</li>

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- What happens when 2 nodes concurrently mine a block? Fork

## Mining Details

• Find a nonce that results in SHA256(block) < Difficulty

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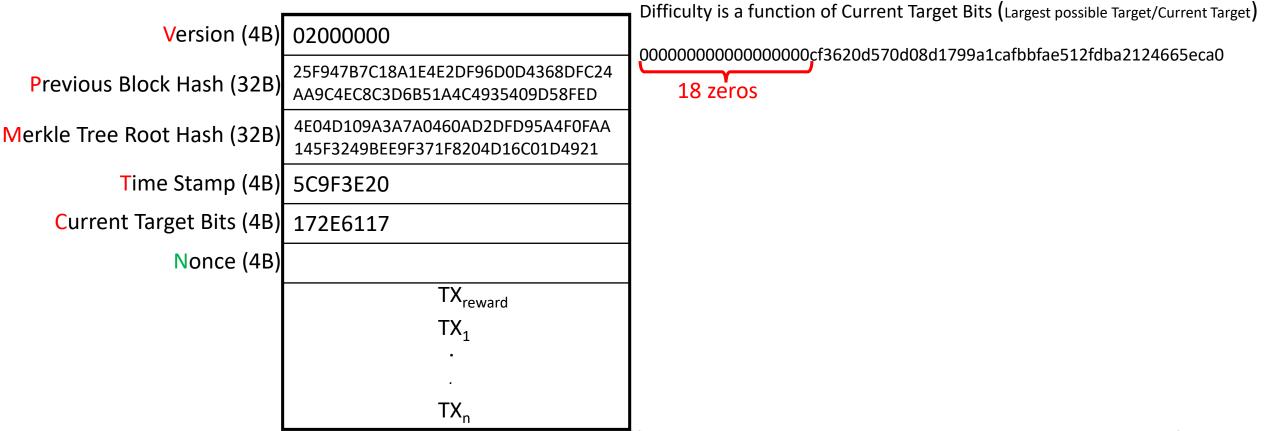
• Find a nonce that results in SHA256(block) < Difficulty

Version (4B)	02000000
Previous Block Hash (32B)	25F947B7C18A1E4E2DF96D0D4368DFC24 AA9C4EC8C3D6B51A4C4935409D58FED
Merkle Tree Root Hash (32B)	4E04D109A3A7A0460AD2DFD95A4F0FAA 145F3249BEE9F371F8204D16C01D4921
Time Stamp (4B)	5C9F3E20
Current Target Bits (4B)	172E6117
Nonce (4B)	
	TX <sub>reward</sub>
	$TX_1$
	•
	•
	TX <sub>n</sub>

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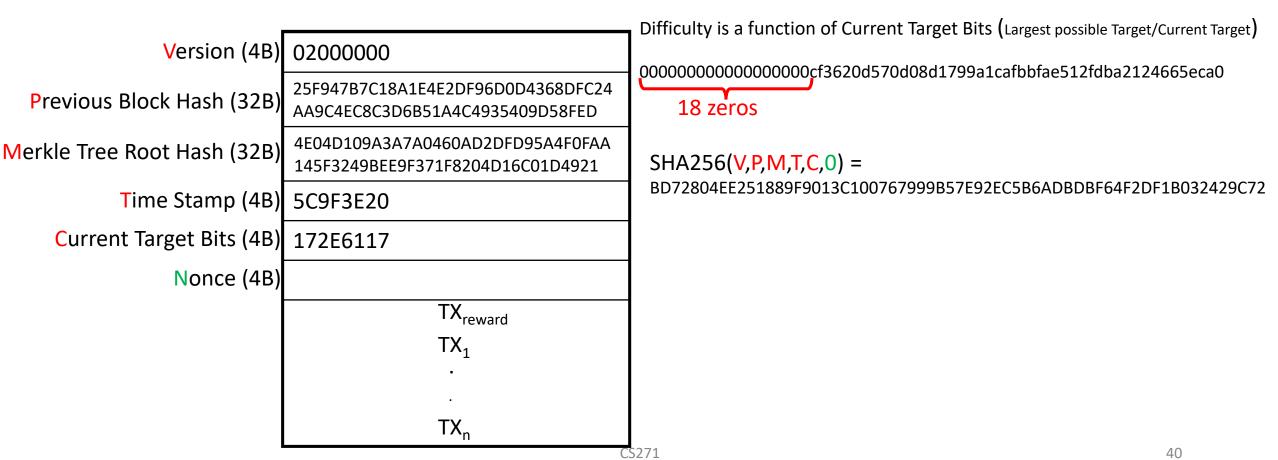


Find a nonce that results in SHA256(block) < Difficulty</li>



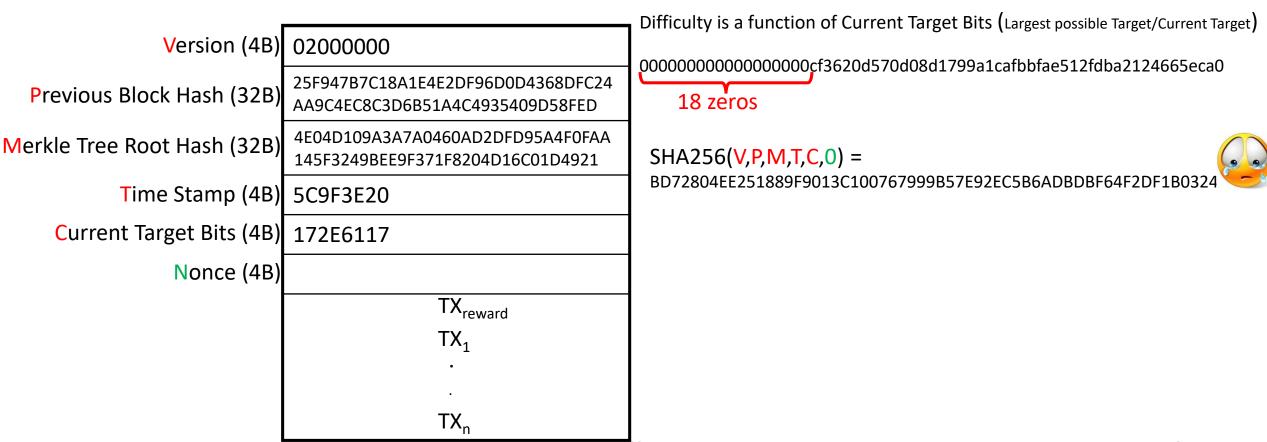


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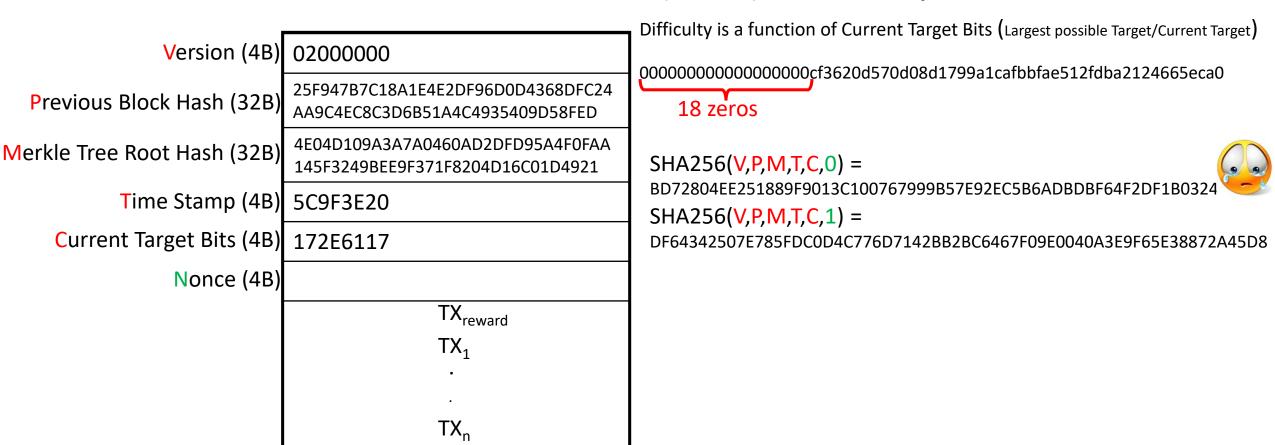


Find a nonce that results in SHA256(block) < Difficulty</li>





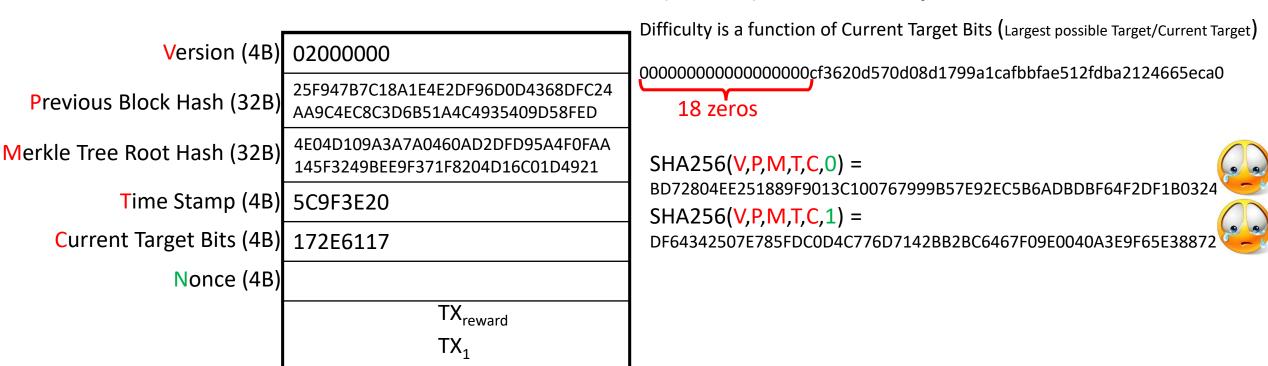
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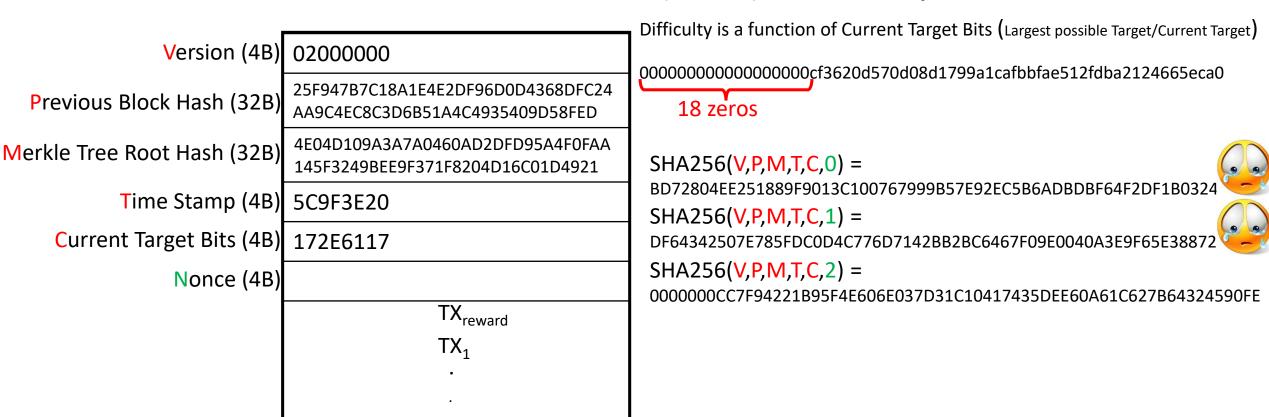
 $TX_n$ 





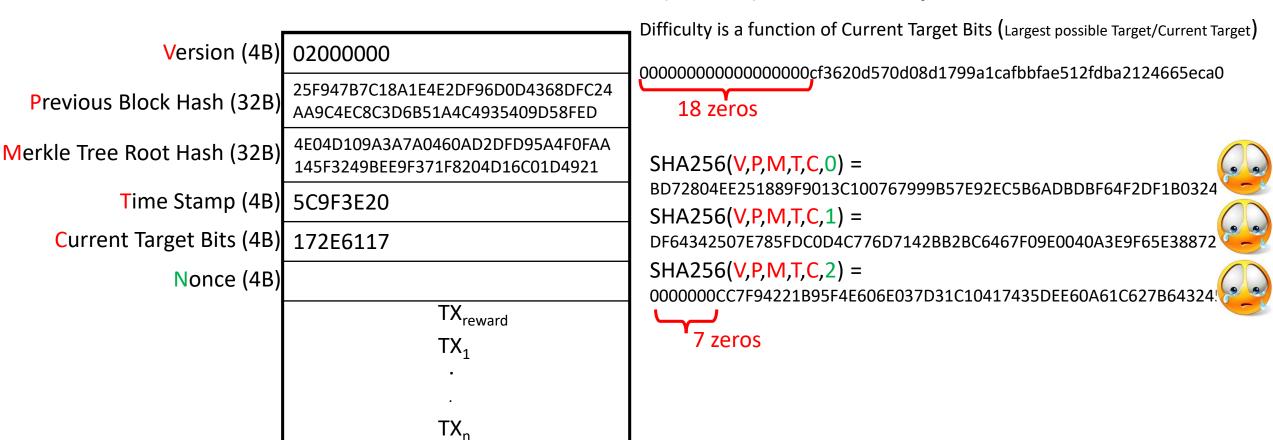
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 $TX_n$ 



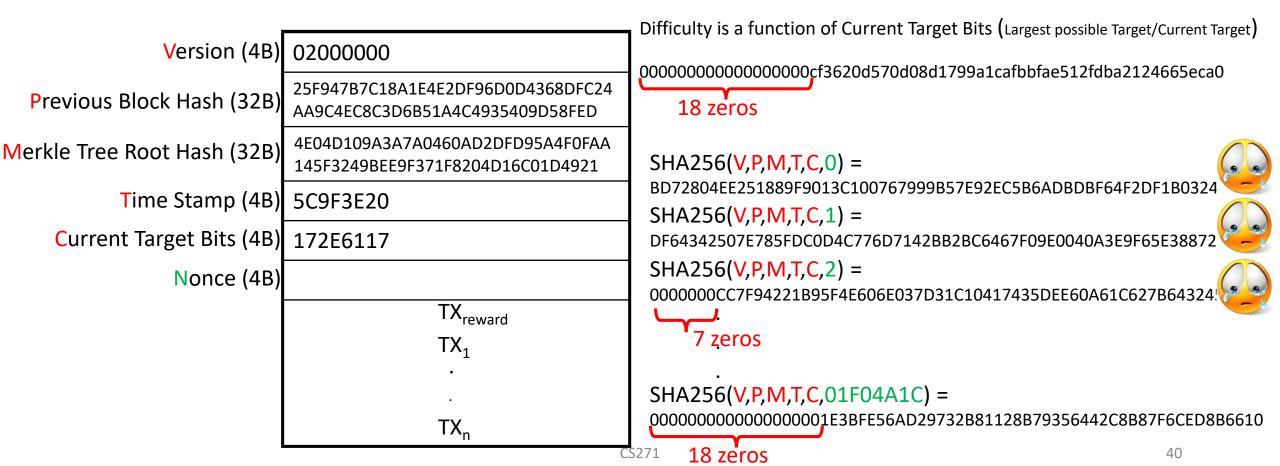


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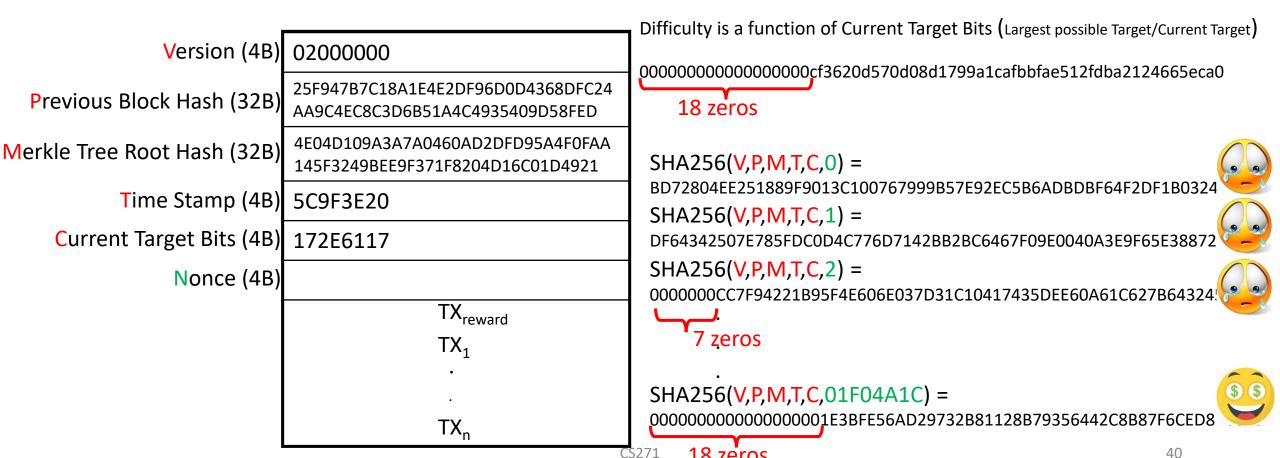
### Mining Details

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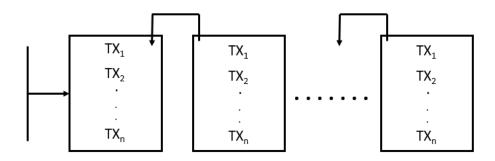


### Mining Details

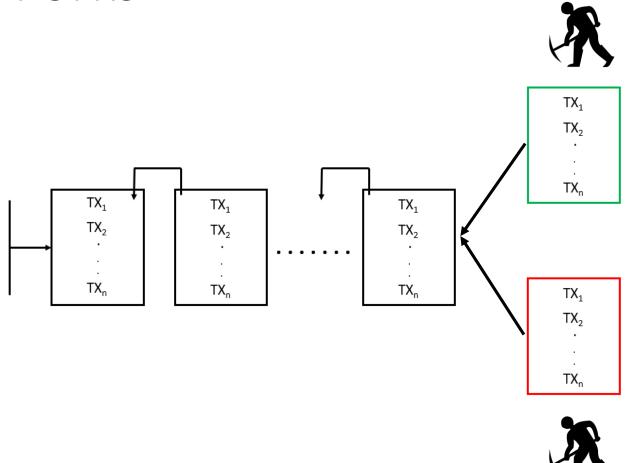
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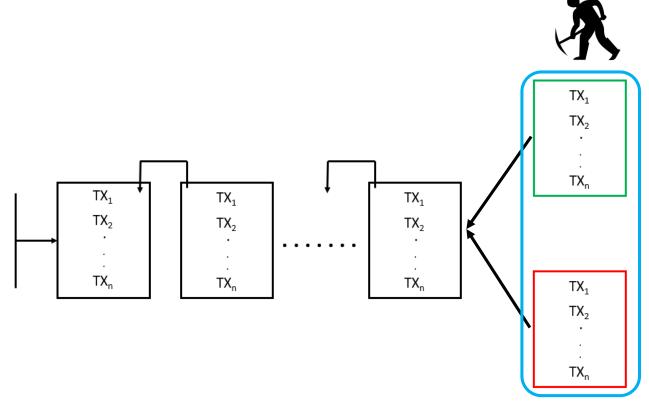
# Forks



# Forks



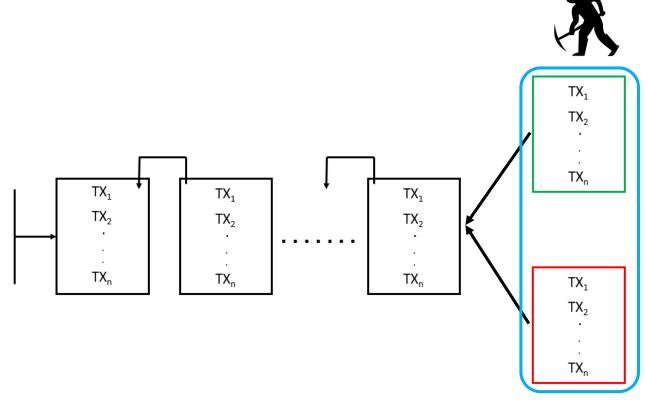
## Forks



Transactions in the forked blocks might have conflicts

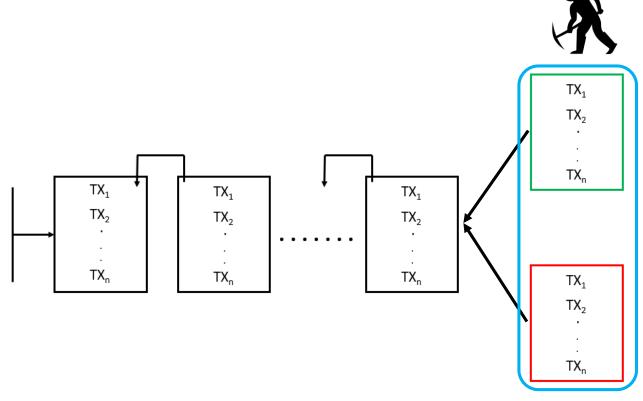


## Forks



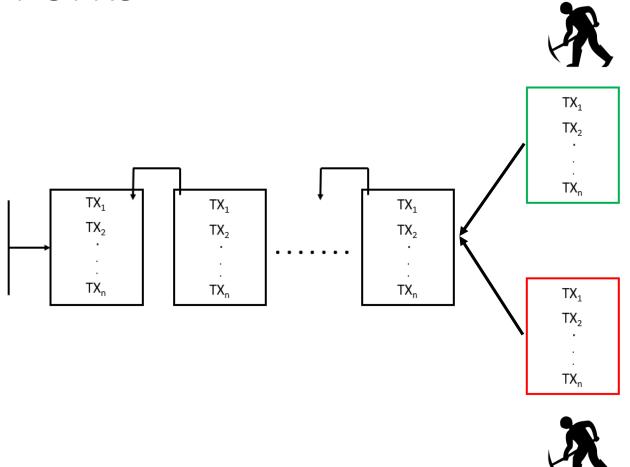
- Transactions in the forked blocks might have conflicts
- Could lead to double spending

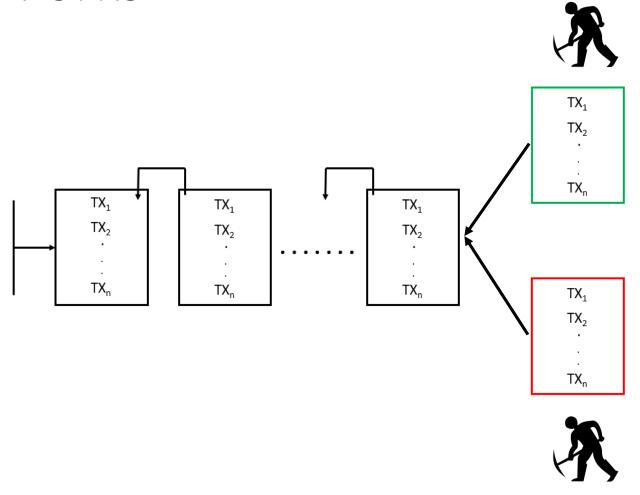


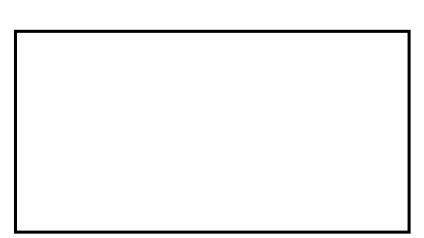


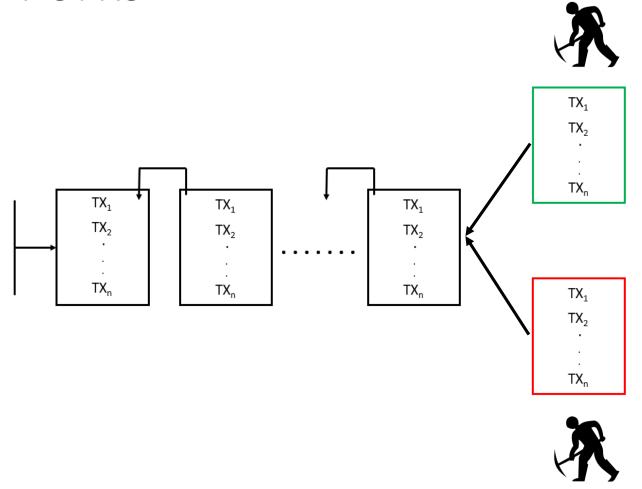
- Transactions in the forked blocks might have conflicts
- Could lead to double spending
- Forks have to be eliminated

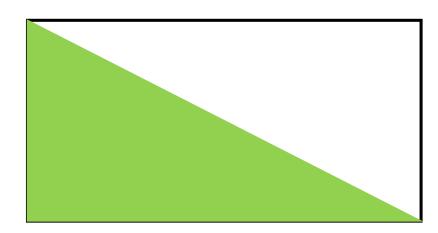


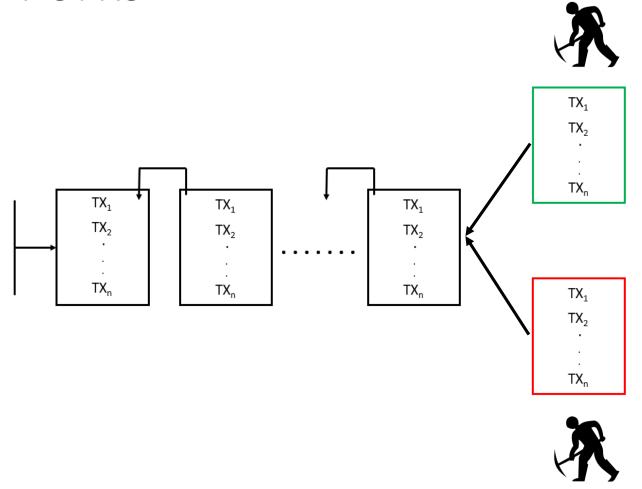


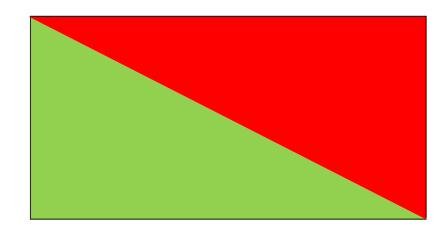


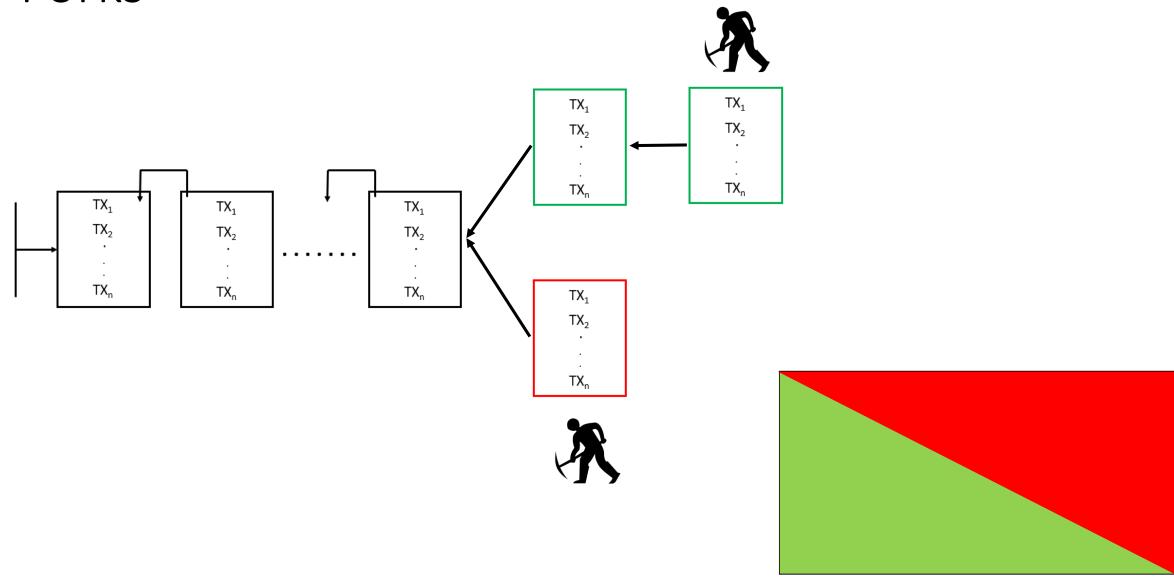


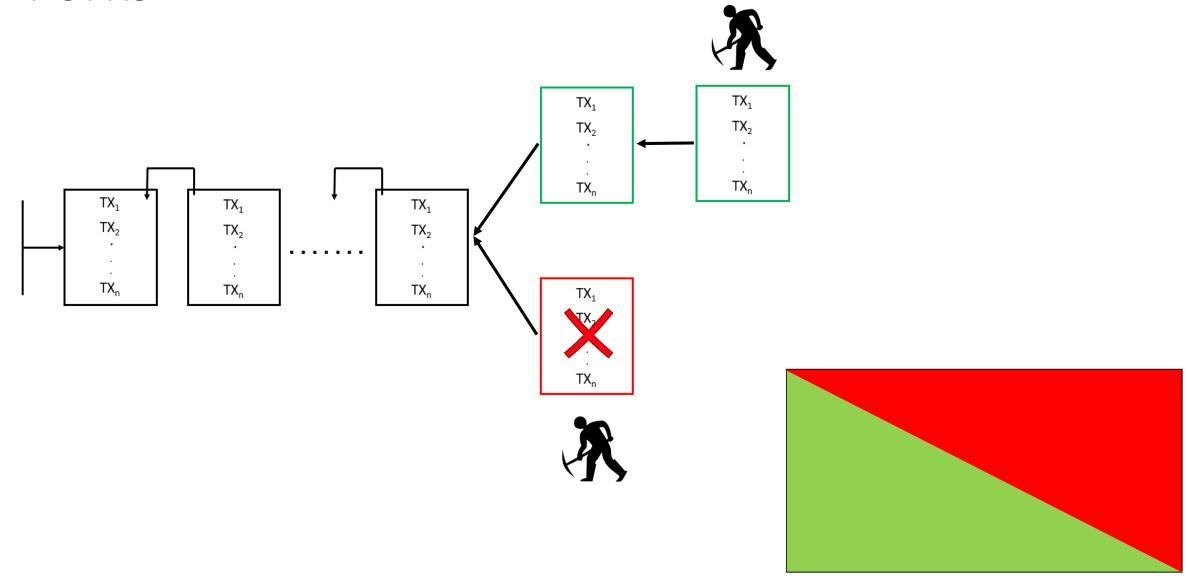




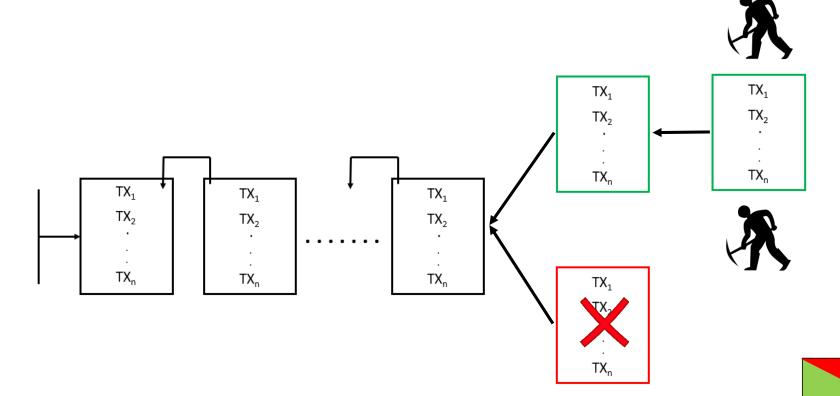




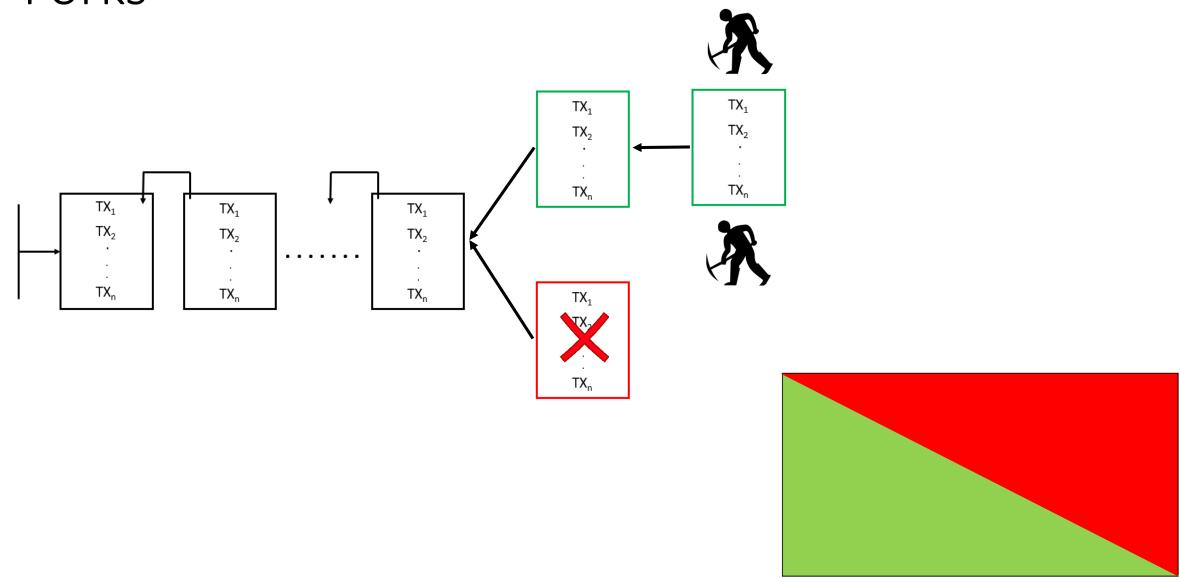




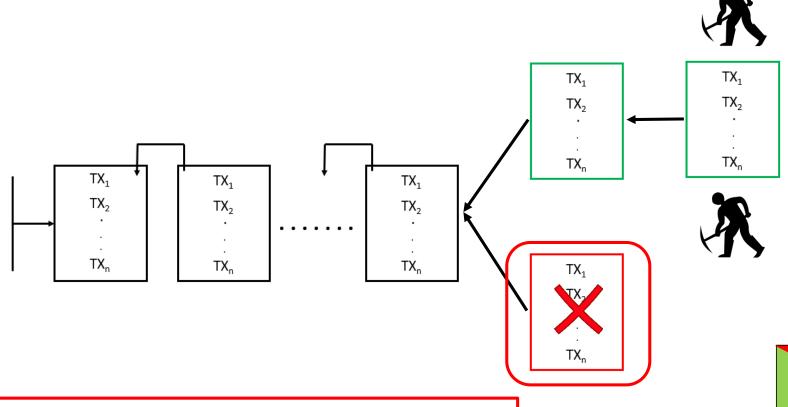
# Forks



Miners join the longest chain to resolve forks

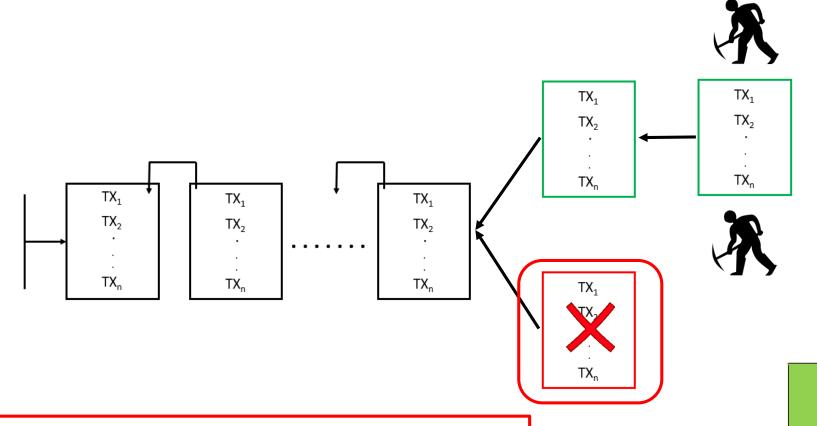


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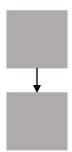


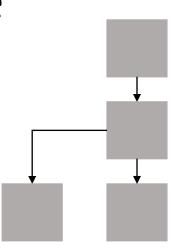
Transactions in this block have to be resubmitted

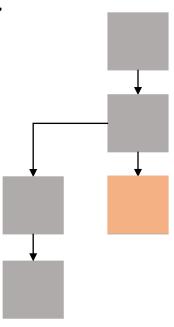
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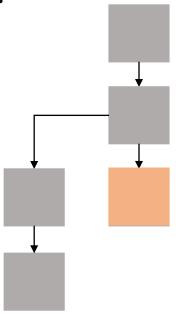
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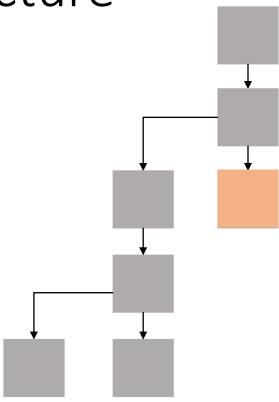


Forks: The Big Picture



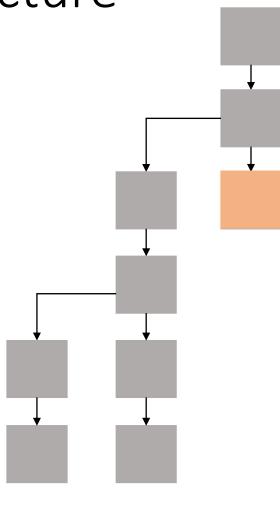
Abandoned

Forks: The Big Picture



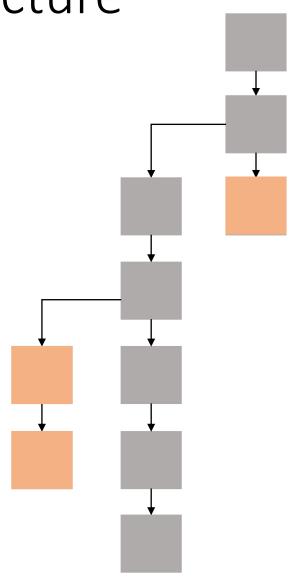
Abandoned

Forks: The Big Picture



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Forks: The Big Picture



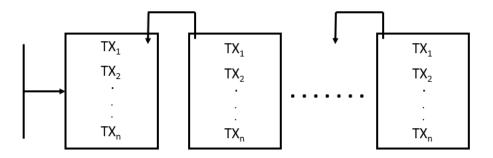
Abandoned

## 51% Attack

- If 51% of the computation (hash) power are malicious:
  - They can cooperate to fork the chain at any block
- Can lead to double spending

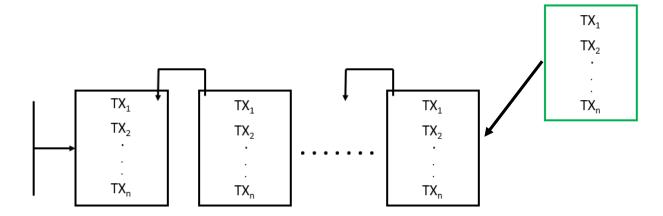
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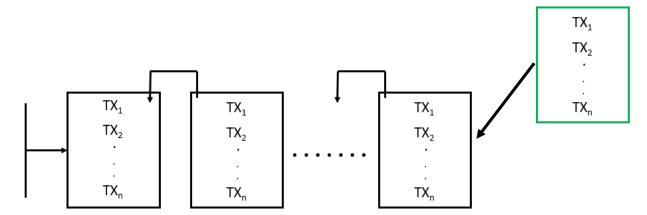
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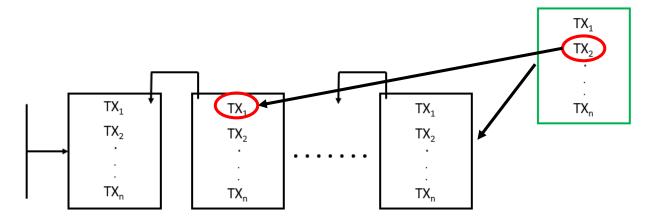




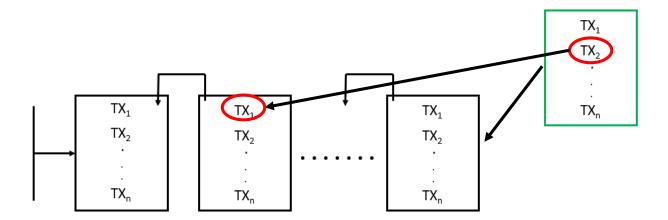
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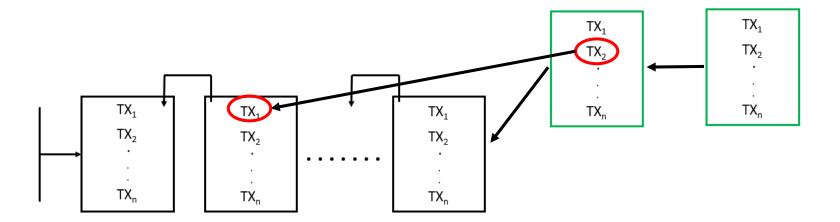


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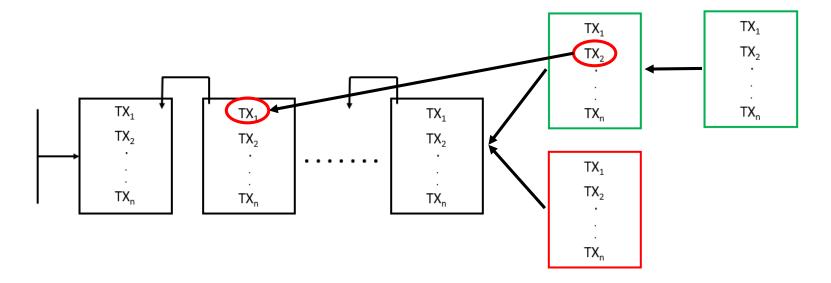


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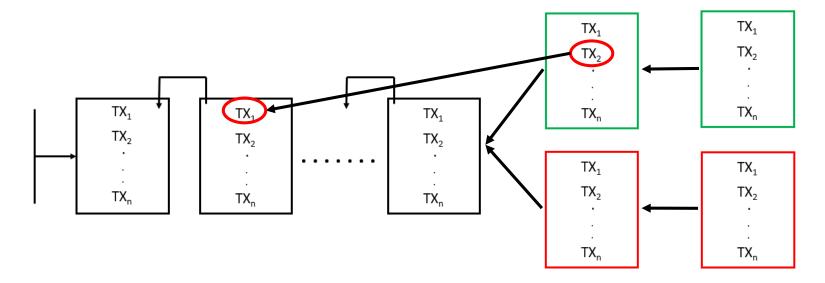


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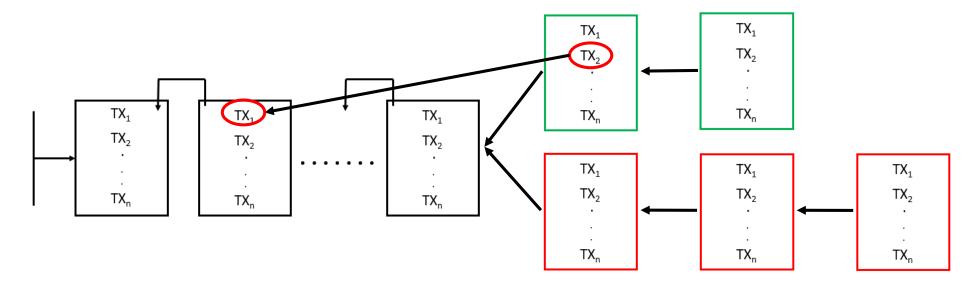


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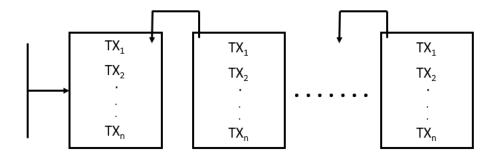
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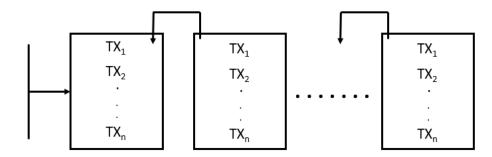








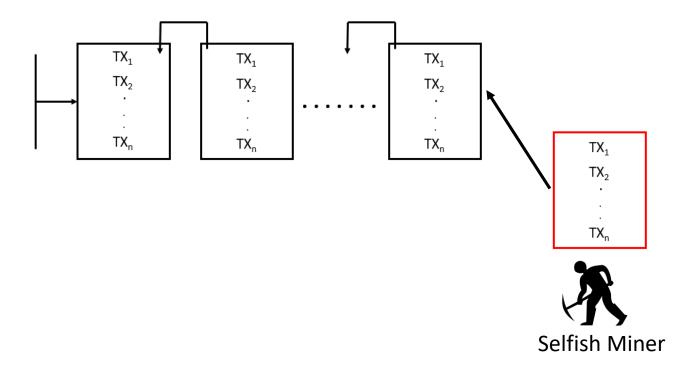






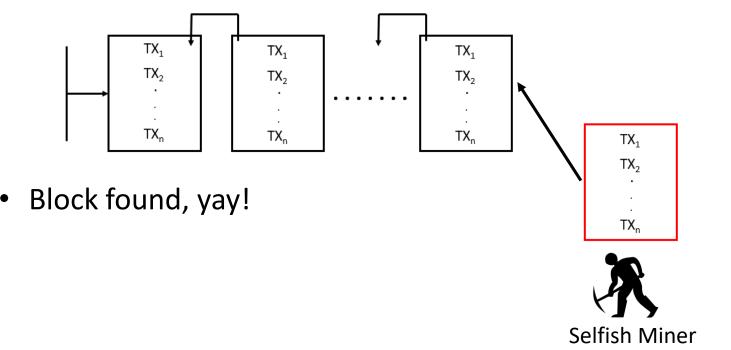






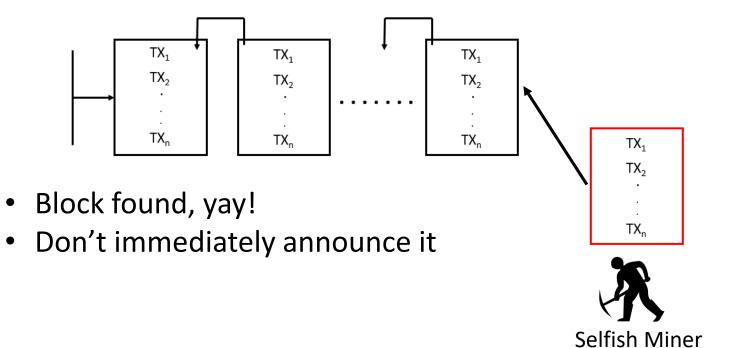








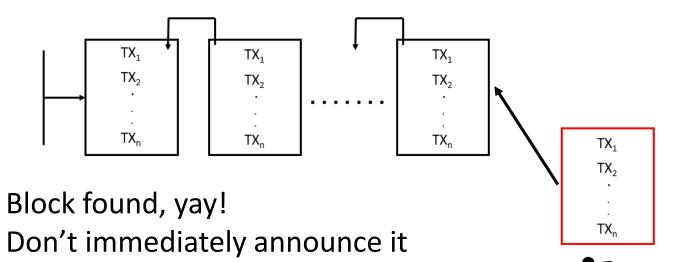






**Honest Miner** 

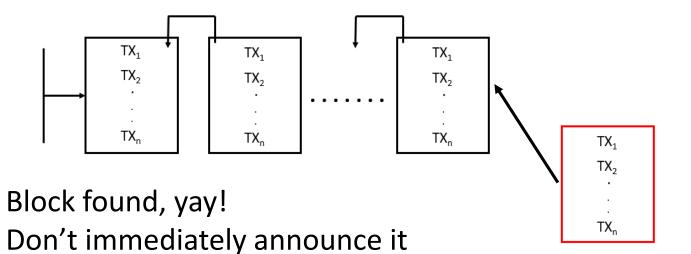






#### **Honest Miner**







Start mining the next block (Advantage)

Let honest miners waste their mining

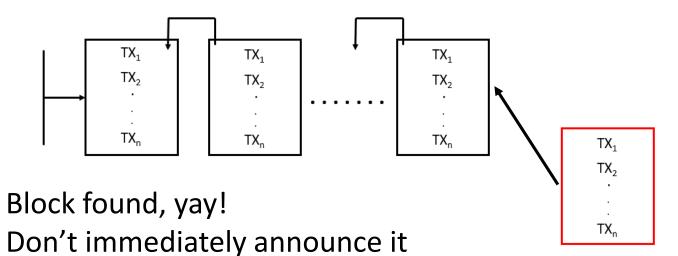
power on an obsolete block

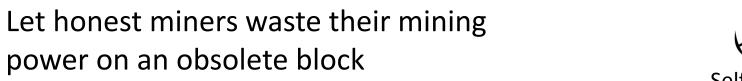








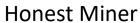




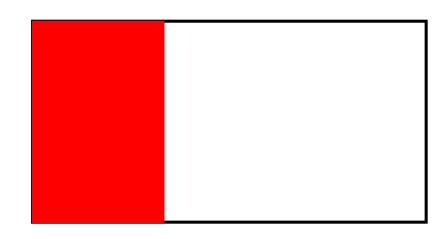
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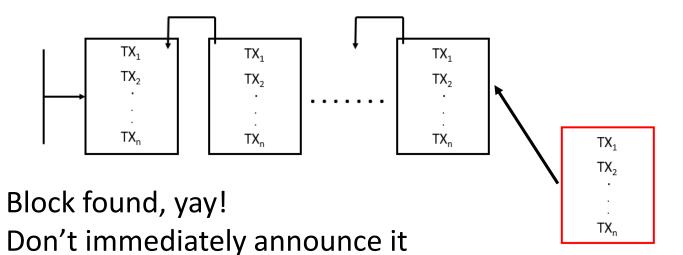


## Selfish Mining











Start mining the next block (Advantage)

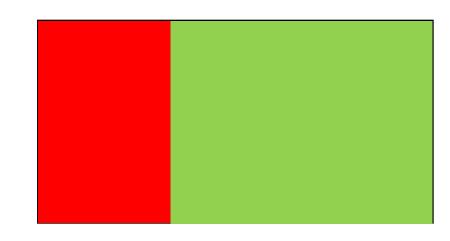
Let honest miners waste their mining

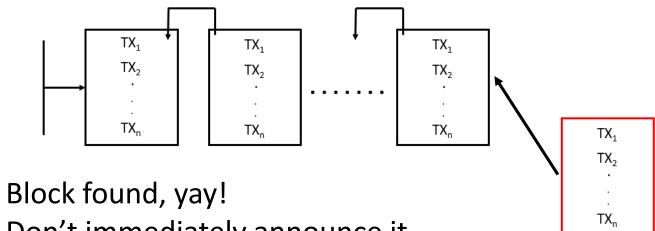
power on an obsolete block

## Selfish Mining









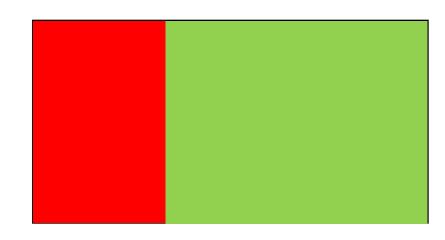
- Don't immediately announce it
- Let honest miners waste their mining power on an obsolete block
- Start mining the next block (Advantage)

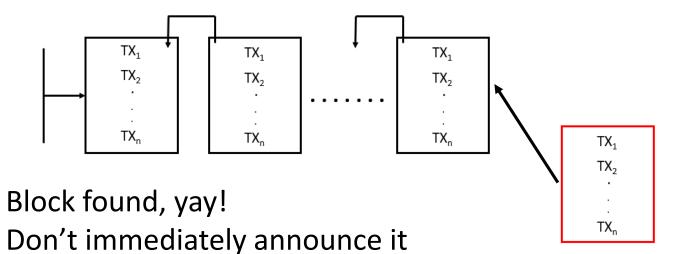


## Selfish Mining











Start mining the next block (Advantage)

Let honest miners waste their mining

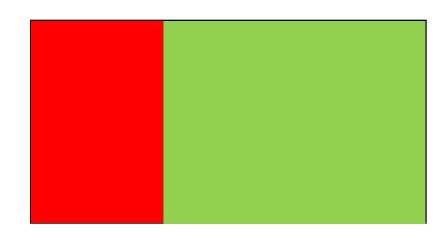
Two possible outcomes

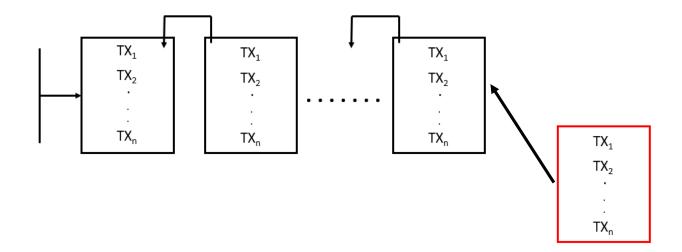
power on an obsolete block









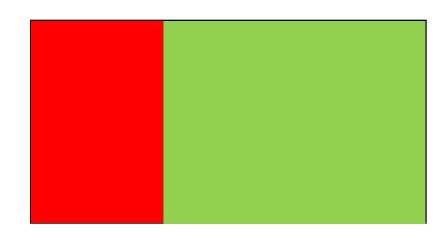


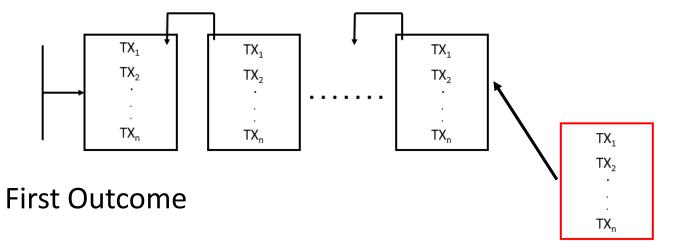




**Honest Miner** 





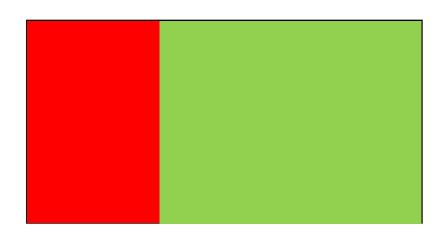


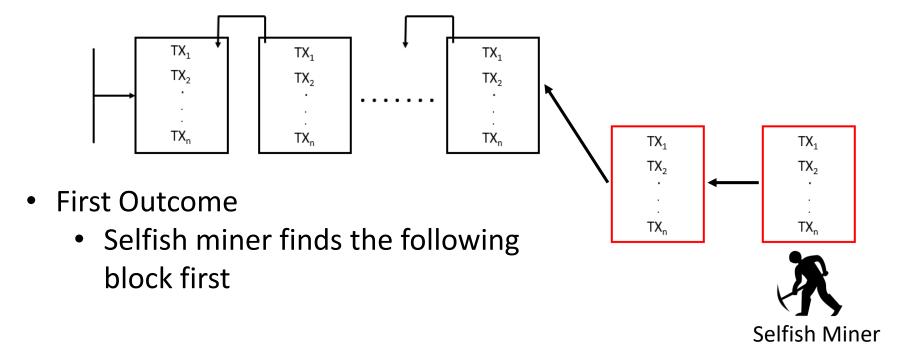










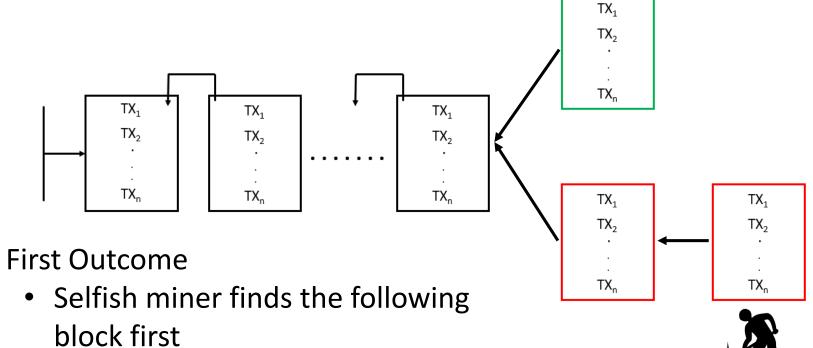


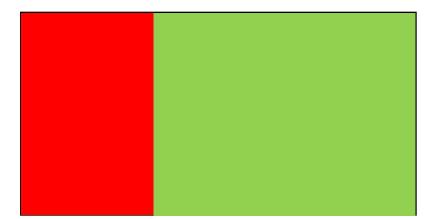
Selfish Mining

Once an honest miner finds a block

**Honest Miner** 







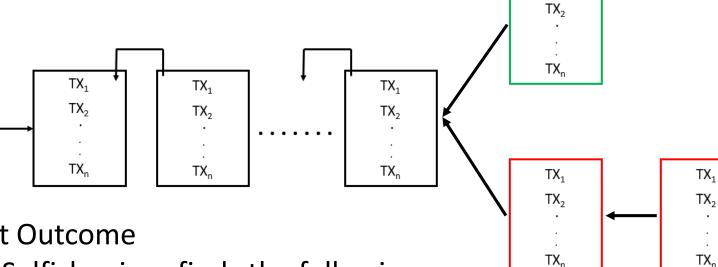
Selfish Miner

Selfish Mining

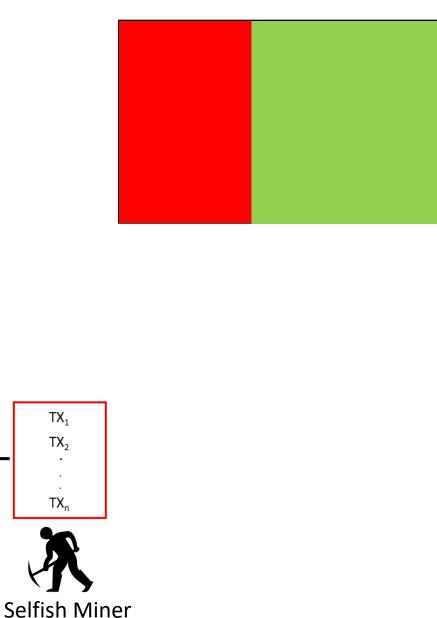
**Honest Miner** 



 $TX_1$ 



- First Outcome
  - Selfish miner finds the following block first
  - Once an honest miner finds a block
    - Selfish miner announces 2 blocks

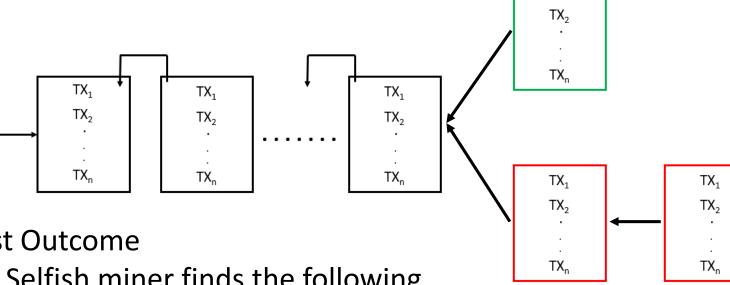


Selfish Mining

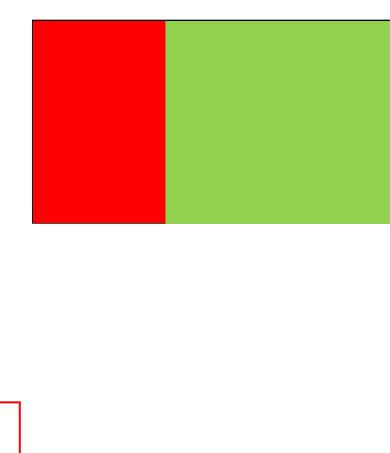
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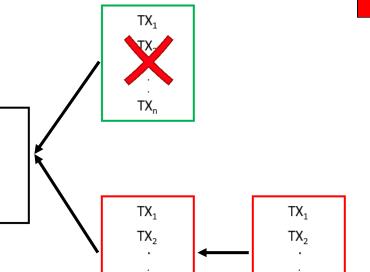


Selfish Miner

Selfish Mining

**Honest Miner** 





 $TX_n$ 

First Outcome

 $TX_1$ 

 $TX_2$ 

 Selfish miner finds the following block first

 $TX_1$ 

 $TX_2$ 

 $TX_1$ 

 $TX_2$ 

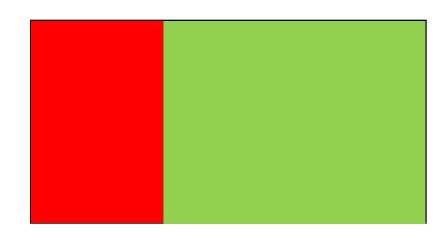
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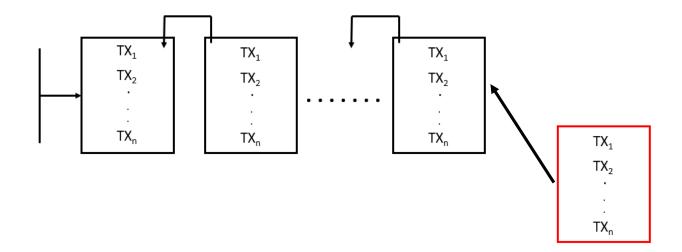










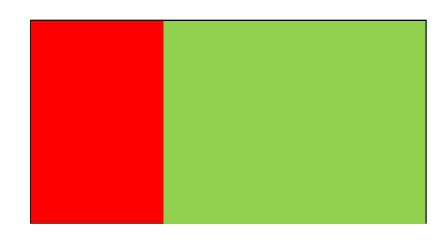


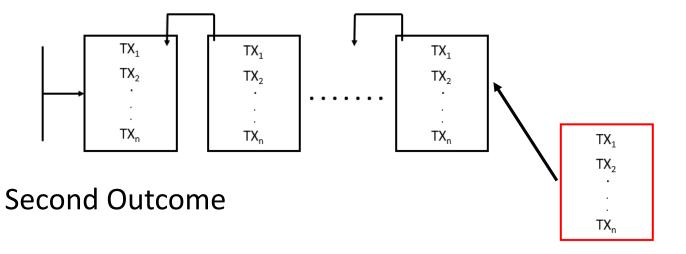












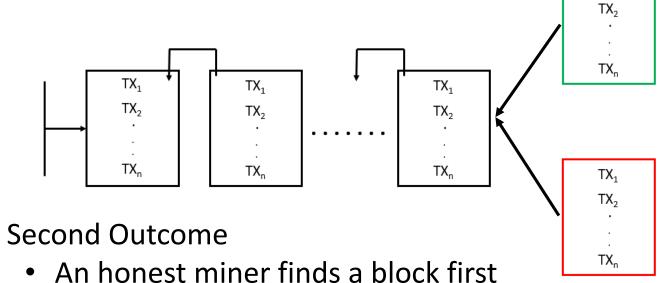


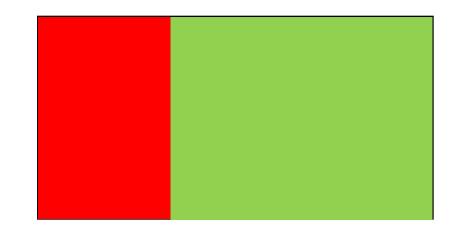
Selfish Mining

**Honest Miner** 



 $\mathsf{TX}_1$ 





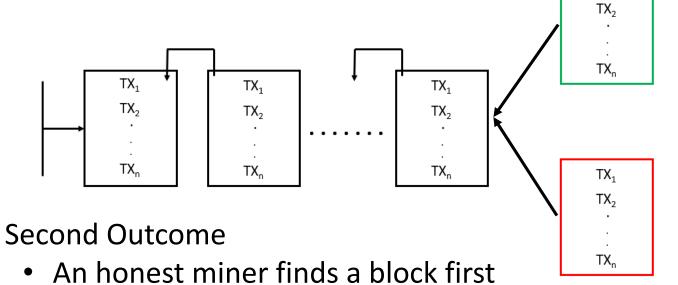


Selfish Mining

**Honest Miner** 



 $TX_1$ 



Selfish miner immediately announces

the previously found block



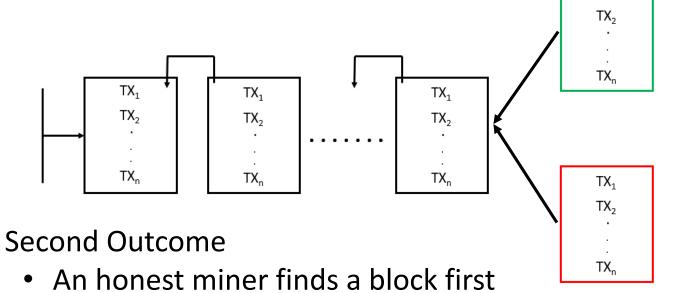


Selfish Mining

**Honest Miner** 



 $\mathsf{TX}_1$ 

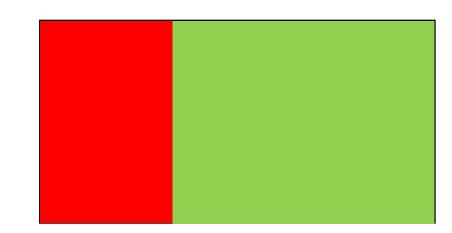




• This splits the power of honest miners

the previously found block

Selfish miner immediately announces

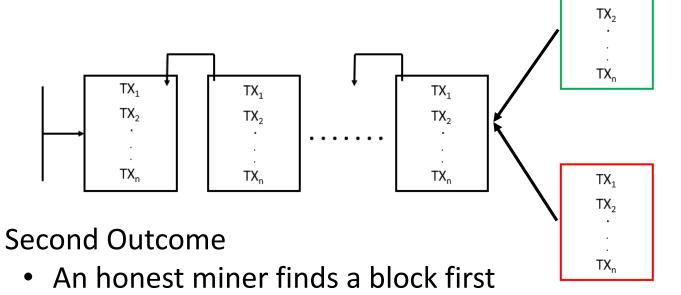


Selfish Mining

**Honest Miner** 



 $\mathsf{TX}_1$ 





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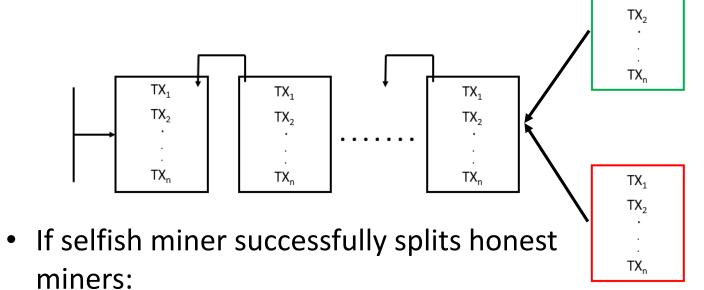


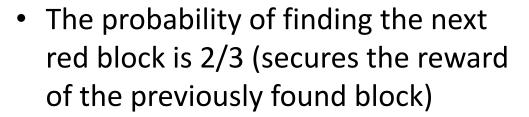
Selfish Mining

**Honest Miner** 

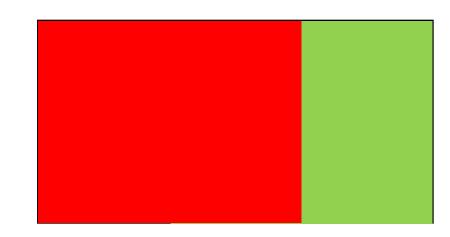


 $\mathsf{TX}_1$ 







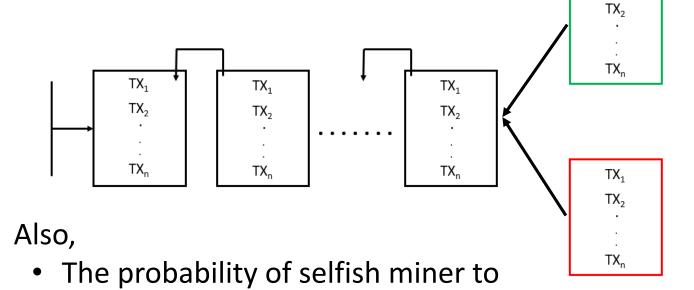


Selfish Mining

**Honest Miner** 



 $\mathsf{TX}_1$ 



find the next red block is 1/2 even if

selfish miner has 1/3 of the mining

resources (Advantage)





### Limitations of Bitcoin

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High transaction-confirmation latency

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New block added every 10 minutes.

How to scale Bitcoin?

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• Two obvious options for increasing Bitcoin's transaction throughput:

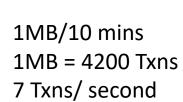
### How to scale Bitcoin?

 Two obvious options for increasing Bitcoin's transaction throughput: increase the size of blocks, or decrease the block interval

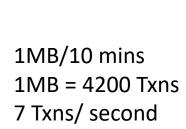
## Increasing Block Size

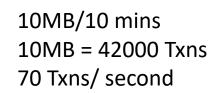
1MB/10 mins 1MB = 4200 Txns 7 Txns/ second

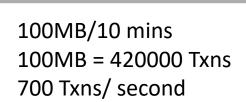
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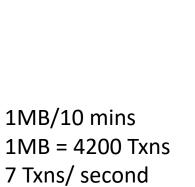


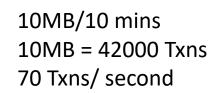
10MB/10 mins 10MB = 42000 Txns 70 Txns/ second

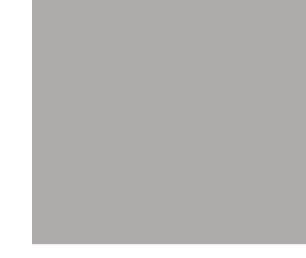












100MB/10 mins 100MB = 420000 Txns 700 Txns/ second

- Why they don't work?
  - Decreases fairness giving large miners an advantage
  - Requires more storage space (1  $\rightarrow$  10  $\rightarrow$  100 MB/ 10 mins)
  - Requires more Network bandwidth
  - Requires more verification time

# Decrease Block Interval

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1MB/10 mins 1MB = 4200 Txns 7 Txns/ second

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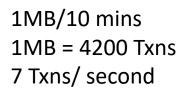
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1MB/1 min 1MB = 4200 Txns 70 Txns/ second

# Decrease Block Interval



1MB/5 mins 1MB = 4200 Txns 14 Txns/ second

1MB/1 min 1MB = 4200 Txns 70 Txns/ second

## Decrease Block Interval

- Requires to mining decrease difficulty
- Leads to more forks
- Results on network instability (many branches)

# **DSL at UCSB**Overview

• Increase throughput by reducing consensus from all nodes to smaller set

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Mine once, publish txns many times

**BitcoinNG** 

• Increase throughput by reducing consensus from all nodes to smaller set

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**BitcoinNG** 

Form a committee to vouch for new block

**ByzCoin** 

Increase throughput by reducing consensus from all nodes to smaller set

Mine once, publish txns many times

Form a committee to vouch for new block

ByzCoin

Shard txns across different committees

Elastico

BitcoinNG (Next Generation)

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Observation: In Bitcoin, blocks provide two purpose:

consensus and txn verification

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#### **Keyblocks**:

Used for Leader Election and created using Proof-of-work

# BitcoinNG (Next Generation)

- Key-block miner → leader till next key-block is mined
- Leader publishes micro-blocks while in tenure

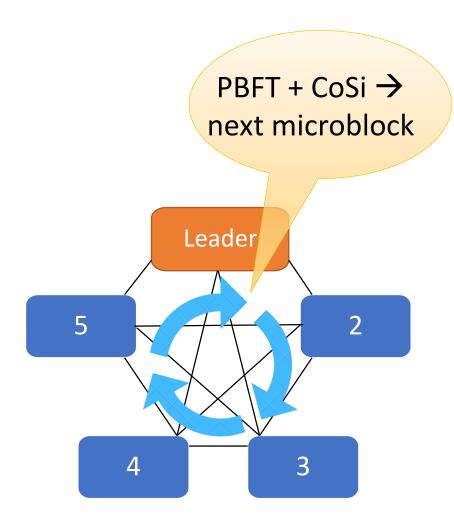
Allowing one miner to be a leader, even for a brief interval, presents many concerns!!

Eyal, Ittay, et al. "Bitcoin-NG: A Scalable Blockchain Protocol." NSDI. 2016.

ByzCoin

# ByzCoin

- Uses key-blocks and micro-blocks
- Key-block miner (PoW) in window becomes a trustee
- Micro-block decided by trustees
- Trustees use PBFT to reach consensus on next micro-block
- Each block is signed using Collective Signing approach



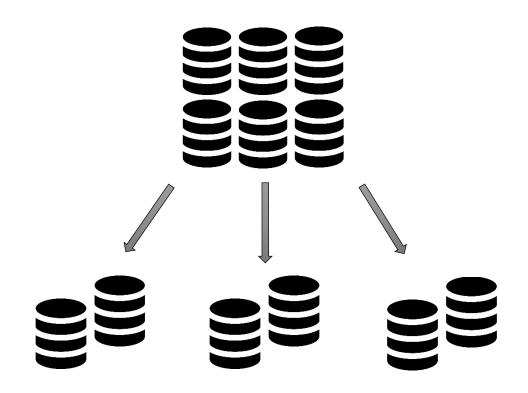
Kogias, Eleftherios Kokoris, et al. "Enhancing bitcoin security and performance with strong consistency via collective signing." 25th USENIX Security Symposium (USENIX Security 16). 2016.

# Elastico

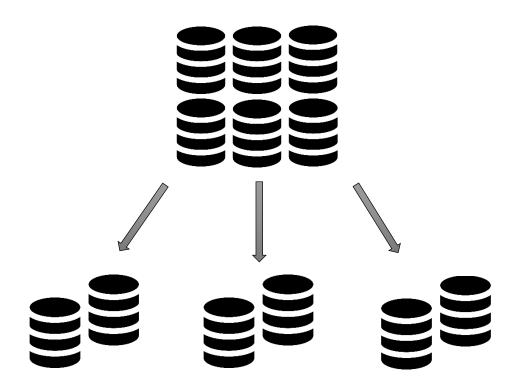
 Key idea: split all servers into smaller sized groups, committees

# Elastico

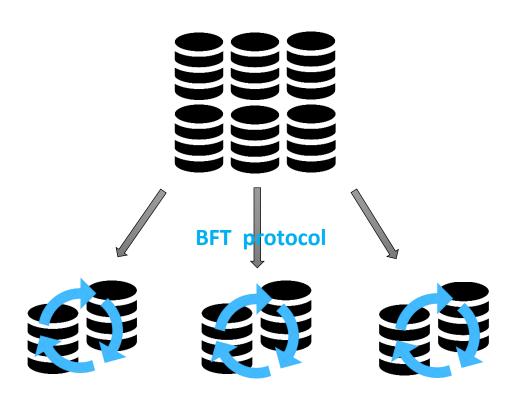
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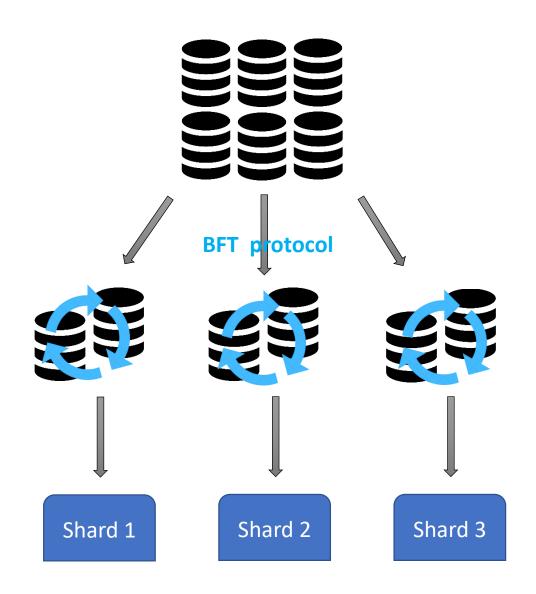
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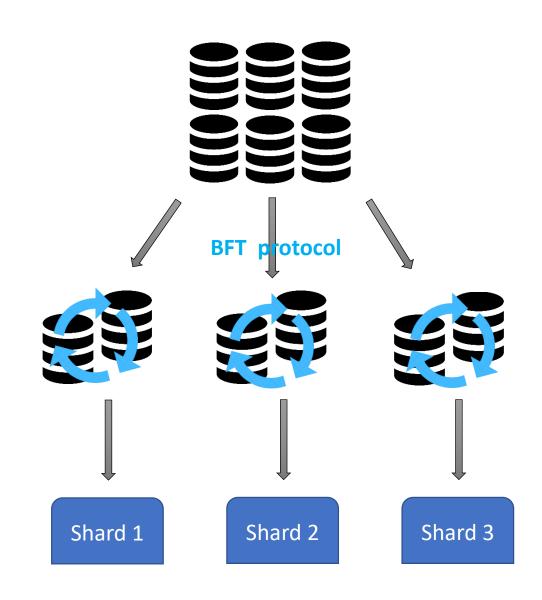
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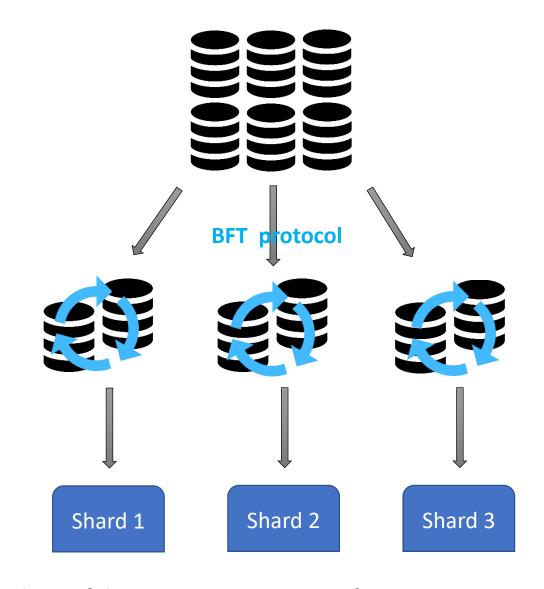


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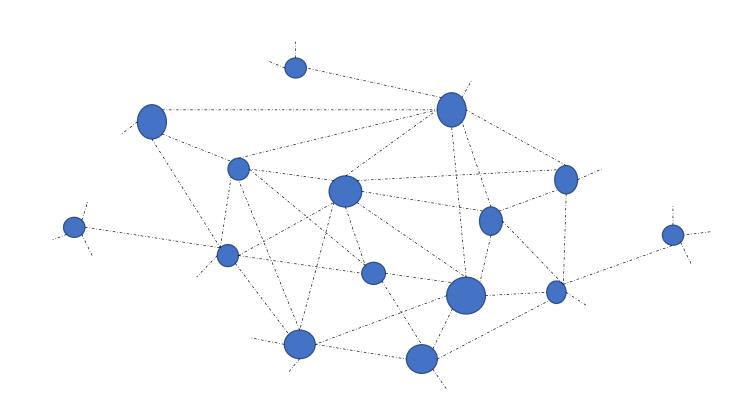


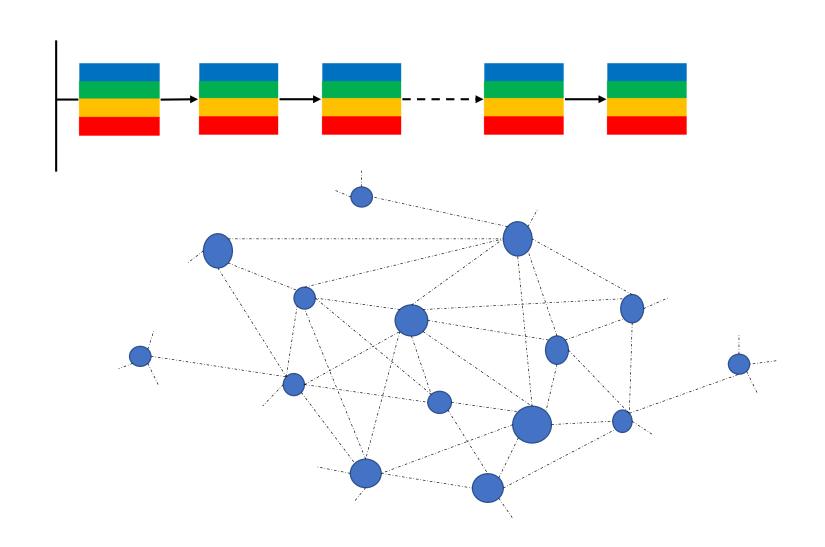
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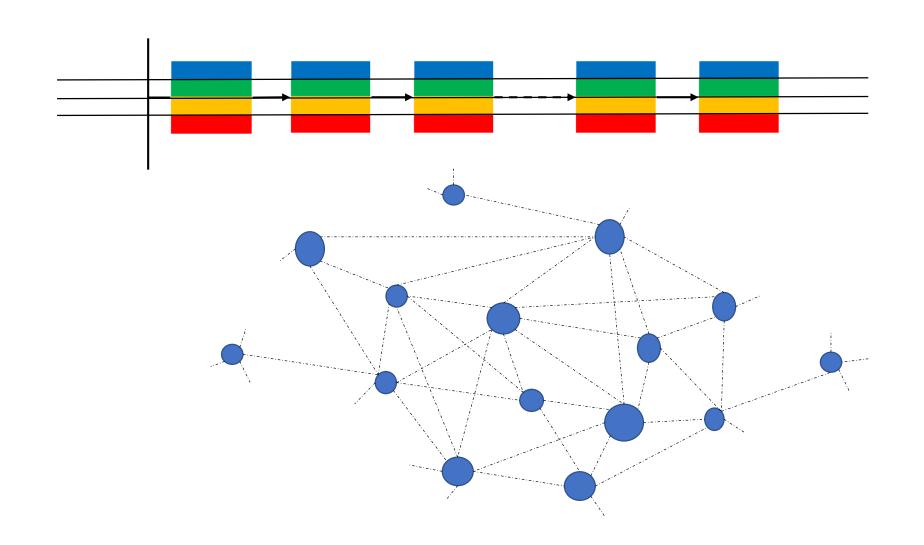


Luu, Loi, et al. "A secure sharding protocol for open blockchains." *Proceedings of the 2016 ACM SIGSAC Conference on Computer and Communications Security*. ACM, 2016.

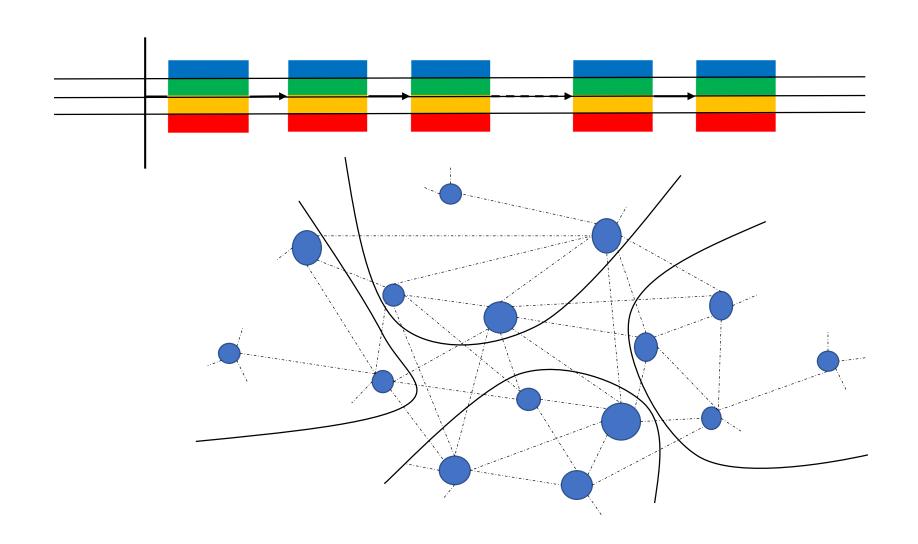




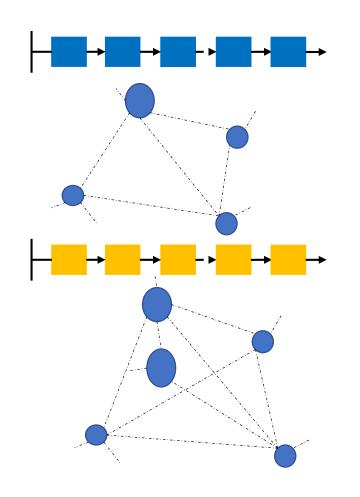


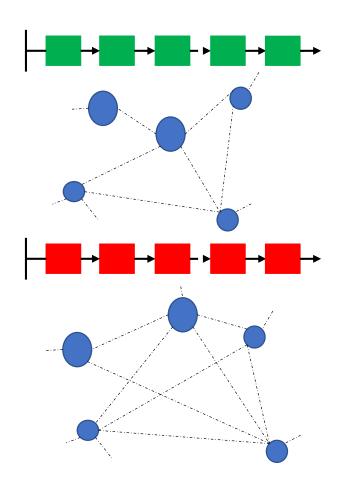






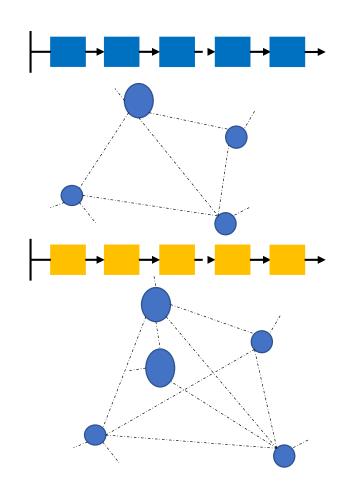


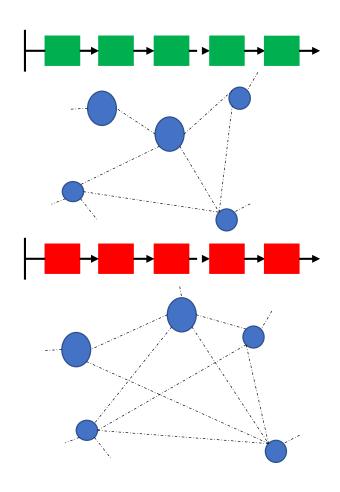






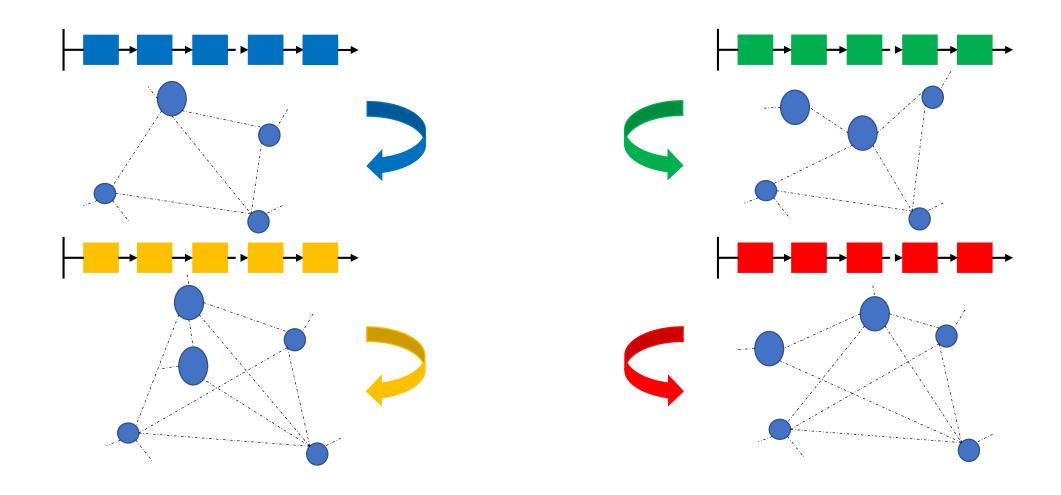
# Classes of Transactions



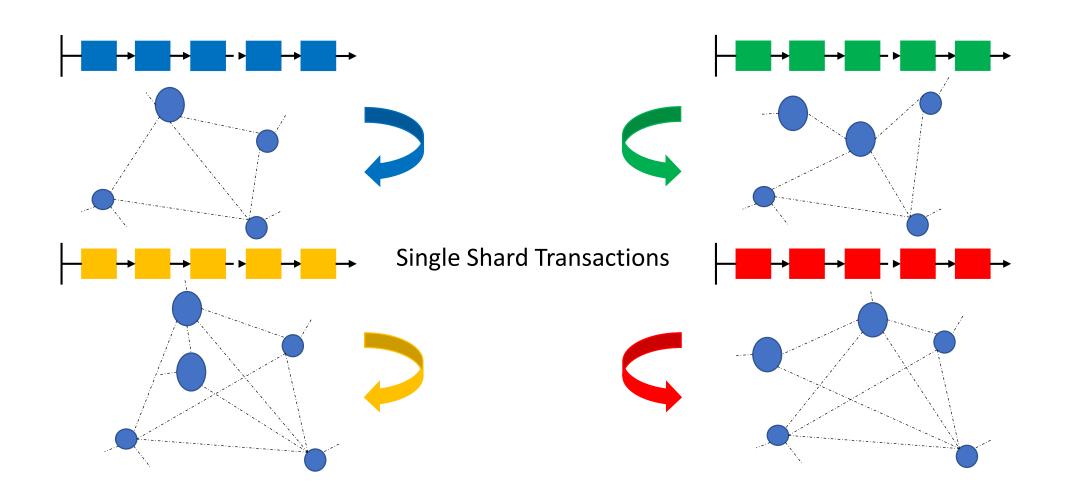




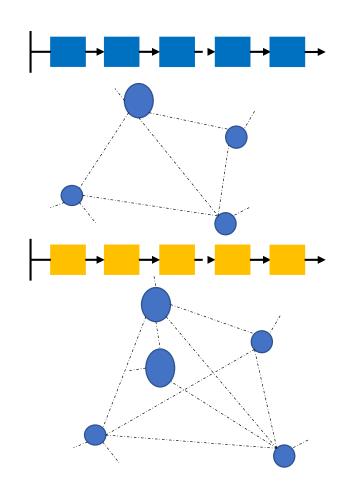
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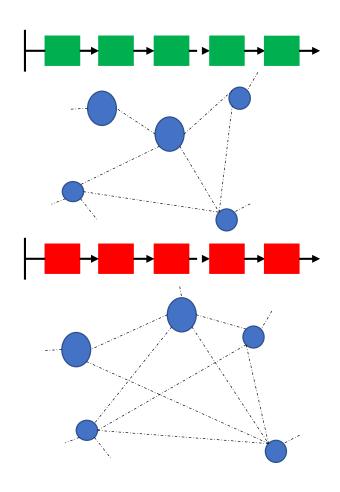




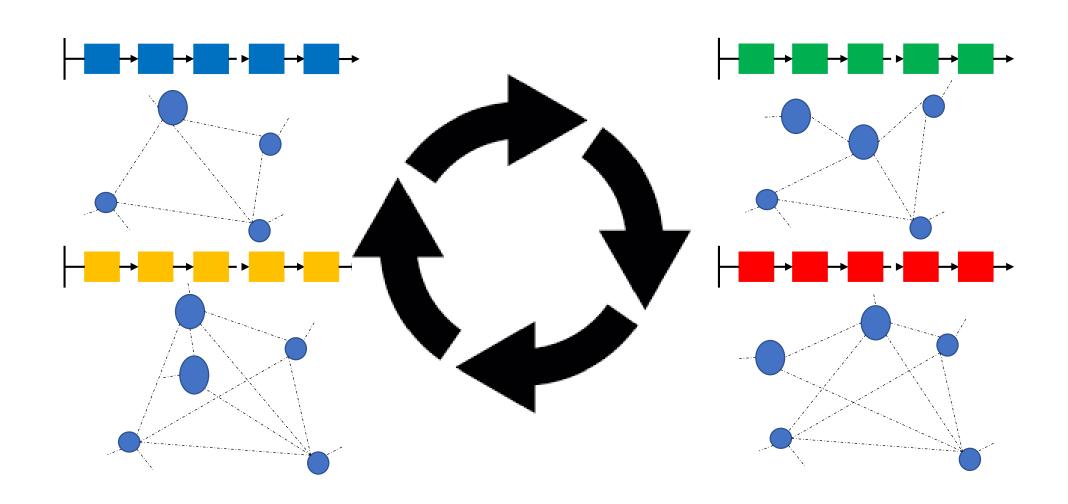




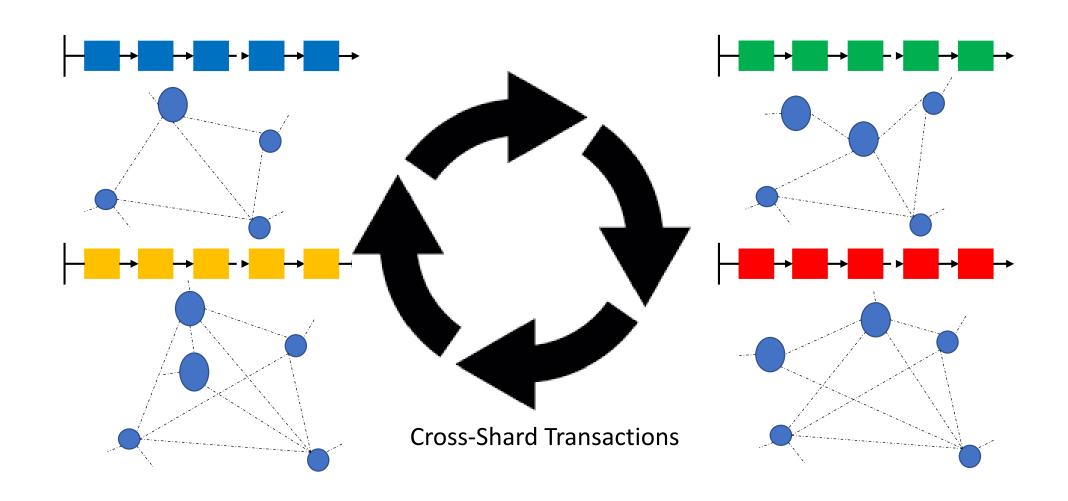






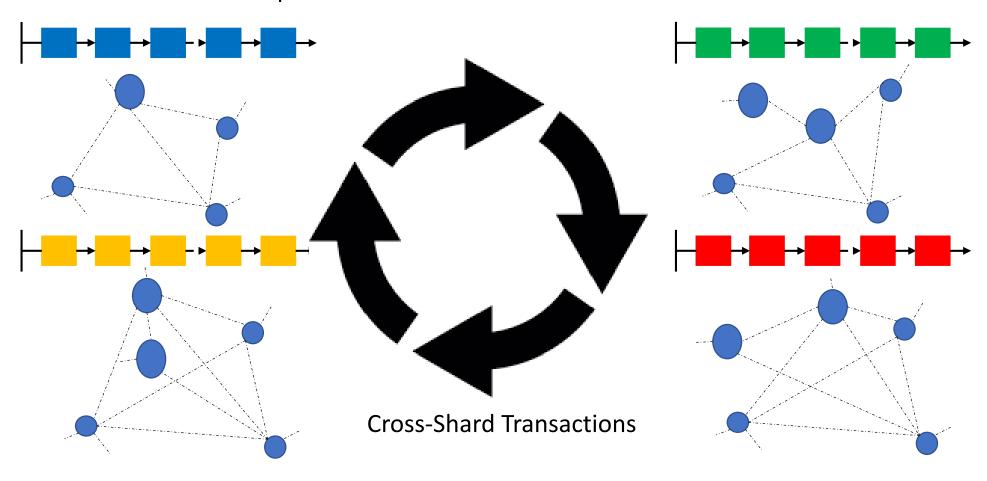








#### Requires Atomic Cross-Shard Commitment Protocol



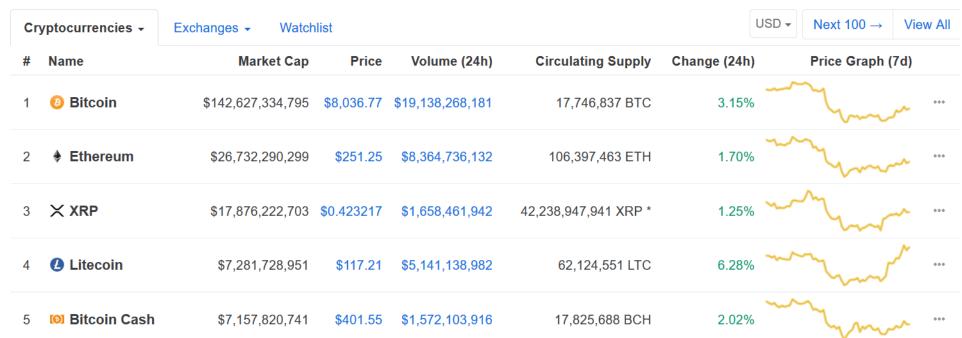


Cryptocurrencies: 2225 • Markets: 18851 • Market Cap: \$257,486,187,861 • 24h Vol: \$66,548,083,112 • BTC Dominance: 55.4%

CoinMarketCap

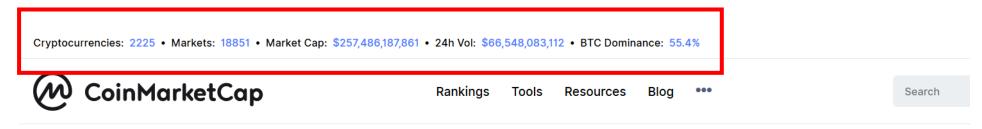
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#### **Top 100 Cryptocurrencies by Market Capitalization**



Source: coinmarketcap.com on June 7th at 5:00pm PST





#### **Top 100 Cryptocurrencies by Market Capitalization**

Cryptocurrencies •	Exchanges - Watch	nlist			USD •	Next 100 → View A
# Name	Market Cap	Price	Volume (24h)	Circulating Supply	Change (24h)	Price Graph (7d)
1 <sup>(3)</sup> Bitcoin	\$142,627,334,795	\$8,036.77	\$19,138,268,181	17,746,837 BTC	3.15%	The "
2 <b>♦ Ethereum</b>	\$26,732,290,299	\$251.25	\$8,364,736,132	106,397,463 ETH	1.70%	my "
3 <b>XRP</b>	\$17,876,222,703	\$0.423217	\$1,658,461,942	42,238,947,941 XRP *	1.25%	-M
4 (2) Litecoin	\$7,281,728,951	\$117.21	\$5,141,138,982	62,124,551 LTC	6.28%	my "
5	sh \$7,157,820,741	\$401.55	\$1,572,103,916	17,825,688 BCH	2.02%	my mu "

Source: coinmarketcap.com on June 7th at 5:00pm PST



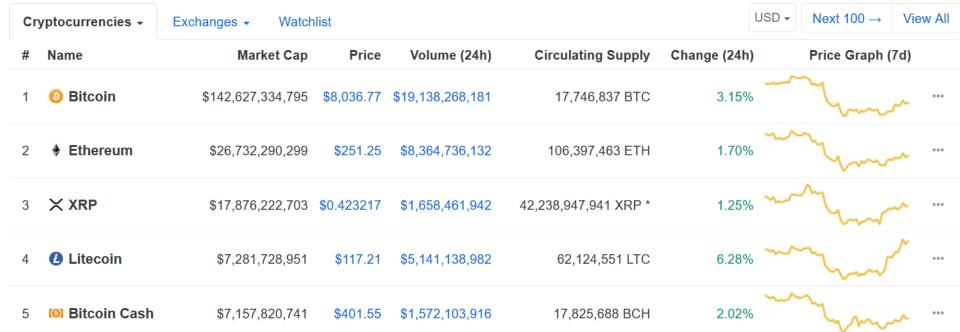
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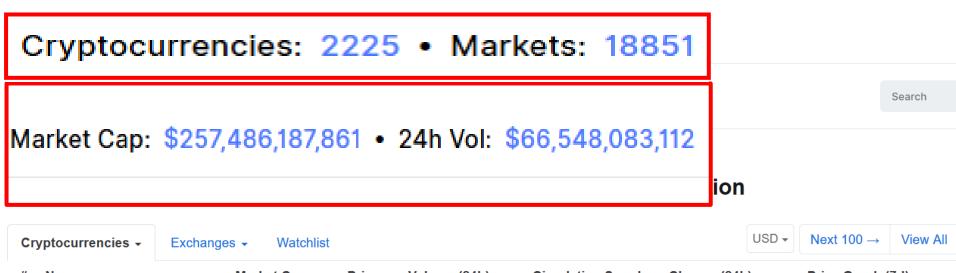
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1 <sup>(3)</sup> Bitcoin	\$142,627,334,795	\$8,036.77	\$19,138,268,181	17,746,837 BTC	3.15%	
2 <b>♦ Ethereum</b>	\$26,732,290,299	\$251.25	\$8,364,736,132	106,397,463 ETH	1.70%	my "
3 <b>XRP</b>	\$17,876,222,703	\$0.423217	\$1,658,461,942	42,238,947,941 XRP *	1.25%	· · · · · · · · · · · · · · · · · · ·
4 Litecoin	\$7,281,728,951	\$117.21	\$5,141,138,982	62,124,551 LTC	6.28%	
5 O Bitcoin Cash	\$7,157,820,741	\$401.55	\$1,572,103,916	17,825,688 BCH	2.02%	my "

Source: coinmarketcap.com on June 7<sup>th</sup> at 5:00pm PST

# The Landscape

• Thousands of Blockchains

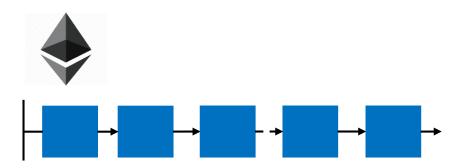
- Thousands of Blockchains
- Tens of thousands of markets

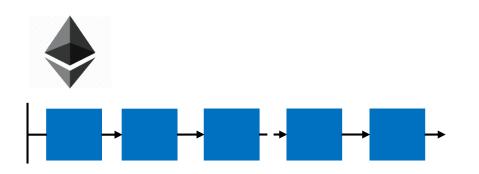
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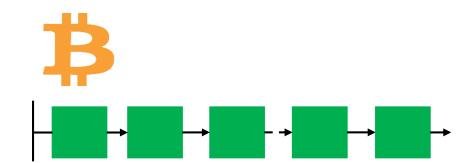
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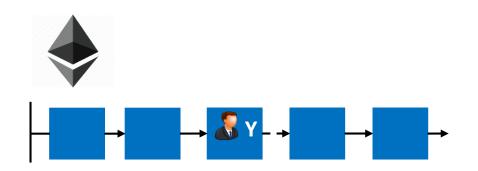
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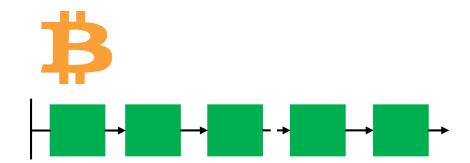
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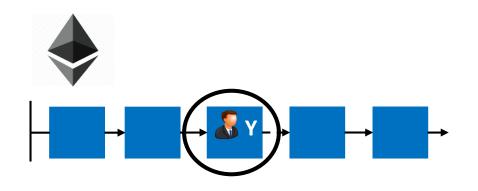


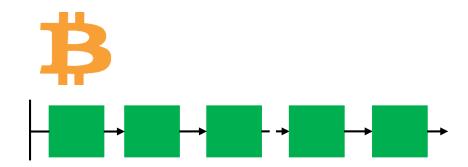


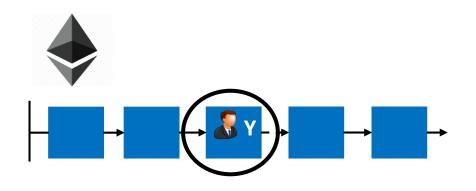


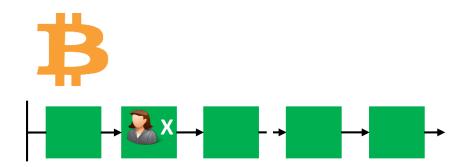


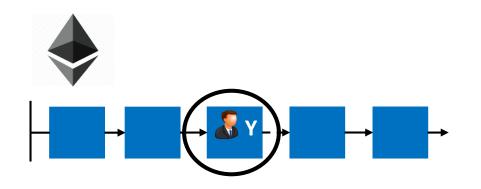


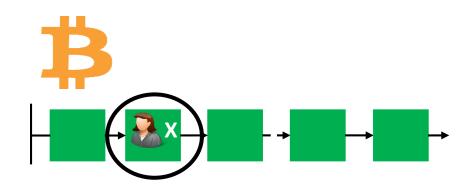


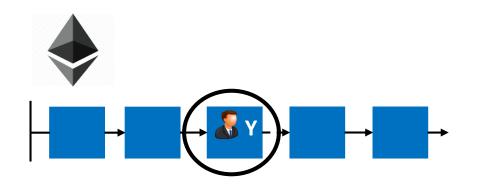


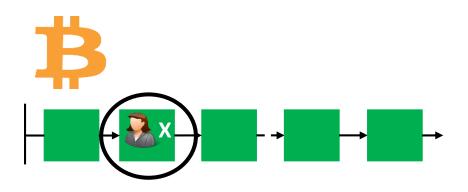






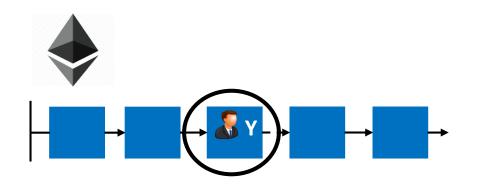


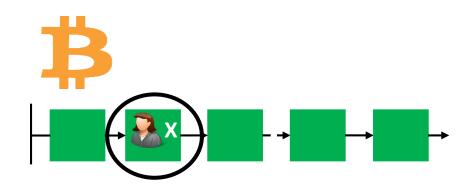




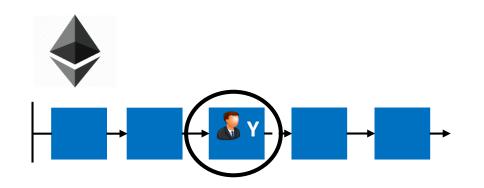


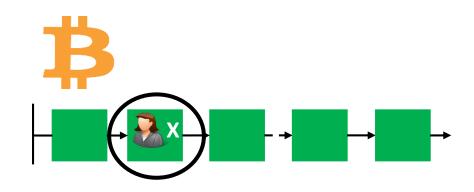


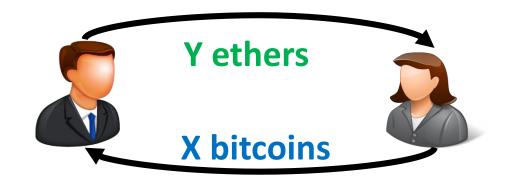




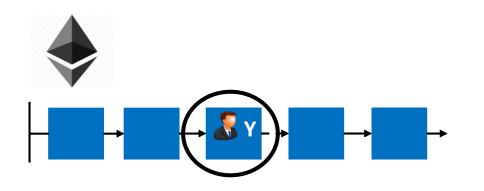


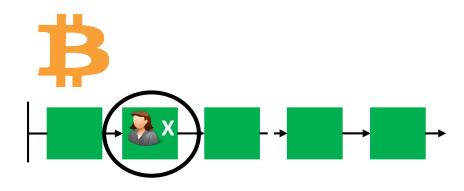




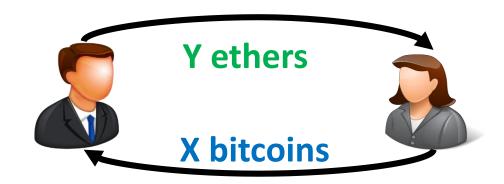


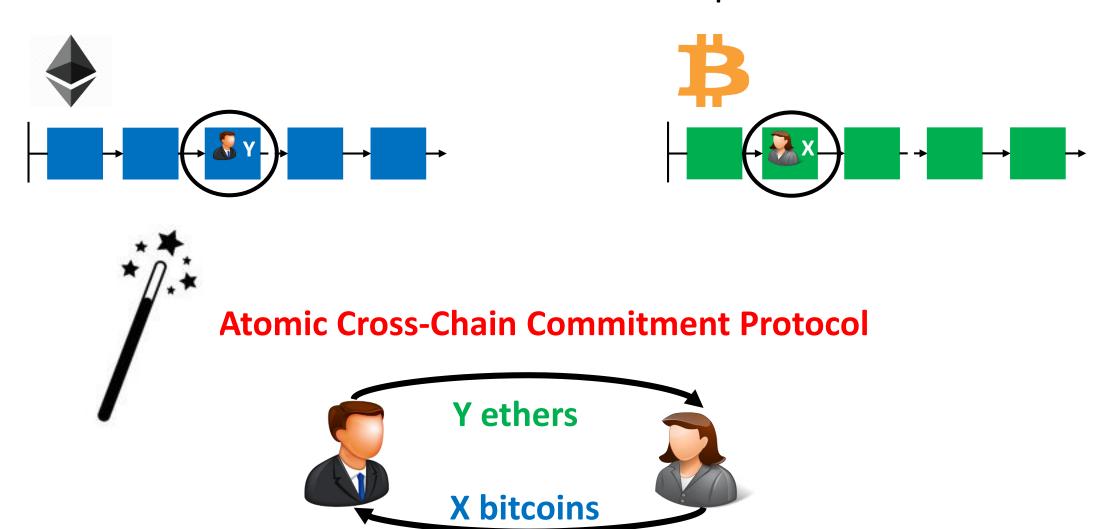
# Cross-ChainTransaction Example

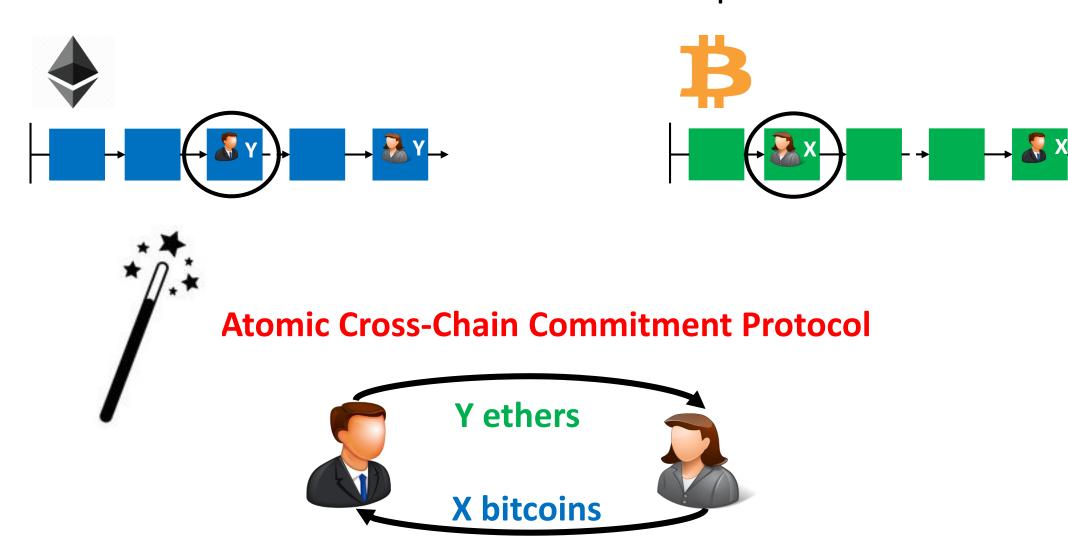


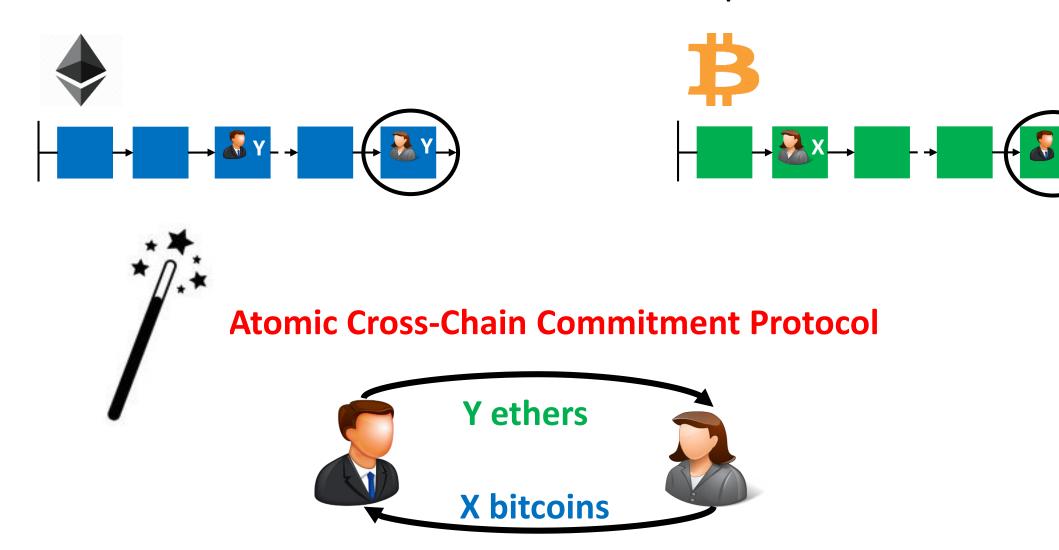


#### **Atomic Cross-Chain Commitment Protocol**









### **Smart Contracts**

• Like classes in Object Oriented Programming Languages

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- Can be used to implement generic transaction logic:
  - Conditionally lock assets in the blockchain
  - Transfer asset ownership on some condition



```
class AtomicSwap {
    sender: s // Alice
    recipient: r // Bob
    asset: a // X bitcoins
    secretHash: h
    constructor() {
    redeem (secret srt) {
            if(hash(srt) == h)
                     transfer a to r
```



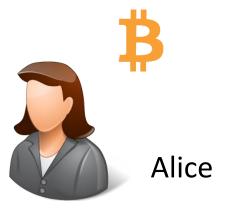
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Alice wants to trade Bitcoin for Ethereum with Bob

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Calculate its hash h = H(s)





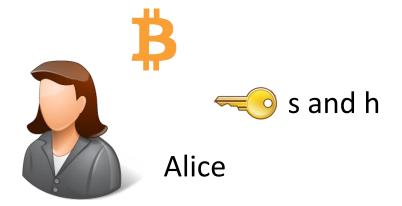
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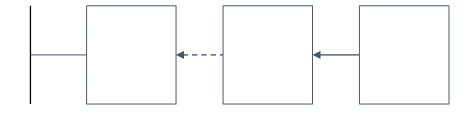
s and h

Alice

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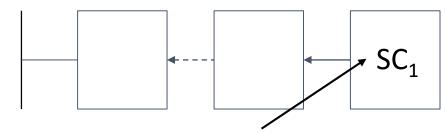




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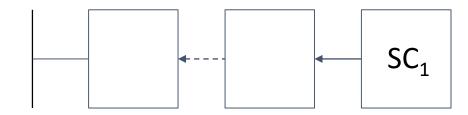




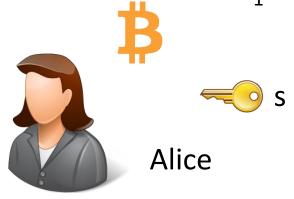
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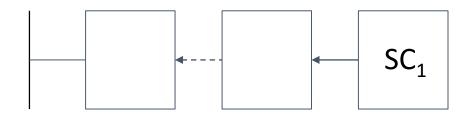


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Bitcoin blockchain

Alice's X bitcoins are locked in



 $SC_2$  Move Y Ethereum to Alice if Alice provides secret s | h = H(s)

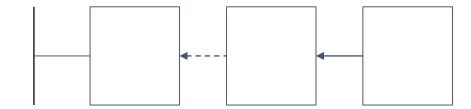




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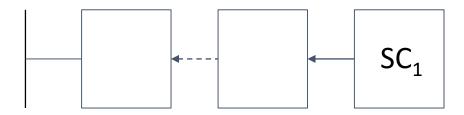
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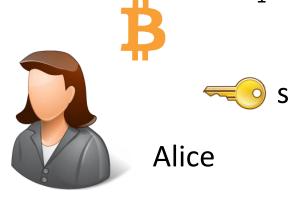


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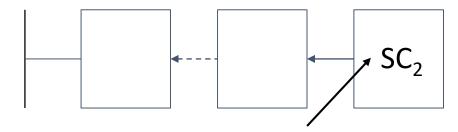




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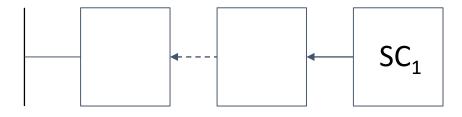
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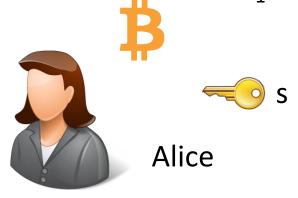


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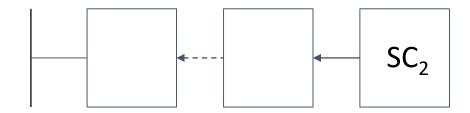




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• Now, for Alice to execute SC<sub>2</sub> and redeem Y Ethereum, she reveals s

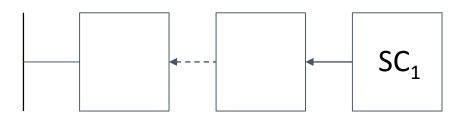
Ethereum blockchain



Bob's Y Ethereum are locked in smart contract SC<sub>2</sub>



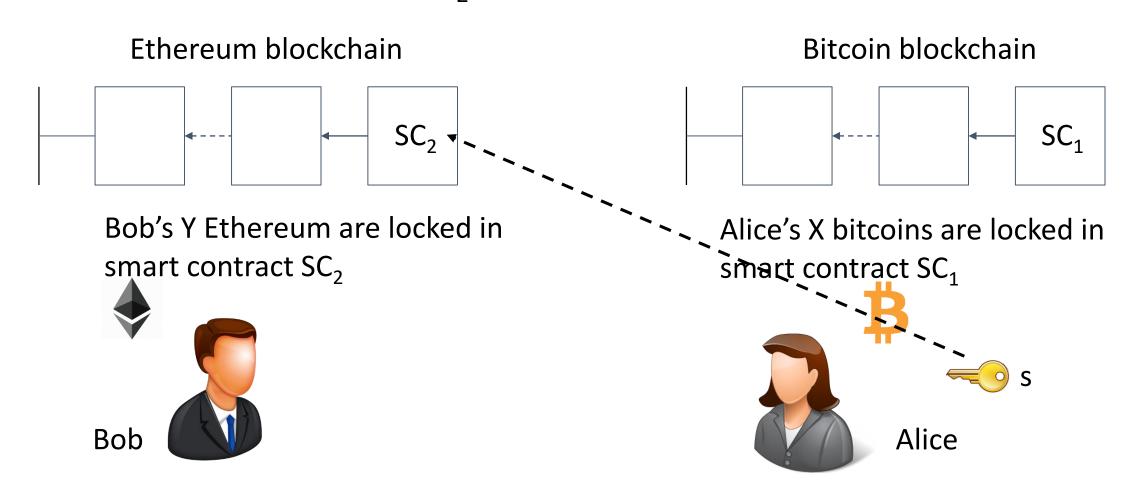
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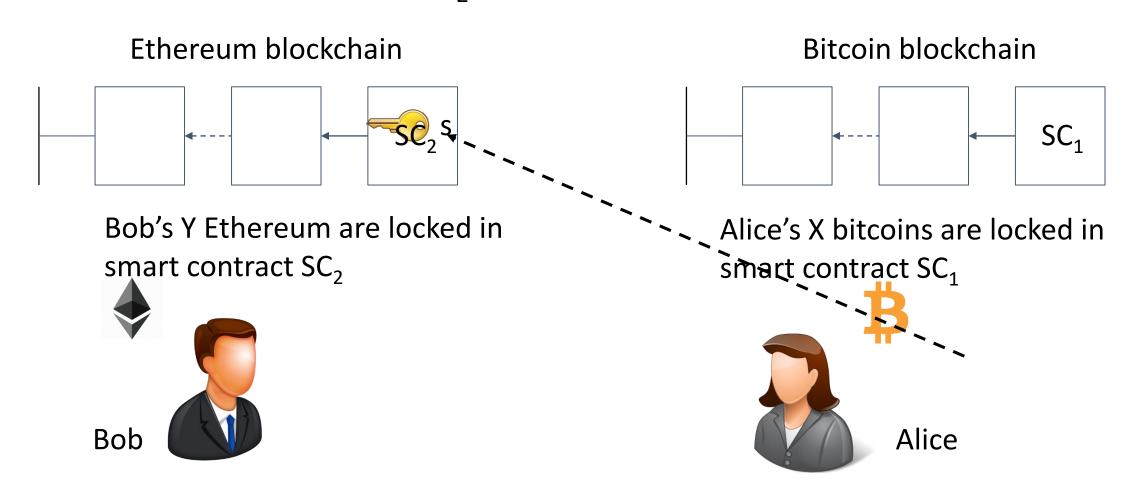
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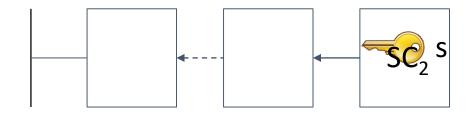
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## Atomic Swap[Nolan'13, Herlihy'18]

• Revealing s, executes SC<sub>2</sub>. Now s is public in Ethereum's blockchain

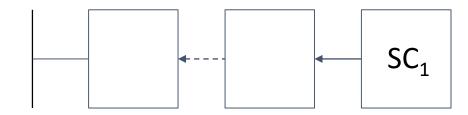
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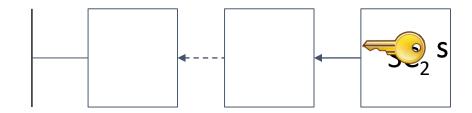




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• Now, Bob uses s to execute SC<sub>1</sub> and redeem his Bitcoins

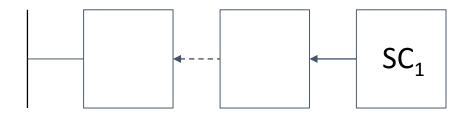
#### Ethereum blockchain



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#### Bitcoin blockchain

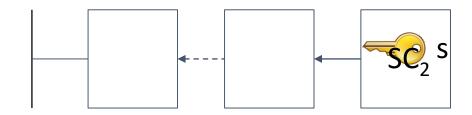




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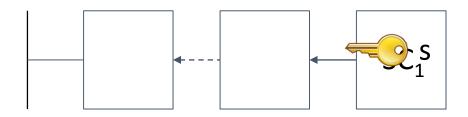
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## Atomic Swap Example: What can go wrong?

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- Now, Alice's Bitcoins are locked for good
  - A conforming party (Alice) ends up worse off because Bob doesn't follow the protocol
- Prevention
  - Use timelocks to expire a contract
  - Specify that an expired contract is refunded to the creator of this contract

Atomic Swap[Nolan'13, Herlihy'18]: Timelocks





## Atomic Swap[Nolan'13, Herlihy'18]: Timelocks



Refund SC<sub>1</sub> to Alice if Bob does not execute SC<sub>1</sub> before 48 hours

 $SC_1$ : Move X bitcoins to Bob if Bob provides secret s | h = H(s)



# Atomic Swap[Nolan'13, Herlihy'18]: Timelocks

Refund SC<sub>2</sub> to Bob if Alice does not execute SC<sub>2</sub> before **24** hours

 $SC_2$ : Move Y Ethereum to Alice if Alice provides secret s | h = H(s)

Bob

Refund SC<sub>1</sub> to Alice if Bob does not execute SC<sub>1</sub> before **48** hours

 $SC_1$ : Move X bitcoins to Bob if Bob provides secret s | h = H(s)



## Atomic Swap[Nolan'13, Herlihy'18]: Timelocks

Refund SC<sub>2</sub> to Bob if Alice does not execute SC<sub>2</sub> before 4 hours

 $SC_2$ : Move Y Ethereum to Alice if Alice provides secret s | h = H(s)

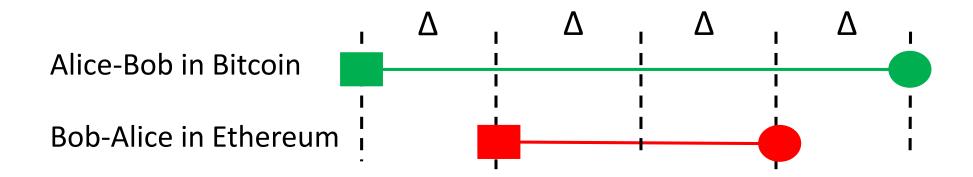
Bob

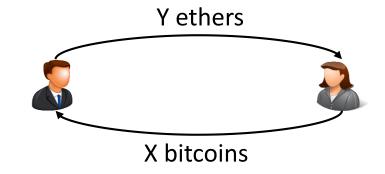
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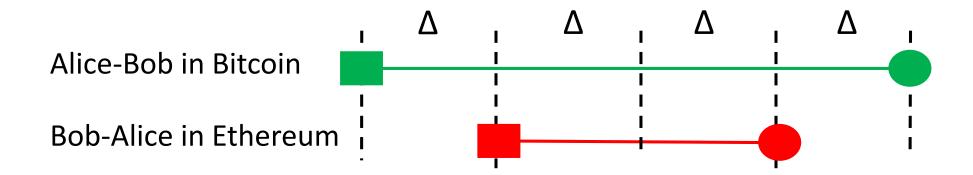


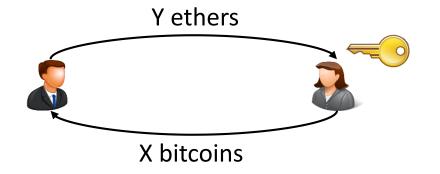
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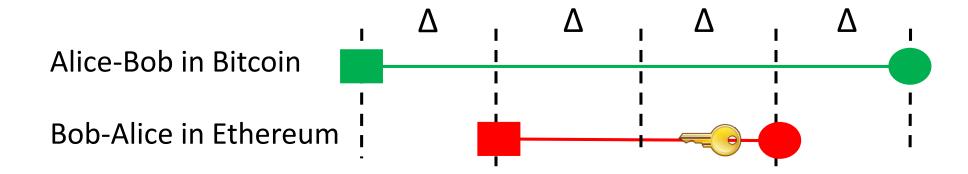


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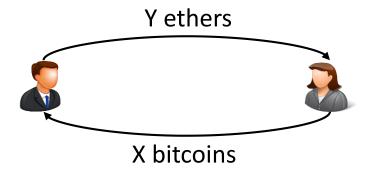




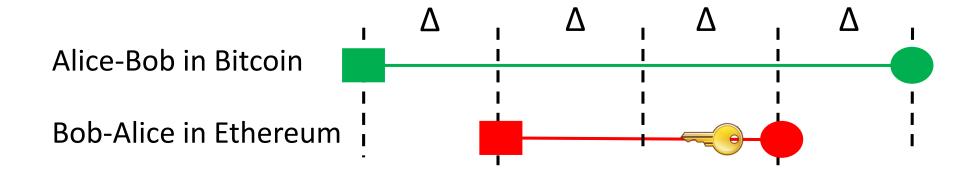
### Atomic Swap Example [Nolan'13, Herlihy'18]



Alice reveals the secret to Bob's contract and claims the Y ether

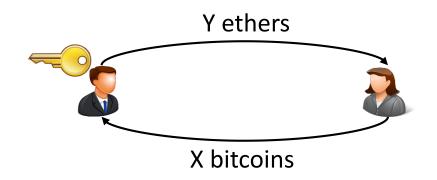


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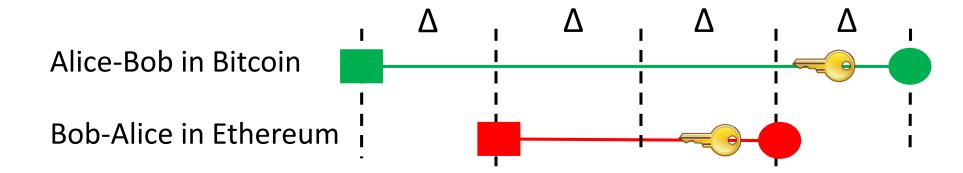


Alice reveals the secret to Bob's contract and claims the Y ether

Supposedly, Bob takes the secret, reveals it to Alice's contract and claims the X bitcoins

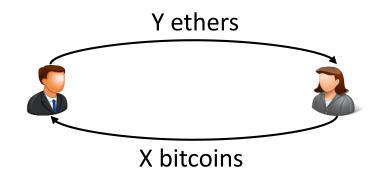


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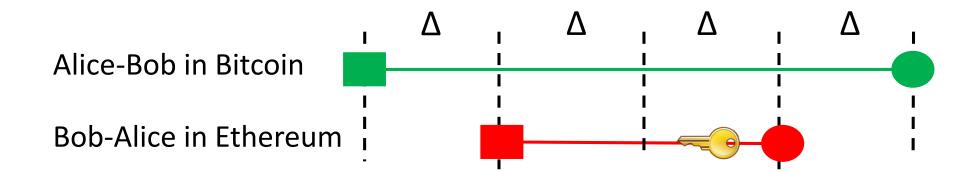


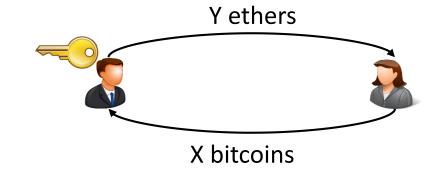
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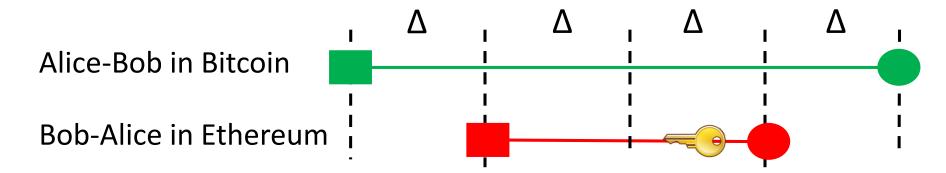


## What can go wrong?

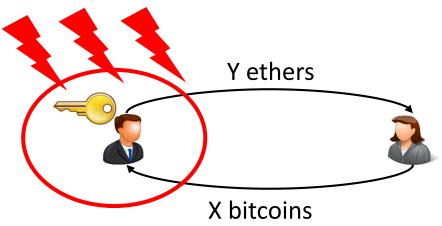




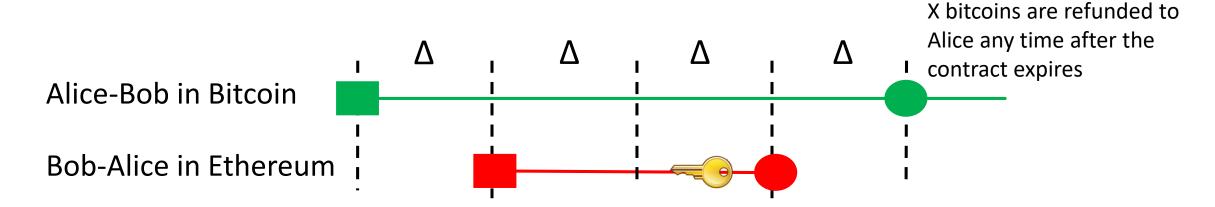
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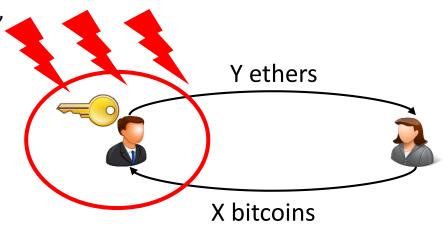
If Bob fails or suffers a network denial of service attack for a Δ, Alice's contract will expire and Bob will lose his X bitcoins



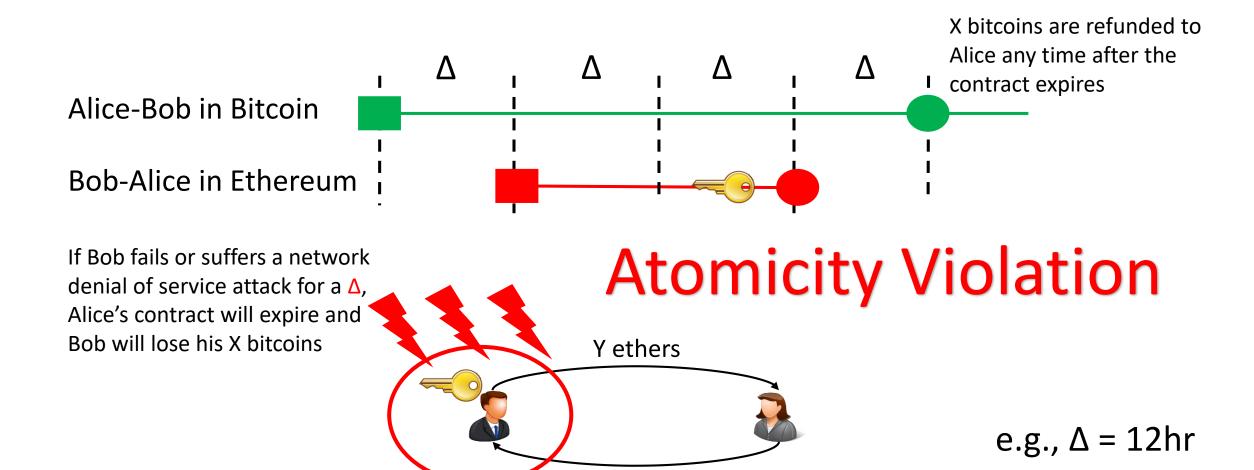
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  - A transaction cannot commit unless a commit decision is reached
  - A transaction cannot abort unless an abort decision is reached



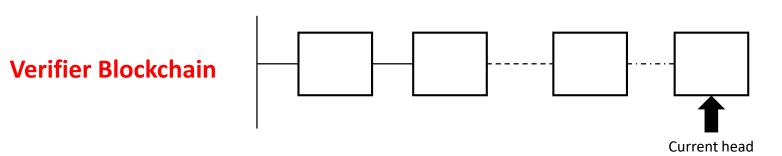
# Atomic Commitment Across Blockchains

Victor Zakhary, Divyakant Agrawal, Amr El Abbadi

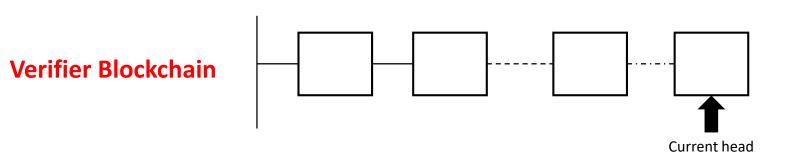


- How can miners of one blockchain:
  - Verify a transaction in another blockchain?

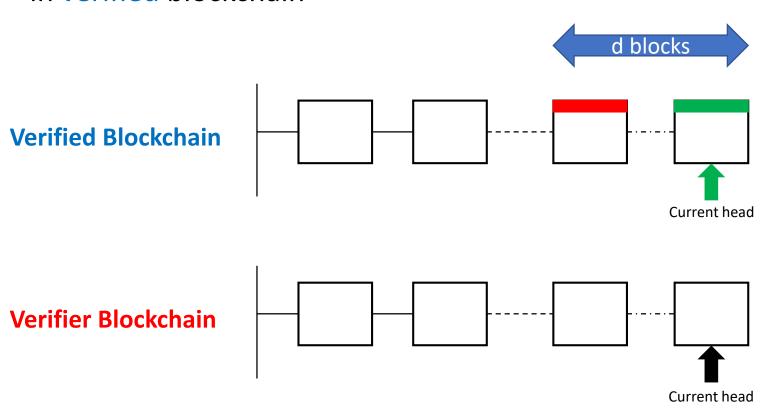
- How can miners of one blockchain:
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  - Without maintaining a copy of this other blockchain.



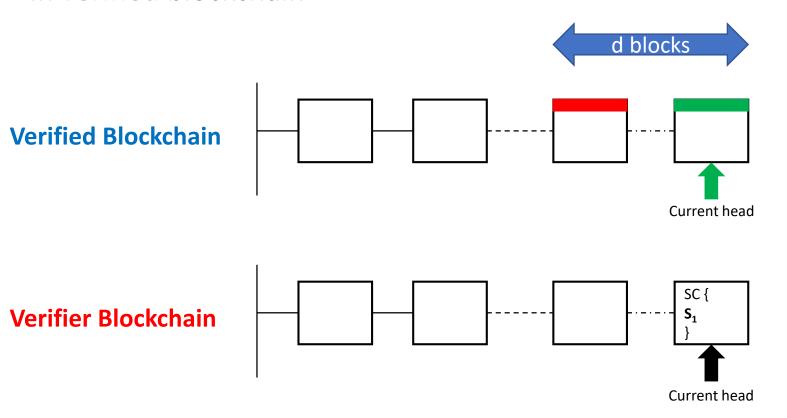
# Building block: Cross-Chain Verification



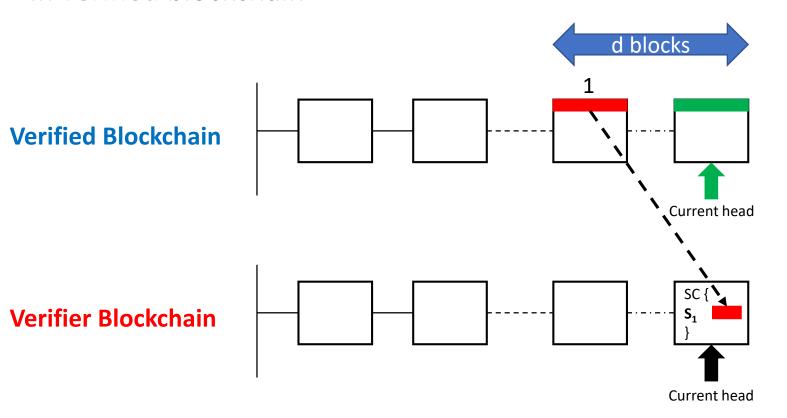




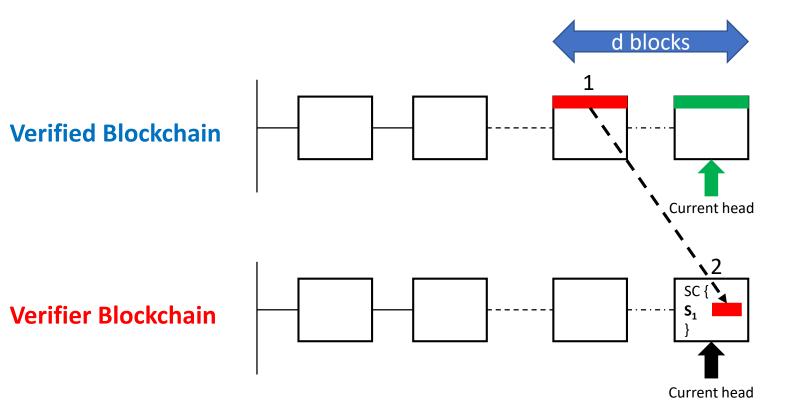




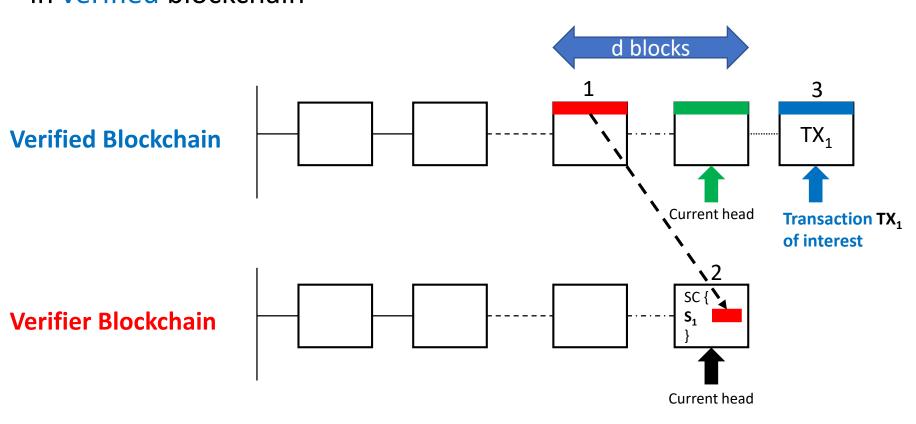




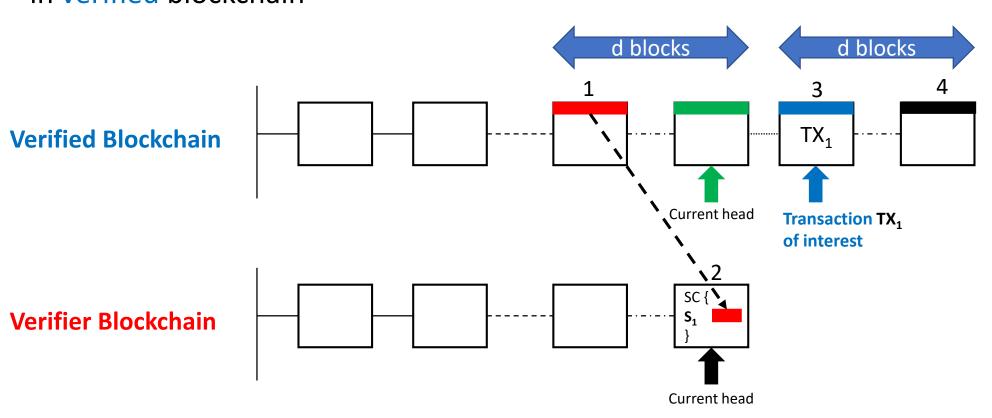




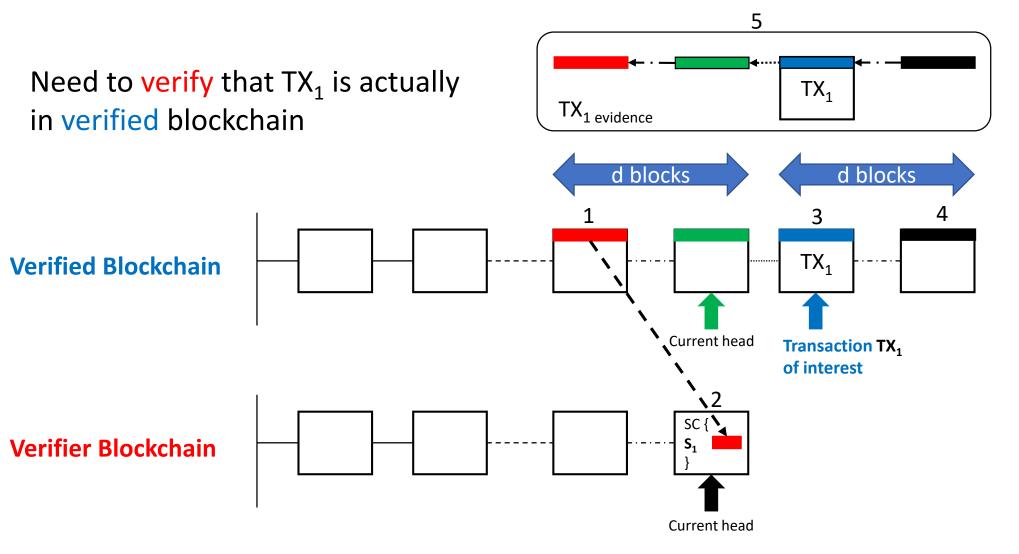




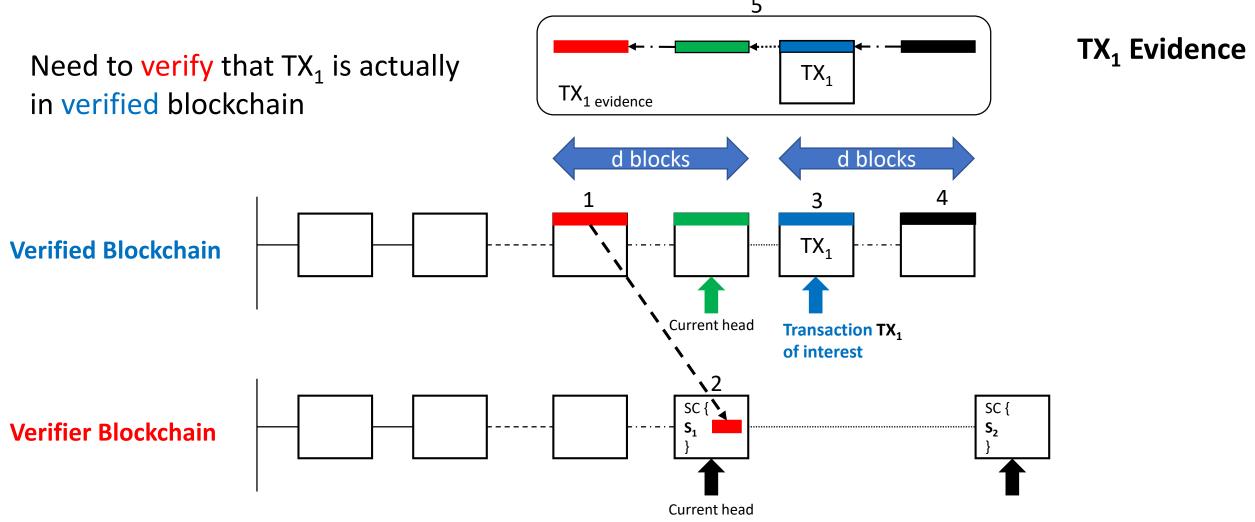


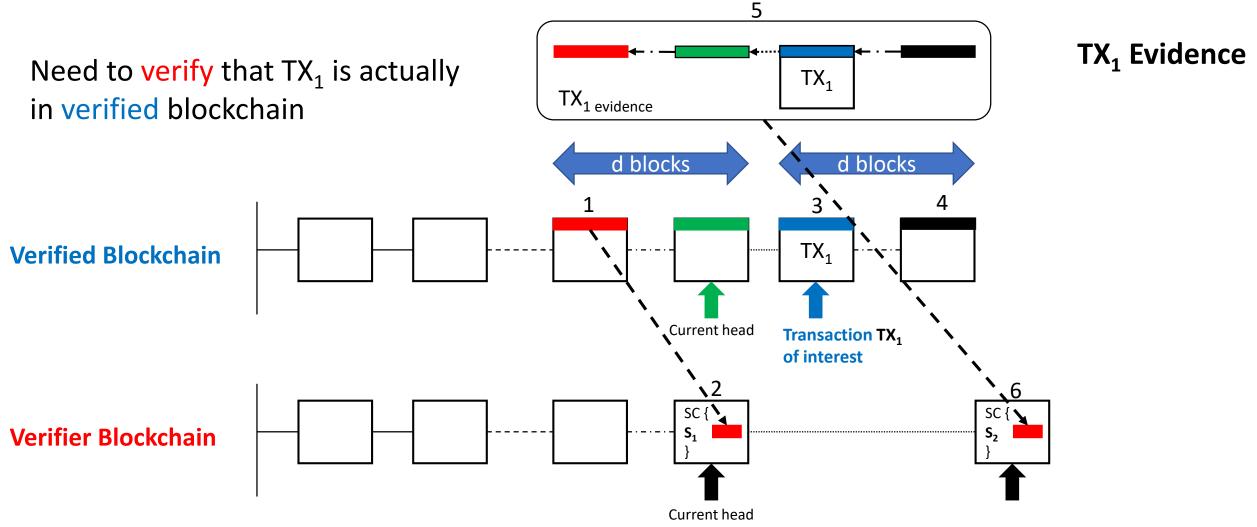


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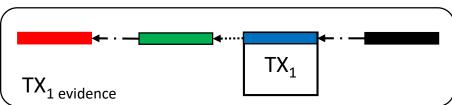


**TX<sub>1</sub> Evidence** 



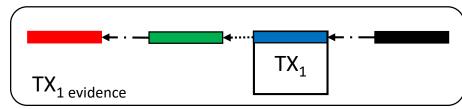






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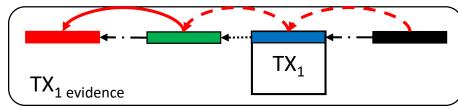
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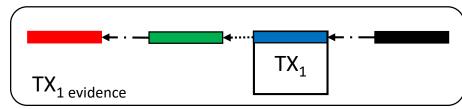
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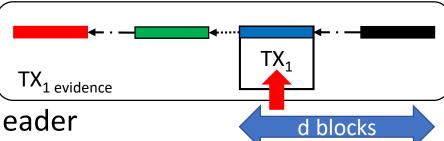
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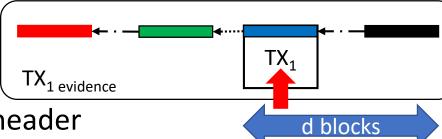
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  - TX<sub>1</sub> is buried under d blocks
- The cost of generating evidence:
  - Choose d to make this cost > the value transacted in TX<sub>1</sub>
  - If true, a malicious user has no incentive to create a fake evidence



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  - Miners of assets' blockchains verify the decision made in the witness network

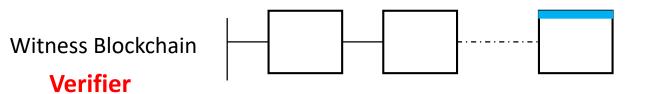
## Protocol Sketch

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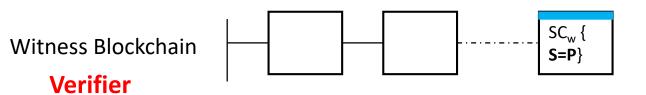
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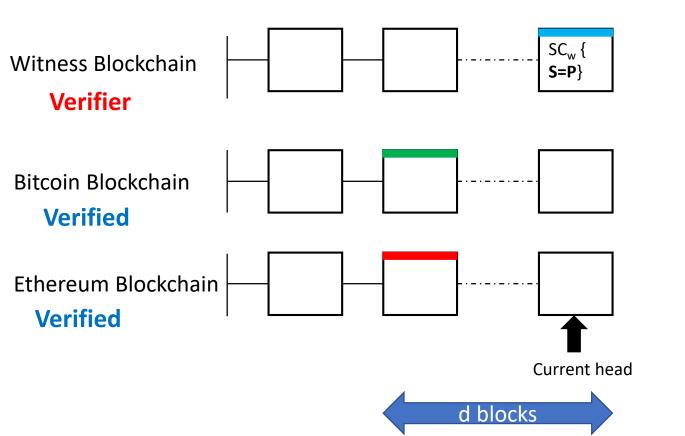
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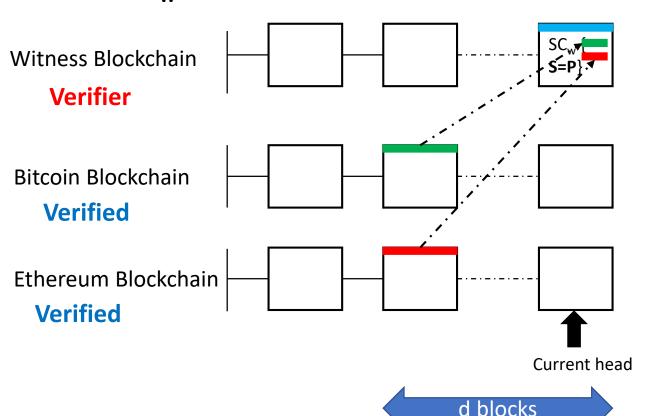
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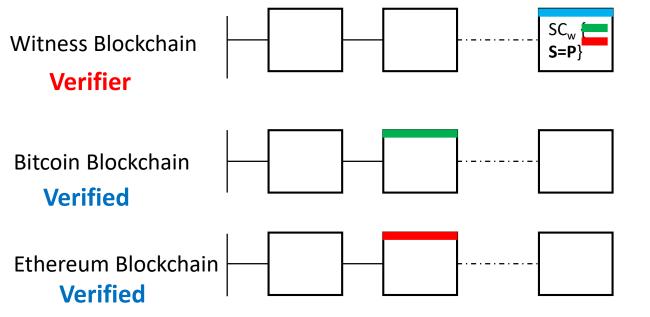
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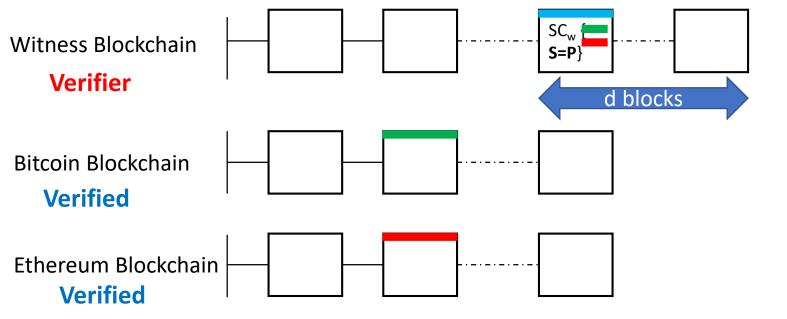


## Protocol Sketch

- Deploy a contract SC<sub>w</sub> in the witness network with state Published (P)
- SC<sub>w</sub> has a header of a block at depth d of all blockchains in the swap

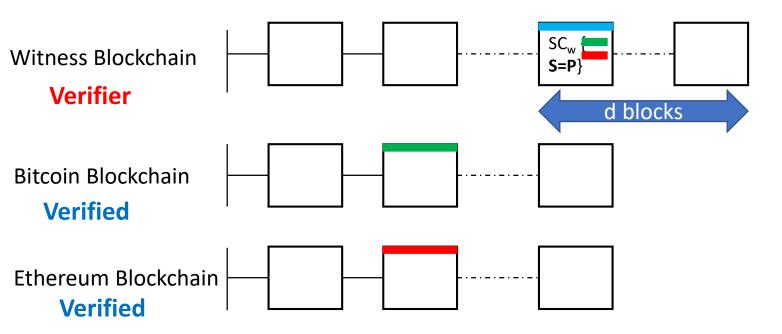






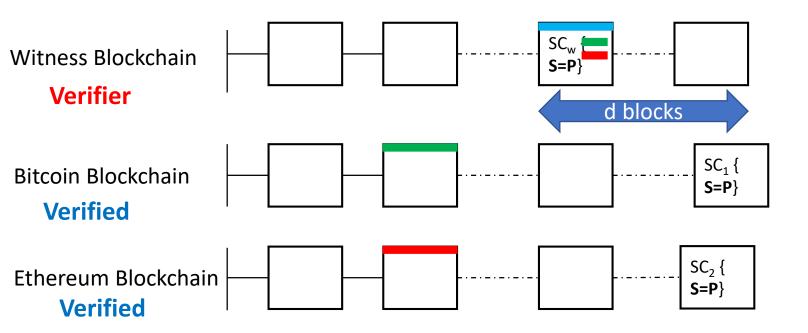


Participants deploy their contracts in the corresponding blockchains



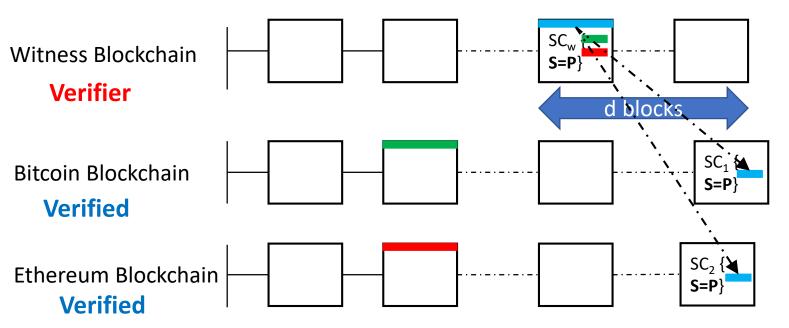


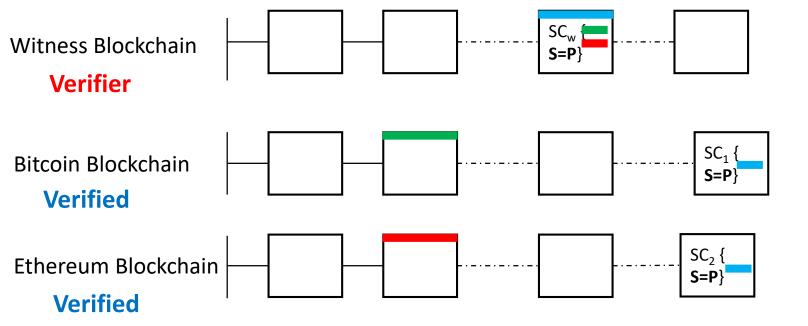
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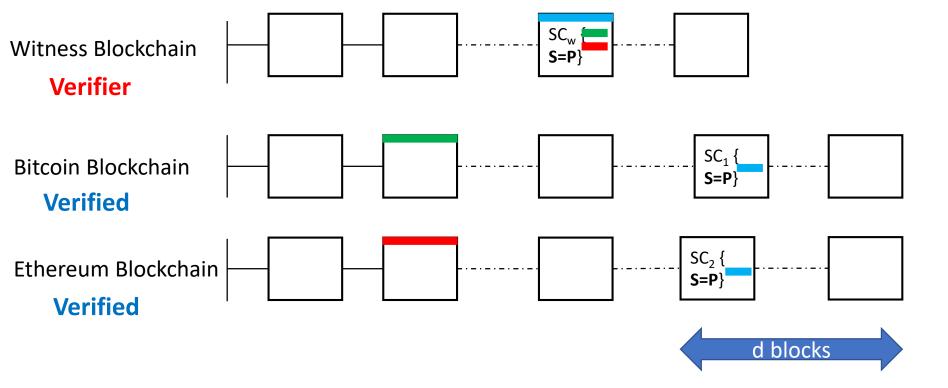


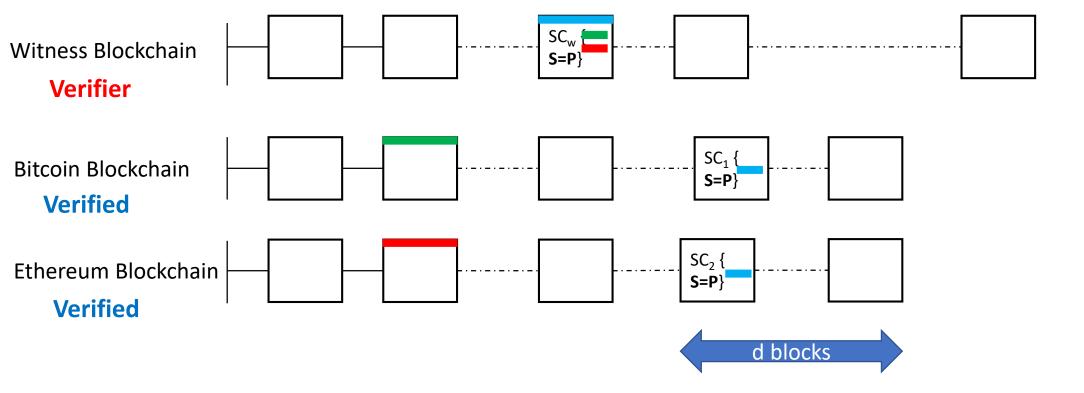


- Participants deploy their contracts in the corresponding blockchains
- Participants add the header of SC<sub>w</sub> to their contracts



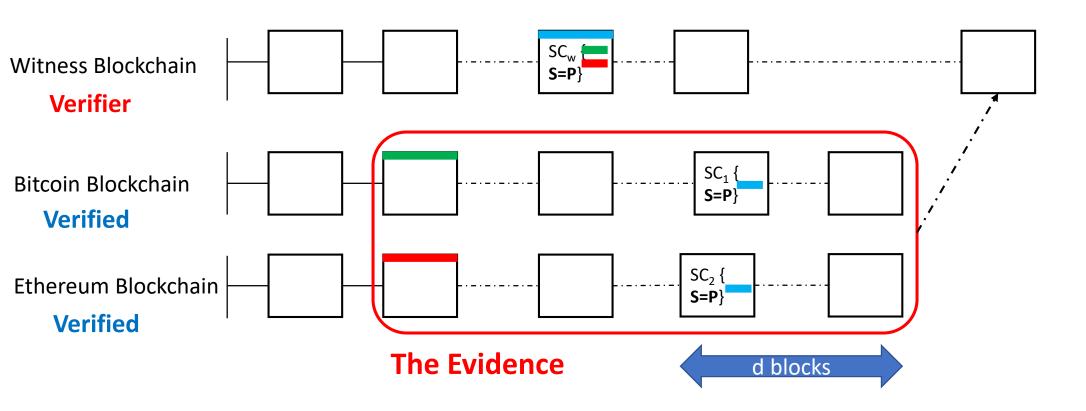






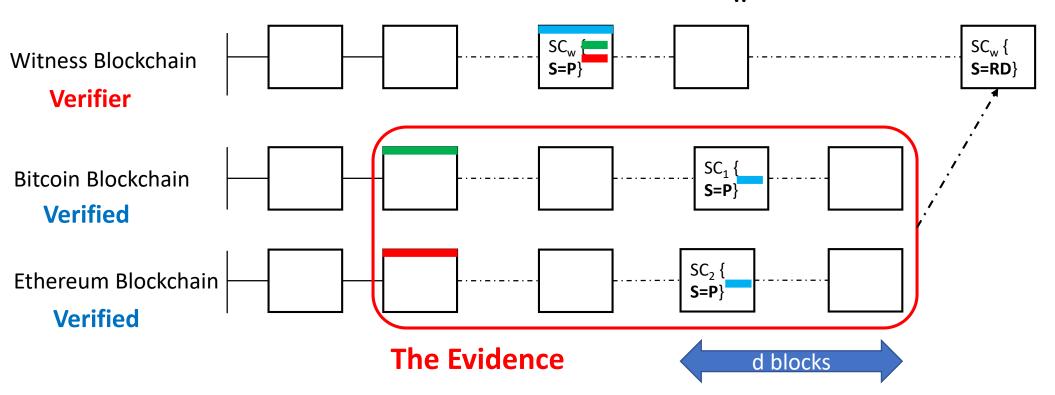


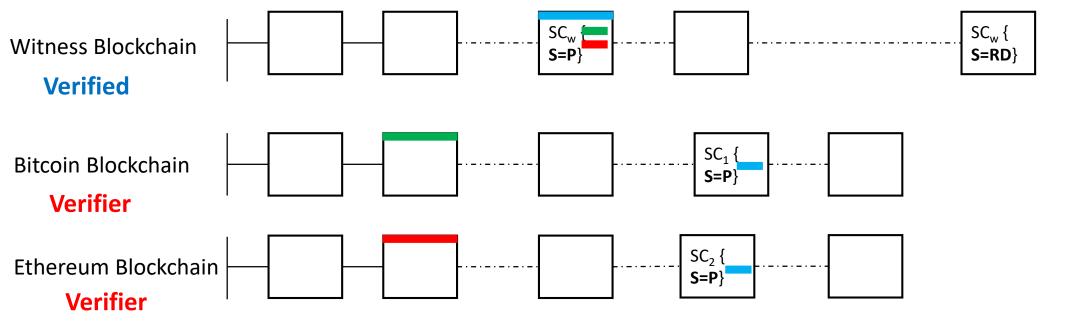
 Participants submit evidence of publishing the smart contracts in Assets Blockchains

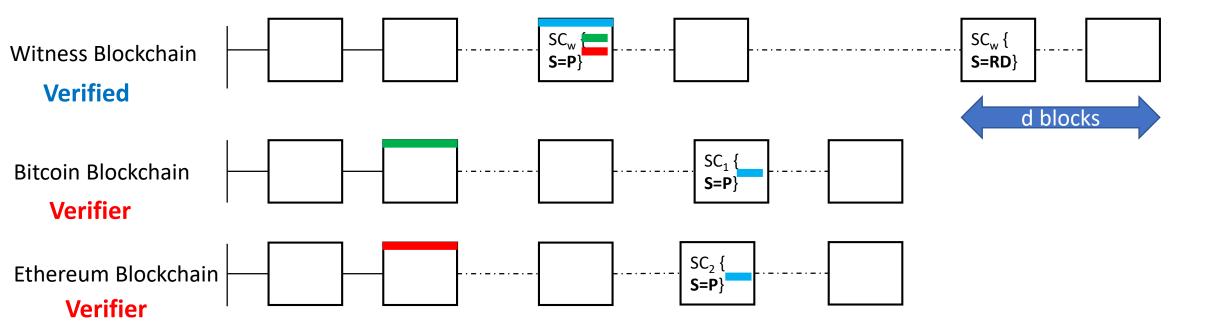


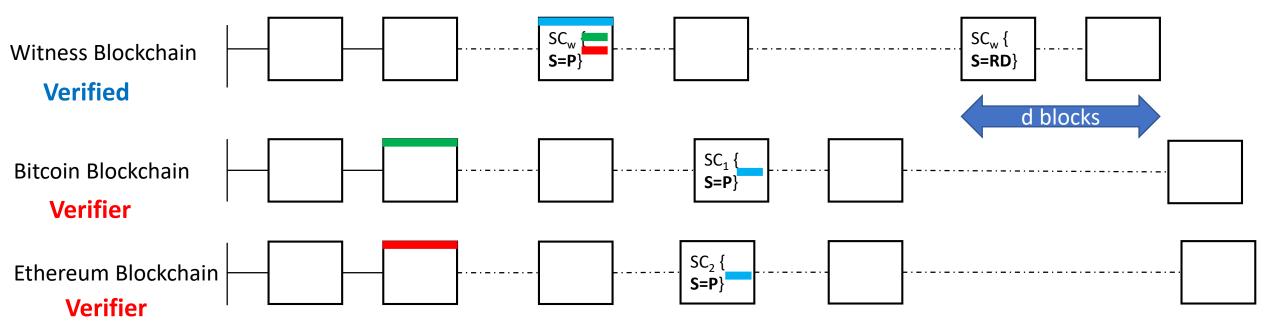


- Participants submit evidence of publishing the smart contracts in Assets Blockchains
- If all contracts are published and correct, SC<sub>w</sub>'s state is altered to redeem (RD)

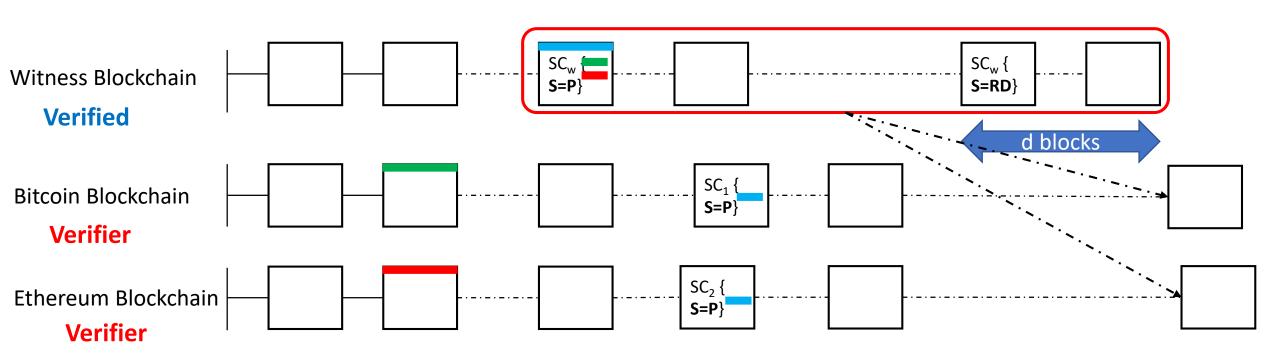




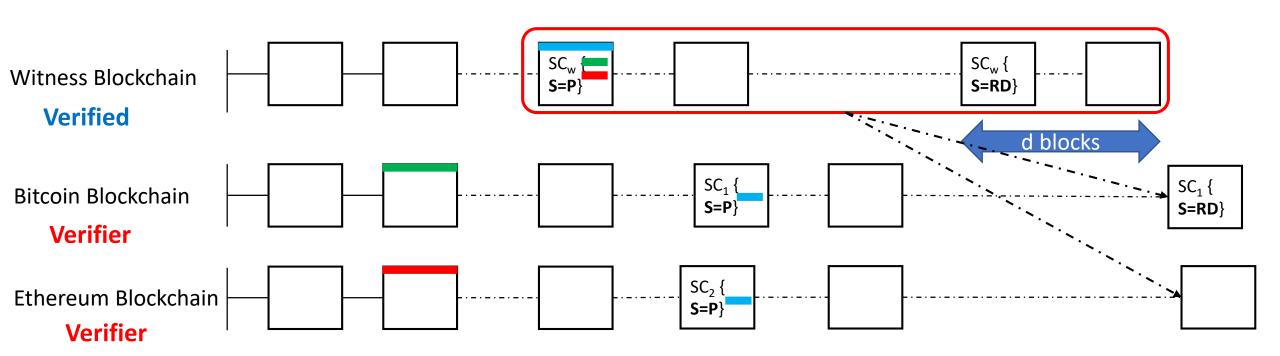




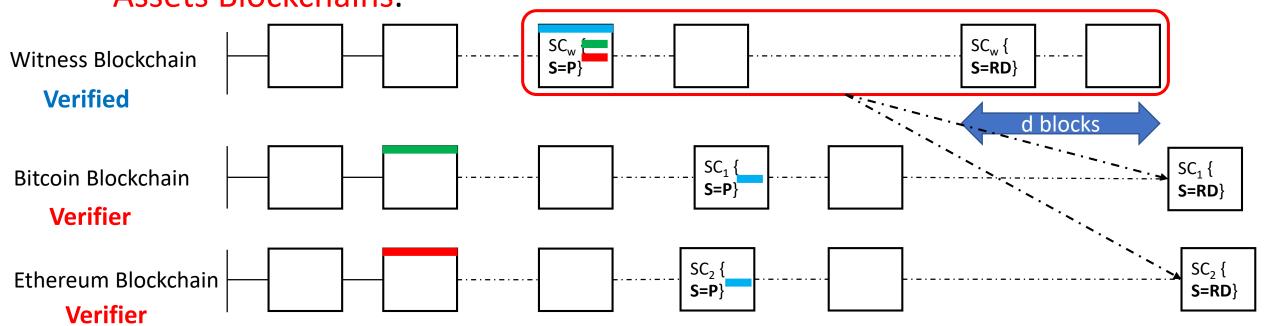
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- After evidence verification, participants redeem their assets from the Assets Blockchains.



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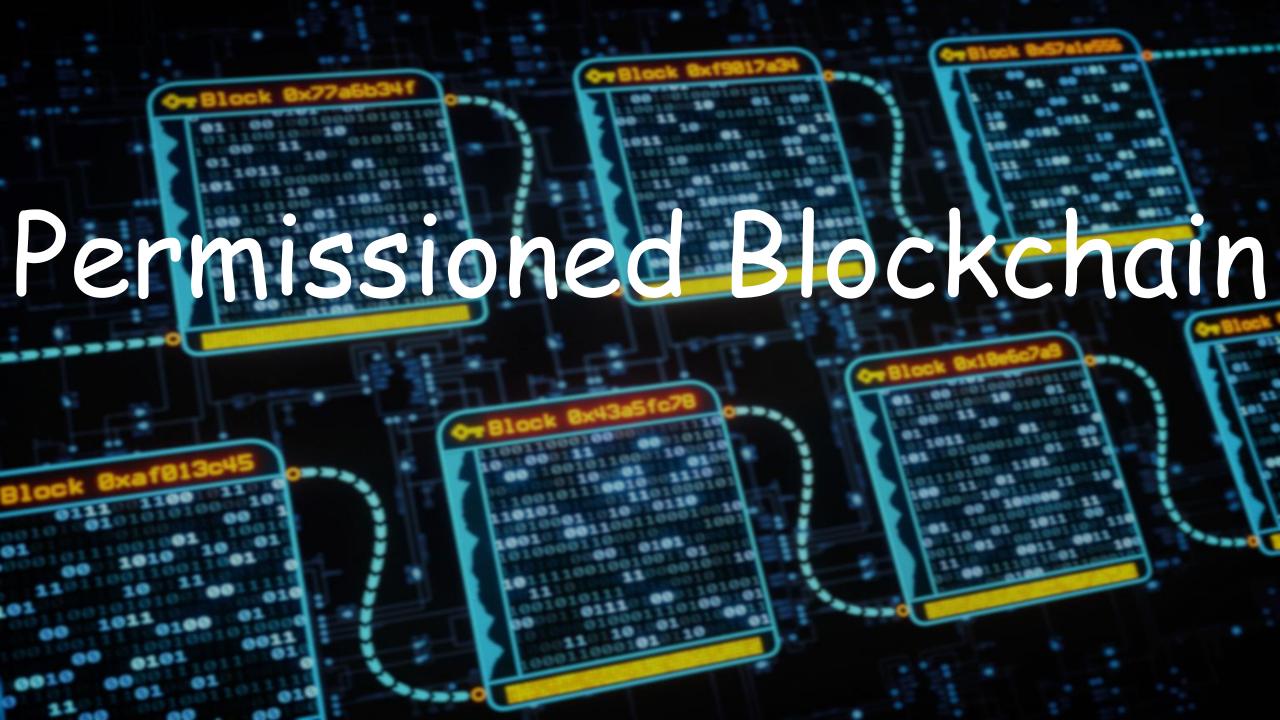
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- Any protocol is prone to fork attacks



# Any applications other than Cryptocurrency?

# Supply Chain Management: Tracking Fish from Ocean to Table

- Ocean fishing represents more than \$70B in worldwide trade<sup>1</sup>
  - Estimates suggest at least 20% of all fish are caught illegally—yet only a tiny fraction are ever inspected<sup>2</sup>.
  - Nearly one in three fish were mislabeled by sellers<sup>3</sup>
  - 87% of snapper and 59% of tuna were mislabelled4
  - 95% of all sushi restaurants were serving mislabeled fish<sup>4</sup>

<sup>&</sup>lt;sup>1</sup> Food and Agriculture Organization, United Nations. 2016. The State of World Fisheries and Aquaculture 2016.

<sup>&</sup>lt;sup>2</sup> Stolen Seafood: The Impact of Pirate Fishing on Our Oceans. Oceana. 2013.

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#### Challenges:

- Many different paths from ocean to table
- Lack of global authority for tracing
- Proprietary tracing systems do not scale
- Most existing processes are paper-based
- The supply chain is extremely complex and includes many participants from different industries

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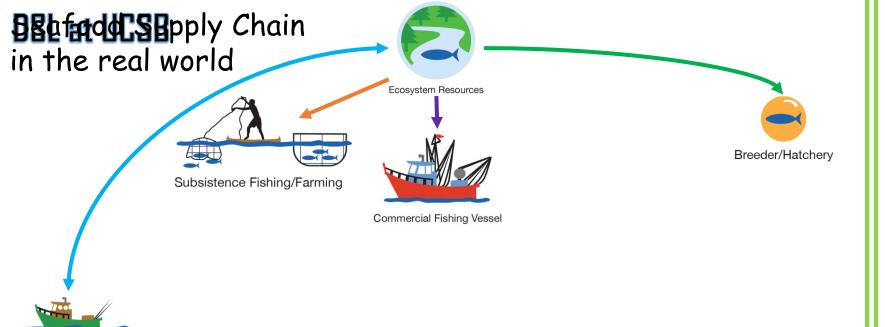
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#### BELIFET CHEST PP Chain in the real world





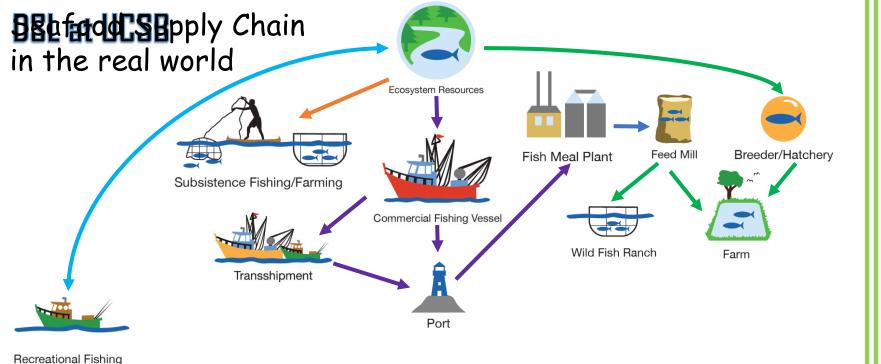


Seafood is caught by fishermen and physically tagged with IOT enabled sensors



Source: Advancing Traceability in the Seafood Industry, FishWise

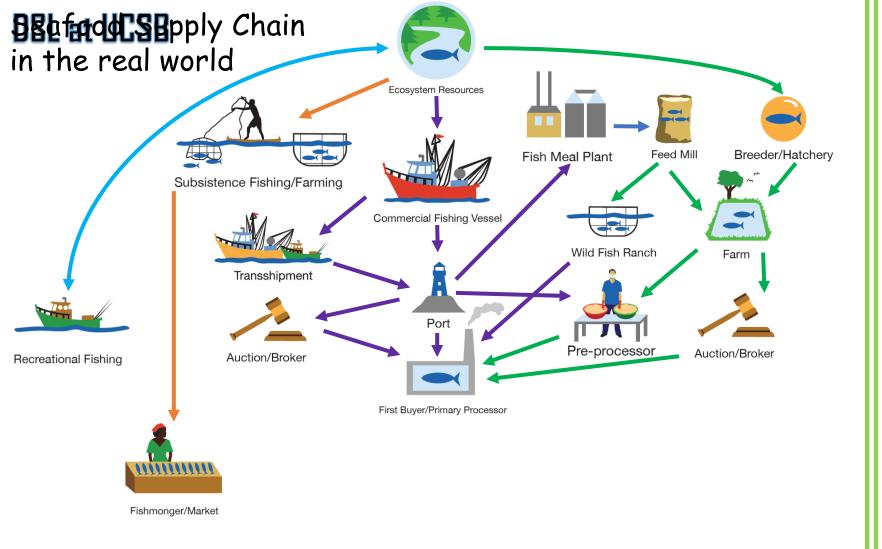
Recreational Fishing



Seafood is caught by fishermen and physically tagged with IOT enabled sensors

Sensors continuously transmit data about time and location to Blockchain

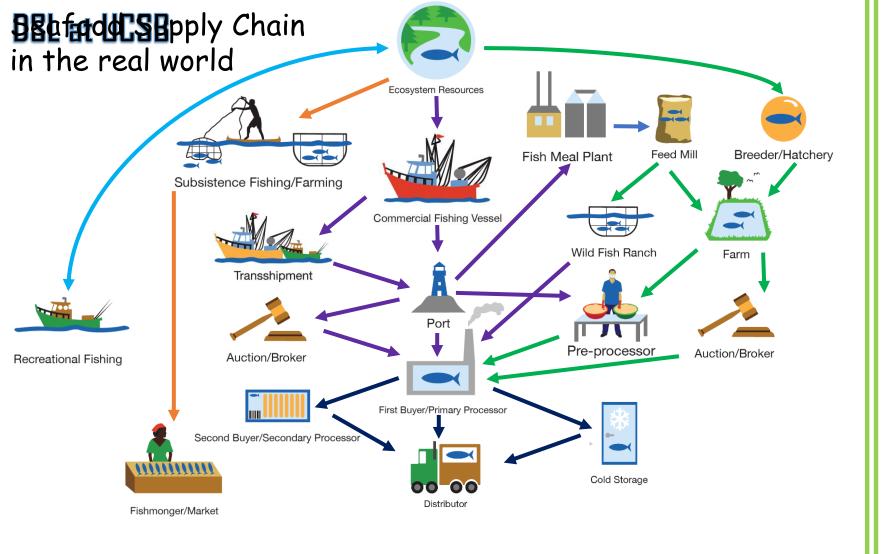




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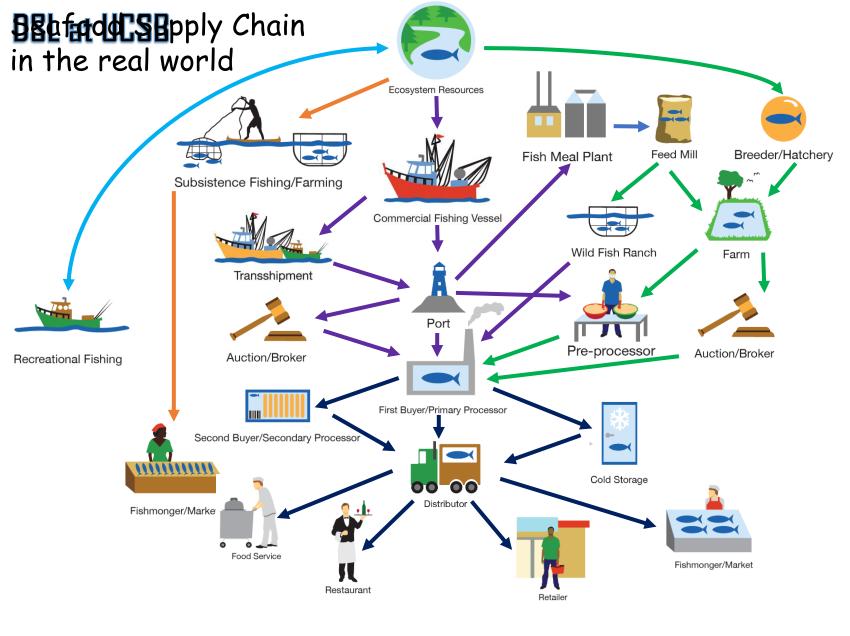


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**Subsistence Fishing** 

Recreational Fishing Aquaculture Wild Capture Fisheries Processing and Distribution

Seafood Supply Chain in Blockchain

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The buyer can access a comprehensive record of the fish's provenance

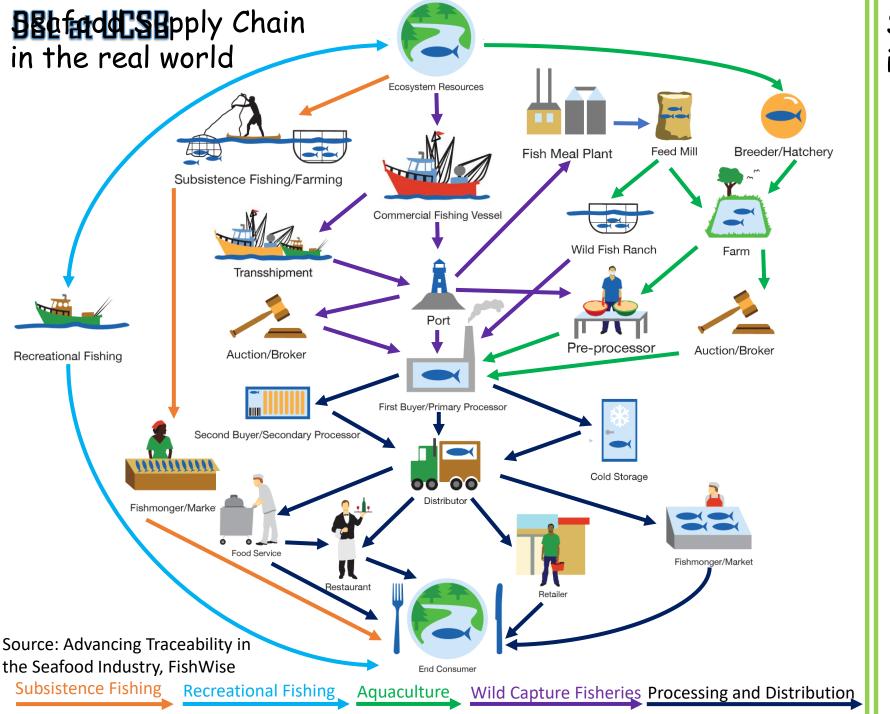












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- Access end-to-end supply chain data instantly and easily with full transparency
- Minimize waste and allocate inventory using insights from real-time demand forecasts



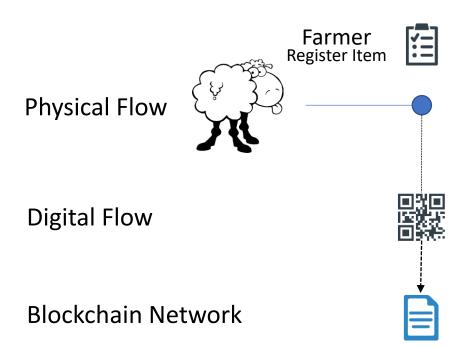
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Digital Flow

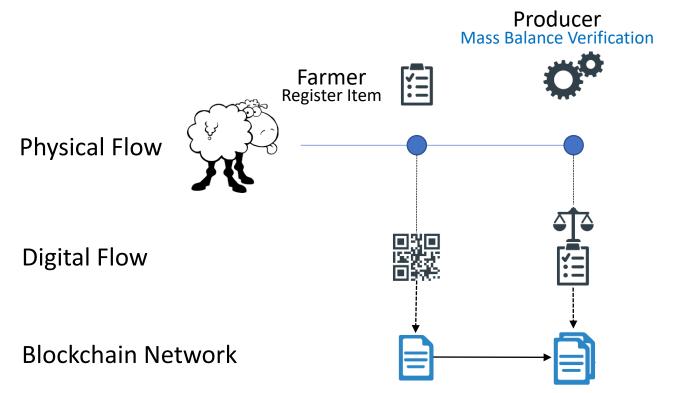
**Blockchain Network** 

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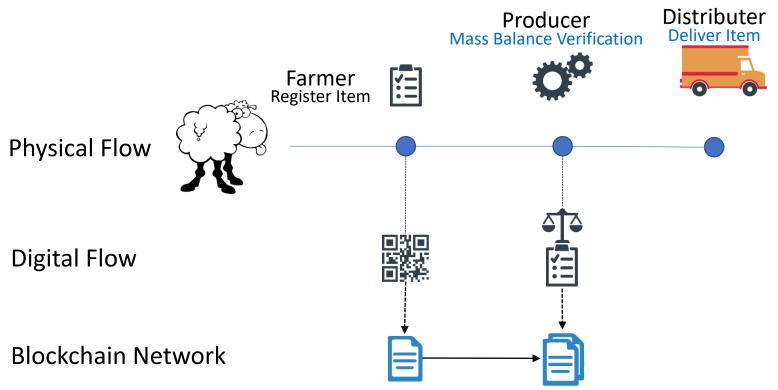


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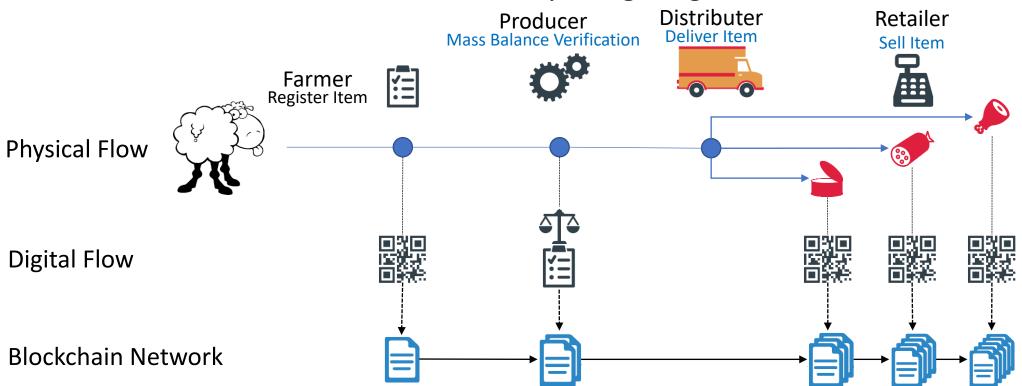


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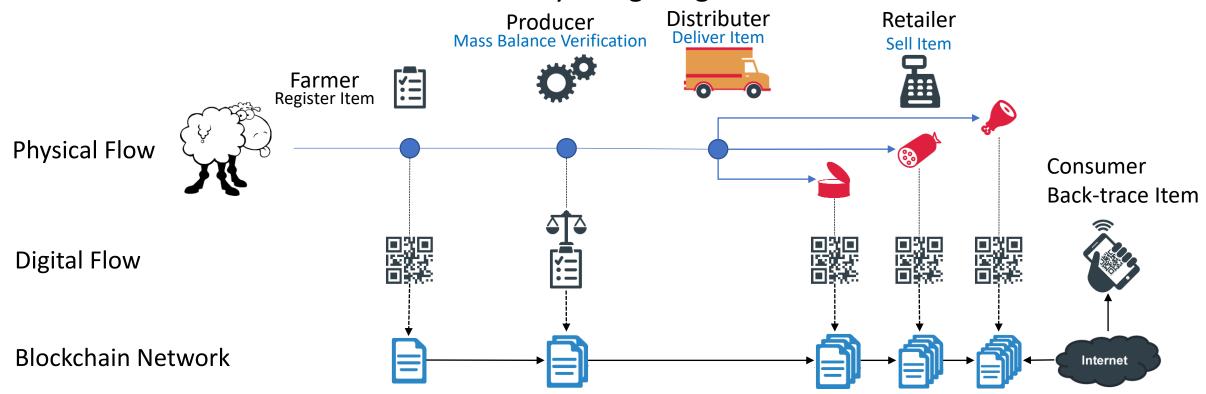


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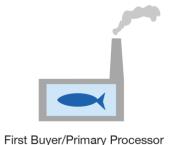


# The difference between Bitcoin and Supply Chain?!

In Supply Chain Participants are known and Identified









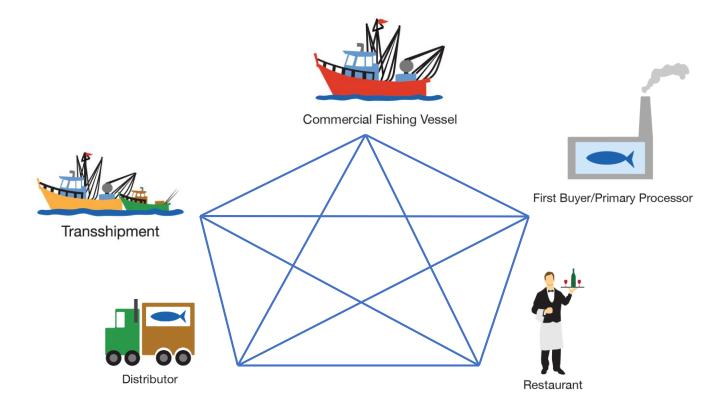


The difference between Bitcoin and Supply Chain?!

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#### Traditional Consensus Protocols can be used





identified nodes that might not fully trust each other.



#### Permissioned Blockchain

- Run a blockchain among a set of known, identified participants
- Provides a way to secure the interactions among a group of entities that have a common goal but which do not fully trust each other
- The ledger is distributed among all the nodes



#### Permissioned Blockchain

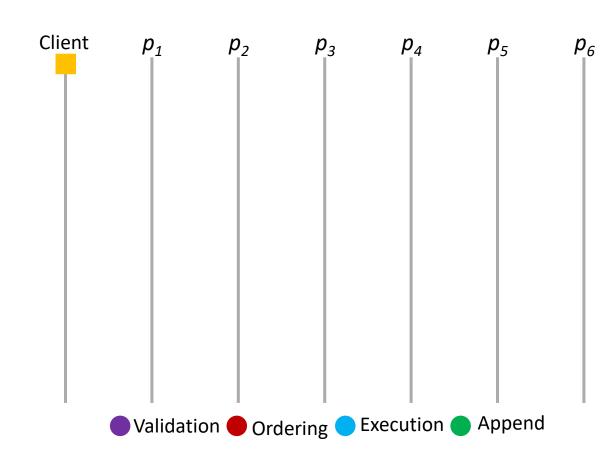
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	Permissionless	Permissioned
Participants	Anonymous, Could be malicious	Known, Identified
Consensus Mechanisms	<ul> <li>Proof of Work, Proof of Stake,</li> <li>Large energy consumption</li> <li>No finality</li> <li>51% attack</li> </ul>	Byzantine fault tolerance Consensus, e.g., PBFT  Lighter Faster Low energy consumption Enable finality
Transaction Approval time	Long (Bitcoin: 10 min or more)	Short (100x msec)

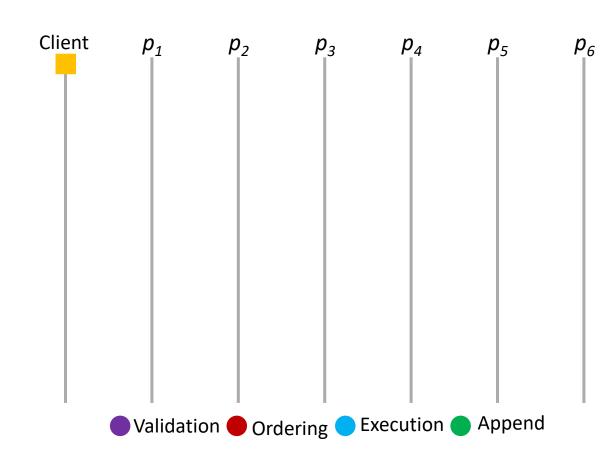
# Consensus Protocols in Permissioned Networks

- Types of systems: synchronous and asynchronous
- Problem statement: given N processes (one of them is the *leader*):
  - Agreement: all correct processes agree on the same value
  - Validity: If initiator does not fail, all correct processes agree on its value
- Types of failure:
  - Crash
  - Malicious (or Byzantine)
- Important *impossibility* result:
  - FLP, in asynchronous systems:
    - With even one crash failure, termination is not guaranteed (no liveness)
  - Synchronous systems:
    - Termination is guaranteed if number of failed malicious processes (f) is at most 1/3 n

# DSL at UCSB Bitcoin review

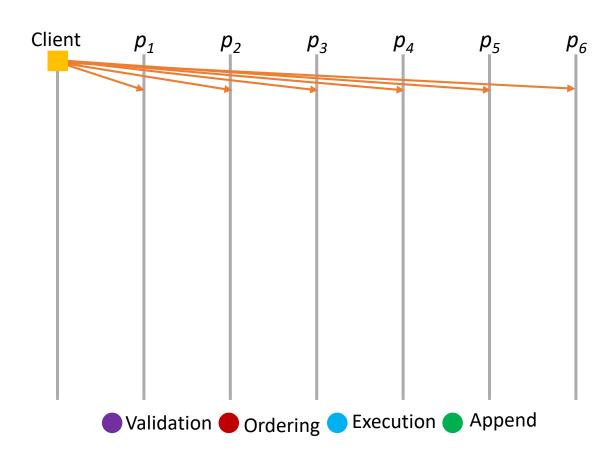


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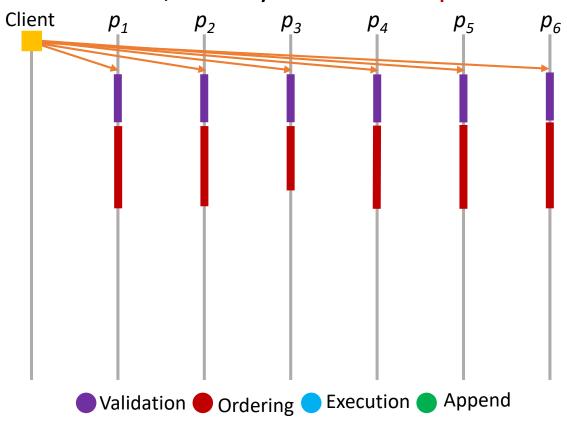
### Bitcoin review

• Clients multicasts their requests



#### Bitcoin review

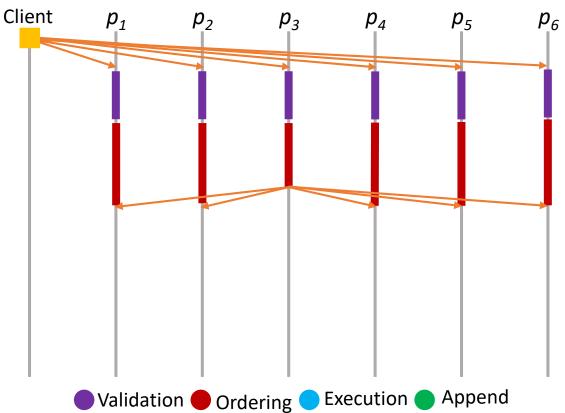
- Clients multicasts their requests
- Nodes validate the transactions, put them into the blocks, and try to solve the puzzle



#### Bitcoin review

WIN

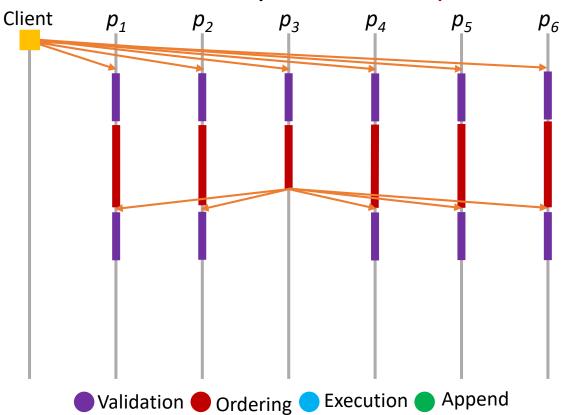
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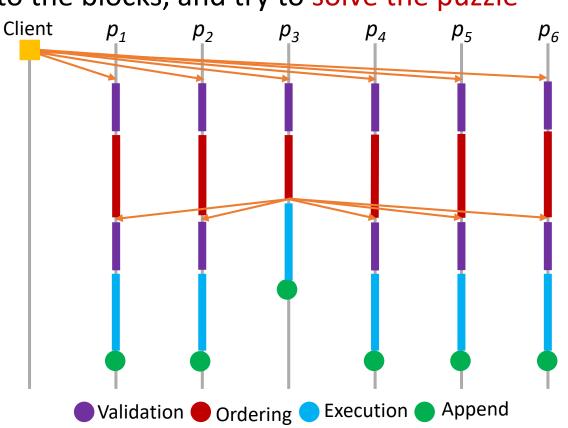
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- Each node validates the transactions within the block



#### Bitcoin review

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- Clients multicasts their requests
- Nodes validate the transactions, put them into the blocks, and try to solve the puzzle
- The lucky node who solves the puzzle first multicasts the block
- Each node validates the transactions within the block
- Transactions are deterministically executed by every node and appended to the ledger

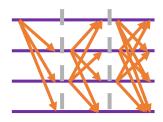


Order-execute Architecture

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  - Confidentiality of execution: all smart contracts run on all peers!

Execute-Order-Validate Architecture

## Execute-Order-Validate Architecture

• Each transaction (of an application) is first executed by a subset of nodes (endorsers of the application)



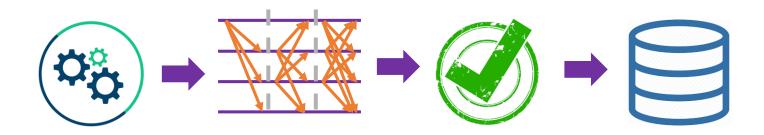
## Execute-Order-Validate Architecture

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**Parallel Execution:** Transactions of different applications can be executed in parallel

# Hyperledger Fabric

• Three types of Nodes: Clients, Endorsers, and Orderers

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    - All endorsers maintain the blockchain ledger
    - Each application has its own set of endorsers





Endorsers (of different applications)

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  - Clients send transactions to be executed.
  - Endorsers execute transaction proposals and validate transactions.
    - All endorsers maintain the blockchain ledger
    - Each application has its own set of endorsers
  - Orderers stablish the total order of all transactions using a consensus protocol
    - Do not maintain the blockchain ledger or smart contracts
    - The consensus protocol is pluggable



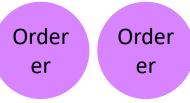






Endors er Endors er

Endorsers (of different applications)



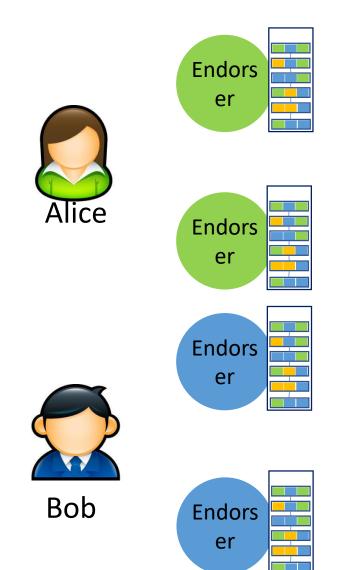
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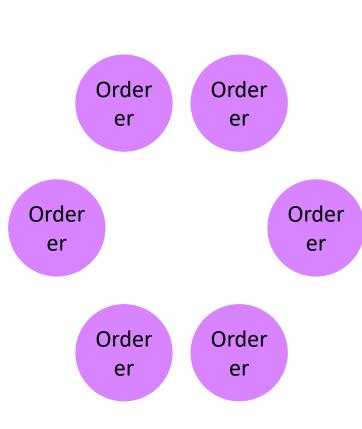
**Orderers** 

Clients (of different applications)

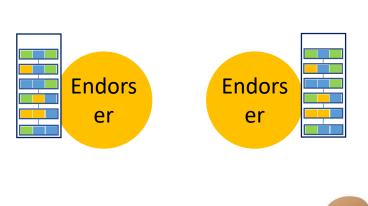
Hyperledger Fabric

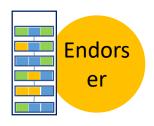
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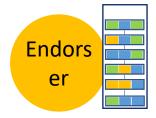




Three Applications (Green, Blue, Yellow)
Three Clients (Alice, Bob, Charlie)
Green and Blue have two Endorsers, Yellow
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There are totally six Orderers

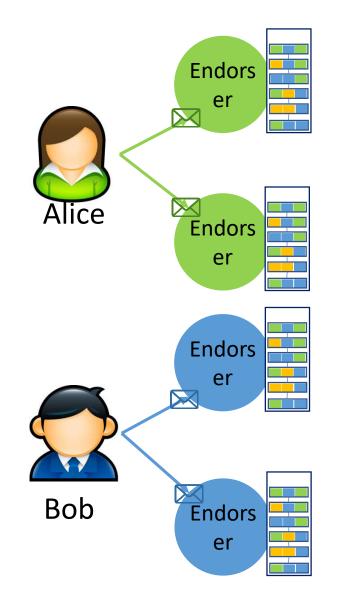


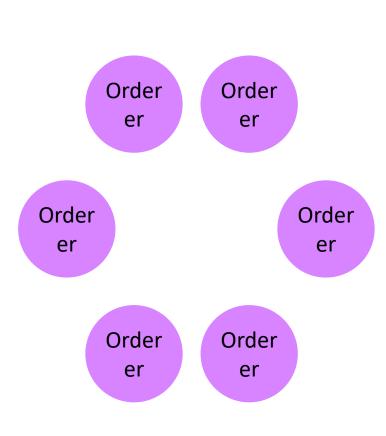


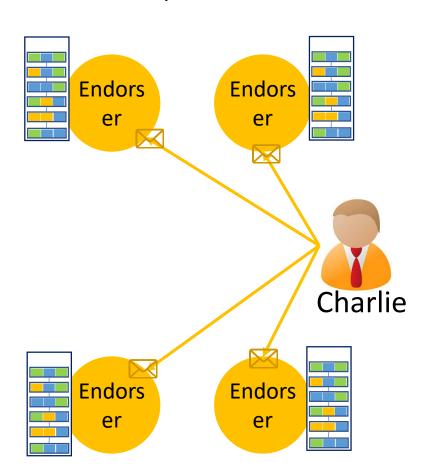


Charlie

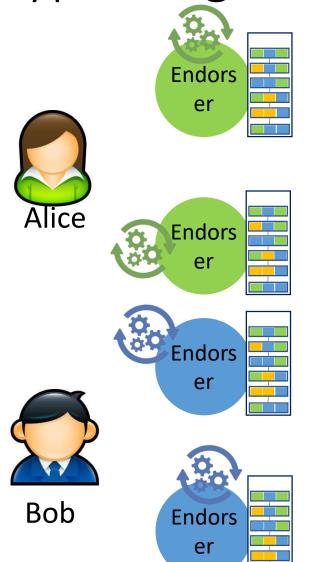








# Hyperledger Fabric



Transactions of different applications are executed in parallel

> Order er

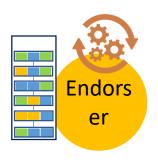
Order er

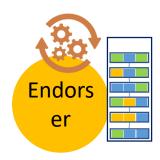
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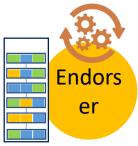
Order er

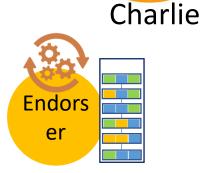
Order er

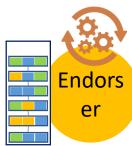
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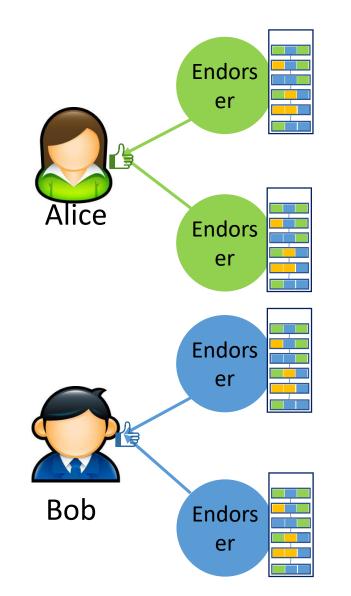


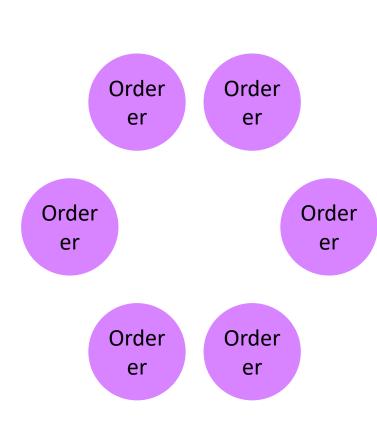


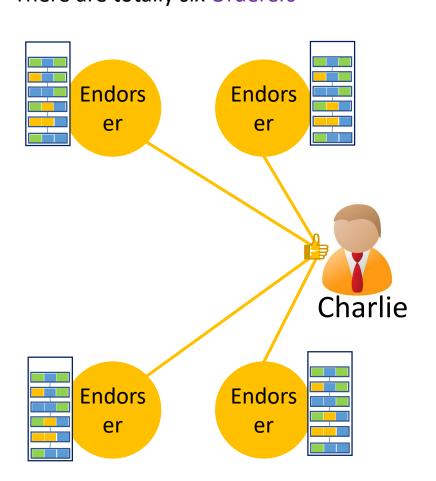










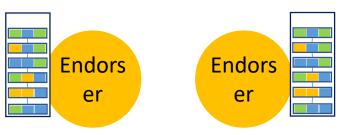


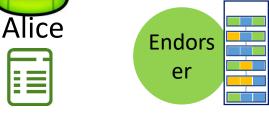
# Hyperledger Fabric

**Endors** er

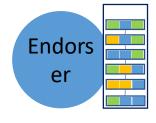
If the results are *identical*, the client put them into a request





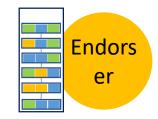


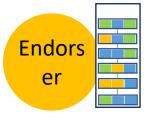






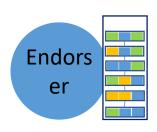


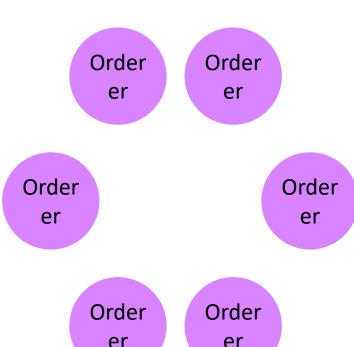




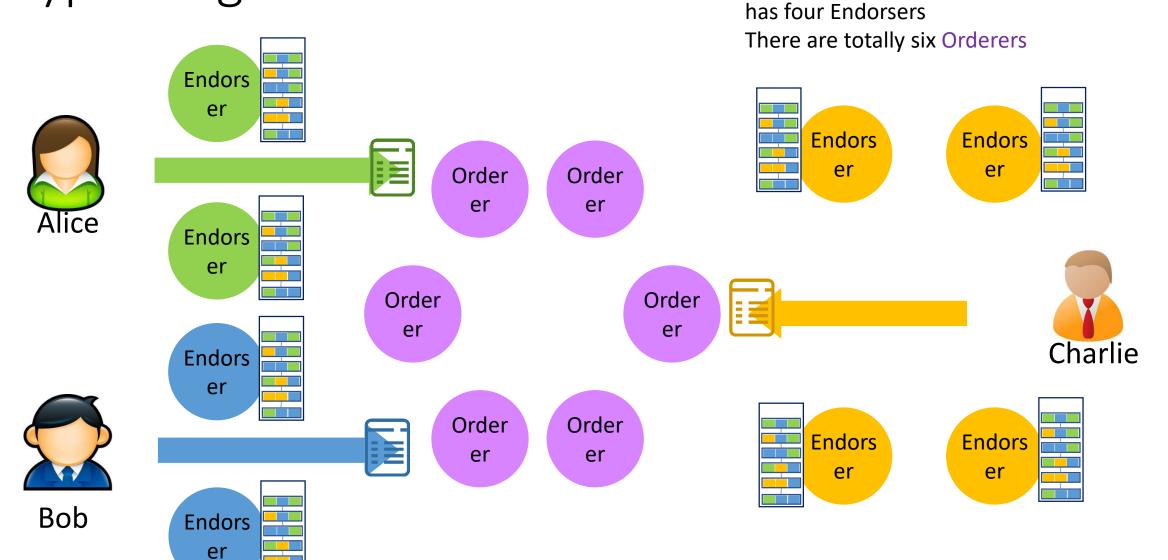








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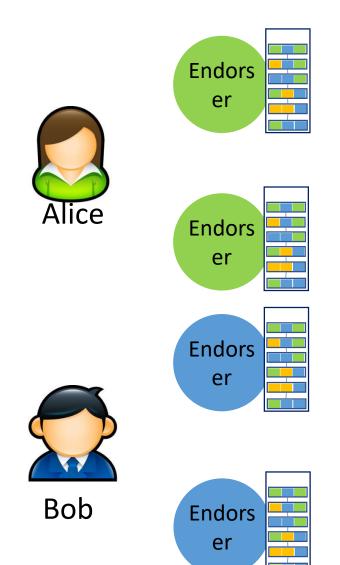


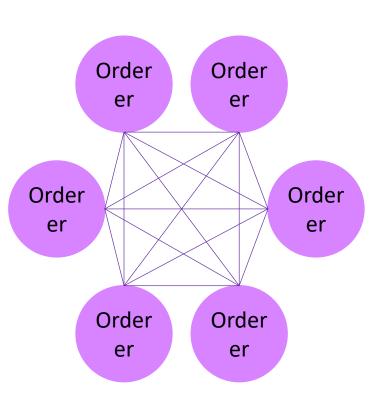
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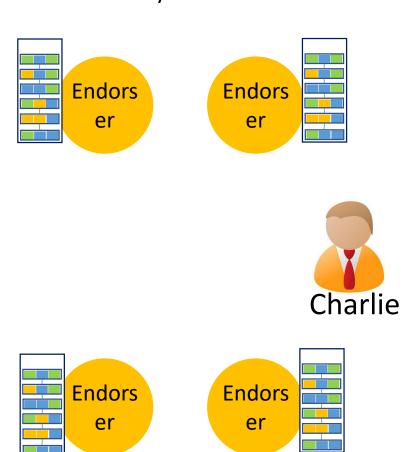
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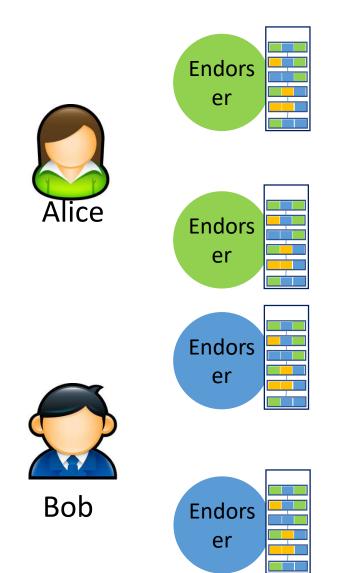
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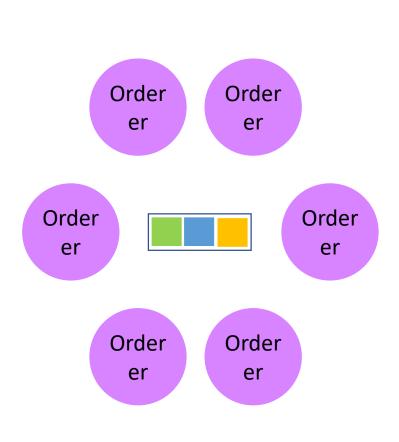




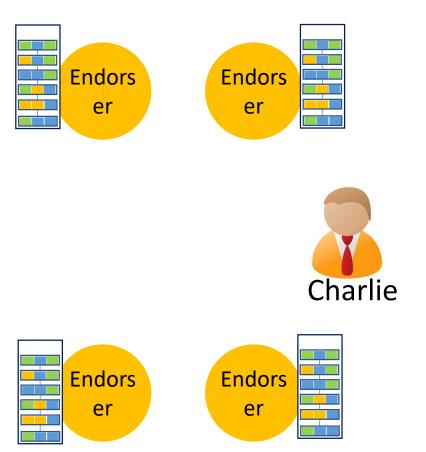


# Hyperledger Fabric



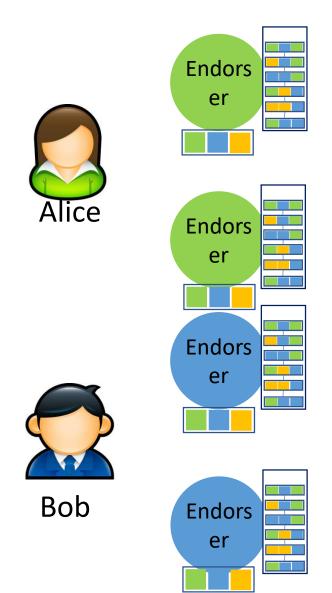


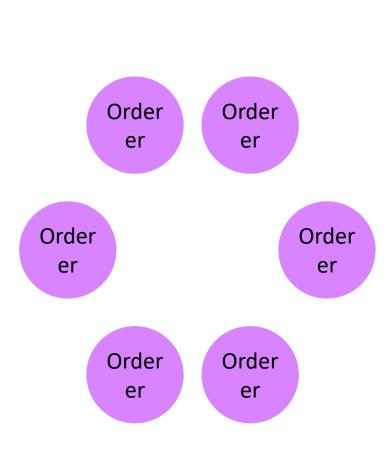
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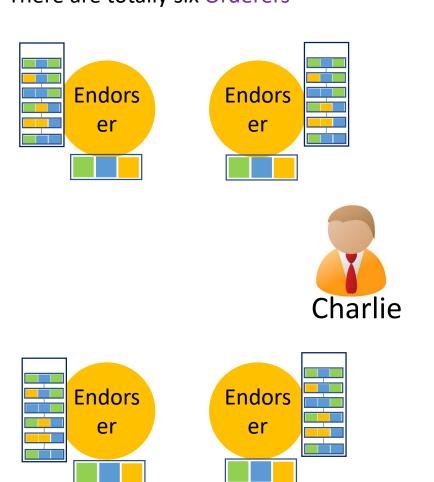


A block might contains multiple transactions from the same application

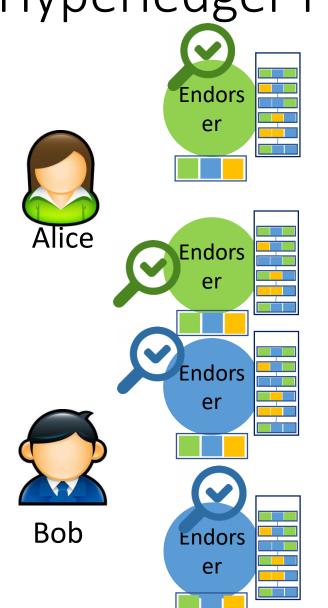
# Hyperledger Fabric

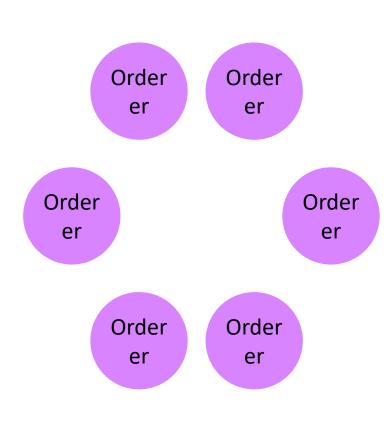




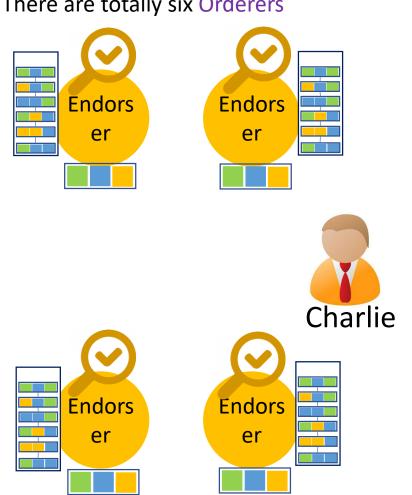


Hyperledger Fabric



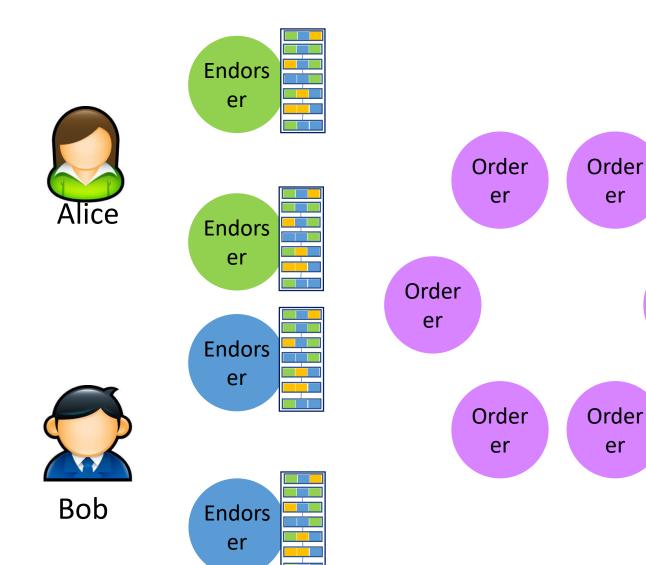


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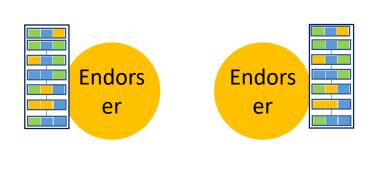


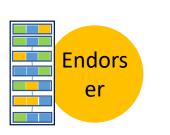
In the validation phase, Endorsers check: (1) *validity* of transactions, (2) read-write conflicts

# Hyperledger Fabric



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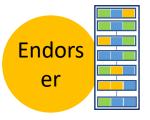


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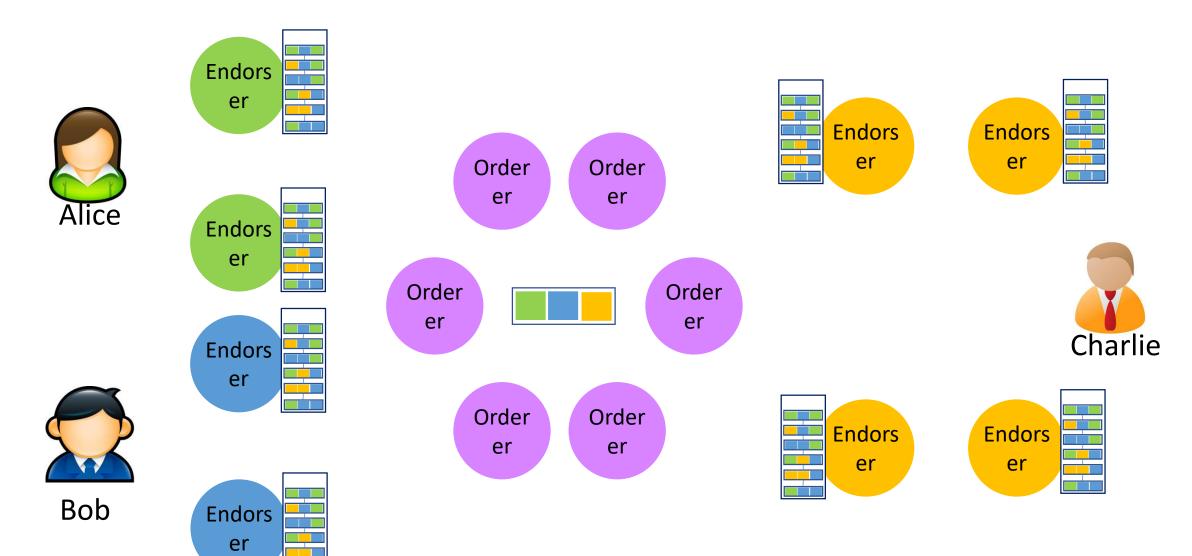
Order

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Charlie

# Hyperledger Fabric



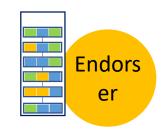
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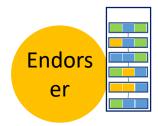
Writes record A
Reads record A
Reads record A

#### What if transactions are conflicting?

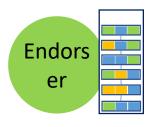
Endors

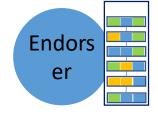
transactions access the same record and one of them is a write operation

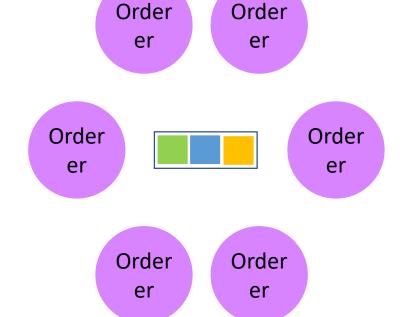






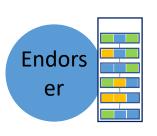


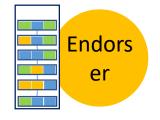


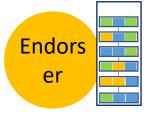










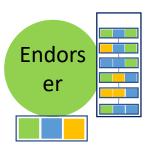


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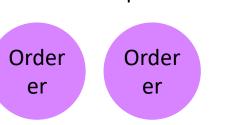
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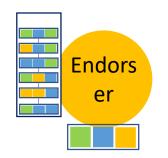
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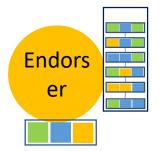


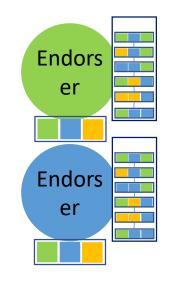


transactions access the same record and one of them is a write operation













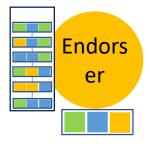


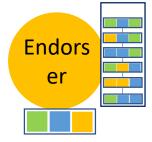
Bob

Endors er









# Hyperledger Fabric

**Endors** 

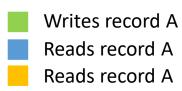
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**Endors** 

er

**Endors** 

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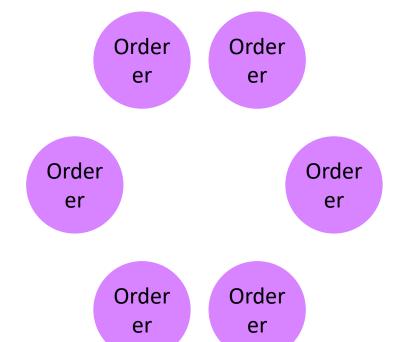


Alice

Bob

#### What if transactions are conflicting?

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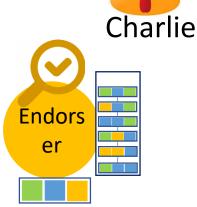




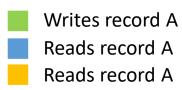


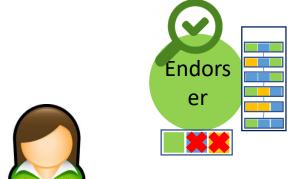
**Endors** 

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# Hyperledger Fabric





Alice

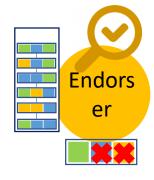
Bob

#### What if transactions are conflicting?

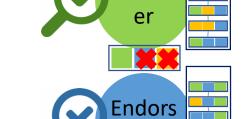
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**Endors** 

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**Endors** 

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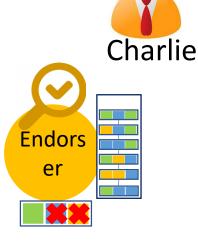








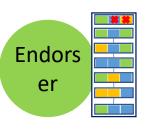




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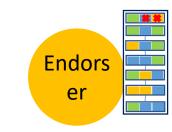
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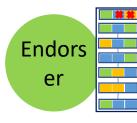
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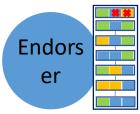


transactions access the same record and one of them is a write operation





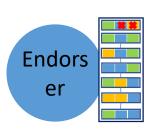






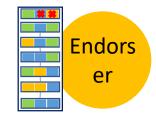


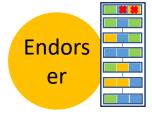












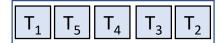
## Dependency Graph

• A dependency graph exposes conflicts between transactions to give a partial order of transactions.

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```
T_1 Read = \{a\} Write = \{a,b\}
```



$$T_2$$
 Read =  $\{f\}$  Write =  $\{d\}$ 

$$T_3$$
 Read =  $\{f\}$  Write =  $\{e\}$ 

$$T_4$$
 Read =  $\{b\}$  Write =  $\{c\}$ 

$$T_5$$
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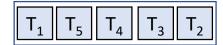
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 $\mathsf{T_1}$ 

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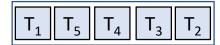
 $\mathsf{T}_1$ 

**T**<sub>5</sub>

## Dependency Graph

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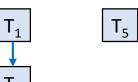
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Write = 
$$\{e\}$$

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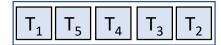


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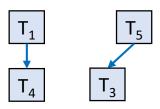


$$T_2$$
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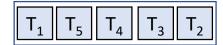
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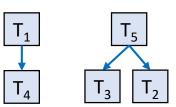


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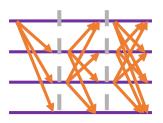
T<sub>3</sub> writes e that is read by T<sub>5</sub>

T<sub>2</sub> writes d that is written by T<sub>5</sub>

Order-Parallel Execute (OXII) Architecture

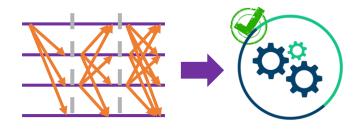
# Order-Parallel Execute (OXII) Architecture

 A separate set of nodes (orderers) orders the transactions, puts them into blocks, generates a dependency graph for the block, and multicasts it to all the nodes.



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- The nodes multicast the results of execution and append the block



### ParBlockchain







Order-Execute Architecture: Transactions are first ordered, and then executed





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**Parallel Execution:** non-conflicting transactions of the same or different applications are executed in parallel



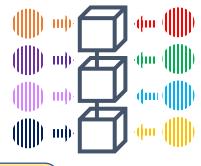


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**Non-deterministic Execution:** inconsistent execution results can be detected in the last phase (results in decreasing the performance)



#### Clients

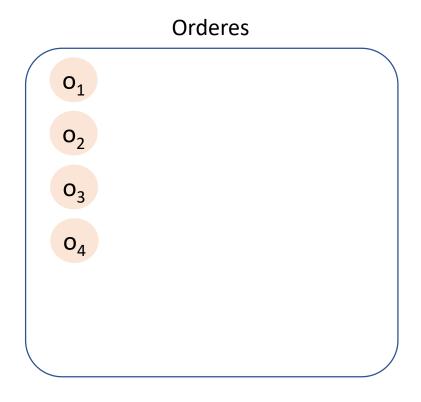


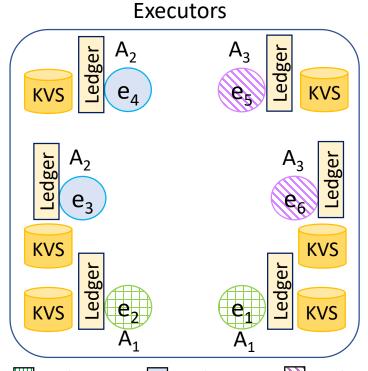












Application  $A_1$  Application  $A_2$  Application  $A_3$ 

Each application has a set of Executors

Each Executor stores a copy of ledger and Data

### ParBlockchain

#### Clients





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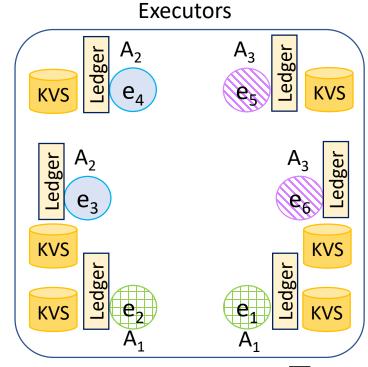
#### Orderes

 $O_1$ 

02

03

 $O_4$ 



 $\blacksquare$  Application  $A_1$  Application  $A_2$  Application  $A_3$ 

Each transaction of an application include records to be read and written

### ParBlockchain

#### Clients







$$T_2$$
 Read = {f} Write = {d}

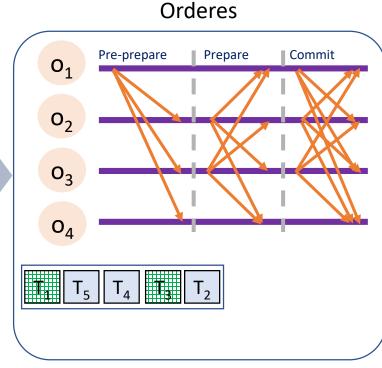




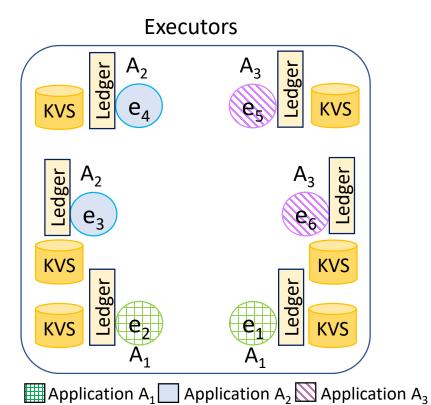
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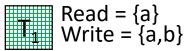
The orderers order transactions using a consensus protocol (e.g. PBFT)



### ParBlockchain

#### Clients







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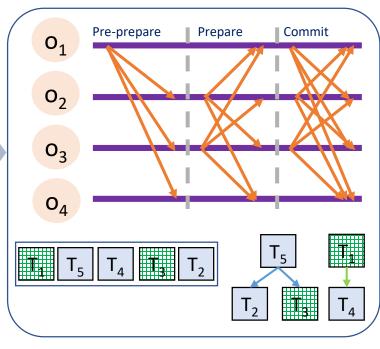




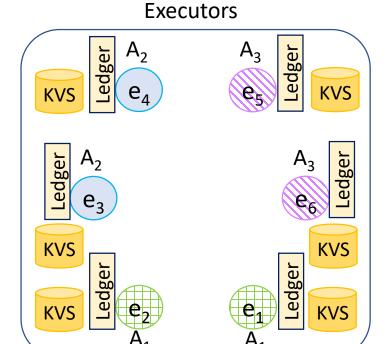
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**Orderes** 



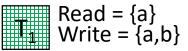
Application  $A_1$  Application  $A_2$  Application  $A_3$ 

Each orderer generates a dependency graph for the block and multicasts it to all Executors

### ParBlockchain

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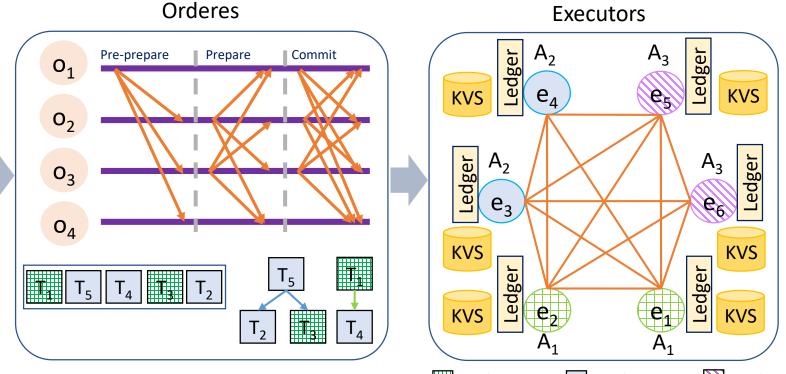




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 $\blacksquare$  Application  $A_1$  Application  $A_2$   $\blacksquare$  Application  $A_3$ 

Executors of each application execute the corresponding transactions following the dependency graph and multicast the results



# Optimistic vs. Pessimistic Execution

Two ways to look at the problem!

Supporting non-deterministic execution

**Supporting High Contention Workloads** 



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Hyperledger

Executes first (does not submit transactions with inconsistent results)

Validates read-write conflicts last (aborts conflicting transactions)



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Supporting non-deterministic execution

Supporting High Contention Workloads

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Validates read-write conflicts last (aborts conflicting transactions)

ParBlcockchain

Validates non-determinist execution last (aborts transactions with inconsistent results)

Checks conflicts first (generates a dependency graph)

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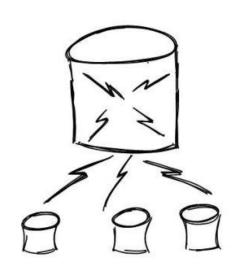
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**Sharding** (as a horizontal technique): Partitioning the data into multiple shards that are maintained by different subsets of nodes



# **Sharding Blockchains**

• Partition the nodes into clusters of *3f+1* nodes (to guarantee safety in each cluster in the presence of f malicious nodes)



















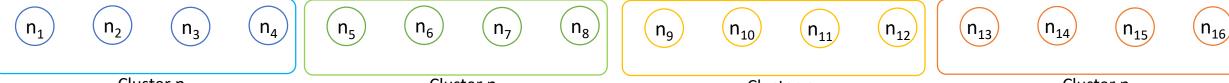




Cluster  $p_1$  Cluster  $p_2$  Cluster  $p_3$  Cluster  $p_4$ 

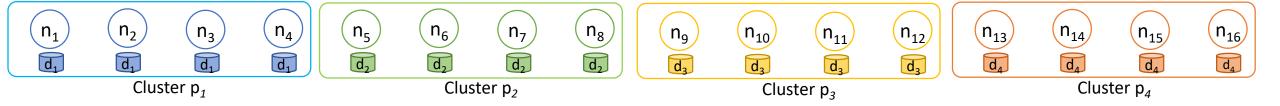
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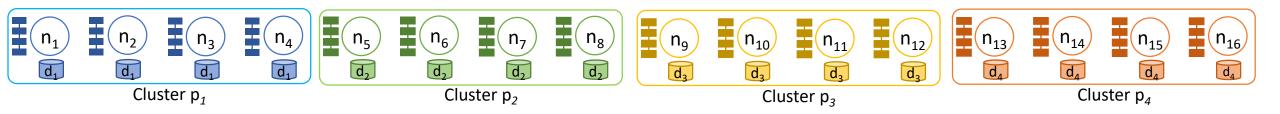


Cluster  $p_1$  Cluster  $p_2$  Cluster  $p_3$  Cluster  $p_4$ 

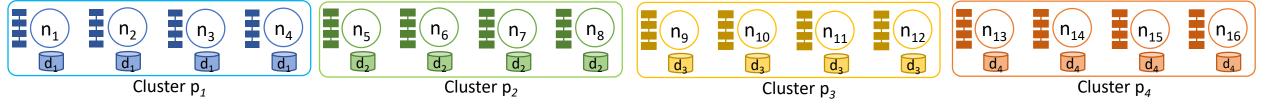
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- Cross-Shard transactions
  - Need the participant of all (and only) involved clusters





Amiri, Mohammad Javad, Divyakant Agrawal, and Amr El Abbadi. Sharding Permissioned Blockchains, IEEE International Conference on Blockchain, 2019



 The blockchain ledger is generalized from a linear chain to a directed acyclic graph (DAG)

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- The entire blockchain ledger is not maintained by any node
- Each node only maintains its own view of the blockchain ledger
  - including the transactions that access the data shard of the cluster

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# SharPer Ledger

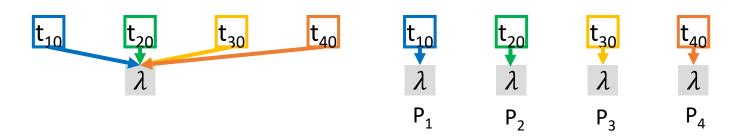
The Blockchain Ledger and the view of clusters P<sub>1</sub>, P<sub>2</sub>, P<sub>3</sub>, and P<sub>4</sub>

 $\lambda$   $\lambda$   $\lambda$   $\lambda$   $\lambda$   $P_1$   $P_2$   $P_3$   $P_4$ 

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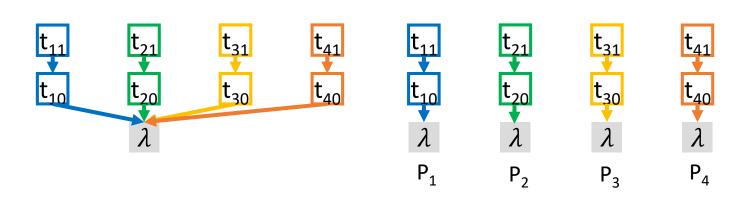
• Intra-shard transactions of different clusters are processed in parallel



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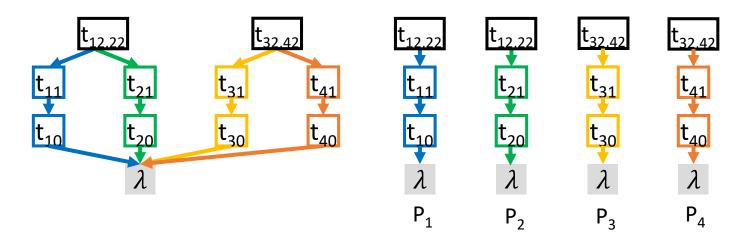
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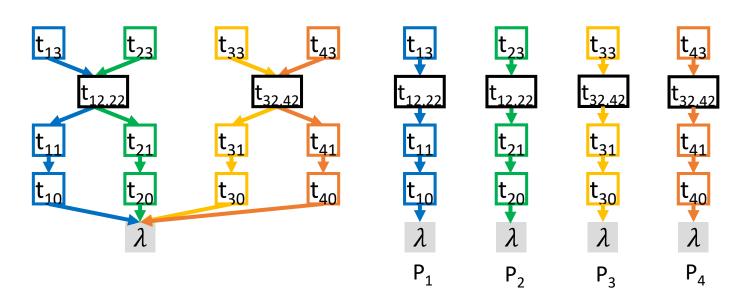


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- A cross-shard transaction includes multiple hash pointers



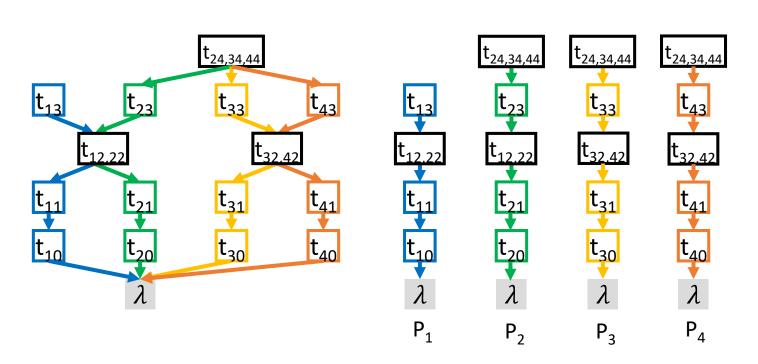


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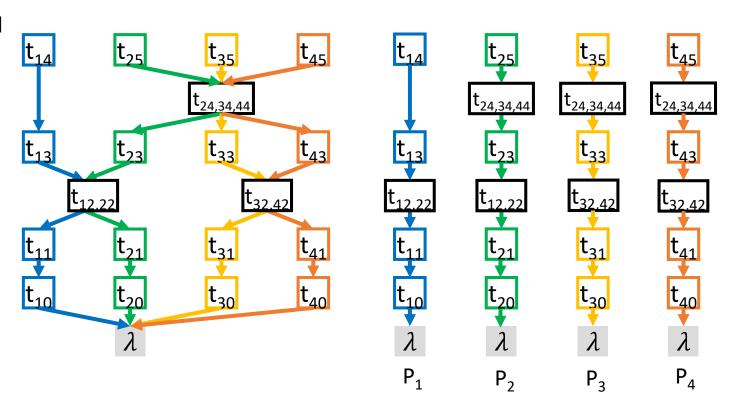
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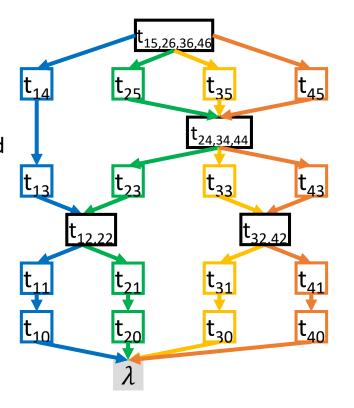


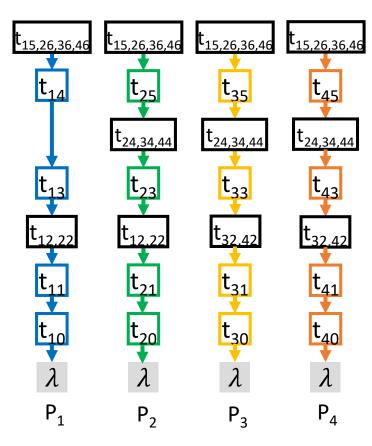
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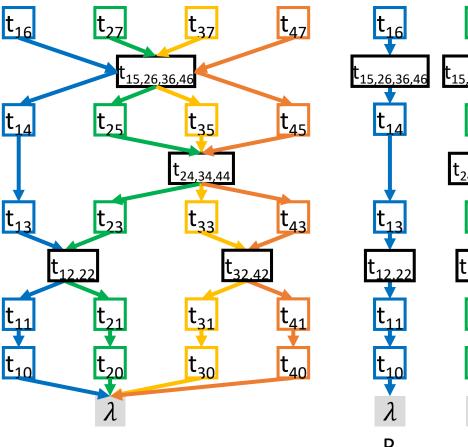
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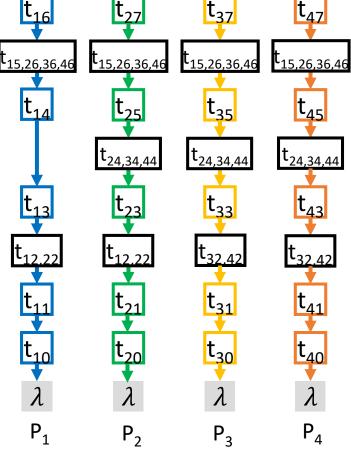






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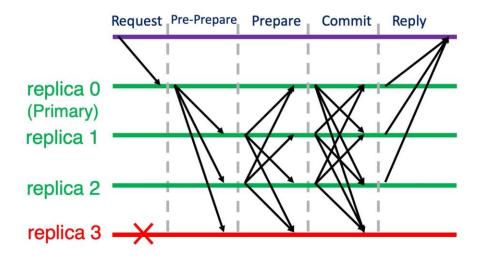


## Consensus in SharPer



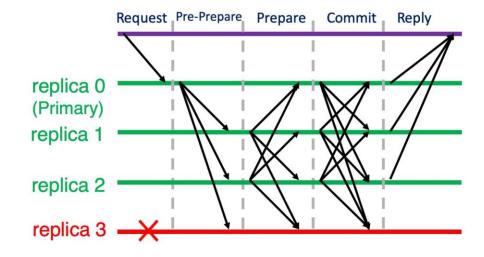
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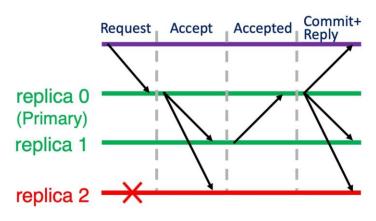
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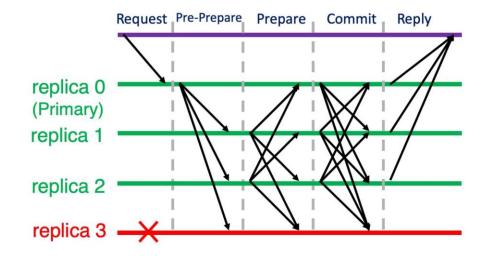
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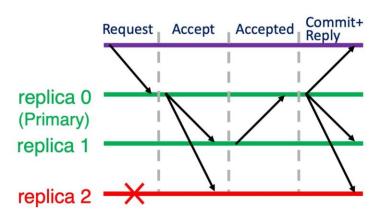




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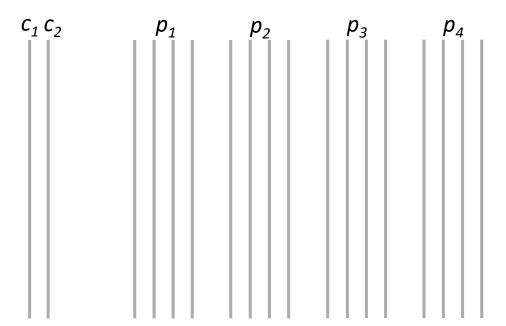
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- If nodes follow crash failure model, use crash fault-tolerant protocol, e.g., Paxos
- Cross-Shard Consensus: needs the participation of all the involved clusters
  - In each step 2f+1 nodes of every involved cluster must participate





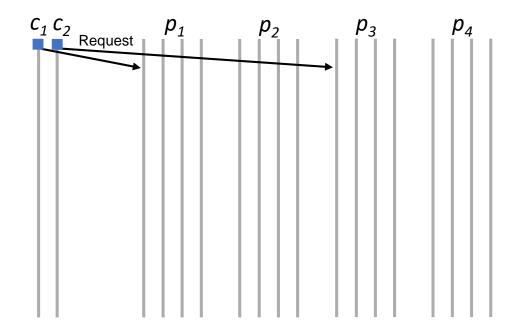


Non-overlapping cross-shard transactions can be processed in parallel



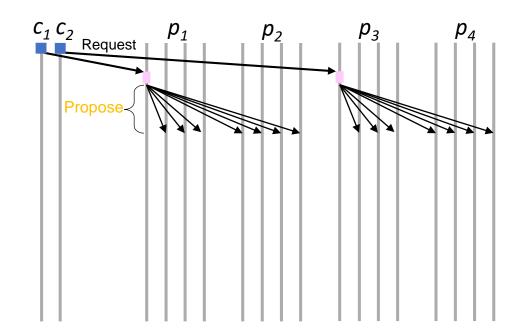


Non-overlapping cross-shard transactions can be processed in parallel Clients ( $c_1$  and  $c_2$ ) send requests to the (pre-elected) primary nodes



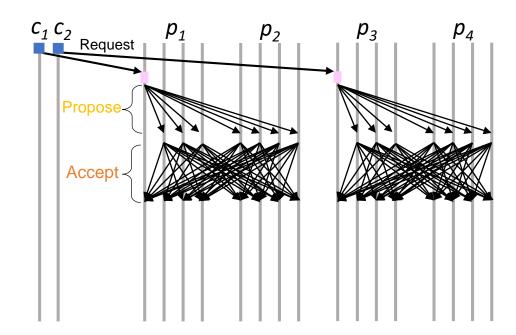


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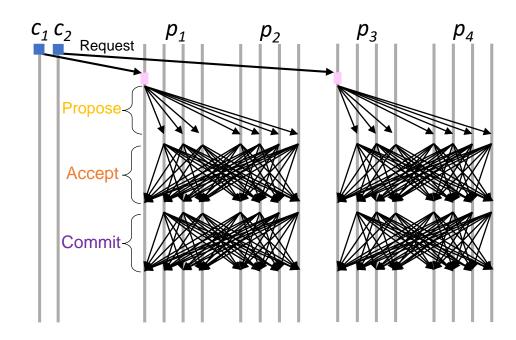


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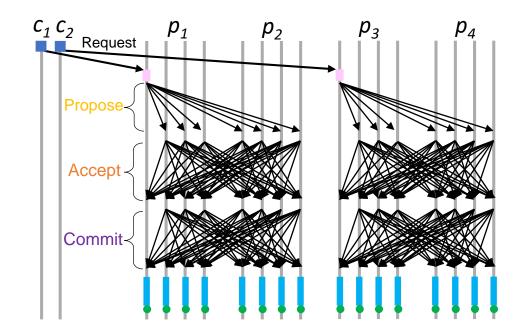


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# Collaborative Workflow: Supply Chain Management

- Different parties (applications) need to communicate across organizations to provide services
- The communication follows Service Level Agreements (agreed upon by all participants)
- They do not trust each other
- The blockchain system should support both cross-application and internal transactions
- Internal data of each party is confidential





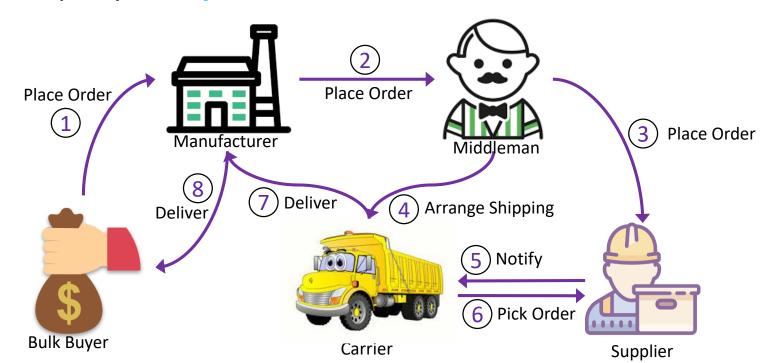






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Performance issue



# CAPER: A Cross-Application Permissioned Blockchain

Distributed applications collaborate with each other following SLAs

Mohammad Javad Amiri, Divyakant Agrawal, Amr El Abbadi, CAPER: A Cross-Application Permissioned Blockchain, The 45th International Conference on Very Large Data Bases (VLDB), 2019.



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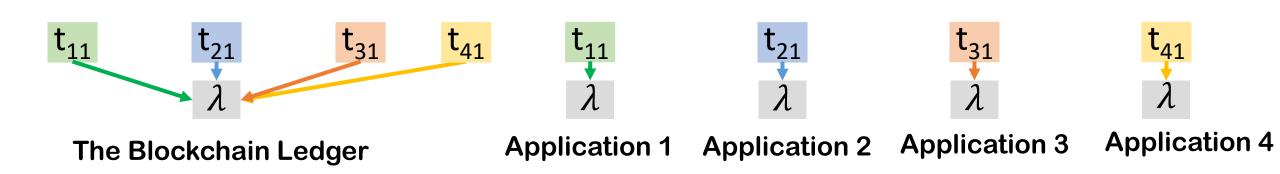
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- Each application maintains only its own view of the ledger
  - including its internal and all cross-application transactions.

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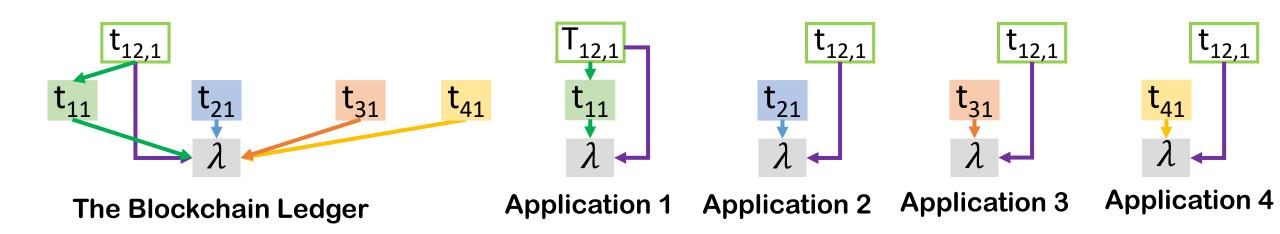
# The Blockchain Ledger of CAPER

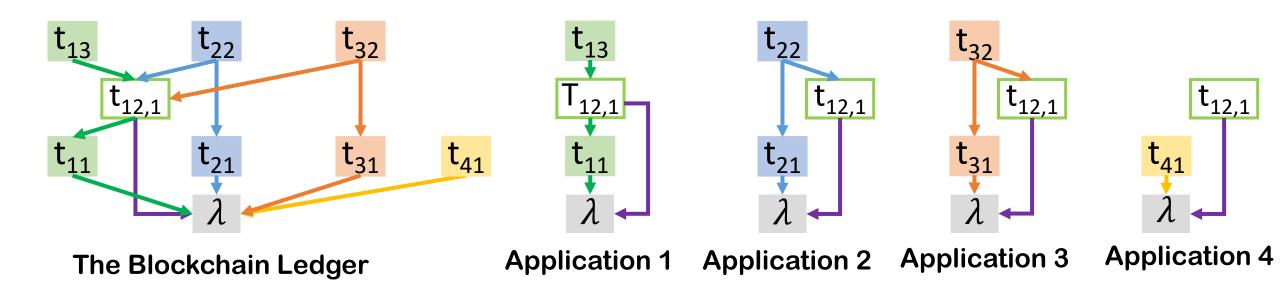
Each application has its own internal transactions

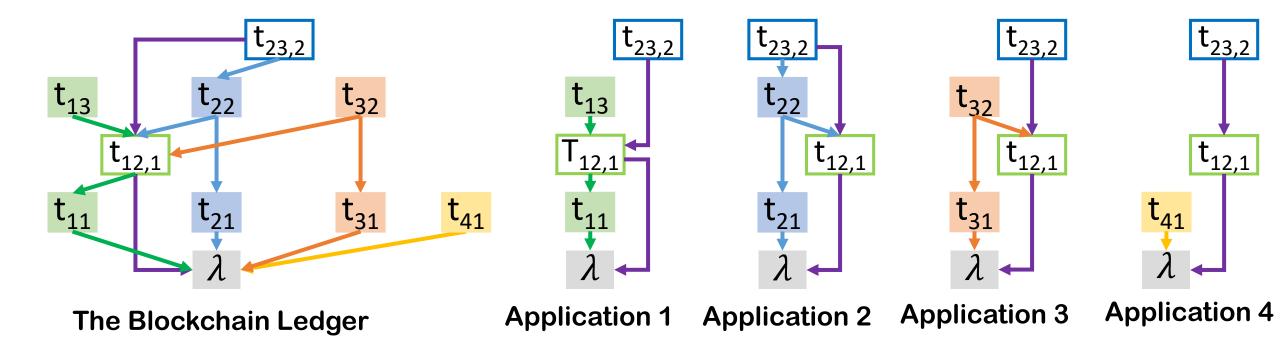


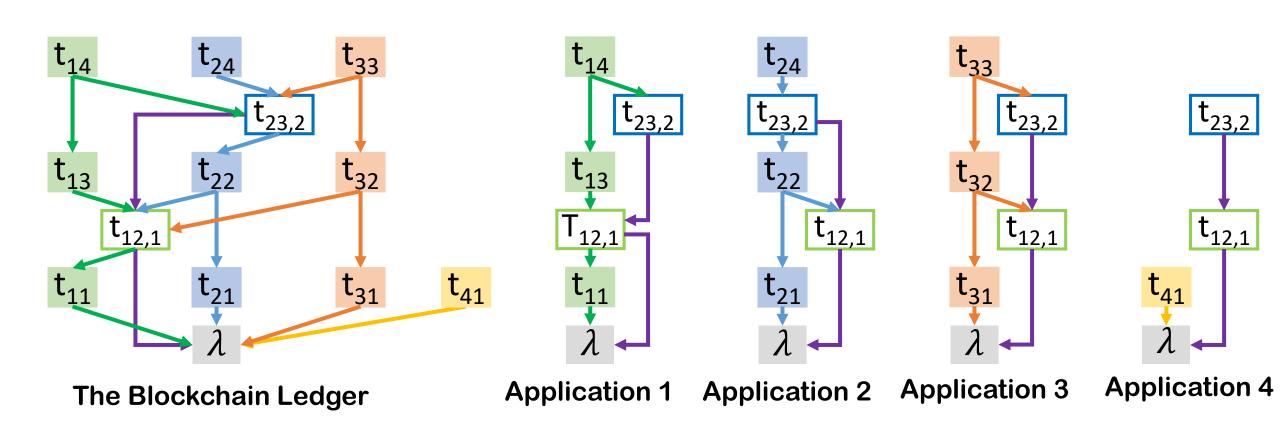
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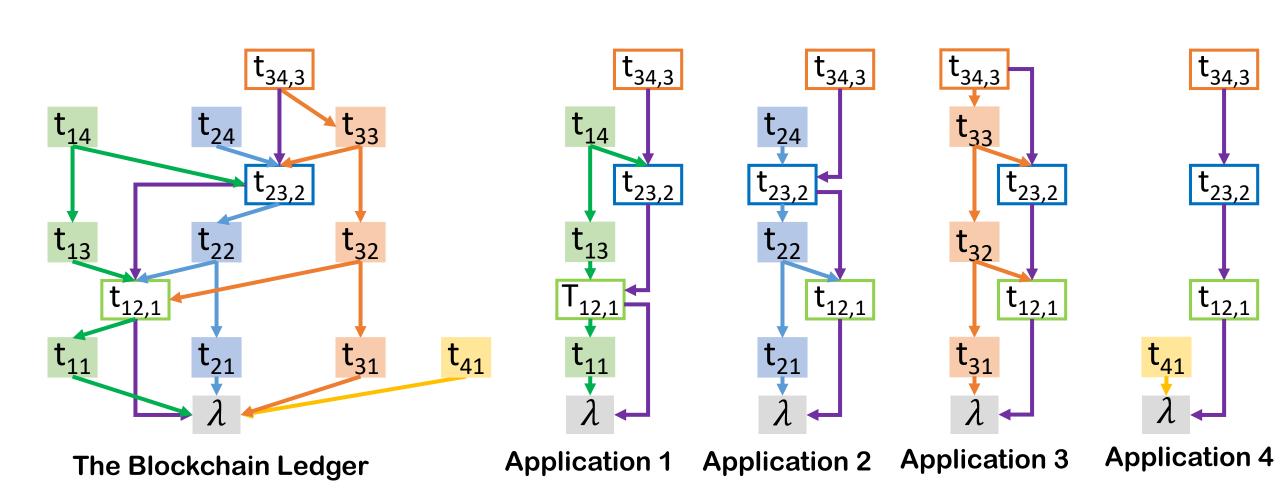
Cross-application transactions are maintained by every application

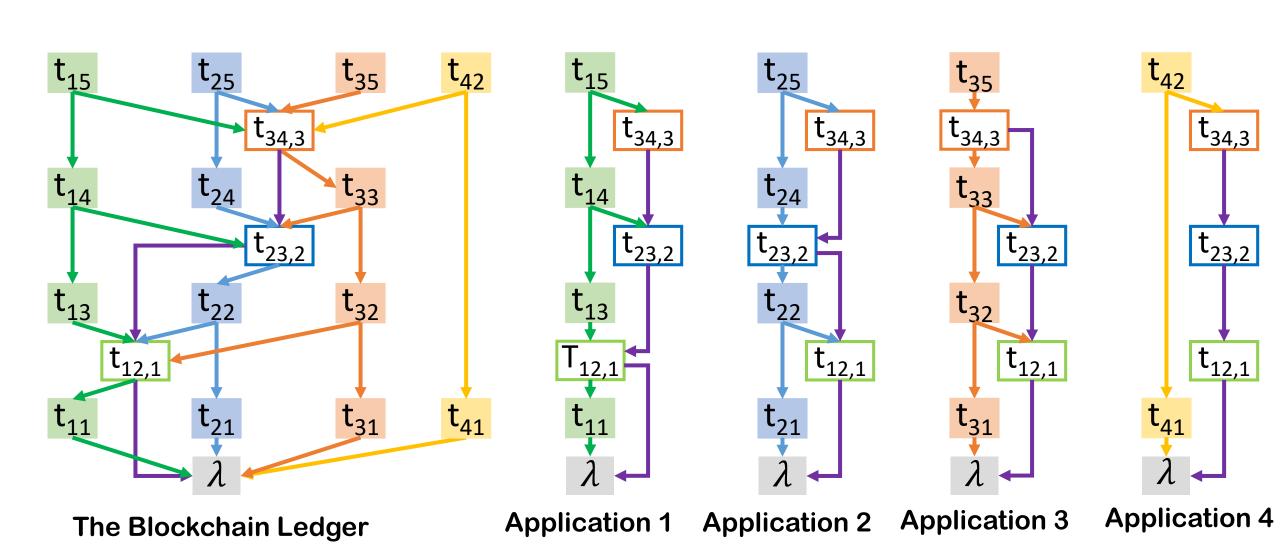












- In CAPER:
  - Internal transactions read both private and public data and write on private data
  - Cross-application transactions read/write only public data

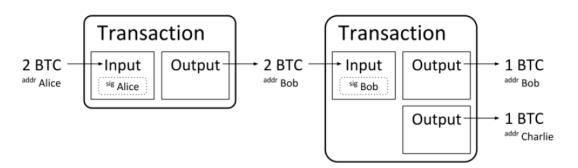
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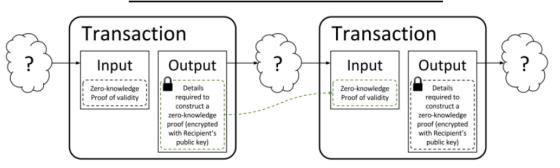
# Confidentiality of Cross-Application Transactions

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- How to validate private transactions without revealing any information?
- Cryptography techniques are needed!
  - Quorum uses zero knowledge proof
  - Fabric defines Private data collections

### Transaction verification in Bitcoin



### Transaction verification in Zcash



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### Case Study on Change Healthcare's use of Hyperledger Fabric

Change Healthcare turned to Hyperledger Fabric to begin blockchain-enabling its Intelligent Healthcare Network, which now processes 50 million transactions a day.

LEARN MORE IN THE BLOG

READ THE CASE STUDY

# Join Hyperledger as a Member

Hyperledger Member Summit is coming up July 30-31 in Tokyo, Japan. Now is a great time to consider joining Hyperledger as a member so you can attend this annual event to discuss the current and future state of Hyperledger technologies.

LEARN MORE

### Hyperledger Transact Now Available

Announcing our latest project to join the Hyperledger Greenhouse. Hyperledger Transact provides a platform-agnostic library that handles the execution of smart contracts, including all aspects of scheduling, transaction dispatch, and state management.

LEARN MORE IN THE BLOG

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Business Blockchain Frameworks & Tools Hosted by Hyperledger













Community Stewardship and Technical, Legal, Marketing, Organizational Infrastructure

### Frameworks













Permissionable smart contract machine (EVM)

Permissioned with channel support

WebAssembly-based project for building supply chain solutions

Decentralized identity

Mobile application focus

Permissioned & permissionless support; EVM transaction family

#### Tools



Infrastructure for

peer-to-peer

interactions



Blockchain framework benchmark platform



As-a-service deployment



Model and build blockchain networks



View and explore data on the blockchain



Ledger interoperability





Advanced transaction execution and state management

Shared Cryptographic Library



# From Cryptocurrencies to Global Asset Management

Victor Zakhary, Mohammad Amiri, Sujaya Maiyya, **Divyakant Agrawal**, **Amr El Abbadi** 

# From Cryptocurrencies to Global Assets

• So far, Mining Node:

- So far, Mining Node:
  - Store cryptocurrency units
  - Store ownership
  - Execute Transactions (transfer ownership of currency units)

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   The new public cloud
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- Transact on:
  - General Assets (e.g., buy a house, rent a car etc)



# **Smart Contracts**

• Alice registers her car





- Alice registers her car
  - Make: Honda
  - Model: Civic
  - Year: ..
  - VIN: ...





- Alice registers her car
  - Make: Honda
  - Model: Civic
  - Year: ..
  - VIN: ...
  - Owner: Alice
  - Price: x ethers

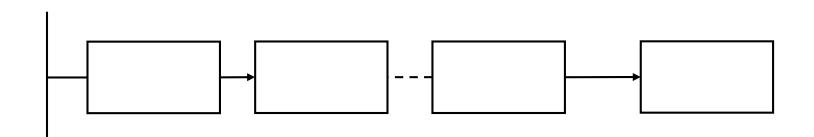


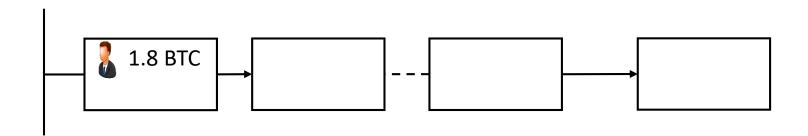


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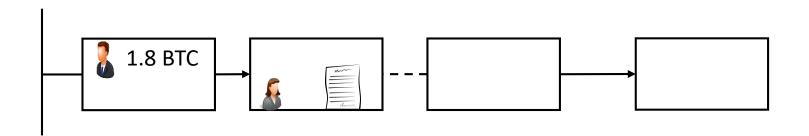
```
Buy () {
    // transfer ownership code
}
```



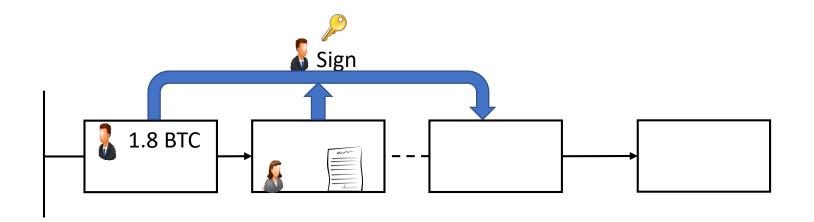




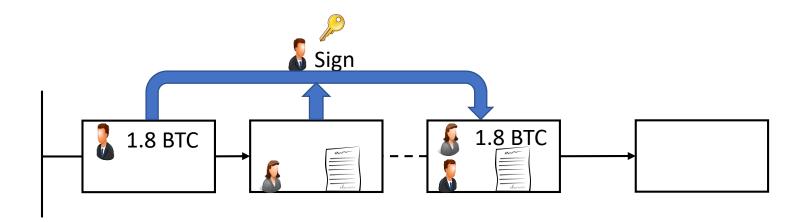
### **Smart Contracts**



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### **Smart Contracts**



## Challenges

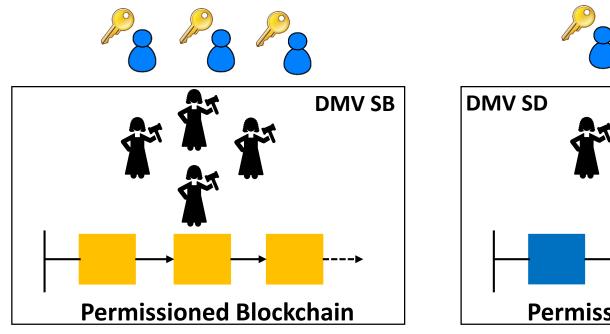
- Asset Authenticity
- Double Spending
  - Deploy two smart contracts for the same car
  - On the same blockchain or different blockchains
- Legality
  - Implementing taxation laws

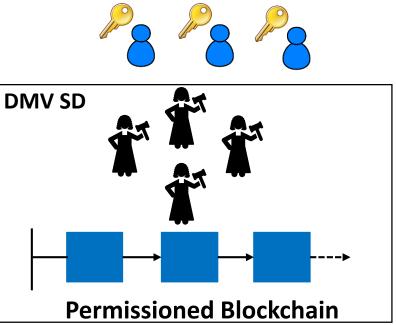
### Permissioned and Permissionless Unite!

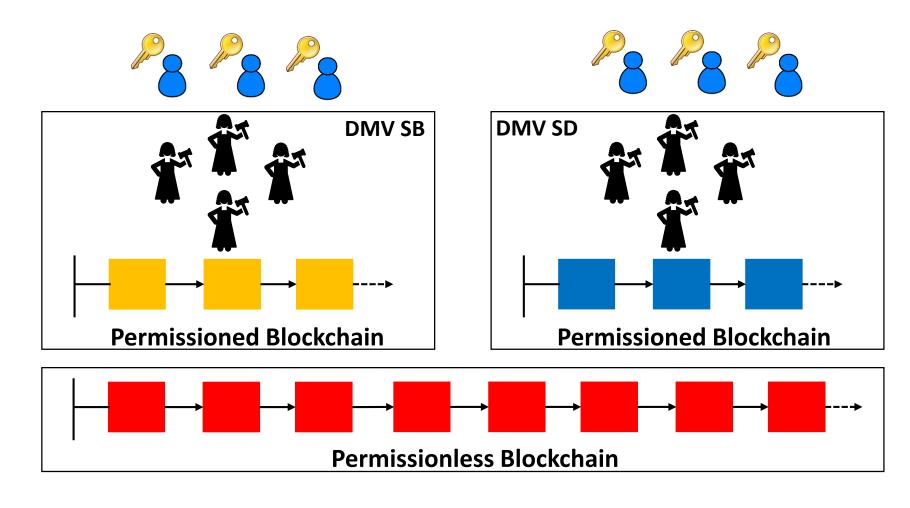
- Permissioned Blockchains
  - Requires trust
  - Trust can be distributed among several organizations
    - Banks
    - Governments
    - NGOs

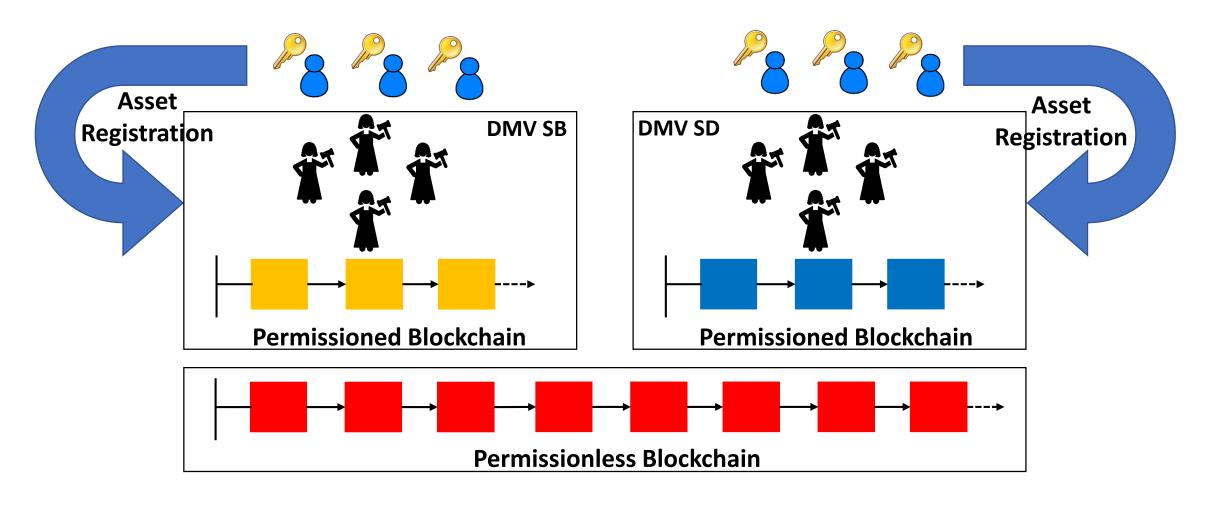


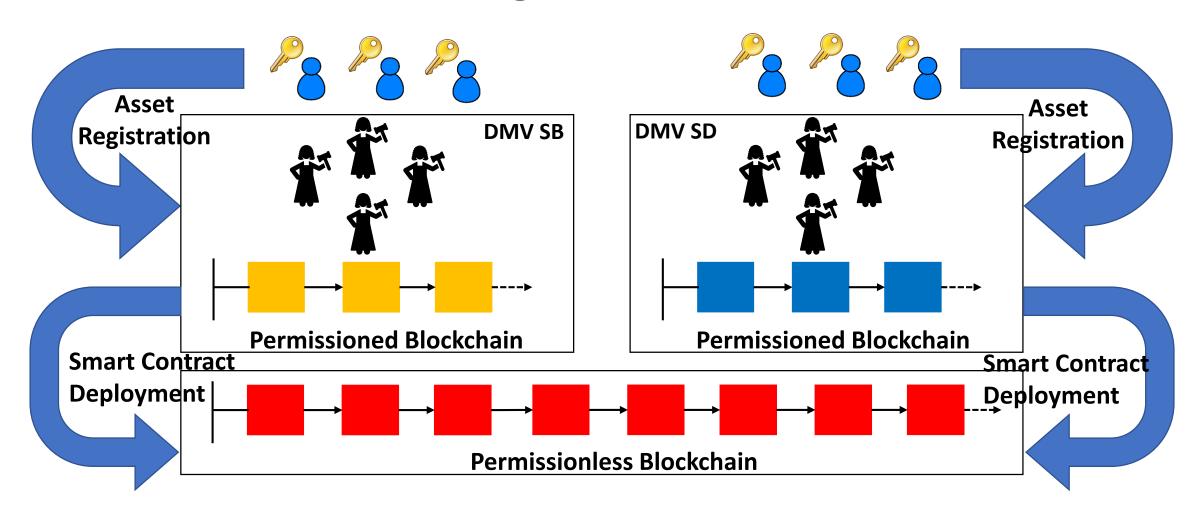


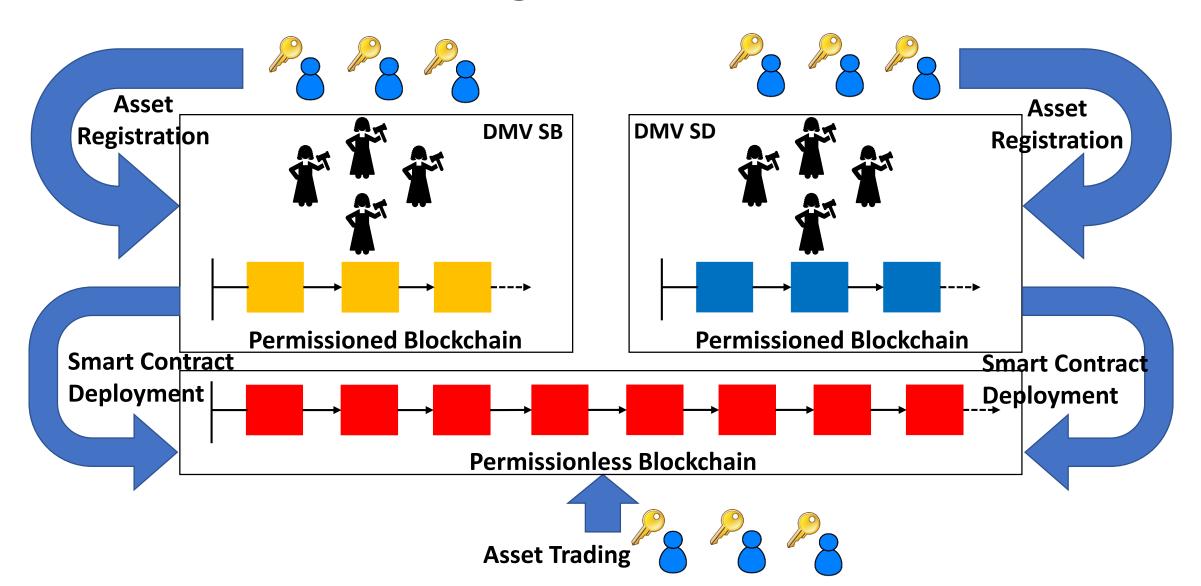












# Challenges Revisited

Asset Authenticity

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- Legality
  - Encode the Taxation law in the smart contract code

## Open research questions

- Scalability
- Identity theft
- Flexibility of asset marketing



# Blockchain: Panacea for all our data problems?

- Resource cost:
  - Proof-of-work consumes resources at the planetary scale
- Mythical notion of democratization:
  - Handful of miners control the progress of Bitcoin blockchain
- False notion of security:
  - An Individual vulnerable to the security of his/her key

- Extreme distribution:
  - is it really worth it?
- Extreme redundancy:
  - is it really necessary?
- Social consequences:
  - Are we comfortable if this technology is used for dark causes?