Lecture slide courtesy of Prof. Steven Skiena



CSE 373 Analysis of Algorithms Fall 2016 Instructor: Prof. Sael Lee

LEC15: HEURISTIC SEARCH (247-)

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- Backtracking gave us a method to find the best of all possible solutions, as scored by a given objective function. However, any algorithm searching all configurations is doomed to be impossible on large instances.
- * Heuristic methods provide an alternate way to approach difficult combinatorial optimization problems.
- **We look at three different heuristic search methods:**
 - + Random sampling,
 - + gradient-descent search, and
 - + simulated annealing.

TWO COMMON COMPONENTS

- Solution space representation: complete yet concise description of the set of possible solutions for the problem.
 - + *TSP:* solution space consists of (n-1)! elements all possible circular permutations of the vertices, represented using an array S of n 1 vertices.
- Cost function: cost or evaluation function to access the quality of each element of the solution space.
 - + TSP: given candidate solution *S, cost function* is just the sum of the costs involved.

RANDOM SAMPLING AKA MONTE CARLO METHOD

- Repeatedly construct random solutions and evaluate t hem, stopping as soon as we get a good enough soluti on, or (more likely) when we are tired of waiting
- Report the best solution found over the course of our sampling.

WHEN MIGHT RANDOM SAMPLING DO WELL?

- When there are a high proportion of acceptable soluti ons
 - + *Ex*> Finding prime numbers
- When there is no coherence in the solution space th ere is no sense of when we are getting closer to a solu tion
 - + Ex>Finding who of your friends has a social security number that ends in 00.

LOCAL SEARCH

- A local search heuristic starts from an arbitrary element of the solution space, and then scans the neighborhood looking for a favorable transition to take.
 - + EX> TSP, this would be *transition*, which lowers the cost of the tour.
- * We want a general transition mechanism that takes us to the next solution by slightly modifying the current one.
 - + Ex> TSP: swap the current tour positions of a random pair of vertices Si and Sj,



}: Improving a TSP tour by swapping vertices 2 and 6

WHEN DOES LOCAL SEARCH DO WELL?

- × When there is great coherence in the solution space
 - + Cost function is either convex or concave
- Whenever the cost of incremental evaluation is much cheaper than global evaluation
 - + EX> TSP: It costs Θ(n) to evaluate the cost of an arbitrary n-vertex candidate TSP solution. Once that is found, however, the cost of the tour after swapping a given pair of vertices can be determined in constant time.

SIMULATED ANNEALING

- Simulated annealing is a heuristic search procedure that allows <u>occasional transitions leading to more</u> <u>expensive (and hence inferior) solutions</u> – which helps keep our search from getting stuck in local optima
- Simulated annealing comes from the physical process of cooling materials down to the solid state.
- * A particle's energy state jumps about randomly, with such transitions governed by the temperature of the system where Transition probability *P*(*ei, ej, T*) from energy *ei* to *ej* at temperature *T*:

$$P(e_i, e_j, T) = e^{\frac{e_i - e_j}{k_b T}}$$

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\begin{array}{l} \mbox{Simulated-Annealing()} \\ \mbox{Create initial solution $S$} \\ \mbox{Initialize temperature $t$} \\ \mbox{repeat} \\ \mbox{for $i=1$ to iteration-length do} \\ \mbox{Generate a random transition from $S$ to $S_i$} \\ \mbox{If $(C(S) \geq C(S_i))$ then $S=S_i$} \\ \mbox{else if $(e^{(C(S)-C(S_i))/(k\cdot t)} > random[0,1)$)$ then $S=S_i$} \\ \mbox{Reduce temperature $t$} \\ \mbox{until $(no change in $C(S)$)} \\ \mbox{Return $S$} \end{array}
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