CSE 548: Analysis of Algorithms

Lectures 7 & 8
(Divide-and-Conquer Algorithms: Akra-Bazzi Recurrences)

Rezaul A. Chowdhury

Department of Computer Science
SUNY Stony Brook
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Deterministic Select

Input: An array $A[\,q : r\,]$ of distinct elements, and integer $k \in [1, r-q+1]$.

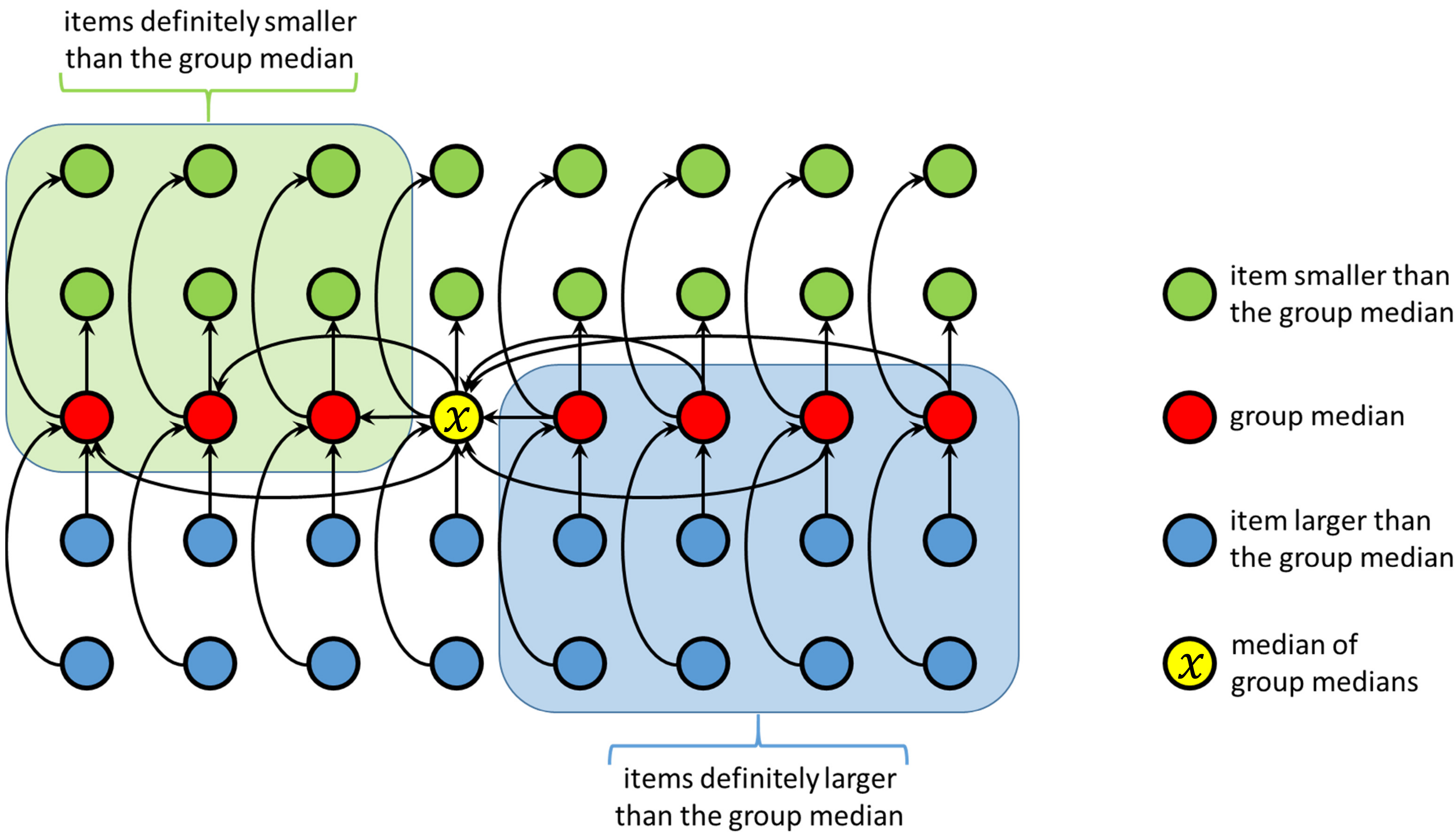
Output: An element x of $A[\,q : r\,]$ such that $rank(x, A[\,q : r\,]) = k$.

Select ($A[\,q : r\,]$, k)

1. $n \leftarrow r - q + 1$
2. *if* $n \leq 140$ *then*
3. sort $A[\,q : r\,]$ and *return* $A[\,q + k - 1\,]$
4. *else*
5. divide $A[\,q : r\,]$ into blocks B_i's each containing 5 consecutive elements (last block may contain fewer than 5 elements)
6. *for* $i \leftarrow 1$ *to* $\lceil n / 5 \rceil$ *do*
7. $M[\,i\,] \leftarrow$ median of B_i using sorting
8. $x \leftarrow$ *Select* ($M[\,1 : \lceil n / 5 \rceil\,], \lfloor (\lceil n / 5 \rceil + 1) / 2 \rfloor$) { median of medians }
9. $t \leftarrow$ *Partition* ($A[\,q : r\,], x$) { partition around x which ends up at $A[\,t\,]$ }
10. *if* $k = t - q + 1$ *then return* $A[\,t\,]$
11. *else if* $k < t - q + 1$ *then return* *Select* ($A[\,q : t - 1\,], k$)
12. *else return* *Select* ($A[\,t + 1 : r\,], k - t + q - 1$)

Deterministic Select

$\text{SELECT}(A, k)$: Given an unsorted set A of n $(= |A|)$ items, find the k^{th} smallest item in the set



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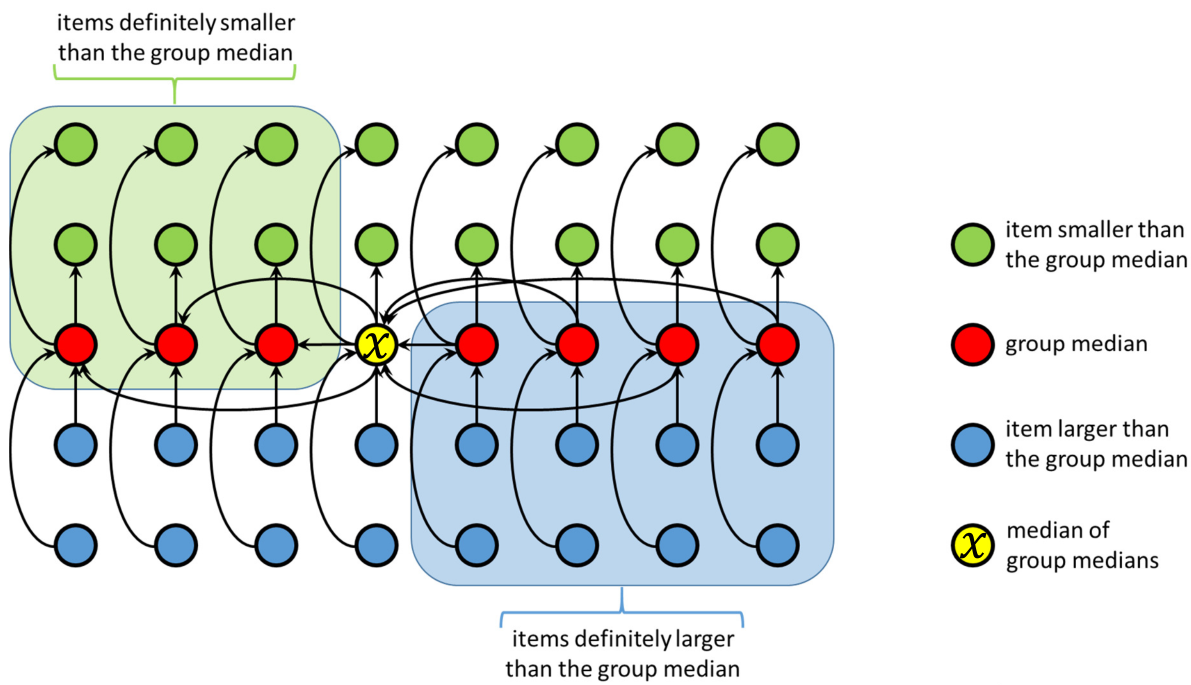

#items definitely smaller than x is $\geq 3\left(\left\lfloor \frac{1}{2}\left\lceil \frac{n}{5} \right\rceil \right\rfloor - 1\right) \geq \frac{3n}{10} - 6$

#items definitely larger than x is $\geq 3\left(\left\lfloor \frac{1}{2}\left\lceil \frac{n}{5} \right\rceil \right\rfloor - 1\right) \geq \frac{3n}{10} - 6$

#items in any recursive call (lines 11/12) $\leq n - \left(\frac{3n}{10} - 6\right) = \frac{7n}{10} + 6$

Deterministic Select

The following recurrence describes the worst-case running time of the deterministic selection algorithm (given in Section 9.3 of CLRS):

$$T(n) \le \begin{cases} \Theta(1), & if\ n < 140, \\ T\left(\left\lceil \frac{n}{5} \right\rceil\right) + T\left(\frac{7n}{10} + 6\right) + \Theta(n), & if\ n \ge 140. \end{cases}$$

Dropping the ceiling for simplicity, and observing that $\frac{7n}{10} + 6 \le \frac{8n}{10}$ when $n \ge 60$, we obtain the following upper bound on $T(n)$.

$$T'(n) = \begin{cases} \Theta(1), & if\ n < 140, \\ T'\left(\frac{n}{5}\right) + T'\left(\frac{4n}{5}\right) + \Theta(n), & if\ n \ge 140. \end{cases}$$

How do you solve for $T'(n)$?

Deterministic Select

The following recurrence describes the worst-case running time of the deterministic selection algorithm (given in Section 9.3 of CLRS):

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Dropping the ceiling for simplicity, and observing that $\frac{7n}{10} + 6 \le \frac{7.5n}{10}$ when $n \ge 120$, we obtain the following upper bound on $T(n)$.

$$T''(n) = \begin{cases} \Theta(1), & if\ n < 140, \\ T''\left(\frac{n}{5}\right) + T''\left(\frac{3n}{4}\right) + \Theta(n), & if\ n \ge 140. \end{cases}$$

How do you solve for $T''(n)$?

Akra-Bazzi Recurrences

Consider the following recurrence:

$$T(x) = \begin{cases} \Theta(1), & if\ 1 \leq x \leq x_0, \\ \sum_{i=1}^{k} a_i T(b_i x) + g(x), & if\ x > x_0; \end{cases}$$

where,

1. $k \geq 1$ is an integer constant
2. $a_i > 0$ is a constant for $1 \leq i \leq k$
3. $b_i \in (0,1)$ is a constant for $1 \leq i \leq k$
4. $x \geq 1$ is a real number
5. $x_0 \geq \max\left\{\frac{1}{b_i}, \frac{1}{1-b_i}\right\}$ is a constant for $1 \leq i \leq k$
6. $g(x)$ is a nonnegative function that satisfies a *polynomial-growth condition* (to be specified soon)

Polynomial-Growth Condition

We say that $g(x)$ satisfies the *polynomial-growth condition* if there exist positive constants c_1 and c_2 such that for all $x \geq 1$, for all $1 \leq i \leq k$, and for all $u \in [b_i x, x]$,

$$c_1 g(x) \leq g(u) \leq c_2 g(x),$$

where x, k, b_i and $g(x)$ are as defined in the previous slide.

The Akra-Bazzi Solution

Consider the recurrence given in the previous two slides under the conditions specified there:

$$T(x) = \begin{cases} \Theta(1), & if\ 1 \le x \le x_0, \\ \sum_{i=1}^{k} a_i T(b_i x) + g(x), & if\ x > x_0. \end{cases}$$

Let p be the unique real number for which $\sum_{i=1}^{k} a_i b_i^p = 1$. Then

$$T(x) = \Theta\left(x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right)\right)$$

Deterministic Select

$$T'(n) = \begin{cases} \Theta(1), & if\ n < 140, \\ T'\left(\frac{n}{5}\right) + T'\left(\frac{4n}{5}\right) + \Theta(n), & if\ n \geq 140. \end{cases}$$

From $\left(\frac{1}{5}\right)^p + \left(\frac{4}{5}\right)^p = 1$ we get $p = 1$.

Hence, $T'(n) = \Theta\left(n^p\left(1 + \int_1^n \frac{u}{u^{p+1}}\,du\right)\right)$

$$= \Theta\left(n\left(1 + \int_1^n \frac{du}{u}\right)\right)$$

$$= \Theta(n \ln n)$$

Deterministic Select

$$T''(n) = \begin{cases} \Theta(1), & if\ n < 140, \\ T''\left(\frac{n}{5}\right) + T''\left(\frac{3n}{4}\right) + \Theta(n), & if\ n \geq 140. \end{cases}$$

From $\left(\frac{1}{5}\right)^p + \left(\frac{3}{4}\right)^p = 1$ we get $p < 1$.

Hence, $T''(n) = \Theta\left(n^p\left(1 + \int_1^n \frac{u}{u^{p+1}}\,du\right)\right)$

$$= \Theta\left(n^p\left(1 + \int_1^n \frac{du}{u^p}\right)\right)$$

$$= \Theta\left(\left(\frac{1}{1-p}\right)n - \left(\frac{p}{1-p}\right)n^p\right)$$

$$= \Theta(n)$$

Examples of Akra-Bazzi Recurrences

Example 1: $T(x) = 2T\left(\frac{x}{4}\right) + 3T\left(\frac{x}{6}\right) + \Theta(x \log x)$

Then $p = 1$ and $T(x) = \Theta\left(x\left(1 + \int_1^x \frac{u \log u}{u^2} du\right)\right) = \Theta(x \log^2 x)$

Example 2: $T(x) = 2T\left(\frac{x}{2}\right) + \frac{8}{9}T\left(\frac{3x}{4}\right) + \Theta\left(\frac{x^2}{\log x}\right)$

Then $p = 2$ and $T(x) = \Theta\left(x^2\left(1 + \int_1^x \frac{u^2/\log u}{u^3} du\right)\right) = \Theta(x^2 \log\log x)$

Example 3: $T(x) = T\left(\frac{x}{2}\right) + \Theta(\log x)$

Then $p = 0$ and $T(x) = \Theta\left(1 + \int_1^x \frac{\log u}{u} du\right) = \Theta(\log^2 x)$

Example 4: $T(x) = \frac{1}{2}T\left(\frac{x}{2}\right) + \Theta\left(\frac{1}{x}\right)$

Then $p = -1$ and $T(x) = \Theta\left(\frac{1}{x}\left(1 + \int_1^x \frac{1}{u} du\right)\right) = \Theta\left(\frac{\log x}{x}\right)$

A Helping Lemma

Lemma: If $g(x)$ is a nonnegative function that satisfies the polynomial-growth condition, then there exist positive constants c_3 and c_4 such that for $1 \le i \le k$ and all $x \ge 1$,

$$c_3 g(x) \le x^p \int_{b_i x}^{x} \frac{g(u)}{u^{p+1}} du \le c_4 g(x).$$

Proof:

$$b_i x \le u \le x$$

$$\Rightarrow \frac{1}{\max\{(b_i x)^{p+1}, x^{p+1}\}} \le \frac{1}{u^{p+1}} \le \frac{1}{\min\{(b_i x)^{p+1}, x^{p+1}\}}$$

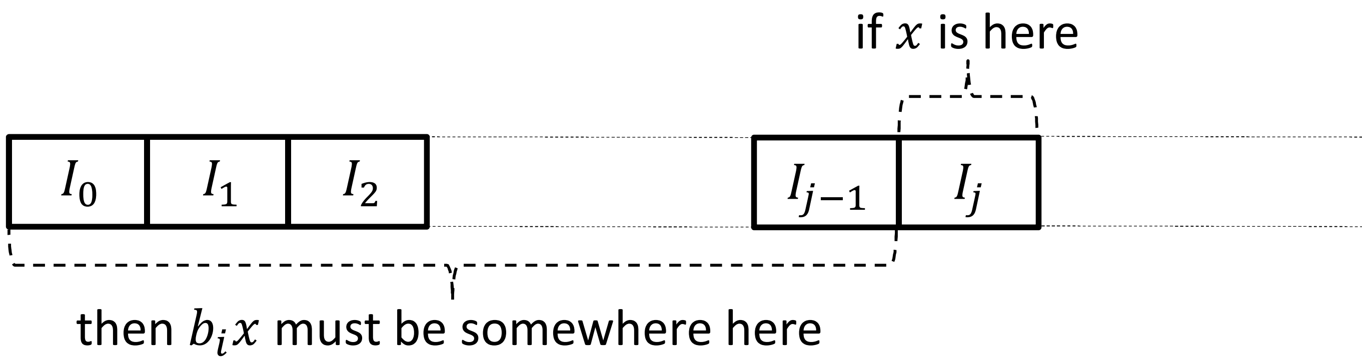
$$\Rightarrow \frac{x^p c_1 g(x)}{\max\{(b_i x)^{p+1}, x^{p+1}\}} \int_{b_i x}^{x} du \le x^p \int_{b_i x}^{x} \frac{g(u)}{u^{p+1}} du \le \frac{x^p c_2 g(x)}{\min\{(b_i x)^{p+1}, x^{p+1}\}} \int_{b_i x}^{x} du$$

$$\Rightarrow \frac{(1-b_i)c_1}{\max\{1, b_i^{p+1}\}} g(x) \le x^p \int_{b_i x}^{x} \frac{g(u)}{u^{p+1}} du \le \frac{(1-b_i)c_2}{\min\{1, b_i^{p+1}\}} g(x)$$

$$\Rightarrow c_3 g(x) \le x^p \int_{b_i x}^{x} \frac{g(u)}{u^{p+1}} du \le c_4 g(x)$$

Partitioning the Domain of x

Let $I_0 = [1, x_0]$ and $I_j = [x_0 + j - 1, x_0 + j]$ for $j \geq 1$.

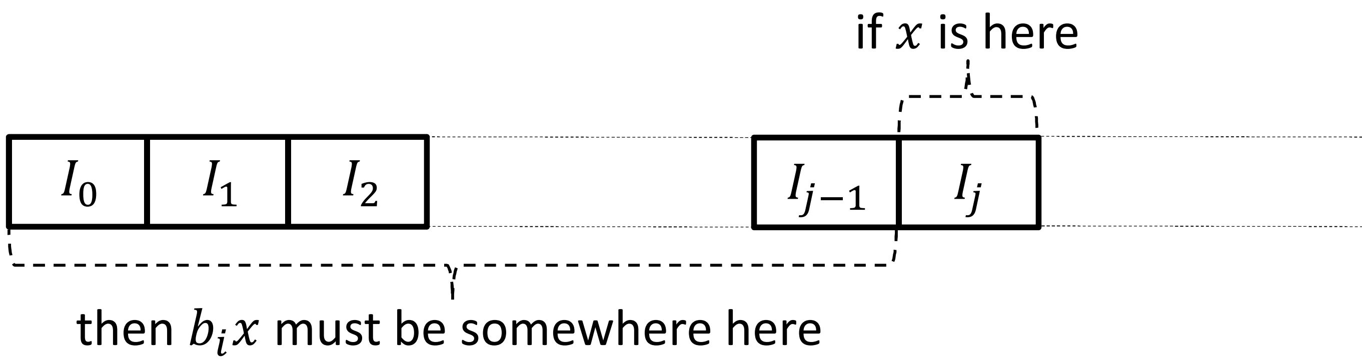


That allows us to use induction in the proof of:

$$T(x) = \begin{cases} \Theta(1), & if\ 1 \leq x \leq x_0, \\ \sum_{i=1}^{k} a_i T(b_i x) + g(x), & if\ x > x_0. \end{cases}$$

Partitioning the Domain of x

Let $I_0 = [1, x_0]$ and $I_j = [x_0 + j - 1, x_0 + j]$ for $j \geq 1$.



Proof:

$$x_0 + j - 1 < x \leq x_0 + j$$

$$\Rightarrow b_i(x_0 + j - 1) < b_i x \leq b_i(x_0 + j)$$

$$\Rightarrow b_i x_0 < b_i x \leq b_i x_0 + j$$

$$\Rightarrow 1 < b_i x \leq x_0 + j - (1 - b_i) x_0$$

$$\Rightarrow 1 < b_i x \leq x_0 + j - 1$$

Derivation of the Akra-Bazzi Solution

Lower Bound: There exists a constant $c_5 > 0$ such that for all $x > x_0$,

$$T(x) \geq c_5 x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right).$$

Proof: By induction on the interval I_j containing x.

Base case ($j = 0$) follows since $T(x) = \Theta(1)$ when $x \in I_0 = [1, x_0]$.

Induction: $$T(x) = \sum_{i=1}^{k} a_i T(b_i x) + g(x) \geq \sum_{i=1}^{k} a_i c_5 (b_i x)^p \left(1 + \int_1^{b_i x} \frac{g(u)}{u^{p+1}} du\right) + g(x)$$

$$= c_5 x^p \sum_{i=1}^{k} a_i b_i^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du - \int_{b_i x}^x \frac{g(u)}{u^{p+1}} du\right) + g(x)$$

$$\geq c_5 x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du - \frac{c_4}{x^p} g(x)\right) \sum_{i=1}^{k} a_i b_i^p + g(x)$$

$$= c_5 x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right) + (1 - c_4 c_5) g(x) \geq c_5 x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right)$$

(assuming $c_4 c_5 \leq 1$)

Derivation of the Akra-Bazzi Solution

Upper Bound: There exists a constant $c_6 > 0$ such that for all $x > x_0$,

$$T(x) \leq c_6 x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right).$$

Proof: Similar to the lower bound proof.