XML and Web Data

CSE 532, Theory of Database Systems
Stony Brook University

http://www.cs.stonybrook.edu/~cse532
XML

- Semistructured data
- XML & DTD
- XML Schema – user-defined data types, integrity constraints
- XPath & XPointer – core query language for XML
- XSLT – document transformation language
- XQuery – full-featured query language for XML
- SQL/XML – XML extensions of SQL
Why XML?

- XML is a standard for data exchange that is taking over the World
- All major database products have been retrofitted with facilities to store and construct XML documents
- There are already database products that are specifically designed to work with XML documents rather than relational or object-oriented data
- XML is closely related to object-oriented and so-called *semistructured* data
Semistructured Data

- An HTML document to be displayed on the Web:

```html
<dt>Name: John Doe
    <dd>Id: s111111111
    <dd>Address:  
        <ul>
            <li>Number: 123</li>
            <li>Street: Main</li>
        </ul>
    </dd>
</dt>
<dt>Name: Joe Public
    <dd>Id: s222222222
        ...
        ...
</dt>
```

*HTML does not distinguish between attributes and values*
Semistructured Data (cont’d.)

- To make the previous student list suitable for machine consumption on the Web, it should have these characteristics:
  - Be *object-like*
  - Be *schemaless* (not guaranteed to conform exactly to any schema, but different objects have some commonality among themselves)
  - Be *self-describing* (some schema-like information, like attribute names, is part of data itself)

- Data with these characteristics are referred to as *semistructured.*
What is Self-describing Data?

- Non-self-describing (relational, object-oriented):

  **Data part:**

  ```
  (#123, [“Students”, [{“John”, s111111111, [123,”Main St”]},
  [“Joe”, s222222222, [321, “Pine St”]}] )
  ```

  **Schema part:**

  ```
  PersonList[ ListName: String,
  Contents: [ Name: String,
  Id: String,
  Address: [Number: Integer, Street: String] ] ]
  ```

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What is Self-Describing Data?

- **Self-describing:**
  - Attribute names embedded in the data itself, *but are distinguished* from values
  - Doesn’t need schema to figure out what is what (but schema might be useful nonetheless)

```plaintext
(#12345,
  [ListName: “Students”,
   Contents: { [ Name: “John Doe”,
                Id: “s111111111”,
                Address: [Number: 123, Street: “Main St.”] ],
   [Name: “Joe Public”,
    Id: “s222222222”,
    Address: [Number: 321, Street: “Pine St.”] } ]
)
```

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XML – The De Facto Standard for Semistructured Data

• **XML:** eXtensible Markup Language
  • Suitable for semistructured data and has become a standard:
    • Easy to describe object-like data
    • Self-describing
    • Doesn’t require a schema (but can be provided optionally)

• **We will study:**
  • DTDs – an older way to specify schema
  • XML Schema – a newer, more powerful (and much more complex!) way of specifying schema
  • Query and transformation languages:
    • XPath
    • XSLT
    • XQuery
    • SQL/XML
Overview of XML

- Like HTML, but any number of different tags can be used (up to the document author) – extensible
- Unlike HTML, no semantics behind the tags
  - For instance, HTML’s `<table>…</table>` means: render contents as a table; in XML: doesn’t mean anything special
  - Some semantics can be specified using XML Schema (types); some using stylesheets (browser rendering)
- Unlike HTML, is intolerant to bugs
  - Browsers will render buggy HTML pages
  - *XML processors* are not supposed to process buggy XML documents
Example

```xml
<?xml version="1.0" ?>

<PersonList Type="Student" Date="2002-02-02">
  <Title Value="Student List" />
  <Person>
    ... ...
  </Person>
  <Person>
    ... ...
  </Person>
  ... ...
</PersonList>
```

- Elements are nested
- Root element contains all others
More Terminology

Opening tag

<Person Name = “John” Id = “s111111111”>

John is a nice fellow

“standalone” text, not very useful as data, non-uniform

<Address>

<Number>21</Number>

<Street>Main St.</Street>

</Address>

Nested element, child of Person

Child of Address, Descendant of Person

Closing tag:
What is open must be closed

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Conversion from XML to Objects

• Straightforward:

  <Person Name="Joe">
    <Age>44</Age>
    <Address><Number>22</Number><Street>Main</Street></Address>
  </Person>

  Becomes:

  (#345, [Name: “Joe”,
    Age: 44,
    Address: [Number: 22, Street: “Main”]
  ] )
Conversion from Objects to XML

- Also straightforward
- Non-unique:
  - Always a question if a particular piece (such as Name) should be an element in its own right or an attribute of an element
  - Example: A reverse translation could give
    
    ```xml
    <Person>
        <Name>Joe</Name>
        <Age>44</Age>
        <Address>
            <Number>22</Number>
            <Street>Main</Street>
        </Address>
    </Person>
    ```

    This or this

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Differences between XML Documents and Objects

- XML’s origin is document processing, not databases
  - Allows things like standalone text (useless for databases)
    
    `<foo> Hello <moo>123</moo> Bye </foo>`
  - XML data is ordered, while database data is not:
    
    `<something><foo>1</foo><bar>2</bar></something>`
    
    is different from
    
    `<something><bar>2</bar><foo>1</foo></something>`
    
    but these two complex values are same:
    
    `[something: [bar:1, foo:2]]`
    
    `[something: [foo:2, bar:1]]`
Attributes aren’t needed – just bloat the number of ways to represent the same thing:

```xml
<foo bar="12">ABC</foo>
```

vs.

```xml
<foobar><foo>ABC</foo><bar>12</bar></foobar>
```

- More concise
- More uniform, database-like
Well-formed XML Documents

- Must have a root element
- Every *opening tag* must have matching *closing tag*
- Elements must be *properly nested*
  - `<foo><bar></foo></bar>` is a no-no
- An *attribute* name can occur *at most once* in an opening tag. If it occurs,
  - It *must have an explicitly specified value* (Boolean attrs, like in HTML, are not allowed)
  - The value *must be quoted* (with “ or ‘)
- *XML processors are not supposed to try and fix ill-formed documents* (unlike HTML browsers)
Identifying and Referencing with Attributes

- An attribute can be declared (in a DTD – see later) to have type:
  - **ID** – unique identifier of an element
    - If attr1 & attr2 are both of type ID, then it is illegal to have `<something attr1="abc"> … <somethingelse attr2="abc"> within the same document
  - **IDREF** – references a unique element with matching ID attribute (in particular, an XML document with IDREFs is not a tree)
    - If attr1 has type ID and attr2 has type IDREF then we can have: `<something attr1="abc"> … <somethingelse attr2="abc">`
  - **IDREFS** – a list of references, if attr1 is ID and attr2 is IDREFS, then we can have
    - `<something attr1="abc"> … <somethingelse attr1="cde"> … <someotherthing attr2="abc  cde">`
Example: Report Document with Cross-References

<?xml version="1.0"?>
<Report Date="2002-12-12">
  <Students>
    <Student StudId="s111111111">
      <Name><First>John</First><Last>Doe</Last></Name>
      <Status>U2</Status>
      <CrsTaken CrsCode="CS308" Semester="F1997"/>
      <CrsTaken CrsCode="MAT123" Semester="F1997"/>
    </Student>
    <Student StudId="s666666666">
      <Name><First>Joe</First><Last>Public</Last></Name>
      <Status>U3</Status>
      <CrsTaken CrsCode="CS308" Semester="F1994"/>
      <CrsTaken CrsCode="MAT123" Semester="F1997"/>
    </Student>
    <Student StudId="s987654321">
      <Name><First>Bart</First><Last>Simpson</Last></Name>
      <Status>U4</Status>
      <CrsTaken CrsCode="CS308" Semester="F1994"/>
    </Student>
  </Students>
</Report>

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<Classes>

<Class>
  <CrsCode>CS308</CrsCode>  <Semester>F1994</Semester>
  <ClassRoster Members="s666666666 987654321" />
</Class>

<Class>
  <CrsCode>CS308</CrsCode>  <Semester>F1997</Semester>
  <ClassRoster Members="s111111111" />
</Class>

<Class>
  <CrsCode>MAT123</CrsCode>  <Semester>F1997</Semester>
  <ClassRoster Members="s111111111 s666666666" />
</Class>

</Classes>

…… continued … …
<Courses>
  <Course CrsCode = “CS038” >
    <CrsName>Market Analysis</CrsName>
  </Course>
  <Course CrsCode = “MAT123” >
    <CrsName>Market Analysis</CrsName>
  </Course>
</Courses>

</Report>
XML Namespaces

- A mechanism to prevent name clashes between components of same or different documents

- Namespace declaration
  - **Namespace** – a symbol, typically a URL *(doesn’t need to point to a real page)*
  - **Prefix** – an abbreviation of the namespace, a convenience; works as an alias
  - Actual name (element or attribute) – *prefix:name*
  - Declarations/prefixes have *scope* similarly to begin/end

- Example:

```xml
<item xmlns="http://www.acmeinc.com/jp#supplies"
      xmlns:toy="http://www.acmeinc.com/jp#toys">
  <name>backpack</name>
  <feature>
    <toy:item><toy:name>cyberpet</toy:name></toy:item>
  </feature>
</item>
```

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Scopes of declarations are color-coded:

```xml
<item xmlns="http://www.foo.org/abc"
     xmlns:cde="http://www.bar.com/cde">
  <name>…</name>
  <feature>
    <cde:item><cde:name>…</cde:name><cde:item>
  </feature>
</item>
```
Namespaces (cont’d.)

- `xmlns=“http://foo.com/bar”` *doesn’t* mean there is a document at this URL: using URLs is just a convenient convention; and a namespace is just an identifier.

- Namespaces aren’t part of XML 1.0, but all XML processors understand this feature now.

- A number of prefixes have become “standard” and some XML processors might understand them without any declaration. E.g.,
  - `<xs>` for `http://www.w3.org/2001/XMLSchema`
  - `<xsl>` for `http://www.w3.org/1999/XSL/Transform`
  - Etc.
Document Type Definition (DTD)

- A *DTD* is a grammar specification for an XML document
- DTDs are optional – don’t need to be specified
  - If specified, DTD can be part of the document (at the top); or it can be given as a URL
- A document that conforms (i.e., parses) w.r.t. its DTD is said to be *valid*
  - XML processors are not required to check validity, even if DTD is specified
  - But they are required to test well-formedness

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DTDs (cont’d)

• DTD specified as part of a document:
  ```xml
  <?xml version="1.0" ?>
  <!DOCTYPE Report [
      … … …
    ]>
  <Report> ……… … </Report>
  ```

• DTD specified as a standalone thing
  ```xml
  <?xml version="1.0" ?>
  <Report> ……… … </Report>
  ```
DTD Components

<!ELEMENT elt-name
    (...contents...)/EMPTY/ANY >
<!ATTLIST elt-name attr-name CDATA/ID/IDREF/IDREFS
    #IMPLIED/#REQUIRED >

• Can define other things, like macros (called entities in the XML jargon)
<!DOCTYPE Report [  
<!ELEMENT Report (Students, Classes, Courses)>  
<!ELEMENT Students (Student*)>  
<!ELEMENT Classes (Class*)>  
<!ELEMENT Courses (Course*)>  
<!ELEMENT Student (Name, Status, CrsTaken*)>  
<!ELEMENT Name (First,Last)>  
<!ELEMENT First (#PCDATA)>  
<!ELEMENT CrsTaken EMPTY>  
<!ELEMENT Class (CrsCode,Semester,ClassRoster)>  
<!ELEMENT Course (CrsName)>  
<!ATTLIST Report Date CDATA #IMPLIED>  
<!ATTLIST Student StudId ID #REQUIRED>  
<!ATTLIST Course CrsCode ID #REQUIRED>  
<!ATTLIST CrsTaken CrsCode IDREF #REQUIRED>  
<!ATTLIST ClassRoster Members IDREFS #IMPLIED>  
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Limitations of DTDs

- Doesn’t understand namespaces
- Very limited assortment of data types (just strings)
- Very weak w.r.t. consistency constraints (ID/IDREF/IDREFS only)
- Can’t express unordered contents conveniently
- All element names are global: can’t have one Name type for people and another for companies:
  
  ```xml
  <!ELEMENT Name (Last, First)>
  <!ELEMENT Name (#PCDATA)>
  ```
  
  both can’t be in the same DTD
XML Schema

• Came to rectify some of the problems with DTDs

• Advantages:
  • Integrated with namespaces
  • Many built-in types
  • User-defined types
  • Has local element names
  • Powerful key and referential constraints

• Disadvantages:
  • Unwieldy – much more complex than DTDs
Schema Document and Namespaces

```xml
<schema
  xmlns="http://www.w3.org/2001/XMLSchema"
  targetNamespace="http://xyz.edu/Admin">
  ...
</schema>
```

- Uses standard XML syntax.
- `http://www.w3.org/2001/XMLSchema` – namespace for keywords used in a schema document (not an instance document), e.g., “schema”, `targetNamespace`, etc.
- `targetNamespace` – names the namespace defined by the above schema.
Instance Document

• Report document whose structure is being defined by the earlier schema document

```xml
<?xml version="1.0"?>
<Report xmlns="http://xyz.edu/Admin"
       xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
       xsi:schemaLocation="http://xyz.edu/Admin
http://xyz.edu/Admin.xsd">

... 

</Report>
```

• `xsi:schemaLocation` says: the schema for the namespace http://xyz.edu/Admin is found in http://xyz.edu/Admin.xsd

• Document schema & its location are not binding on the XML processor; it can decide to use another schema, or none at all.
Building Schemas from Components

<schema xmlns="http://www.w3.org/2001/XMLSchema"
    targetNamespace="http://xyz.edu/Admin">

    <include schemaLocation="http://xyz.edu/StudentTypes.xsd">
    <include schemaLocation="http://xyz.edu/ClassTypes.xsd">
    <include schemaLocation="http://xyz.edu/CourseTypes.xsd">
        ...
    </include>
</include>

• <include…> works like #include in the C language
  • Included schemas must have the same targetNamespace as the including schema

• schemaLocation – tells where to find the piece to be included
Simple Types

- **Primitive types:** decimal, integer, Boolean, string, ID, IDREF, etc. (defined in XMLSchema namespace)

- **Type constructors:** list and union
  
  - A simple way to derive types from primitive types (disregard the namespaces for now):

    ```xml
    <simpleType name="myIntList">
      <list itemType="integer" />
    </simpleType>
    
    <simpleType name="phoneNumber">
      <union memberTypes="phone7digits phone10digits" />
    </simpleType>
    ```
Deriving Simple Types by Restriction

```xml
<simpleType name="phone7digits">
  <restriction base="integer">
    <minInclusive value="1000000"/>
    <maxInclusive value="9999999"/>
  </restriction>
</simpleType>

<simpleType name="emergencyNumbers">
  <restriction base="integer">
    <enumeration value="911"/>
    <enumeration value="333"/>
  </restriction>
</simpleType>

• Has more type-building primitives (see textbook and specs)
```
Some Simple Types Used in the Report Document

```xml
<simpleType name="studentId">
    <restriction base="ID">
        <pattern value="s[0-9]{9}" />
    </restriction>
</simpleType>

<simpleType name="studentIds">
    <list itemType="adm:studentRef" />
</simpleType>

<simpleType name="studentRef">
    <restriction base="IDREF">
        <pattern value="s[0-9]{9}" />
    </restriction>
</simpleType>
```

**targetNamespace =** http://xyz.edu/Admin

**xmlns:adm =** http://xyz.edu/Admin

*XML ID types always start with a letter*

*Prefix for the target namespace*
Simple Types for Report Document (contd.)

```xml
<simpleType name="courseCode">
    <restriction base="ID">
        <pattern value="[A-Z]{3}[0-9]{3}" />
    </restriction>
</simpleType>

<simpleType name="courseRef">
    <restriction base="IDREF">
        <pattern value="[A-Z]{3}[0-9]{3}" />
    </restriction>
</simpleType>

<simpleType name="studentStatus">
    <restriction base="string">
        <enumeration value="U1" />
        … … … …
        <enumeration value="G5" />
    </restriction>
</simpleType>
```

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<schema xmlns="http://www.w3.org/2001/XMLSchema"
    xmlns:adm="http://xyz.edu/Admin"
    targetNamespace="http://xyz.edu/Admin">
    ...
    <element name="CrsName" type="string"/>
    <element name="Status" type="adm:studentStatus"/>
    ...
    <simpleType name="studentStatus">
        ...
    </simpleType>
</schema>
Complex Types

- Allows the definition of element types that have complex internal structure
- Similar to class definitions in object-oriented databases
  - Very verbose syntax
  - Can define both child elements and attributes
  - Supports ordered and unordered collections of elements
Example: studentType

```xml
<element name="Student" type="adm:studentType" />
<complexType name="studentType">
  <sequence>
    <element name="Name" type="adm:personNameType" />
    <element name="Status" type="adm:studentStatus" />
    <element name="CrsTaken" type="adm:courseTakenType"
      minOccurs="0" maxOccurs="unbounded" />
  </sequence>
  <attribute name="StudId" type="adm:studentId" />
</complexType>
<complexType name="personNameType">
  <sequence>
    <element name="First" type="string" />
    <element name="Last" type="string" />
  </sequence>
</complexType>
```

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Compositors: Sequences, Sets, Alternatives

- **Compositors:**
  - *sequence, all, choice* are required when element has at least 1 child element (\(\equiv\) complex content)

- sequence

- all – can specify sets of elements

- choice – can specify alternative types
Sets

- Suppose the order of components in addresses is unimportant:

  <complexType name="addressType">
  <all>
    <element name="StreetName" type="string"/>
    <element name="StreetNumber" type="string"/>
    <element name="City" type="string"/>
  </all>
  </complexType>

- Problem: all comes with a host of awkward restrictions. For instance, cannot occur inside a sequence; only sets of elements, not bags.
Alternative Types

- Assume addresses can have P.O.Box or street name/number:

```xml
<complexType name="addressType">
  <sequence>
    <choice>
      <element name="POBox" type="string"/>
      <sequence>
        <element name="Name" type="string"/>
        <element name="Number" type="string"/>
      </sequence>
    </choice>
    <element name="City" type="string"/>
  </sequence>
</complexType>
```
Local Element Names

- A DTD can define only global element name:
  - Can have at most one `<!ELEMENT foo ...>` statement per DTD
- In XML Schema, names have scope like in programming languages – the nearest containing complexType definition
  - Thus, can have the same element name (e.g., Name), within different types and with different internal structures
Local Element Names: Example

```xml
<complexType name="studentType">
  <sequence>
    <element name="Name" type="adm:personNameType" />
    <element name="Status" type="adm:studentStatus" />
    <element name="CrsTaken" type="adm:courseTakenType"
              minOccurs="0" maxOccurs="unbounded" />
  </sequence>
  <attribute name="StudId" type="adm:studentId" />
</complexType>

<complexType name="courseType">
  <sequence>
    <element name="Name" type="string" />
  </sequence>
  <attribute name="CrsCode" type="adm:courseCode" />
</complexType>
```

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Importing XML Schemas

- Import is used to share schemas developed by different groups at different sites

- Include vs. import:
  - **Include**:
    - Included schemas are usually under the control of the same development group as the including schema
    - Included and including schemas must have the same target namespace (because the text is physically included)
    - `schemaLocation` attribute required
  
  - **Import**:
    - Schemas are under the control of different groups
    - Target namespaces are different
    - The import statement must tell the importing schema what that target namespace is
    - `schemaLocation` attribute optional
<schema xmlns="http://www.w3.org/2001/XMLSchema"
    targetNamespace="http://xyz.edu/Admin"
    xmlns:reg="http://xyz.edu/Registrar"
    xmlns:crs="http://xyz.edu/Courses">

    <import namespace="http://xyz.edu/Registrar"
        schemaLocation="http://xyz.edu/Registrar/StudentType.xsd" />

    <import namespace="http://xyz.edu/Courses" />

    ................
    ................
    ................

    </schema>
Extension and Restriction of Base Types

- Mechanism for modifying the types in imported schemas
- Similar to subclassing in object-oriented languages
- *Extending* an XML Schema type means adding elements or adding attributes to existing elements
- *Restricting* types means tightening the types of the existing elements and attributes (i.e., replacing existing types with subtypes)
Type Extension: Example

```xml
<schema xmlns="http://www.w3.org/2001/XMLSchema"
    xmlns:xyzCrs="http://xyz.edu/Courses"
    xmlns:fooAdm="http://foo.edu/Admin"
    targetNamespace="http://foo.edu/Admin">
    <import namespace="http://xyz.edu/Courses"/>

    <complexType name="courseType">
        <complexContent>
            <extension base="xyzCrs:courseType">
                <element name="syllabus" type="string"/>
            </extension>
        </complexContent>
    </complexType>
    <element name="Course" type="fooAdm:courseType"/>

    ........
</schema>
```

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Type Restriction: Example

```xml
<schema xmlns="http://www.w3.org/2001/XMLSchema"
xmlns:xyzCrs="http://xyz.edu/Courses"
xmlns:fooAdm="http://foo.edu/Admin"
targetNamespace="http://foo.edu/Admin" >
  <import namespace="http://xyz.edu/Courses" />

  <complexType name="studentType">
    <complexContent>
      <restriction base="xyzCrs:studentType">
        <sequence>
          <element name="Name" type="xyzCrs:personNameType" />
          <element name="Status" type="xyzCrs:studentStatus" />
          <element name="CrsTaken" type="xyzCrs:courseTakenType" minOccurs="0" maxOccurs="60" />
        </sequence>
        <attribute name="StudId" type="xyzCrs:studentId" />
      </restriction>
    </complexContent>
  </complexType>

  <element name="Student" type="fooAdm:studentType" />
</schema>
```

**Must repeat the original definition**

**Tightened type: the original was “unbounded”**
Structure of an XML Schema Document

```xml
<schema xmlns="http://www.w3.org/2001/XMLSchema"
    xmlns:adm="http://xyz.edu/Admin"
    targetNamespace="http://xyz.edu/Admin">
  <element name="Report" type="adm:reportType"/>

  <complexType name="reportType">
    .......
  </complexType>
  <complexType name="...">
    .......
  </complexType>
  .......
</schema>
```

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Anonymous Types

• So far all types were named
  • Useful when the same type is used in more than one place
• When a type definition is used exactly once, anonymous types can save space

```
<element name="Report">
  <complexType>
    <sequence>
      <element name="Students" type="adm:studentList" />
      <element name="Classes" type="adm:classOfferings" />
      <element name="Courses" type="adm:courseCatalog" />
    </sequence>
  </complexType>
</element>
```

"element" used to be empty element – now isn’t

No type name

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Integrity Constraints in XML Schema

- A DTD can specify only very simple kinds of key and referential constraint; only using attributes
- XML Schema also has ID, IDREF as primitive data types, but these can also be used to type elements, not just attributes
- In addition, XML Schema can express complex key and foreign key constraints
A key in an XML document is a sequence of components, which might include elements and attributes, which uniquely identifies document components in a source collection of objects in the document.

Issues:

- Need to be able to identify that source collection
- Need to be able to tell which sequences form the key

For this, XML Schema uses XPath – a simple XML query language.
Basic XPath – for Key Specification

<Offerings>  --  current reference point
  <Offering>
    <CrsCode Section="1">CS532</CrsCode>
    <Semester><Term>Spring</Term><Year>2002</Year></Semester>
  </Offering>
  <Offering>
    <CrsCode Section="2">CS305</CrsCode>
    <Semester><Term>Fall</Term><Year>2002</Year></Semester>
  </Offering>
</Offerings>

Offering/CrsCode/@Section  --  selects occurrences of attribute Section within CrsCode within Offerings
Offering/CrsCode  --  selects all CrsCode element occurrences within Offerings
Offering/Semester/Term  --  all Term elements within Semester within Offerings
Offering/Semester/Year  --  all Year elements within Semester within Offerings

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Keys: Example

<complexType name="reportType">
  <sequence>
    <element name="Students"/>
    <element name="Classes">
      <complexType>
        <sequence>
          <element name="Class" minOccurs="0" maxOccurs="unbounded">
            <sequence>
              <element name="CrsCode"/>
              <element name="Semester"/>
              <element name="ClassRoster"/>
            </sequence>
          </element>
        </sequence>
      </complexType>
    </element>
  </sequence>
</complexType>

... key specification goes here – next slide ...

@endelement
<element name="Courses"/>
<sequence>
</complexType>
Example (cont’d)

- A key specification:

  ```xml
  <key name="PrimaryKeyForClass">
    <selector xpath="Class" />
    <field xpath="CrsCode" />
    <field xpath="Semester" />
  </key>
  ```

  Defines **source collection** of objects to which the key applies. The XPath expression is relative to element to which the key is local.

  Fields that form the key.
  The XPath expression is relative to the source collection of objects specified in selector. So, **CrsCode** is actually **Classes/Class/CrsCode**

**field** must return exactly one value per object specified by **selector**
Foreign Keys

• Like the REFERENCES clause in SQL, but more involved
• Need to specify:
  • Foreign key:
    • Source collection of objects
    • Fields that form the foreign key
  • Target key:
    • A previously defined key (or unique) specification, which is comprised of:
      • Target collection of objects
      • Sequence of fields that comprise the key
Foreign Key: Example

```xml
<keyref name="NoEmptyClasses" refer="adm:PrimaryKeyForClass">
  <selector xpath="Student/CrsTaken" />
  <field xpath="@CrsCode" />
  <field xpath="@Semester" />
</keyref>
```

Every class must have at least one student!
XML Query Languages

• XPath – core query language. Very limited, a selection operator. Very useful, though: used in XML Schema, XSLT, XQuery, many other XML standards
• XSLT – a functional style document transformation language. Very powerful, very complicated
• XQuery – W3C standard. Very powerful, fairly intuitive, SQL-style.
• SQL/XML – attempt to marry SQL and XML, part of SQL:2003.
Why Query XML?

- Need to extract parts of XML documents
- Need to transform documents into different forms
- Need to relate – join – parts of the same or different documents
XPath

- Analogous to path expressions in object-oriented languages (e.g., OQL)
- Extends path expressions with query facility
- XPath views an XML document as a tree
  - Root of the tree is a new node, which doesn’t correspond to anything in the document
  - Internal nodes are elements
  - Leaves are either
    - Attributes
    - Text nodes
    - Comments
    - Other: processing instructions, etc.
<?xml version="1.0" ?>
<!-- Some comment -->

<Students>
  <Student StudId="111111111">
    <Name><First>John</First><Last>Doe</Last></Name>
    <Status>U2</Status>
    <CrsTaken CrsCode="CS308" Semester="F1997"/>
    <CrsTaken CrsCode="MAT123" Semester="F1997"/>
  </Student>
  <Student StudId="987654321">
    <Name><First>Bart</First><Last>Simpson</Last></Name>
    <Status>U4</Status>
    <CrsTaken CrsCode="CS308" Semester="F1994"/>
  </Student>
</Students>
XPath Document Tree

Root of XML tree

Root of XML document

Legend:
- Text
- Element
- Attribute
- Comment
- Root

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Terminology

- **Parent/child** nodes, as usual
- Child nodes (that are of interest to us) are of types text, element, attribute
  - We call them \textit{t-children}, \textit{e-children}, \textit{a-children}
  - Also, \textit{et-children} are child-nodes that are either elements or text, \textit{ea-children} are child nodes that are either elements or attributes, etc.
- Ancestor/descendant nodes — as usual in trees
XPath Basics

• An XPath expression takes a document tree as input and returns a multi-set of nodes of the tree
• Expressions that start with `/` are absolute path expressions
  • Expression `/` – returns root node of XPath tree
  • `/Students/Student` – returns all Student-elements that are children of Students elements, which in turn must be children of the root
  • `/Student` – returns empty set (no such children at root)
**XPath Basics (cont’d)**

- **Current** (or context node) – exists during the evaluation of XPath expressions (and in other XML query languages)
- `. ` – denotes the current node; `.. ` – denotes the parent
  - foo/bar – returns all bar-elements that are children of foo nodes, which in turn are children of the current node
  - ./foo/bar – same
  - ../abc/cde – all cde e-children of abc e-children of the parent of the current node
- Expressions that don’t start with `/` are relative (to the current node)
Attributes, Text, etc.

- `/Students/Student/`\text{@}StudentId – returns all StudentId a-children of Student, which are e-children of Students, which are children of the root
- `/Students/Student/Name/Last/`\text{text()} – returns all t-children of Last e-children of …
- `/comment()` – returns comment nodes under root
Overall Idea and Semantics

- An XPath expression is:
  \text{locationStep1/locationStep2/...}

- \textbf{Location step}:
  \text{Axis::nodeSelector[predicate]}

- \textbf{Navigation axis}:
  - \textit{child, parent} – have seen
  - \textit{ancestor, descendant, ancestor-or-self, descendant-or-self} – will see later
  - some other

- \textbf{Node selector}: node name or wildcard; e.g.,
  - \text{./child::Student} (we used ./Student, which is an abbreviation)
  - \text{./child::*} – any e-child (abbreviation: ./*)

- \textbf{Predicate}: a selection condition; e.g.,
  Students/Student[CourseTaken/@CrsCode = “CS532”]

This is called \textit{full} syntax.
We used \textit{abbreviated} syntax before.
Full syntax is better for describing meaning. Abbreviated syntax is better for programming.
XPath Semantics

- The meaning of the expression \( \text{locationStep1/\text{locationStep2/...}} \) is the set of all document nodes obtained as follows:
  - Find all nodes reachable by \( \text{locationStep1} \) from the current node
  - For each node \( N \) in the result, find all nodes reachable from \( N \) by \( \text{locationStep2} \); take the union of all these nodes
  - For each node in the result, find all nodes reachable by \( \text{locationStep3} \), etc.
  - The value of the path expression on a document is the set of all document nodes found after processing the last location step in the expression
Algorithm

- `locationStep1/locationStep2/...`
  - Find all nodes specified by `locationStep1`
  - For each such node `N`:
    - Find all nodes specified by `locationStep2` using `N` as the current node
    - Take union
  - For each node returned by `locationStep2` do the same

- `locationStep = axis::node[predicate]`
  - Find all nodes specified by `axis::node`
  - Select only those that satisfy `predicate`
More on Navigation Primitives

- 2nd CrsTaken child of 1st Student child of Students:
  
  \[ /Students/Student[1]/CrsTaken[2] \]

- All last CourseTaken elements within each Student element:
  
  \[ /Students/Student/CrsTaken[last( )] \]
Wildcards

- Wildcards are useful when the exact structure of document is not known

- *Descendant-or-self* axis, `//` : allows to descend down any number of levels (including 0)
  - `//CrsTaken` – all `CrsTaken` nodes under the root
  - `Students//@Name` – all `Name` attribute nodes under the elements `Students`, who are children of the current node

- *Note:* `.//Last` and `//Last` are different

- The * wildcard:
  - `*` – any element: `Student/*/text()`
  - `@*` – any attribute: `Students//@*`
XPath Queries (selection predicates)

- Location step = Axis::nodeSelector[predicate]

- **Predicate:**
  - XPath expression = const | built-in function | XPath expression
  - XPath expression
  - built-in predicate
  - a Boolean combination thereof

- Axis::nodeSelector[predicate] ⊆ Axis::nodeSelector but contains only the nodes that satisfy predicate

- Built-in predicate: special predicates for string matching, set manipulation, etc.

- Built-in function: large assortment of functions for string manipulation, aggregation, etc.
XPath Queries – Examples

• Students who have taken CSE532:

  //Student[CrsTaken/@CrsCode="CSE532"]

  True if:  “CSE532” ∈ //Student/CrsTaken/@CrsCode

• Complex example:

  //Student[Status="U3" and starts-with(./Last, "A")

  and contains(concat(./@CrsCode), "ESE")

  and not(./Last = ./First) ]

• Aggregation:  sum( ), count( )

  //Student[sum(./@Grade) div count(./@Grade) > 3.5]
Xpath Queries (cont’d)

- Testing whether a subnode exists:
  - `//Student[CrsTaken/@Grade]` — students who have a grade (for some course)
  - `//Student[Name/First or CrsTaken/@Semester or Status/text() = “U4”]` — students who have either a first name or have taken a course in some semester or have status U4

- Union operator, `|`:
  - `//CrsTaken[@Semester=“F2001”] | //Class[Semester=“F1990”]

- union lets us define *heterogeneous* collections of nodes
XPointer

• **XPointer** = URL + XPath

• **Syntax:**

  
  url # xpointer (XPathExpr1) xpointer (XPathExpr2) …

  • Follow url
  • Compute XPathExpr1
    • Result non-empty?  – return result
    • Else:  compute XPathExpr2; and so on

• Example: you might click on a link and run a query against your Registrar’s database

  http://yours.edu/Report.xml#xpointer(
    //Student[CrsTaken/@CrsCode="CS532"
    and  CrsTaken/@Semester="S2012"] )
XSLT: XML Transformation Language

- Powerful programming language, uses *functional programming paradigm*
- Originally designed as a stylesheet language: this is what “S”, “L”, and “T” stand for
  - The idea was to use it to display XML documents by transforming them into HTML
  - For this reason, XSLT programs are often called *stylesheets*
  - Their use is not limited to stylesheets – can be used to query XML documents, transform documents, etc.
- In wide use, but semantics is very complicated
XSLT Basics

- One way to apply an XSLT program to an XML document is to specify the program as a stylesheet in the document preamble using a processing instruction:

```xml
<?xml version="1.0" ?>
......
<xml-stylesheet type="text/xsl"
......
<Report  Date="2002-11-11">
......
</Report>
```

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**XSLT Example**

- *Extract the list of all students from this (hyperlinked) document*
  
  ```xml
  <?xml version="1.0" ?>
  <StudentList xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
  xsl:version="1.0">
    <xsl:copy-of select="//Student/Name"/>
  </StudentList>
  ```

  **Result:**
  
  ```xml
  <StudentList>
    <Name><First>John</First><Last>Doe</Last></Name>
    <Name><First>Bart</First><Last>Simpson</Last></Name>
  </StudentList>
  ```
<StudentList xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
xsl:version="1.0">
  <xsl:for-each select="/Student">
    <xsl:if test="count(CrsTaken) &gt; 1">
      <FullName>
        <xsl:value-of select="*/Last" /> ,
        <xsl:value-of select="*/First" />
      </FullName>
    </xsl:if>
  </xsl:for-each>
</StudentList>

Result:
<StudentList>
  <FullName>
    Doe, John
  </FullName>
</StudentList>

Extracts contents of element, not the element itself (unlike copy-of)
XSLT Pattern-based Templates

- Where the real power lies
- Issue: how to process XML documents by descending into their structure
- Previous syntax was just a shorthand for template syntax – next slide
Full Syntax vs. Simplified Syntax

- **Simplified syntax:**
  ```xml
  <StudentList xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
               xsl:version="1.0">
    <xsl:for-each select="//Student">
      ...
    </xsl:for-each>
  </StudentList>
  ```

- **Full syntax:**
  ```xml
  <xsl:stylesheet xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
                  xsl:version="1.0">
    <xsl:template match="/">
      <StudentList>
        <xsl:for-each select="//Student">
          ...
        </xsl:for-each>
      </StudentList>
    </xsl:template>
  </xsl:stylesheet>
  ```

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Recursive Stylesheets

• A bunch of templates of the form:
  
  `<xsl:template match="XPath-expression">`
  
  … tags, XSLT instructions …
  
  `</xsl:template>`

• Template is applied to the node that is current in the evaluation process

• Template is used if its XPath expression is matched:
  
  • “Matched” means: current node ∈ result set of XPath expression
  
  • If several templates match: use the best matching template – template with the smallest (by inclusion) XPath expression result set
  
  • If several of those: other rules apply (see XSLT specs)
  
  • If no template matches, use the matching default template
    
    • There is one default template for et-children and one for a-children
Recursive Traversal of Document

- `<xsl:apply-templates/>` – XSLT instruction that drives the recursive process of descending into the document tree
- Constructs the list of et-children of the current node
- For each node in the list, applies the best matching template
- A typical initial template:
  
  ```xml
  <xsl:template match="/">
    <StudentList>
      <xsl:apply-templates />
    </StudentList>
  </xsl:template>
  ```
  
  *Start with the root node – typically the first template to be used in a stylesheet*
- Outputs `<StudentList>` – `</StudentList>` tag pair
- Applies templates to the et-children of the current node
- Inserts whatever output is produced in-between `<StudentList>` and `</StudentList>`
Recursive Stylesheet Example

- **List the names of students with > 1 courses:**

```xml
<?xml version="1.0" ?>
<xsl:stylesheet xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
    xmlns:xs="http://www.w3.org/2001/XMLSchema"
    version="1.0">
    <xsl:template match="/">
        <StudentList>
            <xsl:apply-templates/>
        </StudentList>
    </xsl:template>
    <xsl:template match="//Student">
        <xsl:if test="count(Crs Taken) &gt; 1">
            <FullName>
                <xsl:value-of select="*/Last"/>
                <xsl:value-of select="*/First"/>
            </FullName>
        </xsl:if>
    </xsl:template>
    <xsl:template match="text()"></xsl:template>
</xsl:stylesheet>
```

*Initial template*

*The workhorse, does all the job*

*Empty template – no-op. Needed to block default template for text – later.*

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Example Dissected

- *Initial template*: starts off, applies templates to *et*-children. The only *et*-child is *Students* element
- Stylesheet has no matching template for *Students*!
- Use *default template*: For *e*-nodes or root (/) the default is to go down to the *et*-children:
  
  ```xml
  <xsl:template match="* | /">
    <xsl:apply-templates />
  </xsl:template>
  ```

- Children of *Students* node are two *Student* nodes — the “workhorse” template matches!
  - For each such (*Student*) node output:
    
    ```xml
    <FullName>Last, First</FullName>
    ```

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Example (cont’d)

• Consider this expanded document:

```
<Report>
  <Students>
    <Student StudId="111111111"> .... .... </Student>
    <Student StudId="987654321"> .... .... </Student>
  </Students>
  <Courses>
    <Course CrsCode="CS308">
      <CrsName>Software Engineering</CrsName>
    </Course>
    .... ....
  </Courses>
</Report>
```

• Then the previous stylesheet has another branch to explore

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Example (cont’d)

- No stylesheet template applies to *Courses*-element, so use the default template
- No explicit template applies to children, *Course*-elements – use the default again
- Nothing applies to *CrsName* – use the default
- The child of *CrsName* is a text node. If we used the default here:
  
  For text/attribute nodes the XSLT default is
  
  ```xml
  <xsl:template match="text() | @*" >
    <xsl:value-of select="." />
  </xsl:template>
  ```
  
  i.e., output the contents of text/attribute – we don’t want this!

  This is why we provided the empty template for text nodes – to suppress the application of the default template
XSLT Evaluation Algorithm

- Very involved
- Not even properly defined in the official XSLT specification!
- More formally described in a research paper by Wadler – can only hope that vendors read this
- Will describe simplified version – will omit the `for-each` statement
XSLT Evaluation Algorithm (cont’d)

- Create root node, \( \textit{OutRoot} \), for the output document
- Copy root of the input document, \( \textit{InRoot} \), to output document: \( \textit{InRoot}^R \). Make \( \textit{InRoot}^R \) a child of \( \textit{OutRoot} \)
  - Set current node variable: \( \textit{CN} := \textit{InRoot} \)
  - Set current node list: \( \textit{CNL} := <\textit{InRoot}> \)
- \( \textit{CN} \): always the 1st node in \( \textit{CNL} \)
- When a node \( \textit{N} \) is placed on \( \textit{CNL} \), its copy, \( \textit{N}^R \), goes to the output document (becomes a child of some node – see later)
  - \( \textit{N}^R \) is a marker for where subsequent actions apply in the output document
  - Might be deleted or replaced later
- Find the best matching template for \( \textit{CN} \) (or default template, if nothing applies)
- Apply this template to \( \textit{CN} \) – next slide
XSLT Evaluation Algorithm – Application of a Template

• Application of template can cause these changes:

  Case A: \( CN^R \) is replaced by a subtree

  Example: \( CN = \text{Students} \) node in our document. Assume our stylesheet has the following template instead of the initial template (it thus becomes best-matching):

  \[
  \begin{align*}
  <\text{xsl:template}\text{ match}="//\text{Students}" > \\
  <\text{StudentList}>
  <\text{xsl:apply-templates} /> \\
  </\text{StudentList}>
  </\text{xsl:template}>
  \]

  Then:
  - \( CN^R \) is replaced with \( \text{StudentList} \)
  - Each child of \( CN \) (\( \text{Students} \) node) is copied over to the output tree as a child of \( \text{StudentList} \)
Case B:  \( CN^R \) is deleted and its children become children of the parent of \( CN^R \)

Example: The default template, below, deletes \( CN^R \) when applied to any node:

```xml
<xsl:template match="* | /">
  <xsl:apply-templates />
</xsl:template>
```
The Effect of **apply-templates** on Document Tree

Source Tree

- **N**
  - **P**
  - **Q**

Old Result Tree

- **N**
  - Step 4, Case(a)
    - Old **CN**
      - Old **CNL** = \(<N, \ldots, >\)
    - New **CN**
      - New **CNL** = \(<P, Q, \ldots, >\)

(a) New Result Tree

- **StudentList**
  - **P**
  - **Q**

(b) New Result Tree

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XSLT Evaluation Algorithm (cont’d)

- In both cases (A & B):
  - If \( CN \) has no \( et \)-children, \( CNL \) becomes shorter
  - If it does have children, \( CNL \) is longer or stays the same length
  - The order in which \( CN \)’s children are placed on \( CNL \) is their order in the source tree
  - The new 1\(^{st}\) node in \( CNL \) becomes the new \( CN \)
- Algorithm terminates when \( CNL \) is empty
  - Be careful – might not terminate (see next)
XSLT Evaluation Algorithm – Subtleties

- **apply-templates** instruction can have select attribute:
  
  \[
  \langle \text{xsl:apply-templates select=“node()”} \rangle \quad – \text{equivalent to the usual} \\
  \langle \text{xsl:apply-templates} \rangle \\
  \langle \text{xsl:apply-templates select=“@* | text()”} \rangle \quad – \text{instead of the } \text{et-children of } \text{CN, take } \text{at-children} \\
  \langle \text{xsl:apply-templates select=“.”} \rangle \quad – \text{take the parent of } \text{CN} \\
  \langle \text{xsl:apply-templates select=“.”} \rangle \quad – \text{will cause an infinite loop!!} \\
  \]

- **Recipe to guarantee termination:** make sure that **select** in **apply-templates** selects nodes only from a subtree of **CN**

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Advanced Example

- *Example*: take any document and replace attributes with elements. So that

  ```xml
  <Student StudId="111111111">
    <Name>John Doe</Name>
    <CrsTaken CrsCode="CS308" Semester="F1997" />
  </Student>
  ```

  would become:

  ```xml
  <Student>
    <StudId>111111111</StudId>
    <Name>John Doe</Name>
    <CrsTaken>
      <CrsCode>CS308</CrsCode> <Semester>F1997</Semester>
    </CrsTaken>
  </Student>
  ```
Advanced Example (cont’d)

- *Additional requirement*: don’t rely on knowing the names of the attributes and elements in input document – should be completely general. Hence:

1. Need to be able to output elements whose name is not known in advance (we don’t know which nodes we might be visiting)
   - Accomplished with `xsl:element` instruction and Xpath functions `current()` and `name()`:

   ```xml
   <xsl:element name="name(current())">
     Where am I?
   </xsl:element>
   
   If the current node is `foobar`, will output:
   ```
   ```xml
   <foobar>
     Where am I?
   </foobar>
   ```
Advanced Example (cont’d)

2. Need to be able to copy the current element over to the output document

- The *copy-of* instruction won’t do: it copies elements over with all their belongings. But remember: *we don’t want attributes to remain attributes*

- So, use the *copy* instruction
  - Copies the current node to the output document, but without any of its children

```xml
<xsl:copy>
  … XSLT instructions, which fill in the body of the element being copied over …
</xsl:copy>
```
Advanced Example (cont’d)

```xml
<xsl:stylesheet ... ...
  <xsl:template match="node()">
    <xsl:copy>
      <xsl:apply-templates select="@*"/>
    </xsl:copy>
  </xsl:template>

  <xsl:template match="@*">
    <xsl:element name="name(current())">
      <xsl:value-of select="."/>
    </xsl:element>
  </xsl:template>

  <xsl:template>
    <xsl:element name="name(current())">
      <xsl:value-of select="."/>
    </xsl:element>
  </xsl:template>

</xsl:stylesheet>
```

- **Process elements/text**
- **Process a-children of current element**
- **Process et-children of current element**
- **Deal with attributes separately**
- **Convert attribute to element**

```
... Attr="foo" 
becomes 
<Attr>foo</Attr>
```
Limitations of XSLT as a Query Language

- Programming style unfamiliar to people trained on SQL
- Most importantly: Hard to do joins, i.e., real queries
  - Requires the use of variables
  - Even harder than a simple nested loop (which one would use in this case in a language like C or Java)
XQuery – XML Query Language

- Integrates XPath with earlier proposed query languages: XQL, XML-QL
- SQL-style, not functional-style
- Much easier to use as a query language than XSLT
- Can do pretty much the same things as XSLT and more, but typically easier
- 2004: XQuery 1.0
Consider transcript.xml

<Transcripts>
  <Transcript>
    <Student StudId="111111111" Name="John Doe" />
    <CrsTaken CrsCode="CS308" Semester="F1997" Grade="B" />
    <CrsTaken CrsCode="MAT123" Semester="F1997" Grade="B" />
    <CrsTaken CrsCode="EE101" Semester="F1997" Grade="A" />
    <CrsTaken CrsCode="CS305" Semester="F1995" Grade="A" />
  </Transcript>
  <Transcript>
    <Student StudId="987654321" Name="Bart Simpson" />
    <CrsTaken CrsCode="CS305" Semester="F1995" Grade="C" />
    <CrsTaken CrsCode="CS308" Semester="F1994" Grade="B" />
  </Transcript>
  … … cont’d … …
</Transcripts>
transcript.xml (cont’d)

<Transcript>
    <Student StudId="123454321" Name="Joe Blow" />
    <CrsTaken CrsCode="CS315" Semester="S1997" Grade="A" />
    <CrsTaken CrsCode="CS305" Semester="S1996" Grade="A" />
    <CrsTaken CrsCode="MAT123" Semester="S1996" Grade="C" />
</Transcript>

<Transcript>
    <Student StudId="023456789" Name="Homer Simpson" />
    <CrsTaken CrsCode="EE101" Semester="F1995" Grade="B" />
    <CrsTaken CrsCode="CS305" Semester="S1996" Grade="A" />
</Transcript>

</Transcripts>
XQuery Basics

- **General structure:**
  
  \[
  \text{FOR} \quad \text{variable declarations} \\
  \text{WHERE} \quad \text{condition} \\
  \text{RETURN} \quad \text{document}
  \]

- **Example:**

  \[
  (:\text{students who took MAT123}:) \\
  \text{FOR} \quad \$t \quad \text{IN} \quad \text{doc("http://xyz.edu/transcript.xml")//Transcript} \\
  \text{WHERE} \quad \$t/\text{CrsTaken/}@\text{CrsCode} = \text{"MAT123"} \\
  \text{RETURN} \quad \$t/\text{Student}
  \]

- **Result:**

  \[
  <\text{Student} \quad \text{StudId=}\text{"111111111"} \quad \text{Name=}\text{"John Doe"} /> \\
  <\text{Student} \quad \text{StudId=}\text{"123454321"} \quad \text{Name=}\text{"Joe Blow"} /> \\
  \]

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XQuery Basics (cont’d)

• Previous query doesn’t produce a well-formed XML document; the following does:

```xml
<StudentList>
  
  FOR $t IN doc("transcript.xml")/Transcript
  WHERE $t/CrsTaken/@CrsCode = "MAT123"
  RETURN $ /Student

</StudentList>
```

• FOR binds $t to Transcript elements one by one, filters using WHERE, then places Student-children as e-children of StudentList using RETURN
Document Restructuring with XQuery

- **Reconstruct lists of students taking each class using the Transcript records:**

```xquery
FOR $c IN distinct-values(doc("transcript.xml")//CrsTaken)
RETURN
  <ClassRoster CrsCode = {${c}//@CrsCode} Semester = {${c}//@Semester}>
    {FOR $t IN doc("transcript.xml")//Transcript
     WHERE $t/CrsTaken/[@CrsCode = $c/@CrsCode and @Semester = $c/@Semester]
     RETURN $t/Student}
    ORDER BY $t/Student//@StudId
  </ClassRoster>
ORDER BY $c//@CrsCode
```

Query inside RETURN – similar to query inside SELECT in OQL
Document Restructuring (cont’d)

• **Output elements have the form:**

```xml
<ClassRoster  CrsCode="CS305"  Semester="F1995" >
  <Student  StudId="111111111"  Name="John Doe" />  
  <Student  StudId="987654321"  Name="Bart Simpson" />  
</ClassRoster>
```

• **Problem:** the above element will be output twice — once when $c$ is bound to

```xml
<CrsTaken  CrsCode="CS305"  Semester="F1995"  Grade="A" />
```

and once when it is bound to

```xml
<CrsTaken  CrsCode="CS305"  Semester="F1995"  Grade="C" />
```

**Note:** grades are different – distinct-values() won’t eliminate transcript records that refer to same class!
Solution: instead of

\[
\text{FOR } \$c \text{ IN distinct-values(doc("transcript.xml")//CrsTaken)}
\]

\[
\text{use}
\]

\[
\text{FOR } \$c \text{ IN doc("classes.xml")//Class}
\]

where \text{classes.xml} lists course offerings (course code/semester) explicitly (no need to extract them from transcript records).

Then \$c is bound to each class exactly once, so each class roster will be output exactly once.
http://xyz.edu/classes.xml

<Classes>
  <Class  CrsCode="CS308"  Semester="F1997" >
    <CrsName>SE</CrsName>  <Instructor>Adrian Jones</Instructor>
  </Class>

  <Class  CrsCode="EE101"  Semester="F1995" >
    <CrsName>Circuits</CrsName>  <Instructor>David Jones</Instructor>
  </Class>

  <Class  CrsCode="CS305"  Semester="F1995" >
    <CrsName>Databases</CrsName>  <Instructor>Mary Doe</Instructor>
  </Class>

  <Class  CrsCode="CS315"  Semester="S1997" >
    <CrsName>TP</CrsName>  <Instructor>John Smyth</Instructor>
  </Class>

  <Class  CrsCode="MAR123"  Semester="F1997" >

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Document Restructuring (cont’d)

- More problems: the above query will list classes with no students. Reformulation that avoids this:

```xml
FOR $c$ IN doc("classes.xml")//Class
  WHERE doc("transcripts.xml")
    //CrsTaken[@CrsCode = $c/@CrsCode
    and @Semester = $c/@Semester]
  RETURN
    <ClassRoster CrsCode = {$c/@CrsCode}   Semester =
    {$c/@Semester}> { 
      FOR $t$ IN doc("transcript.xml")//Transcript
        WHERE $t$/CrsTaken[@CrsCode = $c/@CrsCode
        and @Semester = $c/@Semester]
        RETURN  $t$/Student  ORDER BY $t$/Student/@StudId
    } </ClassRoster>
ORDER BY $c/@CrsCode
```

Test that classes aren’t empty
XQuery Semantics

- So far the discussion was informal
- XQuery *semantics* defines what the expected result of a query is
- Defined analogously to the semantics of SQL
XQuery Semantics (cont’d)

- **Step 1**: Produce a list of bindings for variables
  - The FOR clause binds each variable to a *list* of nodes specified by an XQuery expression.
    - The expression can be:
      - An XPath expression
      - An XQuery query
      - A function that returns a list of nodes
  - End result of a FOR clause:
    - Ordered list of tuples of document nodes
    - Each tuple is a binding for the variables in the FOR clause
Example (bindings):

- Let FOR declare $A$ and $B$
- Bind $A$ to document nodes \{v,w\}; $B$ to \{x,y,z\}
- Then FOR clause produces the following list of bindings for $A$ and $B$:
  - $A/v, B/x$
  - $A/v, B/y$
  - $A/v, B/z$
  - $A/w, B/x$
  - $A/w, B/y$
  - $A/w, B/z$
XQuery Semantics (cont’d)

• **Step 2**: filter the bindings via the WHERE clause
  • Use each tuple binding to substitute its components for variables; retain those bindings that make WHERE true

• Example:  WHERE  $A/CrsTaken/@CrsCode = $B/Class/@CrsCode

  • Binding:  $A/w, where w = <CrsTaken CrsCode=“CS308” …/>  
  $B/x, where x = <Class  CrsCode=“CS308” … />  
  • Then w/CrsTaken/@CrsCode = x/Class/@CrsCode, so the WHERE condition is satisfied & binding retained
XQuery Semantics (cont’d)

• Step 3: Construct result
  • For each retained tuple of bindings, instantiate the RETURN clause
  • This creates a fragment of the output document
  • Do this for each retained tuple of bindings in sequence
User-defined Functions

- Can define functions, even recursive ones
- Functions can be called from within an XQuery expression
- Body of function is an XQuery expression
- Result of expression is returned
  - Result can be a primitive data type (integer, string), an element, a list of elements, a list of arbitrary document nodes, …
XQuery Functions: Example

- Count the number of e-children recursively:

  DECLARE FUNCTION countNodes($e AS element()) AS integer {
    RETURN
    IF empty($e/*) THEN 0
    ELSE
      sum(FOR $n IN $e/* RETURN countNodes($n))
      + count($e/*)
  }

Built-in functions sum, count, empty
Class Rosters  (again) Using Functions

DECLARE FUNCTION extractClasses($e AS element()) AS element()* { 
   FOR $c IN $e//CrsTaken 
   RETURN <Class CrsCode={$c/@CrsCode} Semester={$c/@Semester} />
}

<Rosters>
   FOR $c IN distinct-values(FOR $d IN doc("transcript.xml") RETURN extractClasses($d) )
   RETURN 
   <ClassRoster CrsCode = {$c/@CrsCode} Semester = {$c/@Semester} > {
   LET $trs := doc("transcript.xml")
   FOR $t IN $trs//Transcript[CrsTaken/@CrsCode=$c/@CrsCode and CrsTaken/@Semester=$c/@Semester]
   RETURN $t/Student 
   ORDER BY $t/Student/@StudId 
   }
   </ClassRoster>
</Rosters>

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Converting Attributes to Elements with XQuery

- An XQuery reformulation of a previous XSLT query – much more straightforward (but ignores text nodes)

```xquery
DECLARE FUNCTION convertAttributes($a AS attribute()) AS element() {
    RETURN element {name($a)} {data($a)}
}
DECLARE FUNCTION convertElement($e AS node()) AS element() {
    RETURN element {name($e)} {
        FOR $a IN $e/@* RETURN convertAttribute($a)
    }
    IF empty($e/*) THEN $e/text() ELSE {
        FOR $n IN $e/* RETURN convertElement($n)
    }
}
RETURN convertElement(doc("my-document")/*)
```

The actual query: Just a RETURN statement!!
Integration with XML Schema and Namespaces

• Let type *sometype* be defined in \( \text{http://types.r.us/types.xsd} \):

  ```
  IMPORT SCHEMA namespace \text{tr} = \text{http://types.r.us}
  AT \text{http://types.r.us/types.xsd};
  
  DECLARE NAMESPACE \text{foo} = \text{http://foo.org/foo};
  
  DECLARE FUNCTION \text{local:doSomething}($x \text{ AS \text{tr:sometype}})
  AS \text{xs:string} 
  
  FOR $i \text{ IN \text{fn:doc}(...)};
  
  \ldots \ldots 
  
  RETURN
  
  <\text{foo:something}>
  
  \ldots \ldots 
  
  </\text{foo:something}>
  ```

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Grouping and Aggregation

- Does not use separate grouping operator
  - Recall that OQL does not need one either
  - Subqueries inside the RETURN clause obviate this need (like subqueries inside SELECT did so in OQL)
- Uses built-in aggregate functions count, avg, sum, etc. (some borrowed from XPath)
Aggregation Example

• Produce a list of students along with the number of courses each student took:

  FOR $t$ IN fn:doc(“transcripts.xml”)//Transcript,
    $s$ IN $t$/Student
  LET $c :=$t/CrsTaken
  RETURN
  <StudentSummary StudId = {$s/@StudId} Name = {$s/@Name} TotalCourses = {fn:count(fn:distinct-values($c))} />
  ORDER BY StudentSummary/@TotalCourses

• The grouping effect is achieved because $c$ is bound to a new set of nodes for each binding of $t$
Quantification in XQuery

- XQuery supports explicit quantification: SOME (\( \exists \)) and EVERY (\( \forall \))

- Example:

```xml
FOR $t$ IN fn:doc("transcript.xml")//Transcript
WHERE SOME $ct$ IN $t$/CrsTaken
    SATISFIES $ct$/@CrsCode = "MAT123"
RETURN $t$/Student
```

“Almost” equivalent to:

```xml
FOR $t$ IN fn:doc("transcript.xml")//Transcript,
    $ct$ IN $t$/CrsTaken
WHERE $ct$/@CrsCode = "MAT123"
RETURN $t$/Student
```

- Not equivalent, if students can take same course twice!
Implicit Quantification

- Note: in SQL, variables that occur in FROM, but not SELECT are implicitly quantified with $\exists$
- In XQuery, variables that occur in FOR, but not RETURN are similar to those in SQL. However:
  - In XQuery variables are bound to document nodes
    - Two nodes may look textually the same (e.g., two different instances of the same course element), but they are still different nodes and thus different variable bindings
    - Instantiations of the RETURN expression produced by binding variables to different nodes are output even if these instantiations are textually identical
  - In SQL a variable can be bound to the same value only once; identical tuples are not output twice (in theory)

- *This is why the two queries in the previous slide are not equivalent*
Quantification (cont’d)

- Retrieve all classes (from classes.xml) where each student took MAT123
  - Hard to do in SQL (before SQL-99) because of the lack of explicit quantification

For $c$ IN fn:doc(classes.xml)//Class
Let $g := \{
\text{Transcript records that correspond to class } c
\}
For $t$ IN fn:doc("transcript.xml")//Transcript
Where $t/CrsTaken/@Semester = c/@Semester
And $t/CrsTaken/@CrsCode = c/@CrsCode
Return $t
}
Where EVERY $tr$ IN $g$ SATISFIES

NOT fn:empty($tr[CrsTaken/@CrsCode="MAT123"])

RETURN $c$ ORDER BY $c/@CrsCode

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In the past, SQL was extended for O-O:

- added values for reference, tuple(row type), and collection(arrays), …
- took over ODL and OQL standards of ODMG

Currently, SQL is being extended for XML:

- adding data types and functions to handle XML
- will it bring the demise of XQuery?
Why SQL/XML

- Publish contents of SQL tables or entire DB as XML documents – need convention for translating primitive SQL data types
- Create XML documents out of SQL query results – need extension of SQL queries to create XML elements
- Store XML documents in relational DBs and query them – need extension of SQL to use XPath to access the elements of tree structures
Publishing Relations as XML Documents

- Current proposal:
  - no built-in functions to convert tables to XML
  - but can create arbitrary XML documents using extended SELECT statements

- Encoding relational data in XML:
  - Entire relation: an element named after the relation
  - Each row: an element named row
  - Each attribute: an element named after the attribute
### Publishing Relations as XML Doc: Tables

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>DeptId</th>
</tr>
</thead>
<tbody>
<tr>
<td>1024</td>
<td>Bob Smith</td>
<td>CS</td>
</tr>
<tr>
<td>3093</td>
<td>Amy Doe</td>
<td>EE</td>
</tr>
</tbody>
</table>

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Publishing Relations as XML Documents

**SQL:**

```
CREATE TABLE Professor

  Id: INTEGER,
  Name: CHAR(50),
  DeptId: CHAR(3)
```

**XML Schema:**

```
<schema xmlns="http://www.w3.org/2001/XMLSchema"
          targetNamespace="http://xyz.edu/Admin">
  <element name="Professor">
    <complexType>
      <sequence>
        <element name="row" minOccurs="0" maxOccurs="unbounded">
          <complexType>
            <sequence>
              <element name="Id" type="integer"/>
              <element name="Name" type="CHAR_50"/>
              <element name="DeptId" type="CHAR_3"/>
            </sequence>
          </complexType>
        </element>
      </sequence>
    </complexType>
  </element>
</schema>
```
• CHAR\_len: standard conventions in SQL/XML for CHAR(len) in SQL.
  • For instance, CHAR\_50 is defined as
    
    ```
    <simpleType>
      <restriction base="string">
        <length value="50">
      </restriction>
    </simpleType>
    ```

• A lot of the SQL/XML standard deals with such data conversion, and with user-defined types of XML, which are defined in SQL using CREATE DOMAIN.
Creating XML from Queries: Functions

**XMLELEMENT, XMLATTRIBUTES**

- An SQL query does not return XML directly. Produces *tables* that can have columns of type **XML**.

```sql
SELECT p.Id, XMLELEMENT (  
    NAME  "Prof",  -- element name
    XMLATTRIBUTES (P.DeptId AS "Dept"), -- attributes
    P.Name  -- content
  ) AS Info
FROM  Professor P
```

**Produces tuples of the form**

- 1024, <Prof Dept="CS">Bob Smith</Prof>
- 3093, <Prof Dept="EE">Amy Doe</Prof>
Creating XML Using Queries: Functions XMLELEMENT, XMLATTRIBUTES

- XMLELEMENT can be nested:

```sql
SELECT XMLELEMENT (NAME "Prof"
    XMLELEMENT(NAME "Id", P.Id),
    XMLELEMENT(NAME "Name", P.Name),
    XMLELEMENT(NAME "DeptId", P.DeptId),
) AS ProfElement
FROM Professor P
```

Produce tuples of the form

```
<Prof>
    <Id>1024</Id><Name>Bob Smith</Name><DeptId>CS</DeptId>
</Prof>
<Prof>
    <Id>3093</Id><Name>Amy Doe</Name><DeptId>EE</DeptId>
</Prof>
```
Creating XML Using Queries: Function XMLQUERY

SELECT XMLQUERY ('<Prof>
  <Id>{$I}</Id><Name>{$N}</Name><DeptId>{$D}</DeptId>
  </Prof>' -- template with placeholder variables
PASSING BY VALUE
  P.Id AS I,          -- values of I substitute for placeholders
  P.Name AS N,
  P.DeptId AS D
RETURNING SEQUENCE
  ) AS ProfElement
FROM Professor P

- Placeholder can occur in positions of XML elements and attributes
- Expressions can be XML-generating expressions or SELECT statements
  - In the example above, could have
    SELECT QUERY(''<Prof>
      {$I} <Name>{$N}</Name> …
    </Prof>'
    PASSING BY VALUE XMLELEMENT(NAME "Id", P.Id) AS I
    …
- In general, the argument to XMLQUERY can include any XQuery expression (XPath or a full query)
Creating XML from Queries: Grouping without GROUP BY

- In XQuery: group elements as children of another element by putting a subquery in RETURN clause of parent query.
- In SQL/XML: group by putting SELECT inside XMLELEMENT in the SELECT clause of parent.
- Example: group the CrsTaken by student Ids

```sql
SELECT XMLELEMENT ( 
    NAME "Student", 
    XMLELEMENT(NAME "CrsTaken", 
        XMLATTRIBUTES(T.CrsCode AS "CrsCode", 
            T.Semester AS "Semester")) 
FROM Transcript T 
WHERE S.Id=T.StudId) 
FROM Student S
```

Returns a set of 1-tuples, not list of elements. Waiting for the standard to resolve how to convert.
Creating XML from Queries: Grouping and XMLAGG

• Same example: group CrsTaken by student ids

```sql
SELECT xmlelement (  
    name "Student",  
    xmlattributes(s.id AS "Id"),  
    xmlagg(xmlelement(name "CrsTaken",  
        xmlattributes(t.crscode AS "CrsCode",  
            t.semester AS "Semester") )  
    )  
ORDER BY t.crscode) )  
FROM Student S, Transcript T  
WHERE S.Id = T.StudId  
GROUP BY S.Id
```
Storing XML in Relational DB:
Data Type XML

- Not stored as a string, but natively as a tree structure. Supports navigation via efficient storage and indexing.

    CREATE TABLE StudentXML ( 
        Id INTEGER, 
        Details XML 
    )

where Details attribute contains things of the form

    <Student>
    <Name><First>Amy</First><Last>Doe</Last></Name>
    <Status>U4</Status>
    <CrsTaken CrsCode="305" Semester="F2003"/>
    <CrsTaken CrsCode="336" Semester="F2003"/>
    </Student>
To validate, use

```sql
CREATE TABLE StudentXML (  
    Id INTEGER,  
    Details XML,  
    CHECK(Details IS VALID ACCORDING TO SCHEMA 'http://xyz.edu/student.xsd')  
)
```

assuming the schema is stored at http://xyz.edu/student.xsd
Querying XML Stored in Relations: XMLEXISTS

- Tells whether the set of nodes returned by XPath expression is empty.
- Example: return Ids and names of students who have taken a course

```sql
SELECT S.Id, XMLEXTRACT(S.D details, '//Name')
FROM StudentXML S
WHERE XMLEXISTS(XMLQUERY('D//CrsTaken'
    PASSING BY REF S.D details AS D
    RETURNING SEQUENCE))
```
Querying XML Stored in Relations (cont’d)

- Use XQuery expressions and XMLEXISTS function.
- XMLQUERY can be both in SELECT and WHERE clauses.
- Example: return Ids and names of students who have status U3 and took MAT123:

```sql
SELECT S.Id, XMLQUERY(S.Details, '$D//Name'  
  PASSING BY REF S.Details AS D  
  RETURNING SEQUENCE)  
FROM StudentXML S  
WHERE XMLEXISTS(XMLQUERY(  
  'WHERE $D//Status/text() = "U3" AND  
  $D//CrsTaken/@CrsCode = "MAT124"  
  RETURN $D'  
  PASSING BY REF S.Details AS D  
  RETURNING SEQUENCE ))
```

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XML stored as appropriately indexed tree structure, but in SQL is specified as a sequence of characters – so need to parse:

```
INSERT INTO StudentXML(Id, Details)
VALUES(12343,
    XMLPARSE(''<Student>'
        '<Name>'<First>Bob</First><Last>Smith</Last>'
        '<Status>U4</Status>
        <CrsTake CrsCode="CS305" Semester="F2003"/>
        <CrsTake CrsCode="CS339" Semester="S2004"/>
    </Student>' )
```

To validate inserted document:

```
INSERT INTO StudentXML(Id, Details)
VALUES(12343,
    XMLPARSE( 'Student>
        <Name><First>Bob</First><Last>Smith</Last></Name>
        <Status>U4</Status>
        <CrsTake CrsCode="CS305" Semester="F2003"/>
        <CrsTake CrsCode="CS339" Semester="S2004"/>
    </Student>' )
IS VALID ACCORDING TO SCHEMA 'http://xyz.edu/Students.xsd'
```
**XMLSERIALIZE: Reverse of XMLPARSE**

- To convert XML back to a string.
  - Typically used to talk to a host language that does not understand XML
- **XMLSERIALIZE** is often used in embedded SQL in conjunction with *cursors*
  - Example: return Ids and names of professors. Professors’ names are returned as ‘<Prof>Joe</Prof>’.

  ```sql
  EXEC SQL DECLARE GetProfessor CURSOR FOR
  SELECT P.Id, XMLSERIALIZE(XMLELEMENT(Name “Prof”, P.Name))
  FROM Professor P
  ```

  This can then be processed by

  ```sql
  EXEC SQL GetProfessor INTO :profId, :name
  ```