Tabled Resolution

CSE 505 – Computing with Logic Stony Brook University <u>http://www.cs.stonybrook.edu/~cse505</u>

Recap: OLD Resolution

- Prolog follows *OLD resolution*, which is SLD (*Selective Linear Definite*) resolution, but with with left-to-right literal selection & Prolog also uses the order in which the clauses are enumerated in the database to determine the order in which the branches of the search space are investigated
 - Consider a goal G₁, G₂,..., G_n as a "procedure stack" with G₁, the selected literal on top
 - Call **G**₁
 - If and when G₁ returns, continue with the rest of the computation: call G₂, and upon its return call G₃, etc. until nothing is left
 - Note: G₂ is "opened up" only when G₁ returns, not after executing only some part of G₁

OLD Resolution

Depth-first expansion, however, contributes to the *incompleteness* of Prolog's evaluation, which may **not terminate** even when the least model is finite (see the next example!)

Example:Reachability in Directed Graphs

- Determining <u>whether there is</u> a <u>path between two vertices</u> in a directed graph is an important and widespread problem
- For instance, consider checking whether (or not) a program accesses a shared resource before obtaining a lock:
 - A program itself can be considered as a graph with vertices representing program states!
 - A state may be characterized by the program counter value, and values of variables
 - There are richer models for representing program evaluation, but a directed graph is most basic
 - If we can go from state s by executing one <u>instruction</u> to s', then we can place an edge from s to s'
 - The <u>reachability question</u> may be whether we <u>can reach from the start</u> <u>state to a state accessing a shared resource</u>, <u>without going through a state</u> <u>that obtained a lock</u>

Graph Reachability as a Logic Program

- A finite directed graph can be represented by a set of binary facts representing an "<u>edge</u>" relation
 - Predicate "**q**" on right is an example
- Reachability can then be written as a "*transitive closure*" over the edge relation
 - Observe the predicate "**r**" defined on right using two clauses:
 - The first clause: there is a *path* from **X** to **Y** if there is an edge from **X** to **Y**
 - The second clause: there is a *path* from **X** to **Y** if there an intermediate vertex **Z** such that:
 - there is an edge from **X** to **Z**, and
 - there is a path from **Z** to **Y**



q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).

r(X,Y) := q(X,Y).r(X,Y) :=q(X,Z), r(Z,Y).

Bottom-Up Evaluation

- Note that the program on the right is a *Datalog* program, i.e., no function symbols
 - Its Herbrand Universe is **finite**, and its least model computation using the bottom up evaluation will **terminate**:

$$\begin{split} M_0 &= \emptyset \\ M_1 &= T_p(M_0) = M_0 \cup \{q(a,a), q(a,b), q(a,c), q(b,b), q(b,c), q(b,b), q(b,d), q(b,c), q(c,d) \} \\ M_2 &= T_p(M_1) = M_1 \cup \{r(a,a), r(a,b), r(a,c), r(b,b), r(a,c), r(b,b), r(b,d), r(b,c), r(c,d) \} \\ M_3 &= T_p(M_2) = M_2 \cup \{r(a,d), r(a,c), r(a,c) \} \\ M_4 &= T_p(M_3) = M_3 = M_p \end{split}$$



Bottom-Up Evaluation

$$M_{p} = \{q(a,a), q(a,b), q(a,c), q(b,b), q(b,d), q(b,e), q(b,d), q(b,e), q(e,d), r(a,a), r(a,b), r(a,c), r(b,b), r(b,d), r(a,c), r(b,e), r(e,d), r(a,d), r(a,e)\}$$



q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).

r(X,Y) := q(X,Y).r(X,Y) :=q(X,Z), r(Z,Y).

OLD Resolution with Depth-First Expansion ?- r(a,N). а Consider initial query "?- r(a,N)." Let's construct the **<u>SLD tree</u>** for this query С b е d q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d). r(X,Y) := q(X,Y).r(X,Y) :=q(X,Z), r(Z,Y).8



Resolving this goal with the first clause of **r**, we get a new goal "?- q(a,N)."



q(X,Z), r(Z,Y).

r(X,Y) :=



Resolving "?- q(a,N)." results in the empty goal, under three answer substitutions: a, b, and c. There are no more ways to resolve "?- q(a,N)."



r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



Resolving "?- r(a,N)." with the second clause of r results in goal "?- q(a,Z), r(Z,N)."

Note: We will use same variable names as in the program clause when possible, instead of mechanically inventing new variable names in every step.

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r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



The selected literal is q(a,Z) which unifies with fact q(a,a) with Z=a. Thus we get the goal "?-r(a,N)."

Note: This is the same goal that we had at the beginning.











Depth-First Expansion of the OLD tree

- If the underlying graph is acyclic, all branches in the OLD tree will be finite
- 2. If the graph is cyclic, nothing to the right of an infinite branch is expanded
 - This renders the evaluation <u>incomplete</u>: goals for which <u>there</u> <u>are OLD derivations</u>, but they are not found
- Moreover, in both cases, the same answer may be returned multiple times (even <u>infinitely</u>!)
 - Even if the underlying graph is acyclic, this evaluation is not efficient
 - For our example, for the query "?- r(a,N)." we will return
 N=b for each path from "a" to "b"

Expansion of the OLD tree

- <u>Breadth-First</u> expansion does have the completeness property for definite programs:
 - Every OLD derivation will be eventually constructed
 - If something is a logical consequence, we will eventually confirm it in a <u>finite</u> number of levels
 - But we may not be able to conclude *negative information*
 - If something is not a logical consequence, we may never be able to identify it because we don't know when to stop?

Expansion of the OLD tree

- Moreover, Breadth-First expansion does not give a natural operational understanding:
 - If we view predicates as being defined by "procedures", then breadth first expansion switches contexts at the end of each step
 - As in procedural programming, <u>context switching</u> <u>is expensive</u> (in this case, we've to <u>switch</u> <u>substitutions</u>)

Programming our way around the problem

Check for any vertex if it was not already seen in the path **L** • Start from **L=[]** to look for reachable vertices r(X, Y) :p([], X, Y). p(L, X, Y) :q(X, Y), not member(Y, L). p(L, X, Y) :q(X, Z), not member(Z, L), p([Z|L], Z, Y).• In **L** we remember the path so far, and use this to avoid loops

Programming our way around the problem

- We are assured termination for reachability queries
 - We stop if a node has been seen before on the same branch
- Still, this is inefficient
 - We re-execute queries on different branches of the SLD/OLD tree
 - May take exponential time

What is *Tabled Resolution*? *Memoize* calls and results to avoid repeated subcomputations

- Properties:
 - Termination: Avoid performing computations that repeat infinitely often
 Complete for Datalog programs
 Efficiency: Dynamically share common subexpressions



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23















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30












Depth-First Expansion of OLD tree with a little twist: Stop if a goal has been seen before.



Depth-First Expansion of OLD tree with a little twist: Stop if a goal has been seen before.



Depth-First Expansion of OLD tree with a little twist: Stop if a goal has been seen before.





40

Rationale for goal-based stopping

- The OLD tree is a representation of search for successful derivations
 - which are finite sequences of goals terminating in an empty goal
- If there is a successful derivation, then there is an equivalent one that does not repeat the same goal (compare to reachability via loop-free paths in a graph)
- Hence ignoring paths with repeated goals is *sound*: the derivations pruned away by stopping have equivalent ones that will not be ignored
- Unfortunately, this scheme still does not fix the problem of infinite derivations



q(a, a). q(a, b).

q(a, c).

q(b, b).

q(b, d).

q(b, e).

q(e, d).

%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :p(X,Z), q(Z,Y). ?- p(a, N)



q(a, a). q(a, b). q(a, c).

- q(a, c). q(b, b).
- q(b, d).

q(b, e).

q(e, d).

%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :p(X,Z), q(Z,Y).



Expand tree as usual



- q(a, a). q(a, b).
- q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :p(X,Z), q(Z,Y).



Expand tree as usual



- q(a, a). q(a, b).
- q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :p(X,Z), q(Z,Y).



Note that the right-most branch has ever-growing goals!

OLD Resolution with Tabling (OLDT)

- The selected <u>literal</u> at a step in a derivation is known as a call (only individual literals are tabled)
- OLDT maintains a <u>table of calls</u> (initially empty)
- With each call, it maintains a table of computed answers (initially empty)
- Start resolution as in OLD
- When a literal is selected, check the call table.
 - If the literal is in the table, <u>resolve it with its answers</u> <u>in its answer table</u>
 - If the literal is not in the table, resolve with program clauses (as in OLD), and <u>add computed answers to</u>

its answer table

OLDT	Example
------	---------



- q(a, a). q(a, b).
- q(a, c).
- q(a, c). q(b, b).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

:- table p/2.

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).

Calls	Answers

Start with empty tables

OLDT Examp	ole	
$ \begin{array}{c} $?- p(a, N)
q(b, e). q(e, d).		
:- table p/2.	Calls	Answers
% of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y).		
p(X,Y) :- p(X,Z), q(Z,Y).	Pick sele	ected literal. Is it in call table?

OLDT Exampl	е		
		?- p(a, N)	
q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).			
:- table p/2. %Alternative formulation % of reachability: Note	Calls p(a, W)	Answers	
<pre>% use of LEF 1 recursion p(X,Y) :- q(X,Y). p(X,Y) :- p(X,Z), q(Z,Y).</pre>			
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OLDT Exam	ple		
	?- q(a, N) <mark>N=a</mark>	?- p(a, N)	
q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d). :- table p/2.	Calls p(a, W)	Answers {p(a,a)	-
<pre>%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :- p(X,Z), q(Z,Y).</pre>	Add com there)	puted answer to table (if not already

OLDT Exam	ple	
	?- q(a, N)	?- p(a, N)
q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d). :- table p/2.	Calls p(a, W)	Answers {p(a,a), p(a,b), p(a,c)
<pre>%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :- p(X,Z), q(Z,Y).</pre>	Add com there)	puted answer to table (if not already















OLDT Forest

- When we get new answer, we will have to return to previous queries to continue their execution
- When a literal is selected, mark it as a *consumer*
- Check the call table:
 - If the literal is not in the table, **start a new tree for that literal** with its root marked as *generator*
 - **Resolve generator with program clauses** (as in OLD), and add computed answers to its answer table
- Resolve consumer with answers in its generator's table

Calls and answers in tables

- Calls in table are standardized apart: i.e. their variables are renamed so that they are not identical to any other variable
- Answers in a call's computed answer table share variables with their call
- When checking if a literal **1** is in call table:
 - We can check for *variance*: for a call **C** that is identical to the given literal **1**, modulo names of variables

• All answers to **C** are answers to **1**, and vice versa

- We can check for *subsumption*: for a call **c** that is more general than a given literal **1**, . i.e. if there is a substitution $\boldsymbol{\theta}$ such that $\mathbf{c}\boldsymbol{\theta} = \mathbf{1}$
 - Not all answers to ${f c}$ may be answers to ${f l}$, but every answer to ${f l}$ is an answer to ${f c}$

Notes on OLDT

- We can selectively mark which predicates we want to maintain tables for (e.g. "p" in the previous example)
 - In general, no need to maintain tables for predicates defined solely by facts (i.e. clauses with empty bodies)
- For a Datalog program, there can be only finitely many distinct calls and answers
 - So the size of tables is bounded
- The number of literals in each goal is limited by the largest clause in the program (or original goal)
- Hence for Datalog, the OLDT forest as well as table sizes are bounded



q(a,	a).	
q(a,	b).	

- q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

- Construct a forest: one tree for each call
- Root of each tree (blue) is a **generator**
- Selected literal that matches a tabled call (green) is a **consumer**





- q(a, a). q(a, b).
- q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).





- q(a, c).
- q(b, b).
- q(b, d).

q(b, e).

q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).





%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).



q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
%ORIGINAL formulation
% of reachability: Note

% of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).





q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).

% ORIGINAL formulation % of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).







- q(b, b). q(b, d).
- q(b, e).

q(e, d).

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



r(a, N)

?- q(a, Z), r(Z, N)

?- r(b, N)

?-

?- q(a, N)

?- r(a, N)

N=c

N=a

N=a N=b

N=b N=c

?- r(b, N)





- q(b, d).
- q(b, e).
- q(e, d).

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

r(a, N)

?-q(a, Z), r(Z, N) ?-q(b, N)

?- r(b, N)

?- r(b, N)

?-

?- q(a, N)

N=a N=b N=c

?- r(a, N)

N=c

N=a N=b





- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).





- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).




- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).





- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).





- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).





- q(a, a). q(a, b). q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).





- q(a, c).
- q(b, b).
- q(b, d).
- q(b, e).
- q(e, d).









- q(a, a). q(a, b). q(a, c).
- q(a, c)
- q(b, b). q(b, d).
- q(b, d). q(b, e).
- q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).



% of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).

















%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).







q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).











%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).







q(a, a). q(a, b). q(a, c). q(b, b). q(b, d). q(b, e). q(e, d).



?- r(d, N)

%ORIGINAL formulation % of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



%ORIGINAL formulation % of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```





%ORIGINAL formulation % of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



% use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).



%ORIGINAL formulation % of reachability: Note % use of RIGHT recursion r(X,Y) :- q(X,Y). r(X,Y) :q(X,Z), r(Z,Y).





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r(X,Y) := q(X,Y).

q(X,Z), r(Z,Y).

r(X,Y) :=



 OLD Resolution with Tabling (OLDT)
 OLDT evaluation can be used to infer <u>negative</u> <u>answers</u>: e.g. a vertex is not reachable from another

• Note that Breadth-First evaluation, or even the evaluation with goal-based stopping condition cannot do this

```
Tabled Resolution in XSB
 edge(a,a).
 edge(a,b).
 edge(b,c).
 :- table(reach/2).
 %OR :- auto table.
 reach(X,Y) := edge(X,Y).
 reach(X,Y) := reach(X,Z), edge(Z,Y).
 • Call:
```

```
?- reach(a,V).
```







Tabled Resolution



Tabled Resolution



Tabled Resolution



