

## CSE303 Q2 SOLUTIONS Spring 2018

### PART 1: YES/NO QUESTIONS

1. The set  $F$  of final states of any deterministic finite automaton is always non-empty.  
**Justify:** the definition says that  $F$  is a finite set, i.e. can be empty, hence for some  $M$ ,  $L(M) = \emptyset$ . **n**
2. A configuration of a deterministic finite automaton  $M = (K, \Sigma, \delta, s, F)$  is any element of  $K \times \Sigma^* \times K$   
**Justify:** it is any element of  $K \times \Sigma^*$  **n**
3. Given an automaton  $M = (K, \Sigma, \delta, s, F)$ , a binary relation  $\vdash_M \subseteq (K \times \Sigma^*) \times (K \times \Sigma^*)$  is a **transition relation** iff the following condition holds  
 $(q, aw) \vdash_M (q', w)$  iff  $\delta(q', a) = q$ .  
**Justify:** Proper condition is:  $(q, aw) \vdash_M (q', w)$  iff  $\delta(q, a) = q'$  **n**
4. For any  $M = (K, \Sigma, \delta, s, F)$ ,  $L(M) \neq \emptyset$   
**Justify:** for any  $M$ , such that  $F = \emptyset$  we have that  $L(M) = \emptyset$  **n**
5. If  $M = (K, \Sigma, \Delta, s, F)$  is a non-deterministic as defined in the book, then  $M$  is also non-deterministic, as defined in the lecture.  
**Justify:**  $\Sigma \cup \{e\} \subseteq \Sigma^*$  **y**

### PART 2: Few Very Short Questions

**Q1:** M1 has components:  $K = \{q\}$ ,  $s = q$ ,  $\Sigma = \emptyset$ ,  $\delta = \emptyset$ ,  $F = \emptyset$ .

#### Solution

1. M1 is deterministic
2.  $L(M1) = \emptyset$

**Q2:** M2 has components:  $K = \{q_0, q_1, q_2\}$ ,  $s = q_0$ ,  $\Sigma = \{a, b\}$ ,  $F = \{q_1\}$   
 $\delta = \{(q_0, a, q_1), (q_1, a, q_1), (q_0, b, q_2)\}$

#### Solution

1. M2 is non-deterministic;  $\delta$  is not a function with the domain  $K \times \Sigma$ . It can be completed to a function by adding some trap states. But the trap states information was not stated in the problem - so **M2** is N DFA
2.  $L(M2) = aa^*$

**Q3** M3 has components:  $K = \{q_0, q_1, q_2\}$ ,  $\Sigma = \{a, b\}$ ,  $F = \{q_1, q_2\}$   
 $\Delta = \{(q_0, a, q_1), (q_0, b, q_2), (q_1, a, q_1), (q_1, b, q_2), (q_2, ab, q_2)\}$

#### Solution

M3 is NOT an automaton. It does not have the INITIAL state!

### PART 3: PROBLEMS

**QUESTION 1** The components of  $M$  represent a deterministic finite automaton  $M$  with some trap states.

**Components of  $M$  are:**

$$K = \{q_0, q_1, q_2, q_3\}, s = q_0, \Sigma = \{a, b\}, F = \{q_1, q_3\}$$
$$\delta = \{(q_0, a, q_1), (q_1, b, q_2), (q_2, a, q_2), (q_2, b, q_3)\}$$

**Complete** the components to a full definition of a deterministic  $M$  by adding trap state(s)

**Solution** We add a new state  $q_4$  and extend  $\delta$  to function  $\delta_1$  such that

$$\delta_1 : (K \cup \{q_4\}) \times \Sigma \rightarrow K \cup \{q_4\} \text{ as follows}$$

$$\delta_1 = \delta \cup \{(q_0, b, q_4), (q_1, a, q_4), (q_3, a, q_4), (q_3, b, q_4)\}$$

**Write** a regular expression or a property defining  $L(M)$ .

**Solution**  $L(M) = a \cup aba^*b$

**QUESTION 2** Construct a non-deterministic finite automaton  $M$ , such that

$$L(M) = (ba \cup b)^* \cup (bb \cup a)^*.$$

**Solution** Some elements of  $L(M)$  are:  $bab, bba, babb, bbaa, bbaabb, babbbba, \dots$

**Components of  $M$  are:**

$$K = \{q_0, q_1, q_2\}, s = q_0, \Sigma = \{a, b\},$$
$$\Delta = \{(q_0, e, q_1), (q_0, e, q_2), (q_1, ba, q_1), (q_1, b, q_1), (q_2, bb, q_2), (q_2, a, q_2)\},$$
$$F = \{q_1, q_2\}$$

**QUESTION 3**

Let  $M = (K, \Sigma, s, \Delta, F)$  for  $K = \{q_0, q_1, q_2\}$ ,  $s = q_0$ ,  $\Sigma = \{a, b, c\}$ ,  $F = \{q_1, q_2\}$  and  $\Delta = \{(q_0, abc, q_0), (q_0, ab, q_1), (q_0, b, q_2)\}$

**Draw the diagram** an automaton  $M'$  such that  $M' \equiv M$  and  $M'$  is defined by the BOOK definition.

**Solution**

We apply the "stretching" technique to  $M$  and the new  $M'$  is as follows.

$$M' = (K \cup \{p_1, p_2, p_3\}, \Sigma, s = q_0, \Delta', F = F)$$

$$\Delta' = \{(q_0, b, q_2), (q_0, a, p_1), (p_1, b, p_2), (p_2, c, q_0), (q_0, b, q_2), (q_0, a, p_3), (p_3, b, q_1)\}$$