Untyped Lambda Calculus

Principles of Programming Languages

CSE 526

1 Syntax
2 Variables and Substitution
3 Reductions
4 Nameless Representation
Lambda Calculus

- A formal notation to study computability and programming.
- Can be considered as the smallest universal programming language.
  - **Universal**: Can be used to express any computation that can be performed on a Turing Machine
  - **Small**: Has only two constructs: abstraction and application.
- Brief History:
  - Introduced by Church and Kleene in 1930s.
  - Used by Church to study problems in computability.
  - Concepts have heavily influenced *functional programming*.
  - Used to study *types* and type systems in programming languages.
Lambda Terms

Syntax of the $\lambda$-calculus

$$t ::= \text{Terms}$$
Lambda Terms

Syntax of the λ-calculus

\[ t ::= x \quad \text{Variable} \]

\[ t \quad \text{Terms} \]

\[ \lambda x. t \quad \text{Abstraction} \]

\[ t \; t \quad \text{Application} \]
Lambda Terms

Syntax of the $\lambda$-calculus

\[
t ::= \begin{array}{ll}
  \text{Terms} & \text{Variable} \\
  x & \lambda x. t
\end{array}
\]
Lambda Terms

Syntax of the $\lambda$-calculus

$$t ::= \begin{align*}
\ & x & \text{Variable} \\
| \ & \lambda x. \ t & \text{Abstraction} \\
| \ & t \ t & \text{Application}
\end{align*}$$
Lambda Terms

Syntax of the $\lambda$-calculus

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t ::= \begin{array}{ll}
  x & \text{Variable} \\
  \lambda x. t & \text{Abstraction} \\
  t \ t & \text{Application}
\end{array}
\]

Textual Representation:

Use parentheses to represent trees using linear text
Informal Semantics

$\lambda$-expressions can be considered as expressions in a functional language

- **Abstraction**: $(\lambda x. \ t)$ is a “function” with formal parameter $x$ that returns (the value of) term $t$. 

Example 1: $(\lambda x. \ x)$ is the identity function: one that returns the argument value itself.

Example 2: $(\lambda x. \ (\lambda y. \ x))$ is a function that takes two arguments $x$ and $y$ and returns the first argument.

Application: $(t_1 \ t_2)$ is a “function call” where $t_1$ is a function and $t_2$ is the supplied argument.

Example: $((\lambda x. \ x) \ y)$ supplies $y$ as the argument to the identity function.
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  - Example: \(((\lambda x. \ x) \ y)\) supplies \(y\) as the argument to the identity function.
Parentheses can be dropped using the following conventions:

- application is left associative
  e.g. \((f \ f) \ x\) is same as \(f \ f \ x\)
- a \(\lambda\) binds as much as possible to its right.
  e.g. \(\lambda f. \lambda x. f \ (f \ x)\) is same as \(\lambda f.(\lambda x. f \ (f \ x))\)

Multiple consecutive abstractions can be combined:

  e.g. \(\lambda f.\lambda x. f \ (f \ x)\) is same as \(\lambda f \ x. f \ (f \ x)\)
The Meaning of Lambda Expressions

- Recall: $\lambda x. t$ stands for a function with $x$ as the parameter and (the value of) $t$ as the return value.
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- Example: Consider the expression

$$(((\lambda wyx. y (w y x)) \ (\lambda sz. z))$$

This is an instance of an application. The expression in blue is passed as an argument to the function in red.
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  1. $\lambda yx. y ((\lambda sz. z) y x)$
  2. $\lambda yx. y ((\lambda z. z) x)$
The Meaning of Lambda Expressions

- Recall: $\lambda x. \ t$ stands for a function with $x$ as the parameter and (the value of) $t$ as the return value.
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This is an instance of an application. The expression in blue is passed as an argument to the function in red.

- The meaning of an application: replace every occurrence of the formal parameter in the body of the function with the given argument.

In the above example

1. $\lambda \ yx. \ y \ ((\lambda \ sz. \ z) \ y \ x)$
2. $\lambda \ yx. \ y \ ((\lambda \ z. \ z) \ x)$
3. $\lambda \ yx. \ y \ x$
Encoding Booleans in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th>B</th>
<th>$\lambda$-calculus</th>
</tr>
</thead>
<tbody>
<tr>
<td>true</td>
<td>$\lambda x. \lambda y. x$</td>
</tr>
<tr>
<td>false</td>
<td>$\lambda x. \lambda y. y$</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>$\lambda x. \lambda y. ((x \ y) \ false)$</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td>!</td>
<td>$\lambda x. ((x \ false) \ true)$</td>
</tr>
<tr>
<td>if</td>
<td>$\lambda c. \lambda t. \lambda e. ((c \ t) \ e)$</td>
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</table>

This is known as the
Church encoding of Booleans,
or simply $Church Booleans$. 
## Encoding Booleans in the λ-Calculus

<table>
<thead>
<tr>
<th>Boolean</th>
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<tbody>
<tr>
<td>true</td>
<td>( \lambda x. \lambda y. x )</td>
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Example:

\[
(\text{true} \&\& \text{false})
\]

This is evaluated as:

\[
\begin{align*}
(\lambda x. \lambda y. x) \ (\lambda x. \lambda y. y) \\
\big( (\lambda x. \lambda y. ((x \ y) \text{false}) \big) \ (\lambda x. \lambda y. (x \text{true}) \ y) \\
\big( (\lambda x. ((x \text{false}) \text{true}) \big) \text{false} \\
(\lambda x. \lambda y. y) \\
\text{false}
\end{align*}
\]
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Example:
$$(true \land false)$$

$$\equiv (\lambda x. \lambda y. ((x \ y) \ false)))$$

$$(\lambda x. \lambda y. \ x)$$

$$(\lambda x. \lambda y. \ y)$$

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Example:

$(true \ && \ false)$

$\equiv (\lambda x. \lambda y. ((x \ y) \ false))$

$(\lambda x. \lambda y. x)$

$(\lambda x. \lambda y. y)$

$\rightarrow (\lambda y. (((\lambda x. \lambda y. x) \ y) \ false))$

$(\lambda x. \lambda y. y)$
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$$(true \&\& \ false)$$

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$$(\lambda x. \lambda y. x) \quad (\lambda x. \lambda y. y)$$

$$\rightarrow \quad (\lambda y. (((\lambda x. \lambda y. x) \ y) \ false))$$

$$(\lambda x. \lambda y. y)$$

$$\rightarrow \quad (((\lambda x. \lambda y. x) \ (\lambda x. \lambda y. y)) \ false)$$

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Example:

$$(\text{true} \&\& \text{false})$$

\[\equiv (\lambda x. \lambda y. ((x \ y) \text{false})) \quad (\lambda x. \lambda y. x) \quad (\lambda x. \lambda y. y)\]

\[\rightarrow (\lambda y. (((\lambda x. \lambda y. x) \ y) \text{false})) \quad (\lambda x. \lambda y. y)\]

\[\rightarrow ((\lambda x. \lambda y. x) (\lambda x. \lambda y. y) \text{false})\]

\[\rightarrow ((\lambda y. (\lambda x. \lambda y. y)) \text{false})\]
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Example:

$$(true \ &\ & \ false)$$

\[\equiv (\lambda x. \lambda y. ((x \ y) \ false))) (\lambda x. \lambda y. x) (\lambda x. \lambda y. y)\]

\[\rightarrow (\lambda y. (((\lambda x. \lambda y. x) \ y) \ false))) (\lambda x. \lambda y. y)\]

\[\rightarrow (((\lambda x. \lambda y. x) \ (\lambda x. \lambda y. y)) \ false)\]

\[\rightarrow (\lambda y. ((\lambda x. \lambda y. y)) \ false)\]

\[\rightarrow (\lambda x. \lambda y. y)\]

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Example:

$(true \ && \ false) 
\equiv \ ((\lambda x. \lambda y. ((x \ y) \ false)) \ ((\lambda x. \lambda y. x) \ ((\lambda x. \lambda y. y) \ false))) 
\rightarrow \ ((\lambda y. (((\lambda x. \lambda y. x) \ y) \ false)) \ ((\lambda x. \lambda y. y) \ false)) 
\rightarrow \ (((\lambda x. \lambda y. x) \ ((\lambda x. \lambda y. y)) \ false)) \ false) 
\rightarrow \ ((\lambda y. (\lambda x. \lambda y. y)) \ false) \ false) 
\rightarrow \ (\lambda x. \lambda y. y) \ false) 
\equiv \ false$
## Encoding Natural Numbers in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th>$N$</th>
<th>$\lambda$-calculus</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>$\lambda s. \lambda z \cdot z$</td>
</tr>
<tr>
<td>1</td>
<td>$\lambda s. \lambda z \cdot (s \ z)$</td>
</tr>
<tr>
<td>2</td>
<td>$\lambda s. \lambda z \cdot (s \ (s \ z))$</td>
</tr>
<tr>
<td>3</td>
<td>$\lambda s. \lambda z \cdot (s \ (s \ (s \ z)))$</td>
</tr>
</tbody>
</table>

\[ \vdots \]

$inc \quad \lambda n. \lambda s. \lambda z \cdot (s \ ((n \ s) \ z))$

$plus \quad \lambda m. \lambda n. \lambda s. \lambda z \cdot ((m \ s) \ ((n \ s) \ z))$

$times \quad \lambda m. \lambda n. \ ((m \ (plus \ n)) \ 0)$

$iszero \quad \lambda m. \ ((m \ (\lambda x. \ false)) \ true)$

\[ \vdots \]

This is known as the **Church encoding of Naturals**, or simply **Church Numerals**.
Encoding Data Structures in the $\lambda$-Calculus

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pair</code></td>
<td>$\lambda f. \lambda s. \lambda c. ((c f) s)$</td>
</tr>
<tr>
<td><code>fst</code></td>
<td>$\lambda p. (p \text{ true})$</td>
</tr>
<tr>
<td><code>snd</code></td>
<td>$\lambda p. (p \text{ false})$</td>
</tr>
</tbody>
</table>
Encoding Data Structures in the $\lambda$-Calculus

<table>
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</table>

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$\textit{pair} \varphi_1 \varphi_2$
Encoding Data Structures in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th></th>
<th>$\lambda f. \lambda s. \lambda c. ((c \ f) \ s)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$fst$</td>
<td>$\lambda p. (p \ true)$</td>
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<td>$snd$</td>
<td>$\lambda p. (p \ false)$</td>
</tr>
</tbody>
</table>

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$$pair \ \varphi_1 \ \varphi_2$$

$$\equiv \ ( (\lambda f. \lambda s. \lambda c. ((c \ f) \ s) \ \varphi_1) \ \varphi_2)$$
Encoding Data Structures in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th>Function</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>pair</td>
<td>$\lambda f. \lambda s. \lambda c. \left((c , f) , s\right)$</td>
</tr>
<tr>
<td>fst</td>
<td>$\lambda p. \left(p , \text{true}\right)$</td>
</tr>
<tr>
<td>snd</td>
<td>$\lambda p. \left(p , \text{false}\right)$</td>
</tr>
</tbody>
</table>

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$\text{pair } \varphi_1 \, \varphi_2$

$\equiv \left((\lambda f. \lambda s. \lambda c. \left((c \, f) \, s\right) \, \varphi_1) \, \varphi_2\right)$

$\rightarrow^* \lambda c. \left((c \, \varphi_1) \, \varphi_2\right)$
Encoding Data Structures in the $\lambda$-Calculus

| $\text{pair}$ | $\lambda f. \lambda s. \lambda c. ((c f) s)$ |
| $\text{fst}$ | $\lambda p. (p \text{true})$ |
| $\text{snd}$ | $\lambda p. (p \text{false})$ |

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$$\text{pair} \varphi_1 \varphi_2 \equiv (\lambda f. \lambda s. \lambda c. ((c f) s) \varphi_1) \varphi_2 \rightarrow^* \lambda c. ((c \varphi_1) \varphi_2)$$

$$\text{fst} \ (\text{pair} \ \varphi_1 \ \varphi_2)$$
Encoding Data Structures in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th>Syntax</th>
<th>Variables and Substitution</th>
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</table>

| pair       | $\lambda f. \lambda s. \lambda c. ((c \ f) \ s)$ |
| fst        | $\lambda p. (p \ true)$ |
| snd        | $\lambda p. (p \ false)$ |

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$$
\text{pair} \ \varphi_1 \ \varphi_2
\equiv \ (\lambda f. \lambda s. \lambda c. ((c \ f) \ s) \ \varphi_1) \ \varphi_2
\rightarrow^* \ \lambda c. ((c \ \varphi_1) \ \varphi_2)
$$

$$
\text{fst} \ (\text{pair} \ \varphi_1 \ \varphi_2)
\equiv \ (\lambda p. (p \ true)) \ (\text{pair} \ \varphi_1 \ \varphi_2)
$$
Encoding Data Structures in the $\lambda$-Calculus

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<td>$\lambda p. (p, false)$</td>
</tr>
</tbody>
</table>

Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

$$pair\; \varphi_1\; \varphi_2$$

$$\equiv (\lambda f. \lambda s. \lambda c. ((c\, f)\, s)\, \varphi_1)\, \varphi_2)$$

$$\rightarrow^* \lambda c. ((c\, \varphi_1)\, \varphi_2)$$

$$fst\; (pair\; \varphi_1\; \varphi_2)$$

$$\equiv (\lambda p. (p\, true))\, (pair\; \varphi_1\; \varphi_2)$$

$$\rightarrow (pair\; \varphi_1\; \varphi_2)\, true$$
Encoding Data Structures in the $\lambda$-Calculus

<table>
<thead>
<tr>
<th>pair</th>
<th>$\lambda f. \lambda s. \lambda c. ((c\ f)\ s)$</th>
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<tbody>
<tr>
<td>$fst$</td>
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\[
pair\ \varphi_1\ \varphi_2
\equiv\ (\ (\lambda f. \lambda s. \lambda c. \ ((c\ f)\ s)\ \varphi_1)\ \varphi_2)
\rightarrow^*\ \lambda c. ((c\ \varphi_1)\ \varphi_2)
\]

\[
fst\ (pair\ \varphi_1\ \varphi_2)
\equiv\ (\lambda p. (p\ true))\ (pair\ \varphi_1\ \varphi_2)
\rightarrow\ (pair\ \varphi_1\ \varphi_2)\ true
\rightarrow^*\ (\lambda c. ((c\ \varphi_1)\ \varphi_2))\ true
\]
Encoding Data Structures in the \( \lambda \)-Calculus

<table>
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<tr>
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</thead>
<tbody>
<tr>
<td>( fst )</td>
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Example: Let \( \varphi_1 \) and \( \varphi_2 \) be two arbitrary expressions.

\[
\begin{align*}
\text{pair} \ \varphi_1 \ \varphi_2 & \\
\equiv & \quad (\lambda f. \lambda s. \lambda c. ((c f) s) \ \varphi_1) \ \varphi_2 \\
\rightarrow^* & \quad \lambda c. ((c \ \varphi_1) \ \varphi_2)
\end{align*}
\]

\[
\begin{align*}
\text{fst} (\text{pair} \ \varphi_1 \ \varphi_2) & \\
\equiv & \quad (\lambda p. (p \text{ true})) (\text{pair} \ \varphi_1 \ \varphi_2) \\
\rightarrow & \quad (\text{pair} \ \varphi_1 \ \varphi_2) \text{ true} \\
\rightarrow^* & \quad ((\text{true} \ \varphi_1) \ \varphi_2)
\end{align*}
\]
Encoding Data Structures in the $\lambda$-Calculus

| $\text{pair}$ | $\lambda f. \lambda s. \lambda c. ((c f) s)$ |
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\[
\text{pair} \varphi_1 \varphi_2 \\
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\rightarrow^* \lambda c. ((c \varphi_1) \varphi_2)
\]

\[
\text{fst} (\text{pair} \varphi_1 \varphi_2) \\
\equiv (\lambda p. (p \text{ true})) (\text{pair} \varphi_1 \varphi_2) \\
\rightarrow (\text{pair} \varphi_1 \varphi_2) \text{ true} \\
\rightarrow^* (\lambda c. ((c \varphi_1) \varphi_2)) \text{ true} \\
\rightarrow (\text{true} \varphi_1) \varphi_2 \\
\rightarrow \varphi_1
\]
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<th>value</th>
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Example: Let \( \varphi_1 \) and \( \varphi_2 \) be two arbitrary expressions.

\[
\text{pair} \ \varphi_1 \ \varphi_2 \\
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\[
\text{fst} \ (\text{pair} \ \varphi_1 \ \varphi_2) \\
\equiv \ (\lambda p. (p \ \text{true})) \ (\text{pair} \ \varphi_1 \ \varphi_2) \\
\to \ (\text{pair} \ \varphi_1 \ \varphi_2) \ \text{true} \\
\to^* \ (\lambda c. ((c \ \varphi_1) \ \varphi_2)) \ \text{true} \\
\to \ (\text{true} \ \varphi_1) \ \varphi_2 \\
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\text{snd} \ (\text{pair} \ \varphi_1 \ \varphi_2)
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Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

\[
pair \varphi_1 \varphi_2 \\
\equiv (\lambda f. \lambda s. \lambda c. ((c f) s) \varphi_1) \varphi_2 \\
\to^* \lambda c. ((c \varphi_1) \varphi_2)
\]

\[
fst (pair \varphi_1 \varphi_2) \\
\equiv (\lambda p. (p \ true)) (pair \varphi_1 \varphi_2) \\
\to (pair \varphi_1 \varphi_2) \ true \\
\to^* (\lambda c. ((c \varphi_1) \varphi_2)) \ true \\
\to ((true \varphi_1) \varphi_2) \\
\to \varphi_1
\]

\[
snd (pair \varphi_1 \varphi_2) \\
\equiv (\lambda p. (p \ false)) (pair \varphi_1 \varphi_2)
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Encoding Data Structures in the $\lambda$-Calculus

| $\text{pair}$ | $\lambda f. \lambda s. \lambda c. ((c \ f) \ s)$ |
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\end{align*}
\]

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\begin{align*}
\text{fst} \ (\text{pair} \ \varphi_1 \ \varphi_2) & \equiv (\lambda p. (p \ \text{true})) (\text{pair} \ \varphi_1 \ \varphi_2) \\
& \to (\text{pair} \ \varphi_1 \ \varphi_2) \ \text{true} \\
& \to^* (\lambda c. ((c \ \varphi_1) \ \varphi_2)) \ \text{true} \\
& \to ((\text{true} \ \varphi_1) \ \varphi_2) \\
& \to \ \varphi_1
\end{align*}
\]

\[
\begin{align*}
\text{snd} \ (\text{pair} \ \varphi_1 \ \varphi_2) & \equiv (\lambda p. (p \ \text{false})) (\text{pair} \ \varphi_1 \ \varphi_2) \\
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Example: Let $\varphi_1$ and $\varphi_2$ be two arbitrary expressions.

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\]

\[
\text{fst (pair } \varphi_1 \varphi_2) \\
\equiv (\lambda p. (p \text{ true})) (\text{pair } \varphi_1 \varphi_2) \\
\rightarrow (\text{pair } \varphi_1 \varphi_2) \text{ true} \\
\rightarrow^* (\lambda c. ((c \varphi_1) \varphi_2)) \text{ true} \\
\rightarrow ((\text{true } \varphi_1) \varphi_2) \\
\rightarrow \varphi_1
\]

\[
\text{snd (pair } \varphi_1 \varphi_2) \\
\equiv (\lambda p. (p \text{ false})) (\text{pair } \varphi_1 \varphi_2) \\
\rightarrow^* ((\text{false } \varphi_1) \varphi_2) \\
\rightarrow \varphi_2
\]
Basic reduction: \((\lambda x. t_1) t_2 \rightarrow [x \mapsto t_2]t_1\),
where

\([x \mapsto t_2]t_1\) be the term obtained by replacing all
“free” occurrences of \(x\) in \(t_1\) by \(t_2\).
Evaluating Lambda Expressions: An Informal Intro.

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- A sub-term of \(t\) of the form \((\lambda x \ . \ t_1) \ t_2\) is called a \textit{redex} of \(t\).
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- A sub-term of \(t\) of the form \((\lambda x. \ t_1) \ t_2\) is called a redex of \(t\).
- One step in evaluating a \(\lambda\)-term \(t\) is replacing some redex in \(t\) according to the above reduction schema.
Basic reduction: \((\lambda x \cdot t_1) \, t_2 \rightarrow [x \mapsto t_2] t_1\), where

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- A sub-term of \(t\) of the form \((\lambda x \cdot t_1) \, t_2\) is called a **redex** of \(t\).

- One step in evaluating a \(\lambda\)-term \(t\) is replacing some redex in \(t\) according to the above reduction schema.

- In general, there may be many redexes in a term.

**Example:** Let \(id = (\lambda x \cdot x)\) in term \(id \, (id \, (\lambda z \cdot id \, z))\)
## Reduction Strategies

A reduction strategy is used to **choose** a redex where the basic reduction step will be done.
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Reduction Strategies

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- **Full $\beta$-reduction:** Pick a redex non-deterministically

```
  apply
  \lambda x \ x
  \lambda x \ x
  \lambda z
    apply
    \lambda x \ x
    \lambda x \ x
  z
  x
```

Programming Languages

The Untyped Lambda Calculus
A reduction strategy is used to **choose** a redex where the basic reduction step will be done.

- **Full $\beta$-reduction**: Pick a redex *non-deterministically*
- **Normal Order**: choose the left-most, outer-most redex.
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- **Full $\beta$-reduction:** Pick a redex *non-deterministically*
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A reduction strategy is used to choose a redex where the basic reduction step will be done.

- **Full β-reduction:** Pick a redex non-deterministically
- **Normal Order:** choose the left-most, outer-most redex.
- **Call-By-Name:** like normal-order, but ignore redexes inside abstractions.
- **Call-By-Value:** choose the right-most, inner-most redex that is not inside an abstraction.
Evaluating Lambda Expressions

- The key step in evaluating an application then is: *replace every occurrence of a formal parameter with the actual argument.*

Example: \((\lambda x. (\lambda z. x z)) \ y\) \rightarrow (\lambda z. y z)\)
Evaluating Lambda Expressions

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**Example:** \(((\lambda x. (\lambda z. x z)) \ y) \rightarrow (\lambda z. y z)\)

- We can formalize the meaning of application by introducing a function, called *substitution* that maps terms to terms:

\[
(\lambda x. t_1) \ t_2 \rightarrow [x \mapsto t_2] t_1
\]
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- We can formalize the meaning of application by introducing a function, called *substitution* that maps terms to terms:

  \[(\lambda x.t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1\]

- The central problem now is how we define this substitution function.
Substitutions ($1^{st}$ attempt)

\[
\begin{align*}
[x \mapsto s]x & = s \\
[x \mapsto s]y & = y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. t) & = \lambda y. [x \mapsto s]t \\
[x \mapsto s](t_1 \ t_2) & = ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]
Substitutions (1\text{st} attempt)

\[ [x \mapsto s]x = s \]
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\[ [x \mapsto s](\lambda y. t) = \lambda y. [x \mapsto s]t \]
\[ [x \mapsto s](t_1 t_2) = ([x \mapsto s]t_1) ([x \mapsto s]t_2) \]

- Appears to be correct.

**Example:** \( [x \mapsto y](\lambda z. x z) = (\lambda z. y z) \)

**Use:** \( (\lambda x. (\lambda z. x z)) y \) \quad \rightarrow \quad (\lambda z. y z) \)
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- But is incorrect!
  **Example:** \([x \mapsto y](\lambda x. x) = (\lambda x. y)\)
  **Use:** \(((\lambda x.(\lambda x. x)) \ y) \rightarrow (\lambda x. y)\)
**Substitutions (2nd attempt)**

\[
\begin{align*}
[x \mapsto s]x & = s \\
[x \mapsto s]y & = y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. t) & = \begin{cases} 
\lambda y. t & \text{if } x = y \\
\lambda y. [x \mapsto s]t & \text{if } x \neq y 
\end{cases} \\
[x \mapsto s](t_1 t_2) & = ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]

But is still incorrect! e.g. \([x \mapsto y](\lambda y. x y) = (\lambda y. y y)\)

In the result of the above example, one \(y\) is local to the function while the other \(y\) is not local. But going by our definition, there is no way to distinguish between the two \(y\)’s!

Solution: We should get \((\lambda w. y w)\) instead (by suitably renaming “local” variables).
Substitutions (2\textsuperscript{nd} attempt)

\[
[x \mapsto s]x = s \\
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\lambda y. t & \text{if } x = y \\
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\end{cases} \\
[x \mapsto s](t_1 \ t_2) = ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\]

\[
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Bound and Free Variables: An Informal Intro.

Variable \( x \) in \( \lambda \)-expression \( \lambda x. t \) is said to be **bound**.
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- Variable $x$ in $\lambda$-expression $\lambda x. t$ is said to be **bound**.
  - Example 1: $x$ in $\lambda x. x$ is a bound variable.

- Variable $x$ in $\lambda$-expression $\lambda x. \text{(} x \ y \text{)}$, $x$ is bound but $y$ is not bound.
  - Rough meaning: parameters are local to a function definition.

A variable that is not bound is said to be **free**.

- Example 2: in $\lambda x. \text{(} x \ y \text{)}$, $y$ is free.
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Bound and Binding Occurrences

\[ \lambda x. x \]
Bound and Binding Occurrences

\[ \lambda x \cdot x \]

Binding Occurrence (declaration)
Bound and Binding Occurrences

\[ \lambda x. x \]

- Bound Occurrence (use)
- Binding Occurrence (declaration)
Bound and Binding Occurrences

\[ \lambda x. x \]

- **Bound Occurrence (use)**
- **Binding Occurrence (declaration)**
Bound and Binding Occurrences

- Bound Occurrence (use)
- Binding Occurrence (declaration)

\[(\lambda x. x)(\lambda z. (x z))\]
Bound and Binding Occurrences

\[
\lambda x. x
\]

Bound Occurrence (use)
Binding Occurrence (declaration)

\[
(\lambda x. x)(\lambda z. (x z))
\]
Bound and Binding Occurrences

\[ \lambda x. x \]

Bound Occurrence (use)

Binding Occurrence (declaration)

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Bound and Binding Occurrences

Bound Occurrence (use)

Binding Occurrence (declaration)

Free Occurrence

\[ \lambda x. x \]

\[ (\lambda x. x)(\lambda z. (x z)) \]
Bound and Binding Occurrences

- **Bound Occurrence (use)**
  - $\lambda x. x$
  - Binding Occurrence (declaration)

- **Free Occurrence**
  - $(\lambda x. x)(\lambda z. (x z))$

- $(\lambda z. (\lambda x. z (x x)) (\lambda x. z (x x)))$
Bound and Binding Occurrences

- \( \lambda x. x \)
- Bound Occurrence (use)
- Binding Occurrence (declaration)

- \((\lambda x. x)(\lambda z. (x z))\)'
- Free Occurrence

- \((\lambda z. (\lambda x. z (x x)) (\lambda x. z (x x))))\)
Bound and Binding Occurrences

- \( \lambda x. x \)  
  - Bound Occurrence (use)
  - Binding Occurrence (declaration)

- \((\lambda x. x)(\lambda z. (x z))\)  
  - Free Occurrence

- \((\lambda z. (\lambda x. z (x x))) (\lambda x. z (x x)))\)
Bound and Binding Occurrences

- \( \lambda x. x \) (Bound Occurrence (use))
- Binding Occurrence (declaration)

- \((\lambda x. x)(\lambda z. (x z))\)
  - Free Occurrence

- \((\lambda z. (\lambda x. z (x x))) (\lambda x. z (x x)))\)
Bound Variables

**Formal definition**: $bv(t)$, the set of all bound variables of $t$, is such that:

- $t$ is an abstraction of the form $\lambda x.t'$:
Bound Variables

**Formal definition:** $bv(t)$, the set of all bound variables of $t$, is such that:

- $t$ is an abstraction of the form $\lambda x.t'$:
  - $bv(t) = bv(t') \cup \{x\}$
Bound Variables

**Formal definition:** $bv(t)$, the set of all bound variables of $t$, is such that:
- $t$ is an abstraction of the form $\lambda x. t'$:
  - $bv(t) = bv(t') \cup \{x\}$
- $t$ is an application of the form $t_1 t_2$: 
Bound Variables

**Formal definition:** $bv(t)$, the set of all bound variables of $t$, is such that:

- $t$ is an abstraction of the form $\lambda x.t'$:
  - $bv(t) = bv(t') \cup \{x\}$
- $t$ is an application of the form $t_1 t_2$:
  - $bv(t) = bv(t_1) \cup bv(t_2)$
Bound Variables

**Formal definition:** $bv(t)$, the set of all bound variables of $t$, is such that:

- $t$ is an abstraction of the form $\lambda x. t'$:
  - $bv(t) = bv(t') \cup \{x\}$
- $t$ is an application of the form $t_1 \ t_2$:
  - $bv(t) = bv(t_1) \cup bv(t_2)$
- Example:
  - $bv((\lambda x. \ x) \ (\lambda z. (x \ z)))$
Bound Variables

**Formal definition:** \( bv(t) \), the set of all bound variables of \( t \), is such that:

- \( t \) is an abstraction of the form \( \lambda x. t' \):
  - \( bv(t) = bv(t') \cup \{x\} \)
- \( t \) is an application of the form \( t_1 \ t_2 \):
  - \( bv(t) = bv(t_1) \cup bv(t_2) \)
- **Example:**
  
  \[
  bv( (\lambda x. x) (\lambda z. (x \ z)) ) \\
  = bv(\lambda x. x) \cup bv(\lambda z. (x \ z))
  \]
Bound Variables

**Formal definition:** $bv(t)$, the set of all bound variables of $t$, is such that:

- $t$ is an abstraction of the form $\lambda x. t'$:
  - $bv(t) = bv(t') \cup \{x\}$

- $t$ is an application of the form $t_1 \ t_2$:
  - $bv(t) = bv(t_1) \cup bv(t_2)$

- **Example:**
  
  $bv((\lambda x. \ x) \ (\lambda z. \ (x\ z)))$
  
  $= bv(\lambda x. \ x) \cup bv(\lambda z. \ (x\ z))$
  
  $= \{x\} \cup \{z\} = \{x, z\}$
Free Variables

**Formal definition:** \( fv(t) \), the set of all free variables of \( t \), is such that:

- \( t \) is a variable of the form \( x \):
  - \( t \) is a variable of the form \( x \):
    - \( fv(λx.x) = \{x\} \)
    - \( fv(λz.(x z)) = \{x\} \)

Example:

\[
fv((λx.x)(λz.(x z))) = \left(fv(λx.x) \cup fv(\lambda z.(x z))\right) = \{x\}
\]
Free Variables

**Formal definition:** $fv(t)$, the set of all free variables of $t$, is such that:

- $t$ is a variable of the form $x$:
  - $fv(t) = \{x\}$
Free Variables

**Formal definition:** $fv(t)$, the set of all free variables of $t$, is such that:

- $t$ is a variable of the form $x$:
  - $fv(t) = \{x\}$
- $t$ is an abstraction of the form $\lambda x.t'$:
  - $fv(t) = fv(t') - \{x\}$

Example:

$$fv( (\lambda x.x) (\lambda z.(x z)) ) = fv(\lambda x.x) \cup fv(\lambda z.(x z)) = \{x\}$$
Free Variables

**Formal definition:** $fv(t)$, the set of all free variables of $t$, is such that:

- $t$ is a variable of the form $x$:
  - $fv(t) = \{x\}$
- $t$ is an abstraction of the form $\lambda x.t'$:
  - $fv(t) = fv(t') \setminus \{x\}$

Example:

$fv((\lambda x.x)(\lambda z.x z)) = fv(\lambda x.x) \cup fv(\lambda z.x z) = \emptyset \cup \{x\} = \{x\}$
Free Variables

**Formal definition:** $fv(t)$, the set of all free variables of $t$, is such that:

- $t$ is a variable of the form $x$:
  - $fv(t) = \{x\}$
- $t$ is an abstraction of the form $\lambda x.t'$:
  - $fv(t) = fv(t') - \{x\}$
- $t$ is an application of the form $t_1 t_2$: 
  - $fv(t) = fv(t_1) \cup fv(t_2)$
Free Variables

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Free Variables

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- $t$ is an application of the form $t_1 t_2$:
  - $fv(t) = fv(t_1) \cup fv(t_2)$

**Example:**

$$fv\left( (\lambda x. x) (\lambda z. (x z)) \right)$$
Free Variables

Formal definition: $fv(t)$, the set of all free variables of $t$, is such that:

- $t$ is a variable of the form $x$:
  - $fv(t) = \{x\}$
- $t$ is an abstraction of the form $\lambda x.t'$:
  - $fv(t) = fv(t') - \{x\}$
- $t$ is an application of the form $t_1 t_2$:
  - $fv(t) = fv(t_1) \cup fv(t_2)$

Example:

$fv( (\lambda x. x) (\lambda z. (x z)) )$

$= fv(\lambda x. x) \cup fv(\lambda z. (x z))$
**Free Variables**

**Formal definition:** \( \text{fv}(t) \), the set of all free variables of \( t \), is such that:

- \( t \) is a variable of the form \( x \):
  - \( \text{fv}(t) = \{x\} \)
- \( t \) is an abstraction of the form \( \lambda x.t' \):
  - \( \text{fv}(t) = \text{fv}(t') \setminus \{x\} \)
- \( t \) is an application of the form \( t_1 t_2 \):
  - \( \text{fv}(t) = \text{fv}(t_1) \cup \text{fv}(t_2) \)

**Example:**

\[
\text{fv}( (\lambda x. x) (\lambda z. (x z)) ) \\
= \text{fv}(\lambda x. x) \cup \text{fv}(\lambda z. (x z)) \\
= \emptyset \cup \{x\} = \{x\}
\]
**α-Conversion (Renaming)**

- *Intuition*: We can rename a bound variable as long as
\( \alpha \)-Conversion (Renaming)

- **Intuition:** We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
**α-Conversion (Renaming)**

- **Intuition:** We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
  - we replace every occurrence of the bound variable

Two terms $t$ and $t'$ are said to be "α-equivalent" (denoted by $t \equiv_{\alpha} t'$) if they are identical modulo the names of bound variables.
\( \alpha \)-Conversion (Renaming)

- **Intuition:** We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
  - we replace every occurrence of the bound variable.
- **Example 1:** \((\lambda y. x y)\) is equivalent to \((\lambda z. x z)\)
\(\alpha\)-Conversion (Renaming)

- **Intuition**: We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
  - we replace every occurrence of the bound variable

- Example 1: \((\lambda y. x \ y)\) is equivalent to \((\lambda z. x \ z)\)

- Example 2: \((\lambda y. x \ y)\) is *not* equivalent to \((\lambda x. x \ x)\) (the name of new variable is same as that of a free variable)
\( \alpha \)-Conversion (Renaming)

- **Intuition:** We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
  - we replace every occurrence of the bound variable

- Example 1: \((\lambda y. x y)\) is equivalent to \((\lambda z. x z)\)

- Example 2: \((\lambda y. x y)\) is *not* equivalent to \((\lambda x. x x)\) (the name of new variable is same as that of a free variable)

- Example 3: \((\lambda y. x y)\) is *not* equivalent to \((\lambda y. x z)\) (not every occurrence of \(y\) has been replaced).
**$\alpha$-Conversion (Renaming)**

- **Intuition:** We can rename a bound variable as long as
  - the new name is not also the name of a free variable, and
  - we replace every occurrence of the bound variable

- Example 1: $(\lambda y. x \, y)$ is equivalent to $(\lambda z. x \, z)$

- Example 2: $(\lambda y. x \, y)$ is *not* equivalent to $(\lambda x. x \, x)$ (the name of new variable is same as that of a free variable)

- Example 3: $(\lambda y. x \, y)$ is *not* equivalent to $(\lambda y. x \, z)$ (not every occurrence of $y$ has been replaced).

- Two terms $t$ and $t'$ are said to be “$\alpha$-equivalent” (denoted by $t \equiv_\alpha t'$) if they are identical modulo the names of bound variables.
Substitutions (3\textsuperscript{rd} attempt)

\[
\begin{align*}
[x \mapsto s]x &= s \\
[x \mapsto s]y &= y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. t) &= \lambda y. [x \mapsto s]t & \text{if } x \neq y \text{ and } y \not\in \text{fv}(s) \\
[x \mapsto s](t_1 \ t_2) &= ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]
Substitutions (3rd attempt)

\[
\begin{align*}
[x \mapsto s]x &= s \\
[x \mapsto s]y &= y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. \ t) &= \lambda y. [x \mapsto s]t & \text{if } x \neq y \text{ and } y \not\in \text{fv}(s) \\
[x \mapsto s](t_1 \ t_2) &= ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]

- The definition is now incomplete! e.g. \([x \mapsto y](\lambda y. x \ y) = ?\)
Substitutions (3\textsuperscript{rd} attempt)

\[ [x \mapsto s]x = s \]
\[ [x \mapsto s]y = y \quad \text{if } y \neq x \]
\[ [x \mapsto s](\lambda y. t) = \lambda y. [x \mapsto s]t \quad \text{if } x \neq y \text{ and } y \not\in \text{fv}(s) \]
\[ [x \mapsto s](t_1 t_2) = ([x \mapsto s]t_1)([x \mapsto s]t_2) \]

- The definition is now incomplete! e.g. \([x \mapsto y](\lambda y. x \ y) = ??\)
- This drawback is not serious:
Substitutions (3rd attempt)

\[
\begin{align*}
[x \mapsto s]x & = s \\
[x \mapsto s]y & = y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. t) & = \lambda y. [x \mapsto s]t & \text{if } x \neq y \text{ and } y \not\in \text{fv}(s) \\
[x \mapsto s](t_1 t_2) & = ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]

- The definition is now incomplete! e.g. \([x \mapsto y](\lambda y. x y) = ??\)
- This drawback is not serious:
- We can apply a substitution on an \(\alpha\)-equivalent term instead.
Substitutions (3rd attempt)

\[
\begin{align*}
[x \mapsto s]x &= s \\
[x \mapsto s]y &= y & \text{if } y \neq x \\
[x \mapsto s](\lambda y. \ t) &= \lambda y. [x \mapsto s]t & \text{if } x \neq y \text{ and } y \notin \text{fv}(s) \\
[x \mapsto s](t_1 \ t_2) &= ([x \mapsto s]t_1) ([x \mapsto s]t_2)
\end{align*}
\]

- The definition is now incomplete! e.g. \( [x \mapsto y](\lambda y. \ x \ y) = ?? \)
- This drawback is not serious:
- We can apply a substitution on an \( \alpha \)-equivalent term instead.
- E.g. \( [x \mapsto y](\lambda z. \ x \ z) = (\lambda z. \ y \ z) \)
Operational Semantics: Full $\beta$-Reduction

\[
\begin{align*}
& t_1 \rightarrow t'_1 \\
\hline
& t_1 \ t_2 \rightarrow t'_1 \ t_2 & \text{E-APP1} \\
\end{align*}
\]

\[
\begin{align*}
& t_2 \rightarrow t'_2 \\
\hline
& t_1 \ t_2 \rightarrow t_1 \ t'_2 & \text{E-APP2} \\
\end{align*}
\]

\[
\begin{align*}
& t \rightarrow t' \\
\hline
& \lambda x. \ t \rightarrow \lambda x. \ t' & \text{E-ABS} \\
\end{align*}
\]

\[
(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1 & \text{E-APPABS}
\]
Operational Semantics: Call-By-Value

\[ t ::= \ldots \quad \text{Terms (all } \lambda\text{-terms)} \]
\[ v ::= \lambda x. t \quad \text{Values} \]

Evaluation:

\[
\frac{t_1 \rightarrow t'_1}{t_1 t_2 \rightarrow t'_1 t_2} \quad \text{E-APP1}
\]
\[
\frac{t_2 \rightarrow t'_2}{v_1 t_2 \rightarrow v'_1 t'_2} \quad \text{E-APP2}
\]
\[
(\lambda x. t_1) v_2 \rightarrow [x \mapsto v_2] t_1 \quad \text{E-APPABS}
\]
Operational Semantics: Call-By-Value

\[
\begin{align*}
  t & ::= \ldots & \text{Terms (all } \lambda\text{-terms)} \\
  v & ::= \lambda x. t & \text{Values}
\end{align*}
\]

Evaluation:

\[
\frac{t_1 \rightarrow t_1'}{t_1 \, t_2 \rightarrow t_1' \, t_2} \quad \text{E-App1}
\]

\[
\frac{t_2 \rightarrow t_2'}{v_1 \, t_2 \rightarrow v_1 \, t_2'} \quad \text{E-App2}
\]

\[
(\lambda x. \, t_1) \, v_2 \rightarrow [x \mapsto v_2] t_1 \quad \text{E-AppAbs}
\]

- In an application of the form \((t_1 \, t_2)\), if \(t_1\) is a \(\lambda\)-abstraction, then \(t_2\) has to be reduced to a value before the application is done.
Operational Semantics: Call-By-Value

\[
\begin{align*}
  t & ::= \ldots & \text{Terms (all } \lambda\text{-terms)} \\
  v & ::= \lambda x. t & \text{Values}
\end{align*}
\]

Evaluation:

\[
\begin{align*}
  \frac{t_1 \to t'_1}{t_1 \ t_2 \to t'_1 \ t_2} & \quad \text{E-APP1} \\
  \frac{t_2 \to t'_2}{v_1 \ t_2 \to v_1 \ t'_2} & \quad \text{E-APP2} \\
  (\lambda x. \ t_1) \ v_2 \to [x \mapsto v_2]t_1 & \quad \text{E-APPABS}
\end{align*}
\]

- In an application of the form \((t_1 \ t_2)\), if \(t_1\) is a \(\lambda\)-abstraction, then \(t_2\) has to be reduced to a value before the application is done.
- This corresponds to Call-By-Value parameter passing: evaluate the actual arguments first before passing them as parameters to a called function.
Operational Semantics: Call-By-Name

\[
\begin{align*}
  t & ::= \ldots \quad \text{Terms (all } \lambda\text{-terms)} \\
  v & ::= \lambda x. t \quad \text{Values}
\end{align*}
\]

Evaluation:

\[
\frac{t_1 \rightarrow t'_1}{t_1 \ t_2 \rightarrow t'_1 \ t_2} \quad \text{E-APP}
\]

\[
(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2] t_1 \quad \text{E-APPABS}
\]
Operational Semantics: Call-By-Name

\[ t ::= \ldots \quad \text{Terms (all } \lambda\text{-terms)} \]
\[ v ::= \lambda x. t \quad \text{Values} \]

Evaluation:

\[
\frac{t_1 \rightarrow t_1'}{t_1 \; t_2 \rightarrow t_1' \; t_2} \quad \text{E-APP}
\]

\[
(\lambda x. \; t_1) \; t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-APPABS}
\]

- In an application of the form \((t_1 \; t_2)\), if \(t_1\) is a \(\lambda\)-abstraction, then \(t_1\) has to be reduced to a value before the application is done.
Operational Semantics: Call-By-Name

\[ t ::= \ldots \quad \text{(all \( \lambda \)-terms)} \]

\[ v ::= \lambda x.\ t \quad \text{Values} \]

Evaluation:

\[
\frac{t_1 \rightarrow t'_1}{t_1 \ t_2 \rightarrow t'_1 \ t_2} \quad \text{E-App} \\
(\lambda x.\ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}
\]

- In an application of the form \((t_1 \ t_2)\), if \(t_1\) is a \(\lambda\)-abstraction, then \(t_1\) has to be reduced to a value before the application is done.
- In terms of familiar languages, the actual arguments are passed \textit{unevaluated} to the called function. They will be evaluated in the called function if needed.
Consider variables in a $\lambda$-term as named “holes” to be filled in.

Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

Examples:
Nameless Representation of Terms

- Consider variables in a λ-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the λ that binds them.

Examples:

- $\lambda x. \ x$ can be written as

\[
\begin{array}{c}
\lambda x \\
| \uparrow \\
x
\end{array}
\]
Nameless Representation of Terms

- Consider variables in a $\lambda$-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

Examples:

- $\lambda x. x$ can be written as

  \[
  \begin{array}{c|c|c|c|c}
  & \lambda & \lambda & 0 \\
  \hline
  | & \uparrow & \uparrow & \uparrow \\
  x & \lambda & \lambda & \lambda \\
  \end{array}
  \]
Nameless Representation of Terms

- Consider variables in a \( \lambda \)-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the \( \lambda \) that binds them.

**Examples:**

\[
\begin{array}{c}
\lambda x \\
\lambda \\
x \\
0
\end{array}
\]

- \( \lambda x. \ x \) can be written as \( \lambda. \ 0 \)
Nameless Representation of Terms

- Consider variables in a λ-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the λ that binds them.

**Examples:**

\[ \lambda x \]
\[ \lambda y \]
\[ x \]

- \( \lambda x. x \) can be written as \( \lambda. 0 \)
- \( \lambda x. \lambda y. x \) can be written as
Nameless Representation of Terms

- Consider variables in a $\lambda$-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

**Examples:**

\[
\begin{align*}
\lambda x & \quad \lambda \\
\lambda y & \quad \lambda \\
\lambda x & \quad 1
\end{align*}
\]

- $\lambda x. \ x$ can be written as $\lambda. \ 0$
- $\lambda x. \ \lambda y. \ x$ can be written as
Nameless Representation of Terms

- Consider variables in a $\lambda$-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

Examples:

\[
\begin{align*}
\lambda x & \quad \lambda \\
\lambda y & \quad \lambda \\
\lambda x & \quad \lambda 0 \\
\lambda x \; \lambda y \; x & \quad \lambda \lambda 1
\end{align*}
\]
Consider variables in a $\lambda$-term as named “holes” to be filled in.

Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

### Examples:

- $\lambda x. x$ can be written as $\lambda. 0$
- $\lambda x. \lambda y. x$ can be written as $\lambda. \lambda. 1$
- $\lambda x. \lambda y. x (y x)$ can be written as
Nameless Representation of Terms

- Consider variables in a $\lambda$-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

Examples:

- $\lambda x. x$ can be written as $\lambda. 0$
- $\lambda x. \lambda y. x$ can be written as $\lambda. \lambda. 1$
- $\lambda x. \lambda y. x (y x)$ can be written as $\lambda. \lambda. 0 1$
Nameless Representation of Terms

- Consider variables in a $\lambda$-term as named “holes” to be filled in.
- Instead of using symbolic names for variables, one can name the holes w.r.t. the $\lambda$ that binds them.

Examples:

- $\lambda x \cdot x$ can be written as $\lambda. \ 0$
- $\lambda x. \lambda y. \ x$ can be written as $\lambda. \lambda. \ 1$
- $\lambda x. \lambda y. \ x \ (y \ x)$ can be written as $\lambda. \lambda. \ 1 \ (0 \ 1)$
$n$-Terms

De Bruijn terms are defined by a family of sets (each set being a set of terms) $\{T_0, T_1, \ldots\}$ such that $T_n$ represents $\lambda$-terms with $n$ or fewer free variables.

Formally, $T$ is the smallest family of sets $\{T_0, T_1, \ldots\}$ such that:

- $k \in T_n$ whenever $0 \leq k < n$ 
- If $t_1 \in T_n$ then $\lambda t_1 \in T_{n-1}$
- If $t_1, t_2 \in T_n$ then $(t_1 t_2) \in T_n$

$\alpha$-equivalent closed $\lambda$-terms will have the same de Bruijn representation.
Naming Context

- When a $\lambda$-term has free variables, we need information on their relative positions.
- E.g. given $\{v \mapsto 2, w \mapsto 1, x \mapsto 0\}$:
When a $\lambda$-term has free variables, we need information on their relative positions.

E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):
Naming Context

- When a \( \lambda \)-term has free variables, we need information on their relative positions.
- E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):

\[
\lambda v \quad \lambda w \quad \lambda x
\]

\( v \) (\( w \) \( x \)) can be written as

\[
\lambda y \quad \lambda c \quad \lambda.
\]
Naming Context

- When a λ-term has free variables, we need information on their relative positions.
- E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):

\[
\lambda v \lambda w w \lambda x v (w x)
\]

- \( v (w x) \) can be written as \( 2 (1 0) \)
Naming Context

- When a λ-term has free variables, we need information on their relative positions.
- E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):

\[
\begin{align*}
\lambda v \\
\downarrow \\
\lambda w \\
\downarrow \\
\lambda x \\
\downarrow \\
\lambda y \\
\downarrow \\
apply \\
\downarrow \\
w \\
\downarrow \\
y
\end{align*}
\]

\( v (w x) \) can be written as 2 (1 0)

\( \lambda y. w y \) can be written as
When a \( \lambda \)-term has free variables, we need information on their relative positions.

E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):

\[
\begin{array}{c}
\lambda v \\
\downarrow \\
\lambda w \\
\downarrow \\
\lambda x \\
\downarrow \\
\lambda y \\
\downarrow \\
apply \\
\downarrow \\
w \quad y
\end{array}
\]

- \( v \ (w \ x) \) can be written as \( 2 \ (1 \ 0) \)
- \( \lambda y. \ w \ y \) can be written as \( \lambda. \ 2 \ 0 \)
Naming Context

- When a λ-term has free variables, we need information on their relative positions.
- E.g. given \[ \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} : \]

\[
\begin{align*}
\lambda v \\
\lambda w \\
\lambda x \\
\lambda y \\
\lambda c \\
\lambda
\end{align*}
\]

- \( v \ (w \ x) \) can be written as \( 2 \ (1 \ 0) \)
- \( \lambda y. \ w \ y \) can be written as \( \lambda. \ 2 \ 0 \)
- \( \lambda y. \lambda c. \ v \) can be written as
Naming Context

- When a \( \lambda \)-term has free variables, we need information on their relative positions.
- E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):
  
  \[
  \begin{aligned}
  &\lambda v \\
  &\lambda w \\
  &\lambda x \\
  &\lambda y \\
  &\lambda c \\
  &v
  \end{aligned}
  \]
  
  - \( v \ (w \ x) \) can be written as 2 (1 0)
  - \( \lambda y. \ w \ y \) can be written as \( \lambda. \ 2 \ 0 \)
  - \( \lambda y. \lambda c. \ v \) can be written as \( \lambda. \ \lambda. \ 4 \)
Naming Context

- When a \( \lambda \)-term has free variables, we need information on their relative positions.
- E.g. given \( \{ v \mapsto 2, w \mapsto 1, x \mapsto 0 \} \):

\[
\begin{array}{c}
\lambda v \\
\lambda w \\
\lambda x \\
\lambda y \\
\lambda c \\
v
\end{array}
\]

- \( v \ (w \ x) \) can be written as \( 2 \ (1 \ 0) \)
- \( \lambda y. \ w \ y \) can be written as \( \lambda. \ 2 \ 0 \)
- \( \lambda y.\lambda c. \ v \) can be written as \( \lambda. \ \lambda. \ 4 \)

- Naming contexts are often written as a sequence, where \( x_n, x_{n-1}, \ldots, x_1, x_0 \), represents a context where each \( x_i \) has de Bruijn index \( i \).
Substitution

- Term \( (\lambda y. \lambda z. (x \ y) \ (w \ z)) \) under naming context \( v, w, x \) has the following de Bruijn representation:

\[
\lambda. \lambda. (2 \ 1) (3 \ 0)
\]
Substitution

- Term \((\lambda y. \lambda z. (x \ y) \ (w \ z))\) under naming context \(v, w, x\) has the following de Bruijn representation:

\[ \lambda. \lambda. (2 \ 1) \ (3 \ 0) \]

- Term \((v \ w)\) under naming context \(v, w, x\) has the following de Bruijn representation:

\[ (2 \ 1) \]
Substitution

- Term \((\lambda y. \lambda z. (x \ y) (w \ z))\) under naming context \(v, w, x\) has the following de Bruijn representation:
  \[ \lambda. \lambda. (2 \ 1) (3 \ 0) \]

- Term \((v \ w)\) under naming context \(v, w, x\) has the following de Bruijn representation:
  \[ (2 \ 1) \]

- Substitution \([x \mapsto (v \ w)](\lambda y. \lambda z. (x \ y) (w \ z))\) will yield the term
  \[ \lambda y. \lambda z. ((v \ w) \ y) (w \ z) \]
Substitution

- Term \((\lambda y. \lambda z. (x \ y) (w \ z))\) under naming context \(v, w, x\) has the following de Bruijn representation:
  \[
  \lambda. \lambda. (2 \ 1) (3 \ 0)
  \]

- Term \((v \ w)\) under naming context \(v, w, x\) has the following de Bruijn representation:
  \[
  (2 \ 1)
  \]

- Substitution \([x \mapsto (v \ w)](\lambda y. \lambda z. (x \ y) (w \ z))\) will yield the term
  \[
  \lambda y. \lambda z. ((v \ w) \ y) (w \ z)
  \]

- Assuming the naming context is \(v, w, x\), the above term has the following de Bruijn representation:
  \[
  (\lambda. \lambda. ((4 \ 3) \ 1) (3 \ 0))
  \]
Substitution

- Term \((\lambda y. \lambda z. (x \, y) \, (w \, z))\) under naming context \(v, w, x\) has the following de Bruijn representation:

\[
\lambda. \lambda. (2 \, 1) \, (3 \, 0)
\]

- Term \((v \, w)\) under naming context \(v, w, x\) has the following de Bruijn representation:

\[
(2 \, 1)
\]

- Substitution \([x \mapsto (v \, w)](\lambda y. \lambda z. (x \, y) \, (w \, z))\) will yield the term

\[
\lambda y. \lambda z. ((v \, w) \, y) \, (w \, z)
\]

- Assuming the naming context is \(v, w, x\), the above term has the following de Bruijn representation: \((\lambda. \lambda. ((4 \, 3) \, 1) \, (3 \, 0))\)

- Hence, when carrying out substitution, we need to renumber the indices of free variables in the replacement term, and retain the indices of bound variables. This will be done using the shifting operation, defined next.
Shifting

For substitution, we need to

- renumber the indices of free variables (say, by $d$), and
- retain the indices of bound variables (say, those numbered below $c$).

This is done using the \textit{shifting} operation, defined as follows:

\[
\uparrow_c^d (k) = \begin{cases} 
  k & \text{if } k < c \\
  k + d & \text{if } k \geq c
\end{cases}
\]

\[
\uparrow_c^d (\lambda. t_1) = \lambda. \uparrow_{c+1}^d (t_1)
\]

\[
\uparrow_c^d (t_1 \ t_2) = (\uparrow_c^d t_1 \ \uparrow_c^d t_2)
\]

\[
\uparrow^d (t) = \uparrow_0^d (t)
\]
Shifting

For substitution, we need to

- renumber the indices of free variables (say, by $d$), and
- retain the indices of bound variables (say, those numbered below $c$).

This is done using the *shifting* operation, defined as follows:

$$
\uparrow^d_c (k) = \begin{cases} 
  k & \text{if } k < c \\
  k + d & \text{if } k \geq c 
\end{cases}
$$

$$
\uparrow^d_c (\lambda. t_1) = \lambda. \uparrow^d_{c+1} (t_1)
$$

$$
\uparrow^d_c (t_1 t_2) = (\uparrow^d_c t_1 \uparrow^d_c t_2)
$$

$$
\uparrow^d (t) = \uparrow^d_0 (t)
$$

Examples

- $\uparrow^2 (\lambda. \lambda. 1 (0 2)) =$
For substitution, we need to

- renumber the indices of free variables (say, by $d$), and
- retain the indices of bound variables (say, those numbered below $c$).

This is done using the *shifting* operation, defined as follows:

\[
\uparrow^d_c (k) = \begin{cases} 
  k & \text{if } k < c \\
  k + d & \text{if } k \geq c 
\end{cases}
\]

\[
\uparrow^d_c (\lambda. t_1) = \lambda. \uparrow^d_{c+1} (t_1)
\]

\[
\uparrow^d_c (t_1 t_2) = (\uparrow^d_c t_1 \uparrow^d_c t_2)
\]

\[
\uparrow^d (t) = \uparrow^d_0 (t)
\]

**Examples**

- $\uparrow^2 (\lambda. \lambda. 1 (0 2)) = \lambda. \lambda. 1 (0 4)$
Shifting

For substitution, we need to
- renumber the indices of free variables (say, by $d$), and
- retain the indices of bound variables (say, those numbered below $c$).

This is done using the *shifting* operation, defined as follows:

$$\uparrow^d_c (k) = \begin{cases} 
    k & \text{if } k < c \\
    k + d & \text{if } k \geq c
\end{cases}$$

$\uparrow^d_c (\lambda. t_1) = \lambda. \uparrow^{d+1}_{c+1} (t_1)$

$\uparrow^d_c (t_1 \ t_2) = (\uparrow^d_c t_1 \ \uparrow^d_c t_2)$

$\uparrow^d (t) = \uparrow^0 (t)$

Examples
- $\uparrow^2 (\lambda. \lambda. 1 (0 \ 2)) = \lambda. \lambda. 1 (0 \ 4)$
- $\uparrow^2 (\lambda. 0 \ 1 (\lambda. 0 \ 1 \ 2)) =$
Shifting

For substitution, we need to

- renumber the indices of free variables (say, by $d$), and
- retain the indices of bound variables (say, those numbered below $c$).

This is done using the *shifting* operation, defined as follows:

$$
\uparrow^d_c (k) = \begin{cases} 
  k & \text{if } k < c \\
  k + d & \text{if } k \geq c 
\end{cases}
$$

$$
\uparrow^d_c (\lambda. t_1) = \lambda. \uparrow^{d+1}_{c+1} (t_1)
$$

$$
\uparrow^d_c (t_1 t_2) = (\uparrow^d_c t_1 \uparrow^d_c t_2)
$$

$$
\uparrow^d (t) = \uparrow^d_0 (t)
$$

Examples

- $\uparrow^2 (\lambda. \lambda. 1 (0 2)) = \lambda. \lambda. 1 (0 4)$
- $\uparrow^2 (\lambda. 0 1 (\lambda. 0 1 2)) = \lambda. 0 3 (\lambda. 0 1 4)$
Substitution using Shifting

\[
[j \mapsto s]k = \begin{cases} 
  s & \text{if } k = j \\
  k & \text{otherwise}
\end{cases}
\]

\[
[j \mapsto s](\lambda. t_1) = \lambda. [j+1 \mapsto s^1 (s)]t_1
\]

\[
[j \mapsto s](t_1 t_2) = ([j \mapsto s]t_1 [j \mapsto s]t_2)
\]

Examples:

\[ [0 \mapsto 1](0 (\lambda. \lambda. 2)) = \]
Substitution using Shifting

\[ [j \mapsto s]k = \begin{cases} s & \text{if } k = j \\ k & \text{otherwise} \end{cases} \]

\[ [j \mapsto s](\lambda. t_1) = \lambda. [j + 1 \mapsto \uparrow^1 (s)]t_1 \]

\[ [j \mapsto s](t_1 t_2) = ([j \mapsto s]t_1 [j \mapsto s]t_2) \]

Examples:

- \[ [0 \mapsto 1](0 (\lambda. \lambda. 2)) = 1 (\lambda. \lambda. 3) \]
Substitution using Shifting

\[ [j \mapsto s]k = \begin{cases} 
  s & \text{if } k = j \\
  k & \text{otherwise}
\end{cases} \]

\[ [j \mapsto s](\lambda. \ t_1) = \lambda. \ [j+1 \mapsto 1] (s) \ t_1 \]

\[ [j \mapsto s](t_1 \ t_2) = ([j \mapsto s] \ t_1 \ [j \mapsto s] \ t_2) \]

Examples:

- \[ [0 \mapsto 1](0 (\lambda. \ \lambda. \ 2)) = 1 (\lambda. \ \lambda. \ 3) \]
- \[ [0 \mapsto (1 (\lambda. \ 2))](0 (\lambda. \ 1)) = \]
Substitution using Shifting

\[
[j \mapsto s]k = \begin{cases} 
  s & \text{if } k = j \\
  k & \text{otherwise}
\end{cases}
\]

\[
[j \mapsto s](\lambda. t_1) = \lambda. [j + 1 \mapsto \uparrow^1 (s)]t_1
\]

\[
[j \mapsto s](t_1 t_2) = ([j \mapsto s]t_1 [j \mapsto s]t_2)
\]

Examples:

- \([0 \mapsto 1](0 (\lambda. \lambda. 2)) = 1 (\lambda. \lambda. 3)\)
- \([0 \mapsto (1 (\lambda. 2))](0 (\lambda. 1)) = (1 (\lambda. 2)) (\lambda(2 (\lambda. 3)))\)
Substitution using Shifting

\[
[j \mapsto s]k = \begin{cases} 
  s & \text{if } k = j \\
  k & \text{otherwise}
\end{cases}
\]

\[
[j \mapsto s](\lambda. t_1) = \lambda. [j + 1 \mapsto s^1 (s)]t_1
\]

\[
[j \mapsto s](t_1 t_2) = ([j \mapsto s]t_1 [j \mapsto s]t_2)
\]

Examples:

- \([0 \mapsto 1](0 (\lambda. \lambda. 2)) = 1 (\lambda. \lambda. 3)\)
- \([0 \mapsto (1 (\lambda. 2))](0 (\lambda. 1)) = (1 (\lambda. 2)) (\lambda(2 (\lambda. 3)))\)
- \([0 \mapsto 1](\lambda. (0 2)) = \lambda. (0 2)\)
In the calculus with symbolic term representation:

\[(\lambda x.\ t_1)\ t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}\]
### Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1\]  \[\text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \ t_1) \ t_2 \rightarrow \uparrow^{-1} ([0 \mapsto \uparrow^1 (t_2)]t_1)\]  \[\text{E-AppAbs}\]
Evaluation

In the calculus with symbolic term representation:

$$(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}$$

In the calculus with de Bruijn representation:

$$(\lambda. \ t_1) \ t_2 \rightarrow (0 \mapsto 1 (t_2)]t_1) \quad \text{E-AppAbs}$$

- The outer $\lambda$ is removed after application, so the indices have to shift down by 1.
Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1\] \hspace{1cm} E-AppAbs

In the calculus with de Bruijn representation:

\[(\lambda. \ t_1) \ t_2 \rightarrow^{\uparrow -1} ([0 \mapsto^{\uparrow 1} (t_2)]t_1)\] \hspace{1cm} E-AppAbs

- The outer \(\lambda\) is removed after application, so the indices have to shift down by 1.
- Indices in argument \((t_2)\) should not be changed in the end, so we shifting them up by 1 first.
Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \ t_1) \ t_2 \rightarrow \uparrow^{-1} ([0 \mapsto \uparrow^1 (t_2)]t_1) \quad \text{E-AppAbs}\]

- The outer \( \lambda \) is removed after application, so the indices have to shift down by 1.
- Indices in argument \( t_2 \) should not be changed in the end, so we shifting them up by 1 first.
  - Consider \((\lambda x. \ w \times v) (\lambda y. \ (w \ y))\), whose de Bruijn representation is \((\lambda. \ 1 \ 0 \ 2) (\lambda. \ 1 \ 0)\) (assuming naming context \( v, w \)).
Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \ t_1) \ t_2 \rightarrow^{↑-1} ([0 \mapsto^{↑1} (t_2)]t_1) \quad \text{E-AppAbs}\]

- The outer \(\lambda\) is removed after application, so the indices have to shift \textit{down} by 1.
- Indices in argument \((t_2)\) should \textit{not} be changed in the end, so we shifting them \textit{up} by 1 first.
  - Consider \((\lambda x. \ w \ x \ v) \ (\lambda y. \ (w \ y))\), whose de Bruijn representation is \((\lambda. \ 1 \ 0 \ 2) \ (\lambda. \ 1 \ 0)\) \textit{(assuming naming context} \(v, w)\).
  - The result of the application is \(w \ (\lambda y. \ w \ y) \ v\).
Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \ t_1) \ t_2 \rightarrow [x \mapsto t_2] t_1\quad \text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \ t_1) \ t_2 \rightarrow \uparrow^{-1} ([0 \mapsto \uparrow^1 (t_2)] t_1)\quad \text{E-AppAbs}\]

- The outer \(\lambda\) is removed after application, so the indices have to shift down by 1.
- Indices in argument \((t_2)\) should not be changed in the end, so we shifting them up by 1 first.
  - Consider \((\lambda x. \ w \ x \ v) \ (\lambda y. \ (w \ y))\), whose de Bruijn representation is \((\lambda. \ 1 \ 0 \ 2) \ (\lambda. \ 1 \ 0)\) (assuming naming context \(v, w\)).
  - The result of the application is \(w \ (\lambda y. \ w \ y) \ v\).
  - \(\uparrow^1 (\lambda. \ 1 \ 0) = \lambda. \ 2 \ 0\)
Evaluation

In the calculus with symbolic term representation:

\[(\lambda x. \, t_1) \, t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \, t_1) \, t_2 \rightarrow ↑^{-1} ([0 \mapsto ↑^1 (t_2)]t_1) \quad \text{E-AppAbs}\]

- The outer \(\lambda\) is removed after application, so the indices have to shift down by 1.
- Indices in argument \(t_2\) should \textit{not} be changed in the end, so we shifting them up by 1 first.
  - Consider \((\lambda x. \, w \, x \, v) \, (\lambda y. \, (w \, y))\), whose de Bruijn representation is \((\lambda. \, 1 \, 0 \, 2) \, (\lambda. \, 1 \, 0)\) (assuming naming context \(v, w\)).
  - The result of the application is \(w \, (\lambda y. \, w \, y) \, v\).
  - \(↑^1 (\lambda. \, 1 \, 0) = \lambda. \, 2 \, 0\)
  - \([0 \mapsto (\lambda. \, 2 \, 0)](1 \, 0 \, 2) = 1 \, (\lambda. \, 2 \, 0) \, 2\)
**Evaluation**

In the calculus with symbolic term representation:

\[(\lambda x. \; t_1) \; t_2 \rightarrow [x \mapsto t_2]t_1 \quad \text{E-AppAbs}\]

In the calculus with de Bruijn representation:

\[(\lambda. \; t_1) \; t_2 \rightarrow \uparrow^{-1} ([0 \mapsto \uparrow^1 (t_2)]t_1) \quad \text{E-AppAbs}\]

- The outer \(\lambda\) is removed after application, so the indices have to shift *down* by 1.
- Indices in argument \((t_2)\) should *not* be changed in the end, so we shifting them *up* by 1 first.
  - Consider \((\lambda x. \; w \times v) \; (\lambda y. \; (w \; y))\), whose de Bruijn representation is \((\lambda. \; 1 \; 0 \; 2) \; (\lambda. \; 1 \; 0)\) (assuming naming context \(v, w\)).
  - The result of the application is \(w \; (\lambda y. \; w \; y) \; v\).
  - \(\uparrow^1 (\lambda. \; 1 \; 0) = \lambda. \; 2 \; 0\)
  - \([0 \mapsto (\lambda. \; 2 \; 0)](1 \; 0 \; 2) = 1 \; (\lambda. \; 2 \; 0) \; 2\)
  - \(\uparrow^{-1} (1 \; (\lambda. \; 2 \; 0) \; 2) = 0 \; (\lambda. \; 1 \; 0) \; 1\)