Lecture 9
(Parallel Connected Components)

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Symmetry Breaking: List Ranking

1. Flip a coin for each list node
2. If a node \( u \) points to a node \( v \), and \( u \) got a head while \( v \) got a tail, combine \( u \) and \( v \)
3. Recursively solve the problem on the contracted list
4. Project this solution back to the original list
Symmetry Breaking: List Ranking

In every iteration a node gets removed with probability $\frac{1}{4}$

( as a node gets head with probability $\frac{1}{2}$ and the next node gets tail
with probability $\frac{1}{2}$).

Hence, a quarter of the nodes get removed in each iteration
( expected number ).

Thus the expected number of iterations is $\Theta(\log n)$.

In fact, it can be shown that with high probability,

$$T_1(n) = O(n) \text{ and } T_\infty(n) = O(\log n)$$
**Graph Connectivity**

Connected Components: A *connected component* $C$ of an undirected graph $G$ is a maximal subgraph of $G$ such that every vertex in $C$ is reachable from every other vertex in $C$ following a path in $G$.

Problem: Given an undirected graph identify all its connected components.

Suppose $n$ is the number of vertices in the graph, and $m$ is the number of edges.
**Problem:** Identify All connected components of an undirected graph.

Suppose \( n \) is the number of vertices in the graph, and \( m \) is the number of edges.

**Serial Algorithms:** Easy to solve in \( \Theta(m + n) \) time using

- Depth First Search (DFS)
- Breadth First Search (BFS)

**Parallel Algorithms:**

- **DFS:** Inherently sequential
- **BFS:** Depth equal to the diameter of the graph
- **Graph Contraction:** Can reach polylogarithmic depth
Randomized Parallel Connected Components
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Randomized Parallel Connected Components (CC)

**Input:** $n$ is the number of vertices in the graph numbered from 1 to $n$, $E$ is the set of edges, and $L[1:n]$ are vertex labels with $L[v] = v$ initially for all $v$.

**Output:** An array $M[1:n]$ where for all $v$, $M[v]$ is the unique id of the connected component containing $v$.

$$\text{Par-Randomized-CC ( n, E, L )}$$

1. if $|E| = 0$ then return $L$
2. array $C[1:n], M[1:n], S[1:|E|]$
3. parallel for $v \leftarrow 1$ to $n$ do $C[v] \leftarrow \text{RANDOM\{Head, Tail\}}$
4. parallel for each $(u, v) \in E$ do
   5. if $C[u] = \text{Tail}$ and $C[v] = \text{Head}$ then $L[u] \leftarrow L[v]$
6. parallel for $i \leftarrow 1$ to $|E|$ do
   7. if $L[E[i].u] \neq L[E[i].v]$ then $S[i] \leftarrow 1$ else $S[i] \leftarrow 0$
8. $S \leftarrow \text{Par-Prefix-Sum ( S, + )}$
9. array $F[1:S[|E|]]$
10. parallel for $i \leftarrow 1$ to $|E|$ do
11. if $L[E[i].u] \neq L[E[i].v]$ then
    $F[S[i]] \leftarrow (L[E[i].u], L[E[i].v])$
12. $M \leftarrow \text{Par-Randomized-CC ( n, F, L )}$
13. parallel for each $(u, v) \in E$ do
14. if $v = L[u]$ then $M[u] \leftarrow M[v]$
15. return $M$

unbiased coin toss at each vertex

prepare to remove intra-group edges

Map results back to the original graph

find the rank of each inter-group edge among all such edges

find CC in the contracted graph

copy the inter-group edges to $F$
Randomized Parallel Connected Components (CC)

Par-Randomized-CC (n, E, L)
1. if |E| = 0 then return L
2. array C[1 : n], M[1 : n], S[1 : |E|]
3. parallel for v ← 1 to n do
   C[v] ← RANDOM{Head, Tail}
4. parallel for each (u, v) ∈ E do
5.   if C[u] = Tail and C[v] = Head then L[u] ← L[v]
6.   parallel for i ← 1 to |E| do
7.     if L[E[i].u] ≠ L[E[i].v] then S[i] ← 1
     else S[i] ← 0
8.   S ← Par-Prefix-Sum(S, +)
9.   array F[1 : S[|E|]]
10. parallel for i ← 1 to |E| do
11.     if L[E[i].u] ≠ L[E[i].v] then
          F[S[i]] ← (L[E[i].u], L[E[i].v])
12.  M ← Par-Randomized-CC(n, F, L)
13. parallel for each (u, v) ∈ E do
14.     if v = L[u] then M[u] ← M[v]
15. return M

Suppose \( n \) is the number of vertices and \( m \) is the number of edges in the original graph.

Each contraction is expected to reduce #vertices of +ve degree by a factor \( \geq \frac{1}{4} \). [why?]

So, the expected number of contraction steps, \( D = O(\log n) \). [show: the bound holds w.h.p.]

For each contraction step span is \( \Theta(\log^2 n) \), and work is \( \Theta(n + m) \). [why?]

Work: \( T_1(n, m) = \Theta(D(n + m)) = O((n + m)\log n) \) (w.h.p.)

Span: \( T_\infty(n, m) = \Theta(D\log^2 n) = O(\log^3 n) \) (w.h.p.)

Parallelism: \( \frac{T_1(n, m)}{T_\infty(n, m)} = \Theta\left(\frac{n+m}{\log^2 n}\right) \)
The pointer jumping (or path doubling) technique allows fast processing of data stored in the form of a set of rooted directed trees.

For every node $v$ in the set pointer jumping involves replacing $v \rightarrow next$ with $v \rightarrow next \rightarrow next$ at every step.

Some Applications

- Finding the roots of a forest of directed trees
- Parallel prefix on rooted directed trees
- List ranking
Find-Roots (n, P, S)  

{ Input: A forest of rooted directed trees, each with a self-loop at its root, such that each edge is specified by (v, P(v)) for 1 ≤ v ≤ n.  
Output: For each v, the root S(v) of the tree containing v. }

1. parallel for v ← 1 to n do
2. S(v) ← P(v)
3. flag ← true
4. while flag = true do
5. flag ← false
6. parallel for v ← 1 to n do
7. S(v) ← S(S(v))
8. if S(v) ≠ S(S(v)) then flag ← true
Let $h$ be the maximum height of any tree in the forest.
Observe that the distance between $v$ and $S(v)$ doubles after each iteration until $S(S(v))$ is the root of the tree containing $v$.
Hence, the number of iterations is $\log h$. Thus (assuming that each parallel for loop takes $\Theta(1)$ time to execute),

**Work:** $T_1(n) = \Theta(n \log h)$ and **Span:** $T_\infty(n) = \Theta(\log h)$

**Parallelism:** $\frac{T_1(n)}{T_\infty(n)} = O(n)$
Deterministic Parallel Connected Components (CC)

Approach

- Form a set of disjoint subtrees
- Use pointer-jumping to reduce each subtree to a single vertex
- Recursively apply the same trick on the contracted graph

Forming Disjoint Subtrees

- Hook each vertex to a neighbor with larger label (if exists)
- Ensures that no cycles are formed
Deterministic Parallel Connected Components (CC)

Forming Disjoint Subtrees

- Hook each vertex to a neighbor with larger label (if exists)
- Ensures that no cycles are formed

But the number of contraction steps can be as large as \( n - 1 \)!
Deterministic Parallel Connected Components (CC)

Observation:
Let $G = (V, E)$ be an undirected graph with $n$ vertices in which each vertex has at least one neighbor. Then

either $|\{(u, v) \in E \land (u < v)\}| \geq \frac{n}{2}$

or $|\{(u, v) \in E \land (u > v)\}| \geq \frac{n}{2}$

Implication:
Between the two directions of hooking (i.e., smaller to larger label, and larger to smaller label) always choose the one that hooks more vertices.

Then in each contraction step the number of vertices will be reduced by a factor of at least $\frac{1}{2}$. 
Deterministic Parallel Connected Components (CC)

**Input:** $n$ is the number of vertices in the graph numbered from 1 to $n$, $E$ is the set of edges, and $L[1 : n]$ are vertex labels with $L[v] = v$ initially for all $v$.

**Output:** Updated array $L[1: n]$ where for all $v$, $L[v]$ is the unique id of the connected component containing $v$.

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Par-Deterministic-CC (n, E, L)
1. if $|E| = 0$ then return $L$
3. parallel for $v ← 1$ to $n$ do $l2h[v] ← 0, h2l[v] ← 0$
4. parallel for each $(u, v) ∈ E$ do
   if $u < v$ then $l2h[u] ← 1$ else $h2l[u] ← 1$
5. $n_1 ← Par-Sum (l2h, +), n_2 ← Par-Sum (h2l, +)$
6. parallel for each $(u, v) ∈ E$ do
   if $n_1 ≥ n_2$ and $u < v$ then $L[u] ← v$
   else if $n_1 < n_2$ and $u > v$ then $L[u] ← v$
7. Find-Roots (n, L, L)
8. parallel for $i ← 1$ to $|E|$ do $S[i] ← (L[E[i].u] ≠ L[E[i].v]) ? 1 : 0$
9. $S ← Par-Prefix-Sum (S, +)$
10. array $F[1 : S[|E|]]$
11. parallel for $i ← 1$ to $|E|$ do
    if $L[E[i].u] ≠ L[E[i].v]$ then $F[S[i]] ← (L[E[i].u], L[E[i].v])$
12. $L ← Par-Deterministic-CC (n, F, L)$
13. return $L$
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Deterministic Parallel Connected Components (CC)

\textit{Par-Deterministic-CC (n, E, L)}

1. \textit{if} \(|E| = 0\) \textit{then return} \(L\)
2. \textit{array} \(l2h[1 : n], h2l[1 : n], S[1 : |E|]\)
3. \textit{parallel for} \(v \leftarrow 1\) \textit{to} \(n\) \textit{do}
   \hspace{1em} \(l2h[v] \leftarrow 0, h2l[v] \leftarrow 0\)
4. \textit{parallel for each} \((u, v) \in E\) \textit{do}
5. \hspace{1em} \textit{if} \(u < v\) \textit{then} \(l2h[u] \leftarrow 1\) \textit{else} \(h2l[u] \leftarrow 1\)
6. \(n_1 \leftarrow \textit{Par-Sum (l2h, +)}, n_2 \leftarrow \textit{Par-Sum (h2l, +)}\)
7. \textit{parallel for each} \((u, v) \in E\) \textit{do}
8. \hspace{1em} \textit{if} \(n_1 \geq n_2\) \textit{and} \(u < v\) \textit{then} \(L[u] \leftarrow v\)
9. \hspace{1em} \textit{else if} \(n_1 < n_2\) \textit{and} \(u > v\) \textit{then} \(L[u] \leftarrow v\)
10. \textit{Find-Roots (n, L, L)}
11. \textit{parallel for} \(i \leftarrow 1\) \textit{to} \(|E|\) \textit{do}
   \hspace{1em} \(S[i] \leftarrow (L[E[i].u] \neq L[E[i].v]) ? 1 : 0\)
12. \(S \leftarrow \textit{Par-Prefix-Sum (S, +)}\)
13. \textit{array} \(F[1 : S[|E|]]\)
14. \textit{parallel for} \(i \leftarrow 1\) \textit{to} \(|E|\) \textit{do}
15. \hspace{1em} \textit{if} \(L[E[i].u] \neq L[E[i].v]\) \textit{then}
   \hspace{2em} \(F[S[i]] \leftarrow (L[E[i].u], L[E[i].v])\)
16. \(L \leftarrow \textit{Par-Deterministic-CC (n, F, L)}\)
17. \textit{return} \(L\)

Each contraction step reduces the number of vertices by a factor of at least \(\frac{1}{2}\).

So, number of contraction steps, \(D = O(\log n)\).

For contraction step \(k \geq 0\) span is \(O(\log^2 n)\),
and work is \(O(n \log n + m)\). [why?]

\textbf{Work:} \(T_1(n, m) = O(\sum_{0 \leq i < D} (n \log n + m))\)
\hspace{2em} \(= O((n \log n + m)D)\)
\hspace{2em} \(= O((n \log n + m) \log n)\)

\textbf{Span:} \(T_\infty(n, m) = O(D \log^2 n)\)
\hspace{2em} \(= O(\log^3 n)\)

How to get \(T_1(n, m) = O((n + m) \log n)\)?